October 18, 2015

## Literature Search And Expected Results-Jipeng

**Theorem 1.** (Size of Largest Components In A Random Graph) *Consider a random graph*  $G \in G_{n,p}$  *where* p = c/n *for some constant c. Then:* 

- 1. If c < 1, then a.a.s., the largest component of G has size  $O(\log(n))$ , or more exactly, at least  $\frac{3\ln(n)}{(1-c)^2}$ .
- 2. If c > 1, then a.a.s., G has a largest component of size O(n). It is a unique "giant" component of size  $(1 + o(1))\beta n$ , where  $\beta$  is the unique solution in [0,1] to the equation  $\beta + e^{-\beta c} = 1$ . All other components have size  $O(\log(n))$ .
- 3. If c = 1, then a.a.s., G has a component of size  $O(n^{\frac{2}{3}})$ .

The phrase "a.a.s" (asymptotically almost surely) is used to denote an event that holds with probability tending to 1 as  $n \to \infty$ .

According to this theorem, the expected size of the largest component can be plotted. For example, the expected largest component of the result of "graphgen1 1000 0.001 topology" has a size  $O(1000^{2/3}) = O(100)$ . There is no exact expected size value. We can evaluate the correctness by judging whether the size is consistent with log(n), n or  $n^{\frac{2}{3}}$  in probability.

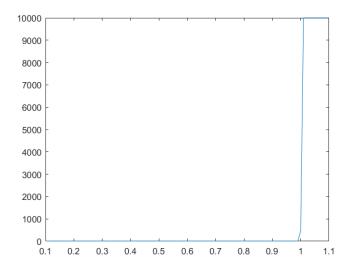


Figure 1: Expected Largest Component Size in Probability

## Graph Analysis-Sam

## References

[1.] latexcompanion S. Janson, T. Luczak, and A. Rucinski. Random Graphs. Wiley, 2000.