

Week 04.2: HDFS

DS-GA 1004: Big Data

Hadoop distributed file system

- HDFS is the storage component of Hadoop
 - Useful beyond map-reduce!
- Provides distributed, redundant storage
- Optimizes for single-write, multiple-read patterns
 - Typical of map-reduce applications

Using HDFS

- HDFS is a "file system", but not like your OS file system
- HDFS sits on top of the operating system's built-in FS
- Better to think of it as an application that stores files for you
 - Kind of like Google Drive or Apple iCloud
 - Data can be accessed through the "hadoop fs" command



Two types of nodes in HDFS



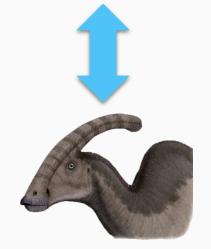


Name node Data nodes

The name node

- Clients talk to the name node to locate data
 - Analogous to the file system (but not storage device) in a standard OS
- Name node knows the mappings of:
 - Files → blocks
 - Blocks → data nodes
- Keeps a journal of transactions
 - Backed up remotely for durability





Data nodes

- Stores each block as two files in the local file system:
 - **Data block** (variable size up to defined max, typically 128MB)
 - Metadata:
 - Checksum
 - Generation stamp
- Checksum: used to detect storage errors
 - Analogous to parity blocks in RAID, but replicated on all nodes
 - Weaker than parity: can detect errors, but not correct them
- **Generation stamp**: used to detect updates

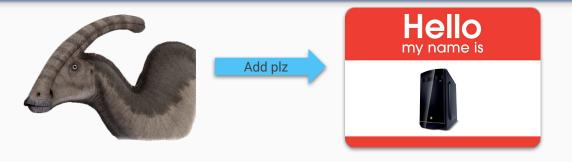


Division of responsibilities

- Name nodes do not store data!
- Data nodes do not store **metadata** (e.g. file names)!
- Name node failure is catastrophic
- Data node failure can be tolerated, up to a point
 - Depends on how much replication you have







1. Client wants to add a block

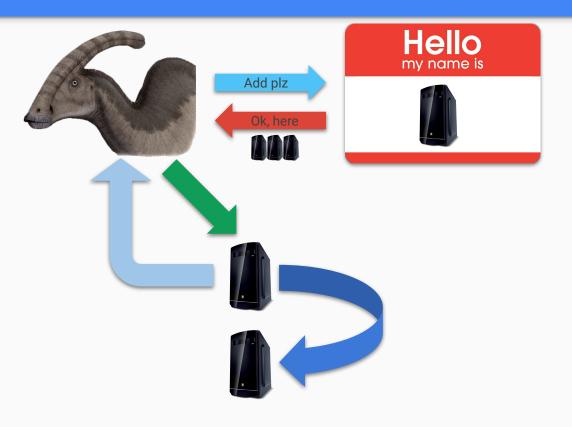




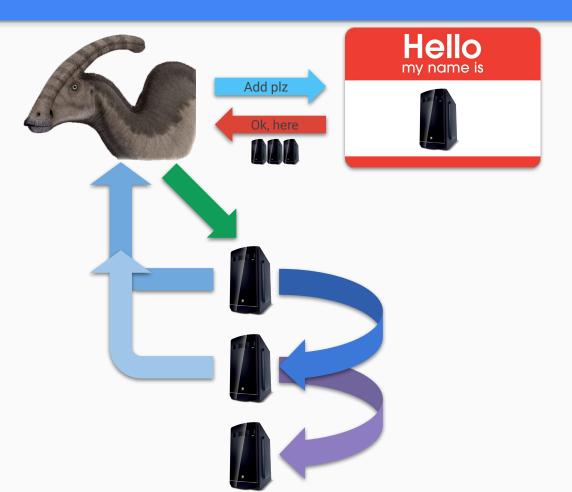
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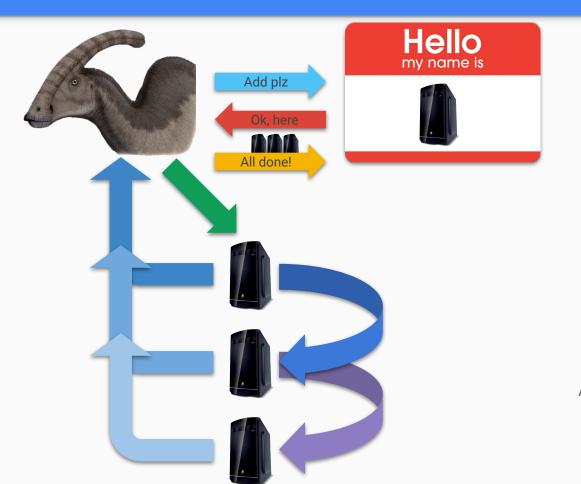
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 - Name node responds with a list of data nodes
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- 3. DN1 stores, acknowledges, and sends block to DN2
- 4. DN2 stores, acknowledges, and sends block to DN3



- Client wants to add a block
 - Name node responds with a list of data nodes
- Client sends block to DN1
- 3. DN1 stores, acknowledges, and sends block to DN2
- 4. DN2 stores, acknowledges, and sends block to DN3
- 5. DN3 stores and acknowledges

Add complete, close file, alert name node

Name and data node communication

- Data nodes periodically signal the name node
 - "Heartbeat"
- The name node always knows which data nodes are alive
 - At least, within the last 3 seconds
 - Name node can infer failures and insufficient replication
- Name node may respond with update messages
 - E.g.: replicate block x from data node y



Recovering from failure: checkpoints

- Checkpoints are snapshots of the current name node's state
 - Directory structure, block maps, and journal
 - Name node keeps all of this information in RAM
- These are created periodically to ensure fast recovery when NN fails
- Checkpoints cannot be updated, only replaced

HDFS isn't quite POSIX-compliant

- Updates are append-only
 - No changing old data!
 - This makes replication logic much simpler
 - What other benefits might immutability offer?
- Not all file modes are supported
 - Not all modes make sense in this limited context anyway
 - E.g., executable

Why doesn't HDFS work like my desktop?

- Desktop computing needs to support all kinds of uses
 - E.g. thousands of small configuration files
 - Files that update frequently (e.g., browser cache)
- Large data analysis jobs have different needs
 - Few, large files
 - Frequent read access (analysis)
 - Infrequent updates (append-only can be okay)

Up next

• Part 3: HDFS + MapReduce