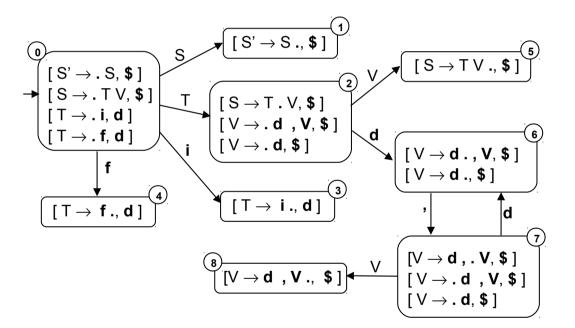
Specify the LR(1) parsing table relevant to the language defined by the following grammar:

```
decl 
ightarrow type \ var-list type 
ightarrow int \mid float var-list 
ightarrow id , var-list \mid id
```

Specify the LR(1) parsing table relevant to the language defined by the following grammar:

$$decl 
ightarrow type \ var-list$$
  
 $type 
ightarrow {
m int} \ | \ {
m float}$   
 $var-list 
ightarrow {
m id}$  ,  $var-list \ | \ {
m id}$ 

$$S' \rightarrow S$$
 $S \rightarrow T V$ 
 $T \rightarrow \mathbf{i} \mid \mathbf{f}$ 
 $V \rightarrow \mathbf{d}, V \mid \mathbf{d}$ 



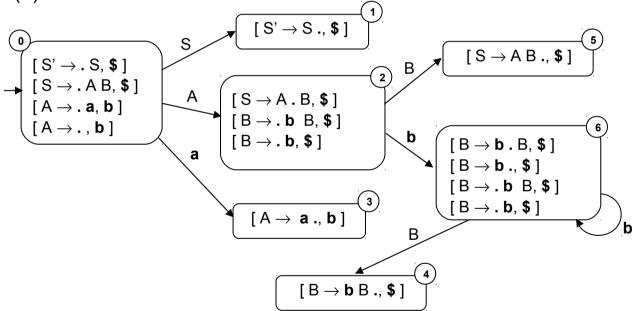
State	Input						Goto		
	i	f	d	,	\$	S	T	V	
0	s3	s4				1	2		
1					accept				
2			s6					5	
3			$T \rightarrow \mathbf{i}$						
4			$T \rightarrow \mathbf{f}$						
5					$S \to T V$				
6				s7	$V \rightarrow \mathbf{d}$				
7			s6					8	
8					$V\! o\!\operatorname{d}$ , $V$				

$$S \rightarrow A B$$
  
 $A \rightarrow \mathbf{a} \mid \varepsilon$   
 $B \rightarrow \mathbf{b} B \mid \mathbf{b}$ 

Check whether the following grammar is LR(1).

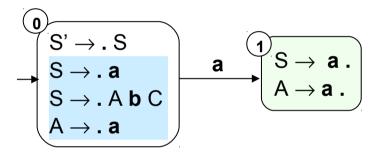
$$S \rightarrow A B$$
  
 $A \rightarrow \mathbf{a} \mid \varepsilon$   
 $B \rightarrow \mathbf{b} B \mid \mathbf{b}$ 

The grammar is LR(1):



$$S \rightarrow \mathbf{a} \mid A \mathbf{b} C$$
  
 $A \rightarrow \mathbf{a}$   
 $C \rightarrow A \mid \mathbf{c}$ 

$$S \rightarrow \mathbf{a} \mid A \mathbf{b} C$$
  
 $A \rightarrow \mathbf{a}$   
 $C \rightarrow A \mid \mathbf{c}$ 



$$FOLLOW(S') = \{ \$ \}$$
  
 $FOLLOW(S) = \{ \$ \}$   
 $FOLLOW(C) = \{ \$ \}$   
 $FOLLOW(A) = \{ \$, b \}$ 

$$FOLLOW(S) \cap FOLLOW(A) = \{ \$ \} \neq \emptyset$$

$$\underbrace{ \text{Not SLR}(1)!}$$

$$S \rightarrow A \ B \mid \mathbf{b}$$
  
 $A \rightarrow A \ \mathbf{a} \mid \mathbf{\epsilon}$   
 $B \rightarrow B \ \mathbf{a} \mid \mathbf{b} \mid \mathbf{\epsilon}$ 

$$S \rightarrow A \ B \mid \mathbf{b}$$
  
 $A \rightarrow A \ \mathbf{a} \mid \mathbf{\epsilon}$   
 $B \rightarrow B \ \mathbf{a} \mid \mathbf{b} \mid \mathbf{\epsilon}$ 

$$\begin{array}{c} \textbf{0} \\ \textbf{S'} \rightarrow \textbf{.} \ \textbf{S} \\ \textbf{S} \rightarrow \textbf{.} \ \textbf{A} \ \textbf{B} \\ \textbf{S} \rightarrow \textbf{.} \ \textbf{b} \\ \textbf{A} \rightarrow \textbf{.} \ \textbf{A} \ \textbf{a} \\ \textbf{A} \rightarrow \textbf{.} \end{array}$$

$$FOLLOW(S') = \{ \$ \}$$
  
 $FOLLOW(S) = \{ \$ \}$   
 $FIRST(B) = \{ a, b, \epsilon \}$   
 $FOLLOW(B) = \{ a, \$ \}$   
 $FOLLOW(A) = \{ a, b, \$ \}$ 

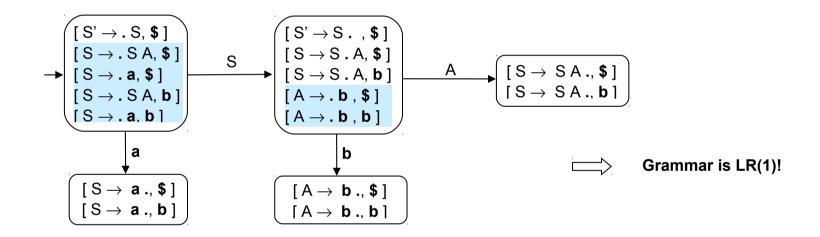
$$\implies \mathbf{b} \in FOLLOW(A) \qquad \stackrel{\frown}{\Longrightarrow} \qquad \underline{\mathbf{Not}} \ \mathbf{SLR}(1)!$$

After generating the relevant parsing automaton, determine (with motivation) whether the following grammar is LR(1):

$$S \to SA \mid \mathbf{a}$$
$$A \to \mathbf{b}$$

After generating the relevant parsing automaton, determine (with motivation) whether the following grammar is LR(1):

$$S \to SA \mid \mathbf{a}$$
$$A \to \mathbf{b}$$



Prove that the following grammar is <u>not</u> LR(1):

$$S \rightarrow A \mathbf{a} \mid \mathbf{b}$$

$$A \rightarrow B \mid \mathbf{a}$$

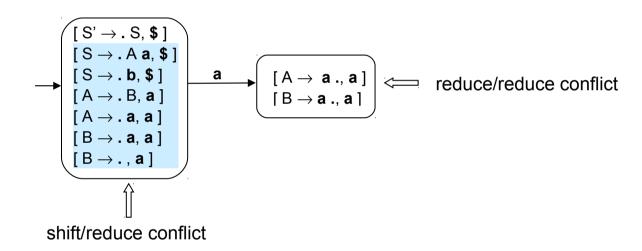
$$B \rightarrow \mathbf{a} \mid \mathbf{\epsilon}$$

Prove that the following grammar is <u>not</u> LR(1):

$$S \rightarrow A \mathbf{a} \mid \mathbf{b}$$

$$A \rightarrow B \mid \mathbf{a}$$

$$B \rightarrow \mathbf{a} \mid \mathbf{\epsilon}$$

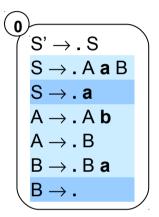


$$S \rightarrow A \mathbf{a} B \mid \mathbf{a}$$
  
 $A \rightarrow A \mathbf{b} \mid B$   
 $B \rightarrow B \mathbf{a} \mid \varepsilon$ 

$$S \rightarrow A \mathbf{a} B \mid \mathbf{a}$$

$$A \rightarrow A \mathbf{b} \mid B$$

$$B \rightarrow B \mathbf{a} \mid \varepsilon$$



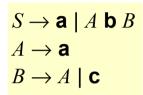
$$FOLLOW(S') = \{ \$ \}$$
  
 $FOLLOW(S) = \{ \$ \}$   
 $FOLLOW(A) = \{ a, b \}$   
 $FOLLOW(B) = \{ a, b, \$ \}$ 

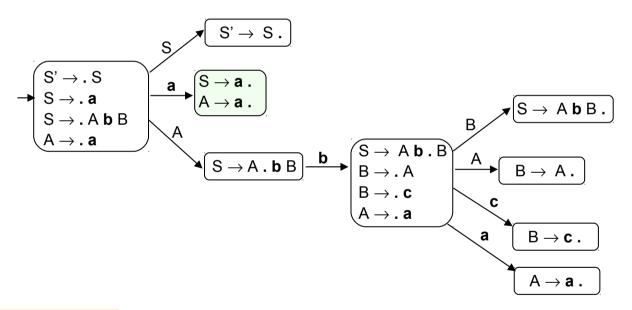
$$a \in FOLLOW(B)$$
  $\longrightarrow$  Not SLR(1)!

After generating the relevant parsing automaton, determine whether the following grammar is SLR(1):

$$S \rightarrow \mathbf{a} \mid A \mathbf{b} B$$
  
 $A \rightarrow \mathbf{a}$   
 $B \rightarrow A \mid \mathbf{c}$ 

After generating the relevant parsing automaton, determine whether the following grammar is SLR(1):





$$FOLLOW(S) = \{ \$ \}$$

$$FOLLOW(B) = \{ \$ \}$$

$$FOLLOW(A) = \{ \mathbf{b}, \$ \}$$

$$FOLLOW(A) \cap FOLLOW(S) = \{ \$ \} \neq \emptyset$$

$$\longrightarrow Not SLR(1)!$$

Generate the initial state of the LR(1) parsing (deterministic) automaton relevant to the following BNF:

$$E \rightarrow T + E \mid T$$
  
 $T \rightarrow F * T \mid F$   
 $F \rightarrow \mathbf{n} \mid (E)$ 

Generate the initial state of the LR(1) parsing (deterministic) automaton relevant to the following BNF:

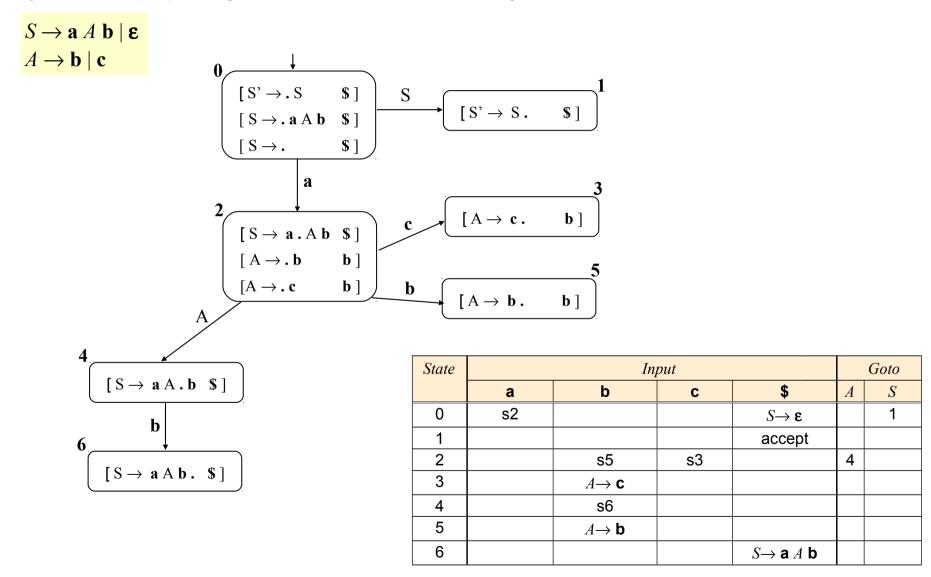
$$E \rightarrow T + E \mid T$$
  
 $T \rightarrow F * T \mid F$   
 $F \rightarrow \mathbf{n} \mid (E)$ 

$(E' \rightarrow . E,$	\$]
[ $E \rightarrow . T + E$ ,	\$]
[ $E \rightarrow .T$ ,	\$]
$[T \rightarrow . F * T,$	+]
$[T \rightarrow . F,$	+]
$[T \rightarrow .F * T,$	\$]
$[T \rightarrow . F,$	\$]
$[F \rightarrow . n,$	* ]
[ $F  o$ . ( $E$ ),	* ]
$[F \rightarrow . n,$	+]
[ $F \to .$ ( $E$ ),	+]
$[F \rightarrow . n,$	\$]
$\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ $	\$]

Specify the LALR(1) parsing table relevant to the following BNF:

$$S \to \mathbf{a} A \mathbf{b} \mid \mathbf{\varepsilon}$$
$$A \to \mathbf{b} \mid \mathbf{c}$$

Specify the LALR(1) parsing table relevant to the following BNF:



$$S \to \mathbf{x} \mid A \mathbf{y} B$$

$$A \to \mathbf{x}$$

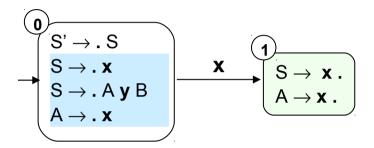
$$B \to A \mid \mathbf{z} \mid \mathbf{w} B$$

Check whether the following grammar is SLR(1):

$$S \rightarrow \mathbf{x} \mid A \mathbf{y} B$$

$$A \rightarrow \mathbf{x}$$

$$B \rightarrow A \mid \mathbf{z} \mid \mathbf{w} B$$



$$FOLLOW(S) \cap FOLLOW(A) = \{ \$ \} \neq \emptyset$$

**Not** SLR(1)!

Generate the initial state of the LR(1) parsing (deterministic) automaton relevant to the following BNF:

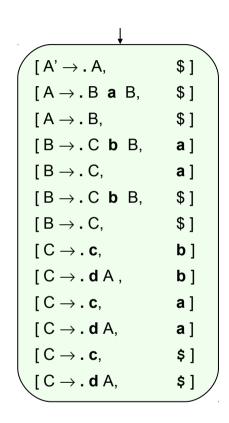
$$A \rightarrow B \mathbf{a} B \mid B$$

$$B \rightarrow C \mathbf{b} B \mid C$$

$$C \rightarrow \mathbf{c} \mid \mathbf{d} A$$

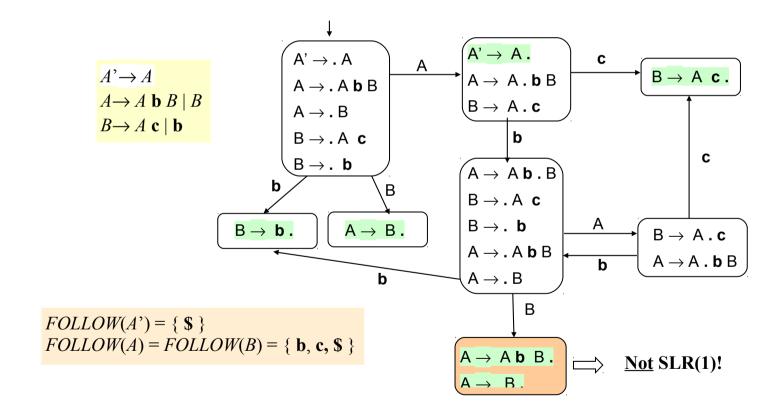
Generate the initial state of the LR(1) parsing (deterministic) automaton relevant to the following BNF:

 $A \rightarrow B \mathbf{a} B \mid B$   $B \rightarrow C \mathbf{b} B \mid C$   $C \rightarrow \mathbf{c} \mid \mathbf{d} A$ 



$$A \rightarrow A \mathbf{b} B \mid B$$
$$B \rightarrow A \mathbf{c} \mid \mathbf{b}$$

$$A \rightarrow A \mathbf{b} B \mid B$$
$$B \rightarrow A \mathbf{c} \mid \mathbf{b}$$



Outline the initial state of the LR(1) parsing (deterministic) automaton relevant to the following BNF:

$$S \rightarrow S \mathbf{a} T \mid T$$
$$T \rightarrow S \mathbf{b} \mid \mathbf{a}$$

Outline the initial state of the LR(1) parsing (deterministic) automaton relevant to the following BNF:

 $S \rightarrow S \mathbf{a} T \mid T$  $T \rightarrow S \mathbf{b} \mid \mathbf{a}$ 

```
 \begin{bmatrix} S' \to . & S, & & S \\ [S \to . & S & a & T, & & S \\ [S \to . & T, & & S \\ [S \to . & T, & & a] \\ [S \to . & T, & & a] \\ [S \to . & T, & & a] \\ [T \to . & S & b, & & S \\ [T \to . & S & b, & & a] \\ [T \to . & S & a & T, & & b] \\ [S \to . & T, & & b] \\ [S \to . & T, & & b] \\ [T \to . & S & b, & & b] \\ [T \to . & S & b, & & b]
```

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After constructing the complete parsing automaton for the following BNF, determine whether the grammar is SLR(1):

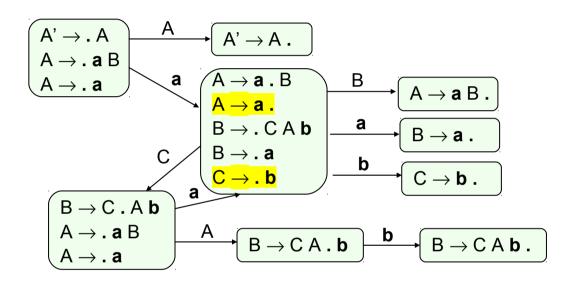
$$A \rightarrow \mathbf{a} \ B \mid \mathbf{a}$$

$$B \rightarrow C \ A \ \mathbf{b} \mid \mathbf{a}$$

$$C \rightarrow \mathbf{b}$$

After constructing the complete parsing automaton for the following BNF, determine whether the grammar is SLR(1):

$$A \rightarrow \mathbf{a} \ B \mid \mathbf{a}$$
  
 $B \rightarrow C A \mathbf{b} \mid \mathbf{a}$   
 $C \rightarrow \mathbf{b}$ 



 $b \in FOLLOW(A) = \{ \$, b \} \Rightarrow BNF \underline{not} SLR(1) !$ 

#### Given the following BNF:

$$A \rightarrow B \mathbf{a} C \mid B$$
$$B \rightarrow C \mathbf{b} B \mid C$$
$$C \rightarrow \mathbf{a}$$

#### We ask to:

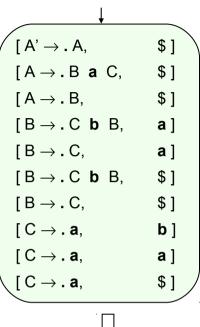
- Generate the initial state of the LR(1) (deterministic) parsing automaton.
- Check whether such a state presents any parsing conflict.

#### Given the following BNF:

$$A \rightarrow B \mathbf{a} C \mid B$$
$$B \rightarrow C \mathbf{b} B \mid C$$
$$C \rightarrow \mathbf{a}$$

#### We ask to:

- Generate the initial state of the LR(1) (deterministic) parsing automaton.
- Check whether such a state presents any parsing conflict.





No parsing conflict

After constructing the parsing automaton for the following grammar,

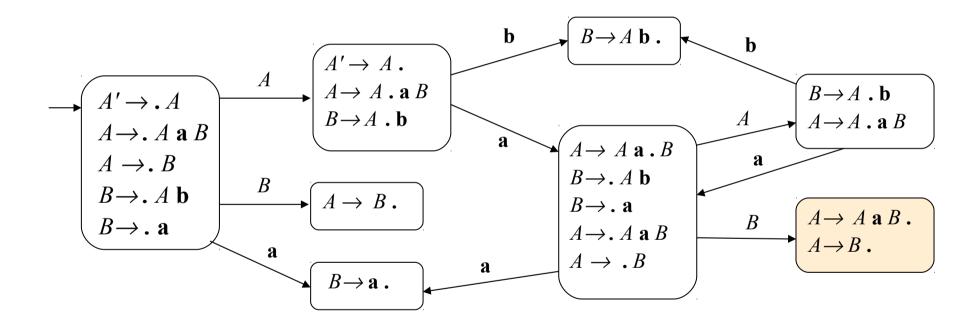
$$A \to A \mathbf{a} B \mid B$$
$$B \to A \mathbf{b} \mid \mathbf{a}$$

determine whether the grammar is SLR(1).

After constructing the parsing automaton for the following grammar,

$$A \rightarrow A \mathbf{a} B \mid B$$
$$B \rightarrow A \mathbf{b} \mid \mathbf{a}$$

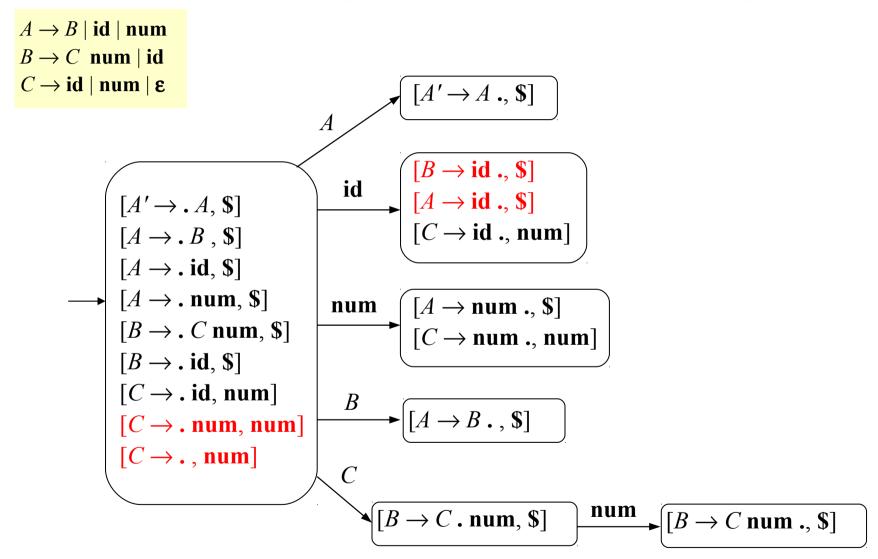
determine whether the grammar is SLR(1).



After constructing the complete parsing automaton, check whether the following grammar is LR(1).

 $A \rightarrow B \mid id \mid num$   $B \rightarrow C \quad num \mid id$  $C \rightarrow id \mid num \mid \epsilon$ 

After constructing the complete parsing automaton, check whether the following grammar is LR(1).



With reference to the following BNF, after constructing the portion of the LR(1) parsing automaton up to the states reached by a <u>single</u> transition from the initial state, check whether this portion of automaton includes any conflict:

$$A \rightarrow \mathbf{a} \ B \mid C$$

$$B \rightarrow \mathbf{b} \ B \mid C$$

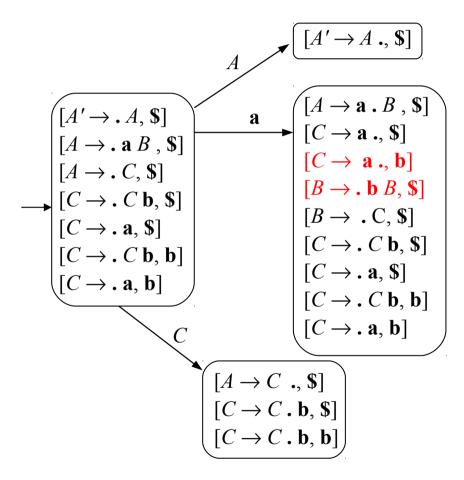
$$C \rightarrow C \ \mathbf{b} \mid \mathbf{a}$$

With reference to the following BNF, after constructing the portion of the LR(1) parsing automaton up to the states reached by a <u>single</u> transition from the initial state, check whether this portion of automaton includes any conflict:

$$A \rightarrow \mathbf{a} \ B \mid C$$

$$B \rightarrow \mathbf{b} \ B \mid C$$

$$C \rightarrow C \ \mathbf{b} \mid \mathbf{a}$$



shift/reduce conflict

After constructing the complete parsing automaton for the following grammar, determine whether it is SLR(1), providing relevant explanation.

$$A' \rightarrow A$$

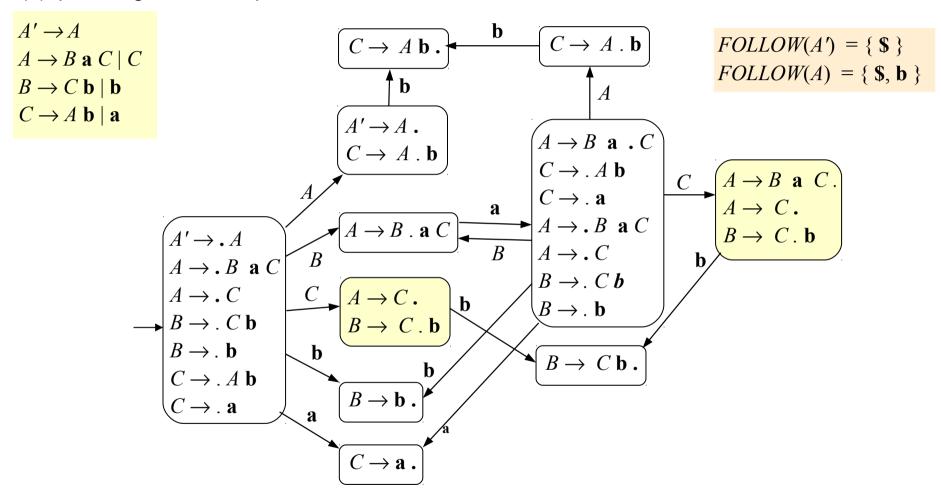
$$A \rightarrow B \mathbf{a} C \mid C$$

$$B \rightarrow C \mathbf{b} \mid \mathbf{b}$$

$$C \rightarrow A \mathbf{b} \mid \mathbf{a}$$

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After constructing the complete parsing automaton for the following grammar, determine whether it is SLR(1), providing relevant explanation.



Outline the LR(1) parsing table relevant to the following BNF.

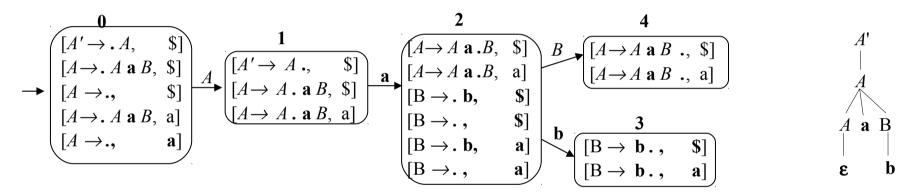
$$A \to A \mathbf{a} B \mid \mathbf{\varepsilon}$$
$$B \to \mathbf{b} \mid \mathbf{\varepsilon}$$

Then, trace the LR(1) parsing of phrase **a b**. Finally, draw the corresponding syntax tree based on the traced parsing actions.

Outline the LR(1) parsing table relevant to the following BNF.

$$A \to A \mathbf{a} B \mid \mathbf{\varepsilon}$$
$$B \to \mathbf{b} \mid \mathbf{\varepsilon}$$

Then, trace the LR(1) parsing of phrase **a b**. Finally, draw the corresponding syntax tree based on the traced parsing actions.



State	Input			Goto	
	а	b	\$	A	В
0	$A \rightarrow \mathbf{\epsilon}$		$A \rightarrow \mathbf{\epsilon}$	1	
1	s2		accept		
2	$B \rightarrow \mathbf{\epsilon}$	s3	$B  o \mathbf{\epsilon}$		4
3	$B \rightarrow \mathbf{b}$		$B \rightarrow \mathbf{b}$		
4	$A \rightarrow A \mathbf{a} B$		$A \rightarrow A \mathbf{a} B$		

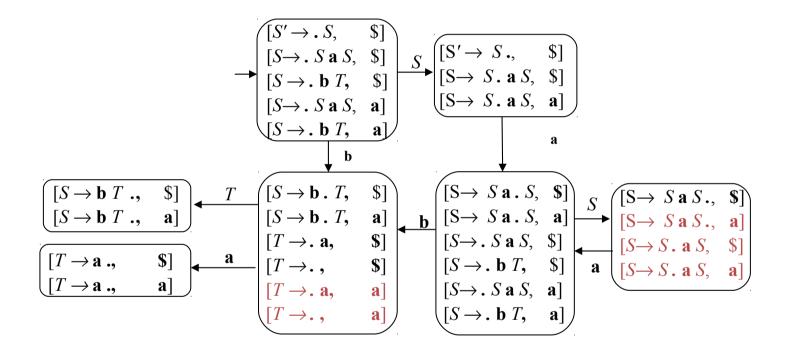
Stack	Input	Action
<b>\$</b> 0	a <b>b \$</b>	$A \rightarrow \mathbf{\epsilon}$
<b>\$</b> 0 A1	ab\$	shift
<b>\$</b> 0 A1 a2	b \$	shift
<b>\$</b> 0 A1 a2 b3	\$	$B \rightarrow \mathbf{b}$
<b>\$</b> 0 A1 a2 B4	\$	$A \rightarrow A \mathbf{a} B$
<b>\$</b> 0 A1	\$	accept

After constructing the parsing automaton for the following BNF, discuss whether the BNF is LR(1).

$$S \rightarrow S \mathbf{a} S \mid \mathbf{b} T$$
  
 $T \rightarrow \mathbf{a} \mid \mathbf{\varepsilon}$ 

After constructing the parsing automaton for the following BNF, discuss whether the BNF is LR(1).

$$S \to S \mathbf{a} S \mid \mathbf{b} T$$
$$T \to \mathbf{a} \mid \mathbf{\epsilon}$$

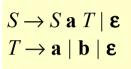


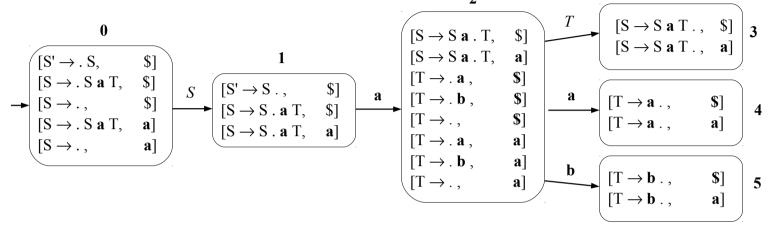
Given the following BNF, we ask to generate the complete LR(1) parsing table and to establish whether the grammar is LR(1).

$$S \rightarrow S \mathbf{a} T \mid \mathbf{\varepsilon}$$
  
 $T \rightarrow \mathbf{a} \mid \mathbf{b} \mid \mathbf{\varepsilon}$ 

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Given the following BNF, we ask to generate the complete LR(1) parsing table and to establish whether the grammar is LR(1).





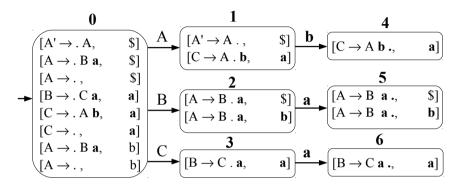
State	Input			Goto	
	а	b	\$	S	T
0	$S \rightarrow \mathbf{\epsilon}$		$S \rightarrow \mathbf{\epsilon}$	1	
1	s2		accept		
2	$T \rightarrow \mathbf{\varepsilon}$ s4	<b>s</b> 5	$T \rightarrow \mathbf{\epsilon}$		3
3	$S \to S \mathbf{a}$ $T$		$S \to S$ a $T$		
4	$T \rightarrow \mathbf{a}$		$T \rightarrow \mathbf{a}$		
5	$T \rightarrow \mathbf{b}$		$T \rightarrow \mathbf{b}$		

Given the following BNF, we ask to (i) outline the corresponding LR(1) parsing table, (ii) trace the LR(1) parsing of phrase **baa**, and (iii) based on the trace, outline the syntax tree of the given phrase.

$$A \rightarrow B \mathbf{a} \mid \mathbf{\epsilon}$$
 $B \rightarrow C \mathbf{a}$ 
 $C \rightarrow A \mathbf{b} \mid \mathbf{\epsilon}$ 

Given the following BNF, we ask to (i) outline the corresponding LR(1) parsing table, (ii) trace the LR(1) parsing of phrase **baa**, and (iii) based on the trace, outline the syntax tree of the given phrase.

$$A \rightarrow B \mathbf{a} \mid \mathbf{\epsilon}$$
 $B \rightarrow C \mathbf{a}$ 
 $C \rightarrow A \mathbf{b} \mid \mathbf{\epsilon}$ 



State	Input			Goto		
	а	b	\$	A	В	C
0	$C \rightarrow \mathbf{\epsilon}$	$A \to \mathbf{\epsilon}$	$A \to \mathbf{\epsilon}$	1	2	3
1		s4	accept			
2	s5					
3	s6					
4	$C \rightarrow A \mathbf{b}$					
5		$A \rightarrow B$ a	$A \rightarrow B$ a			
6	$B \rightarrow C$ a					

Stack	Input	Action
<b>\$</b> 0	baa\$	$A  ightarrow \epsilon$
<b>\$</b> 0 A1	baa\$	shift
<b>\$</b> 0 A1 <b>b</b> 4	a a \$	$C \rightarrow A \mathbf{b}$
<b>\$</b> 0 C3	a a \$	shift
<b>\$</b> 0 C3 <b>a</b> 6	a \$	$B \rightarrow C$ a
<b>\$</b> 0 B2	a \$	shift
<b>\$</b> 0 B2 <b>a</b> 5	\$	$A \rightarrow B$ a
<b>\$</b> 0 A1	\$	accept

