How to calculate $T_{I/O}$ for an HDD given the number of blocks

$$T_{I/O} = \left[(1 - DL) * \left(T_{seek} + T_{latency} \right) + T_{transfer} + T_{CTRL} \right] * N$$

$$T_{latency} = \frac{T_{rotation}}{2} \rightarrow T_{rotation} = \frac{60 * 1000}{RPM}$$

If I don't know RPM (o T_{seek}) $\rightarrow N*(1-DL)*T_{service\ time} + N*DL*(T_{transfer} + T_{CTRL})$

RAID Disks: Reliability Calculation

MTTF = Mean Time To Failure

- RAID 0
 - \circ $MTTF_{RAID,0} = MTTF_{single,disk}/N$
- RAID 1 (per N=2)

$$\circ MTTF_{RAID\ 1} = (MTTF_{single\ disk}/N) * (1/(MTTR/MTTF_{single\ disk})$$

• RAID 1+0

$$OMTTF_{RAID\ 1+0} = (MTTF_{single\ disk}^{2})/(N*MTTR)$$

- RAID 0+1
 - $\circ MTTF_{RAID\ 0+1} = (MTTF_{single\ disk}^{2})/(N*G*MTTR)$
 - G = (N/2) is the number of disks in a stripe group
- RAID 4

$$\circ MTTF_{RAID\ 4} = (MTTF_{single\ disk}^{2})/(N*N-1*MTTR)$$

• RAID 5

$$\circ MTTF_{RAID 5} = (MTTF_{single disk}^{2})/(N*N - 1*MTTR)$$

RAID 6

$$0 \quad MTTF_{RAID 6} = 2 * (MTTF_{single disk}^{3})/(N * N - 1 * N - 2 * MTTR)$$

	RAID 0	RAID 1	RAID 4	RAID 5	RAID 6
Capacity	N	N / 2	N – 1	N – 1	N – 2
Seq. Read	N * S	(N / 2) * S	(N - 1) * S	(N - 1) * S	
Seq. Write	N * S	(N / 2) * S	(N - 1) * S	(N - 1) * S	
Random Read	N * R	N * R	(N – 1) * R	N * R	
Random Write	N * R	(N / 2) * R	R / 2	(N / 4) * R	
Reliability	NO	1 (or N / 2)	1	1	2
(how many					
disk(s) can fail)					

Best performance and most capacity? → RAID 0

Greatest error recovery? → RAID 1 (1+0 or 0+1) or RAID 6

Balance between space, performance, and recoverability? → RAID 5

Performance

Operational laws are based on observable variables:

- T, the length of time we observe the system
- A_k , the number of request arrivals we observe for resource k
- ullet C_k, the number of request completions we observe at resource k
- B_k , the total amount of time during which the system is busy
- N_k , the average number of jobs in the resource k

From these observed values, we can derive the following 4 quantities:

- 1. $\lambda_k = A_k/T$, the arrival rate
- 2. $X_k = C_k/T$, the throughput or completion rate
- 3. $U_k = B_k/T$, the utilisation
- 4. $S_k = B_k/C_k$, the mean service time per completed job

(The upper formulas are also valid if we have total data and not partial ones like X = C/T)

Utilization Law: $U_k = X_k * S_k$ and $X_k = U_k / S_k$

Little's Law: $N_{insystem} = X * R$

Response Time Law: R = N/X - Z (where N number of users, X system throughput and Z think

time)

Forced Law: $X_k = v_k * X$

Other laws:

$V_k = C_k/C$	$R_k(residence\ time) = V_k * R_k\ (response\ time)$			
$D_k = V_k * S_k = B_k / C$	$R = \sum R_k$			
$U_k = X * D_k$	$N = \sum N_k$			
$X_{sys} = U_k/D_k$				

$$R = \sum_{i=1}^{k} \frac{D_k}{1 - \lambda D_k}$$

$$R(\lambda) = \frac{D_{CPU}}{1 - \lambda D_{CPU}} + \frac{D_{disk}}{1 - \frac{\lambda D_{disk}}{n}}$$

Fundamentals laws:

$$D = \sum_{k=1}^{K} D_k$$
 and $D_{max} = \max_k D_k$ $N^* = \frac{D+Z}{D_{max}}$

batch	$\frac{1}{D} \le X(N) \le \min\left(\frac{N}{D}, \frac{1}{D_{max}}\right)$
terminal	$\frac{N}{ND+Z} \le X(N) \le \min\left(\frac{N}{D+Z}, \frac{1}{D_{max}}\right)$
transaction	$X(\lambda) \leq \frac{1}{D_{max}}$

batch $\max(D, ND_{max}) \le R(N) \le ND$ terminal $\max(D, ND_{max} - Z) \le R(N) \le ND$ transaction $D \le R(\lambda)$ Transaction = Open Models
Batch & Terminal = Closed Models

Dependability

MTTF: mean time to (first) failure, the up time before the first failure

MTBF: mean time between failures = $\frac{total\ operating\ time}{number\ of\ failures}$ or $\frac{1}{\lambda}\ where\ \lambda = \frac{number\ of\ failures}{total\ operating\ time}$

Block diagram:

• Component in series

a. 2 blocks, C1 and C2

$$R_S = R_{C1} * R_{C2}$$

In general, if system S is composed by components with a reliability having an exponential distribution (very common case)

$$R_s(t) = e^{-\lambda_s t} \text{ where } \lambda_s = \sum_{i=1}^n \lambda_i \to MTTF_s = \frac{1}{\sum 1/MTTF_i}$$

A special case: when all components are identical

$$R_s(t) = e^{-\lambda_s t}$$



$$R_S(t) = e^{-n\lambda t} = e^{-\frac{nt}{MTTF_1}}$$
 $MTTF_S = \frac{MTTF_1}{n}$

Availability:

$$A_{S} = \prod_{i=1}^{n} \frac{MTTF_{i}}{MTTF_{i} + MTTR_{i}}$$

When all components are the same:

$$A_S(t) = A_1(t)^n$$
 $A = \left(\frac{MTTF_1}{MTTF_1 + MTTR_1}\right)^n$

• Component in parallel

a. 2 blocks, C1 and C2

$$R_S = R_{C1} + R_{C2} - R_{C1} * R_{C2}$$

$$\begin{split} MTTF_{parallel} &= MTTF * \sum_{n=1}^{N} \frac{1}{n} & \text{(if all MTTF are equal)} \\ MTTF_{parallel} &= \sum_{i=1}^{N} MTTF_{i} - \frac{1}{\sum_{i=1}^{N} \frac{1}{MTTF_{i}}} & \text{(if all MTTF are different)} \end{split}$$

System P composed by n components

$$R_P(t) = 1 - \prod_{i=1}^{n} (1 - R_i(t))$$

Availability

$$\begin{split} A_{p}(t) &= 1 - \prod_{i=1}^{n} \left(1 - A_{i}(t) \right) \\ A_{p} &= 1 - \prod_{i=1}^{n} \left(1 - A_{i} \right) = 1 - \prod_{i=1}^{n} \frac{MTTR_{i}}{MTTF_{i} + MTTR_{i}} \end{split}$$