

**Exercise 1** (Summation). Let  $p_i \in \mathbb{F}$ ,  $1 \leq i \leq n$  be a sequence of  $n$  floating-point numbers.

1. Show that the condition number of the computation of the summation satisfies

$$\text{cond}\left(\sum_{i=1}^n p_i\right) = \frac{\sum_{i=1}^n |p_i|}{\left|\sum_{i=1}^n p_i\right|}.$$

We recall that by definition

$$\text{cond}\left(\sum_{i=1}^n p_i\right) := \limsup_{\varepsilon \rightarrow 0} \left\{ \left| \frac{\sum_{i=1}^n \tilde{p}_i - \sum_{i=1}^n p_i}{\varepsilon \sum_{i=1}^n p_i} \right| : |\tilde{p}_i - p_i| \leq \varepsilon |p_i| \text{ for } i = 1, \dots, n \right\}.$$

2. Show that the recursive summation algorithm is *backward-stable*.
3. Derive a bound on the relative error for the summation.
4. Redo all the questions for the dot product.

**Exercise 2** (Polynomial evaluation). Let  $p(x) = \sum_{i=0}^n a_i x^i$  be a polynomial of degree  $n$  with floating-point coefficients.

1. Recall the formula for the condition number  $\text{cond}(p, x)$  of the polynomial evaluation of  $p$  in  $x$ .
2. Show that the Horner scheme for polynomial evaluation is *backward-stable*.
3. Derive a bound on the relative error for the polynomial evaluation.
4. Given a polynomial  $q(x) = \sum_{i=0}^n b_i x^i$ , we define the distance  $d(p, q) = \max_i \{|a_i - b_i|/|a_i|\}$ . Show that given  $p$  and  $z$ ,

$$\min\{d(p, q) : q(z) = 0\} = 1/\text{cond}(p, z).$$

**Exercise 3** (Roots of polynomials). Let  $p(x) = \sum_{i=0}^n a_i x^i$  be a polynomial of degree  $n$  with floating-point coefficients and  $\alpha$  a simple root ( $p(\alpha) = 0$  and  $p'(\alpha) \neq 0$ ).

1. We define the condition number of the simple root  $\alpha$  by

$$K(p, \alpha) := \lim_{\varepsilon \rightarrow 0} \sup_{|\Delta a_i| \leq \varepsilon |a_i|} \left\{ \frac{|\Delta \alpha|}{\varepsilon |\alpha|} \right\}.$$

Show that

$$K(p, \alpha) = \frac{\tilde{p}(|\alpha|)}{|\alpha| |p'(\alpha)|},$$

with  $\tilde{p}(x) := \sum_{i=0}^n |a_i| x^i$ .

2. When is a simple root ill-conditioned?

**Exercise 4** (Conditioning of the inverse of a matrix). In the sequel, we will use the Euclidean  $\|\cdot\|$ . We define the condition number of the computation of the inverse of a matrix by

$$\kappa(A) := \lim_{\varepsilon \rightarrow 0} \sup_{\|\Delta A\| \leq \varepsilon \|A\|} \left( \frac{\|(A + \Delta A)^{-1} - A^{-1}\|}{\varepsilon \|A^{-1}\|} \right).$$

1. Show that  $\kappa(A) = \|A\| \|A^{-1}\|$ .
2. We define the *distance to singularity* of a matrix  $A$  by

$$\text{dist}(A) := \min \left\{ \frac{\|\Delta A\|}{\|A\|} : A + \Delta A \text{ singular} \right\}.$$

Show that  $\text{dist}(A) = \kappa(A)^{-1}$ .

3. Express  $\kappa(A)$  in terms of the singular values of  $A$ .