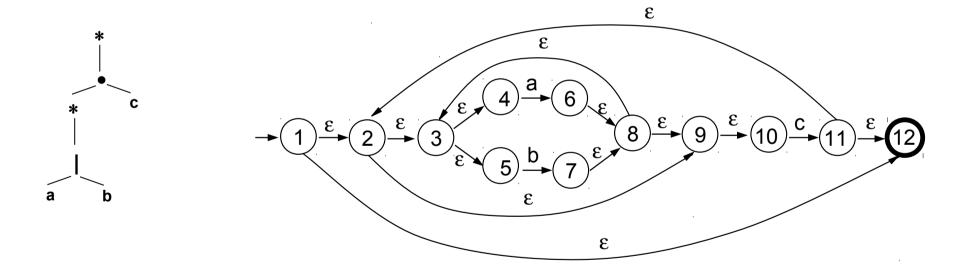
Outline the NFA generated by the construction of Thompson relevant to the following regular expression:

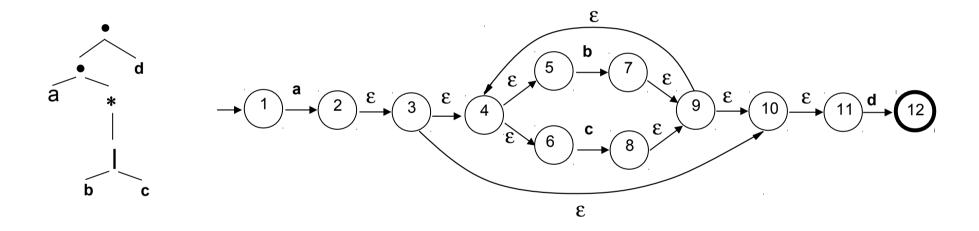
((a | b)\*c)\*

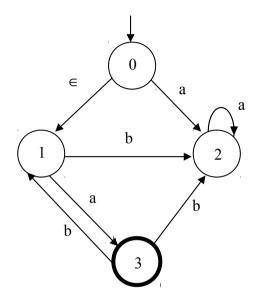
Outline the NFA generated by the construction of Thompson relevant to the following regular expression:

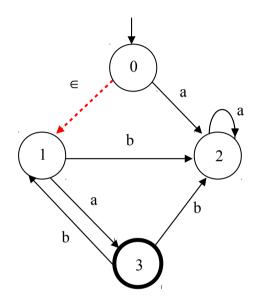


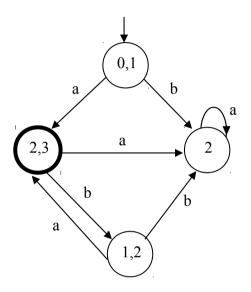
By means of the construction of Thompson, outline the NFA relevant to the following regular expression:

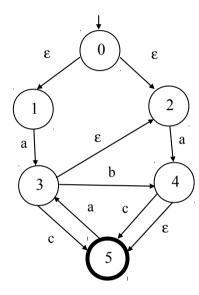
By means of the construction of Thompson, outline the NFA relevant to the following regular expression:

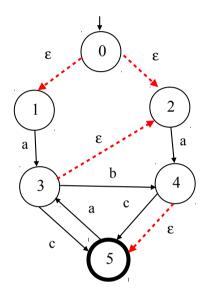


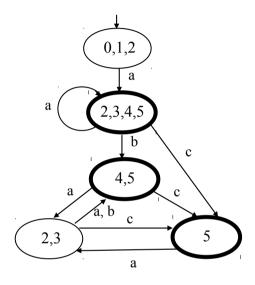


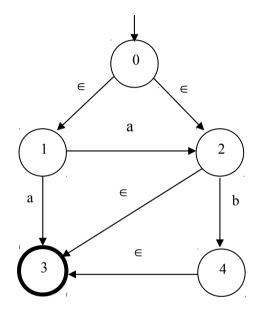


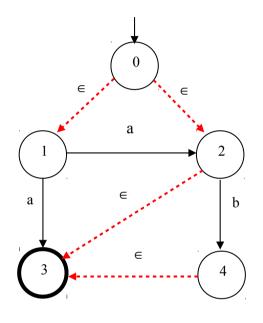


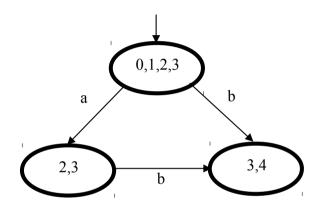


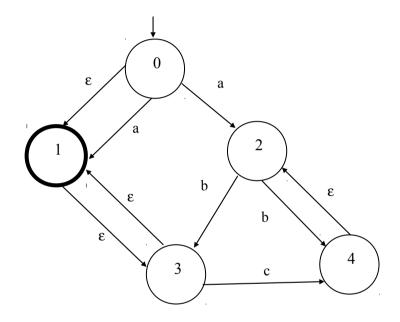


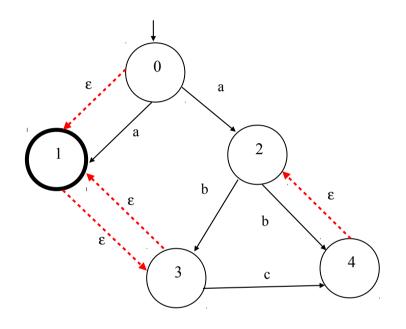


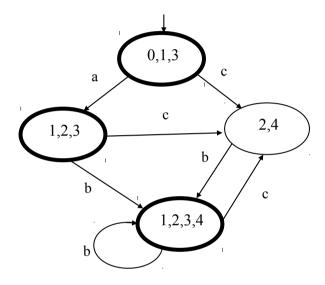


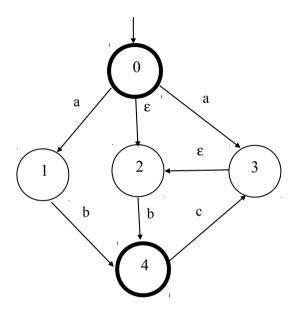


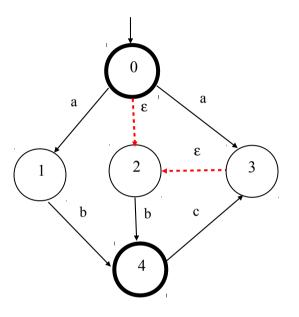


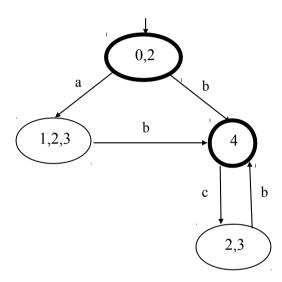


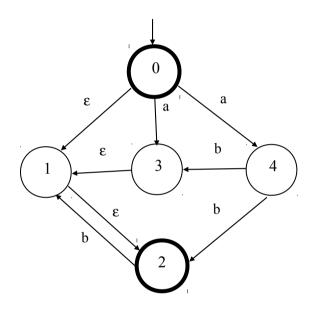


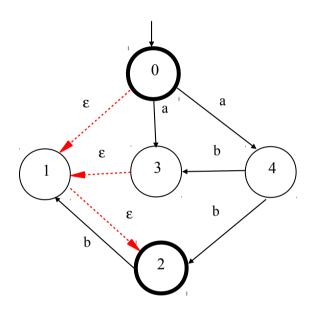


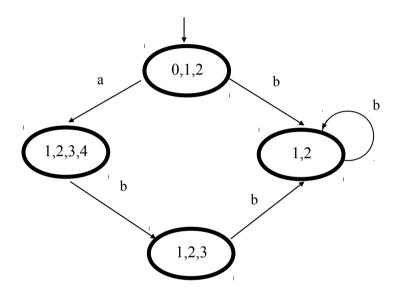










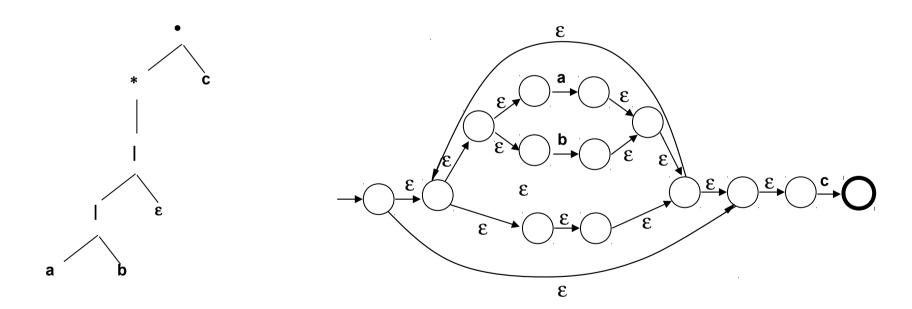


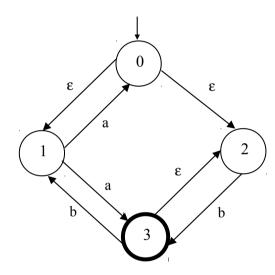
Given the regular expression  $\mathbf{r}$ :  $((\mathbf{a} \mid \mathbf{b}) \mid \mathbf{\epsilon})^* \mathbf{c}$ , we ask to:

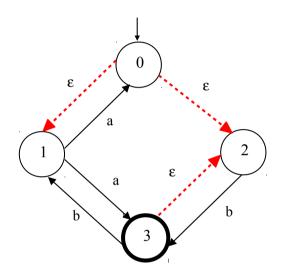
- a) Outline the tree of expression **r**;
- b) Based on the tree of **r**, outline, by means of the construction of Thompson, the NFA recognizing the regular language of **r**.

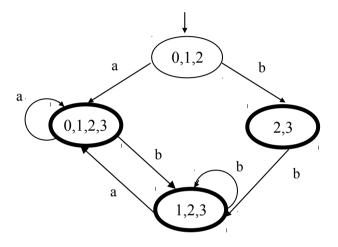
Given the regular expression  $\mathbf{r}$ :  $((\mathbf{a} \mid \mathbf{b}) \mid \mathbf{\epsilon})^* \mathbf{c}$ , we ask to:

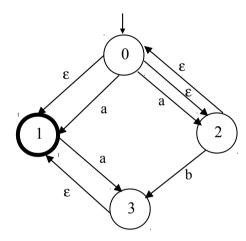
- a) Outline the tree of expression **r**;
- b) Based on the tree of **r**, outline, by means of the construction of Thompson, the NFA recognizing the regular language of **r**.

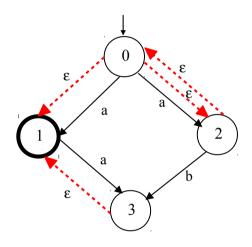


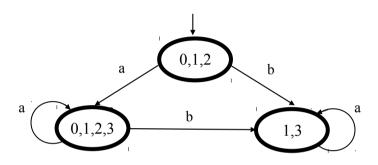










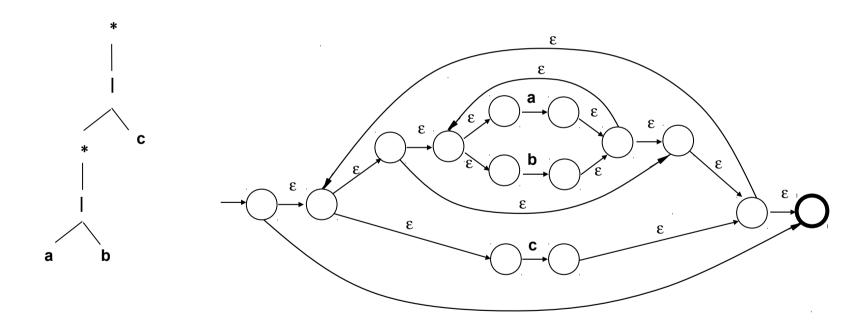


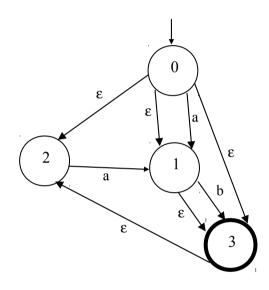
Given the regular expression  $\mathbf{r}$ :  $((\mathbf{a} \mid \mathbf{b})^* \mid \mathbf{c})^*$ , we ask to:

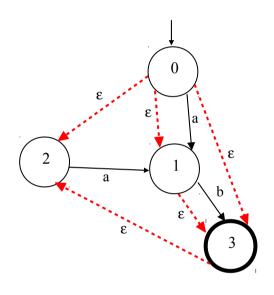
- a) Outline the tree of expression **r**;
- b) Based on the tree of **r**, outline, by means of the canonical construction of Thompson, the DFA recognizing the regular language of **r**.

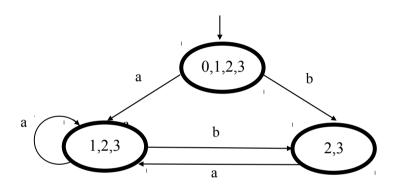
Given the regular expression  $\mathbf{r}$ :  $((\mathbf{a} \mid \mathbf{b})^* \mid \mathbf{c})^*$ , we ask to:

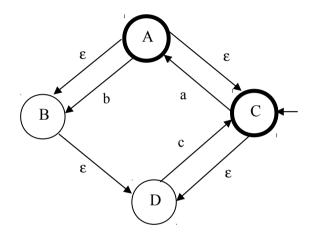
- a) Outline the tree of expression **r**;
- b) Based on the tree of  $\mathbf{r}$ , outline, by means of the canonical construction of Thompson, the DFA recognizing the regular language of  $\mathbf{r}$ .

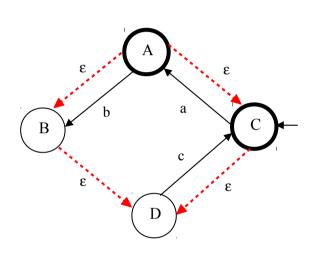


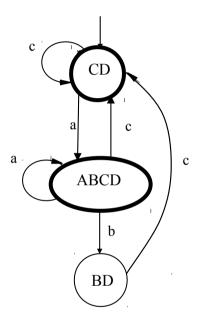




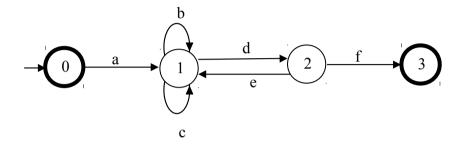




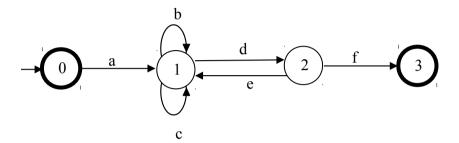




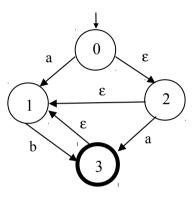
Specify the BNF of the regular language defined by the following DFA:

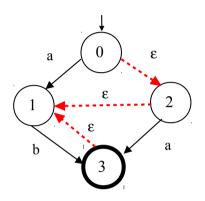


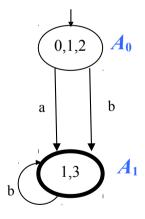
Specify the BNF of the regular language defined by the following DFA:



$$A_0 
ightharpoonup \mathbf{a} A_1 \mid \mathbf{\epsilon}$$
 $A_1 
ightharpoonup \mathbf{b} A_1 \mid \mathbf{c} A_1 \mid \mathbf{d} A_2$ 
 $A_2 
ightharpoonup \mathbf{e} A_1 \mid \mathbf{f} A_3$ 
 $A_3 
ightharpoonup \mathbf{\epsilon}$ 

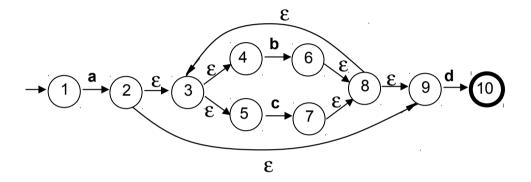


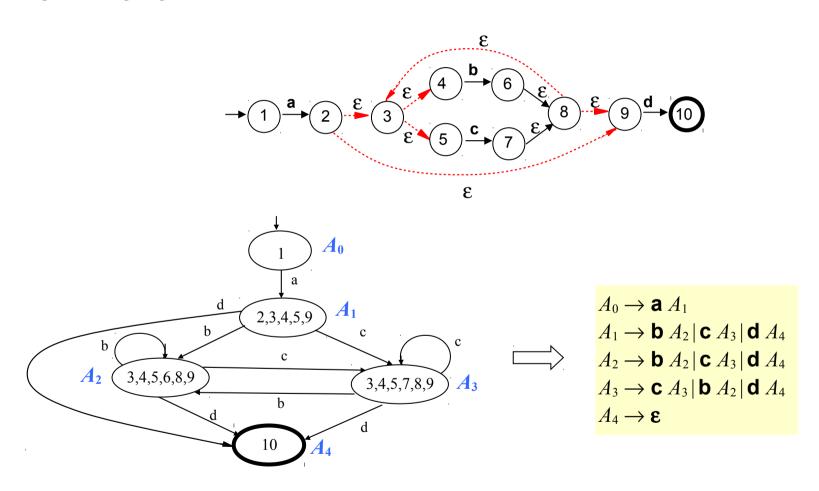


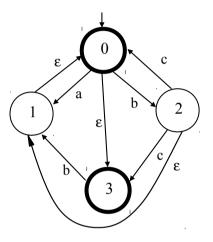


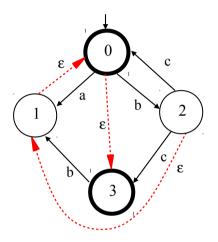


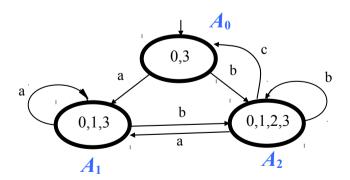
$$A_0 \rightarrow \mathbf{a} A_1 | \mathbf{b} A_1$$
  
 $A_1 \rightarrow \mathbf{b} A_1 | \mathbf{\epsilon}$ 





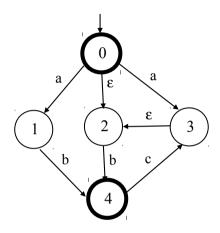




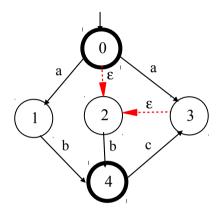


$$A_0 
ightarrow \mathbf{a} A_1 | \mathbf{b} A_2 | \mathbf{\epsilon}$$
 $A_1 
ightarrow \mathbf{a} A_1 | \mathbf{b} A_2 | \mathbf{\epsilon}$ 
 $A_2 
ightarrow \mathbf{a} A_1 | \mathbf{b} A_2 | \mathbf{c} A_0 | \mathbf{\epsilon}$ 

Specify the BNF of the regular language relevant to the following NFA:



Specify the BNF of the regular language relevant to the following NFA:



$$A_0 \rightarrow \mathbf{a} A_1 \mid A_2 \mid \mathbf{a} A_3 \mid \mathbf{\epsilon}$$

$$A_1 \rightarrow \mathbf{b} A_4$$

$$A_2 \rightarrow \mathbf{b}A_4$$

$$A_3 \rightarrow A_2$$

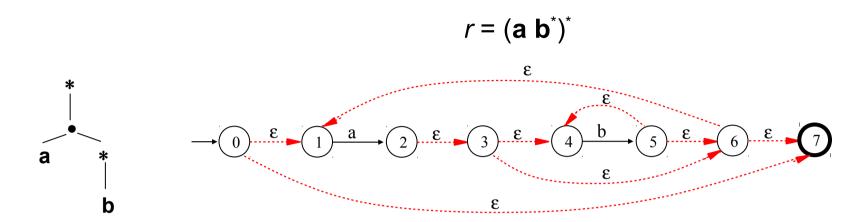
$$A_4 \rightarrow \mathbf{c} A_3 \mid \epsilon$$

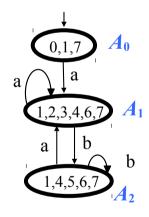
Given the following regular expression r,

$$r = (\mathbf{a} \ \mathbf{b}^*)^*$$

we ask to:

- a) Outline NFA(*r*) generated by the canonical construction of Thompson;
- b) Transform NFA(r) into the equivalent DFA(r);
- c) Based on DFA(r), specify the BNF of the regular language relevant to r.





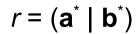
$$A_0 \rightarrow a A_1 \mid \varepsilon$$
 $A_1 \rightarrow a A_1 \mid b A_2 \mid \varepsilon$ 
 $A_2 \rightarrow a A_1 \mid b A_2 \mid \varepsilon$ 

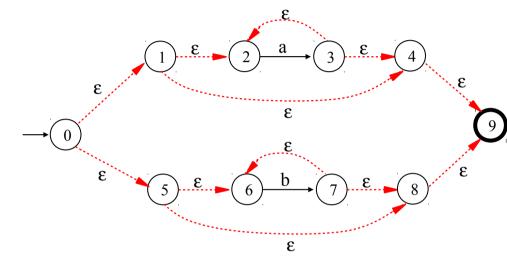
Given the following regular expression:

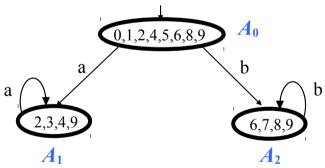
$$r = (\mathbf{a}^* \mid \mathbf{b}^*)$$

we ask to:

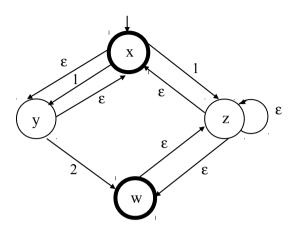
- a) Outline the NFA(r) generated by means of the (canonical) construction of Thompson;
- b) Transform NFA(r) into the equivalent DFA(r);
- c) Based on DFA(r), specifiy the BNF of the regular language relevant to r.



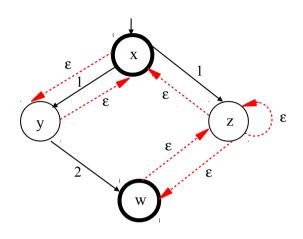


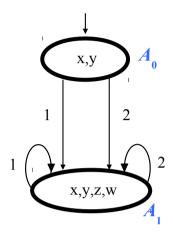


$$A_0 \rightarrow a A_1 \mid b A_2 \mid \epsilon$$
 $A_1 \rightarrow a A_1 \mid \epsilon$ 
 $A_2 \rightarrow b A_2 \mid \epsilon$ 



- a) Generate the equivalent DFA;
- b) Specify the BNF of the regular language of the DFA.





$$A_0 \rightarrow 1 A_1 \mid 2 A_1 \mid \varepsilon$$

$$A_1 \rightarrow 1 A_1 \mid 2 A_1 \mid \varepsilon$$

Given the following regular expression,

$$((a b)^* | (c d)^*)^*$$

#### we ask to:

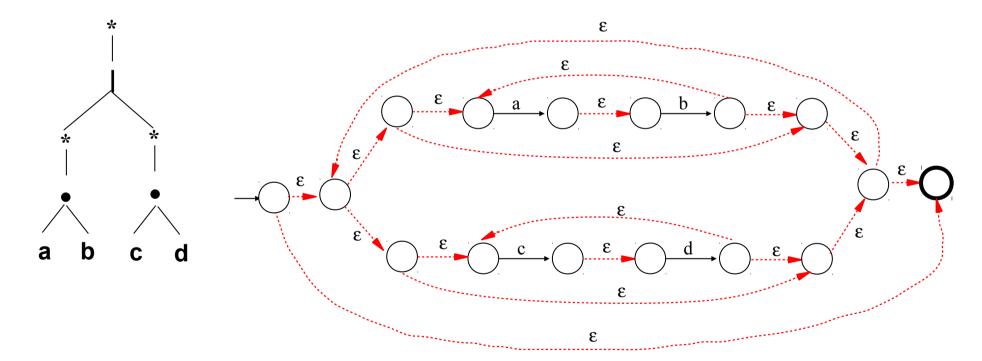
- Outline the tree of the regular expression;
- Outline the DFA generated by the (canonical) construction of Thompson.

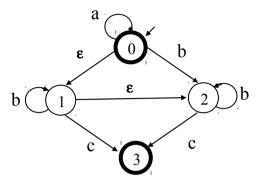
Given the following regular expression,

$$((a b)^* | (c d)^*)^*$$

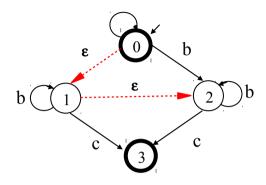
#### we ask to:

- Outline the tree of the regular expression;
- Outline the DFA generated by the (canonical) construction of Thompson.

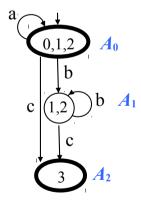




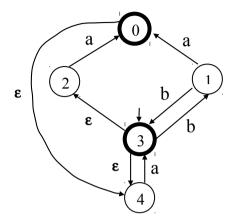
- Generate the equivalent DFA;
- Specify the BNF of the regular language of the DFA.



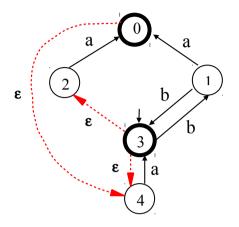
- Generate the equivalent DFA;
- Specify the BNF of the regular language of the DFA.



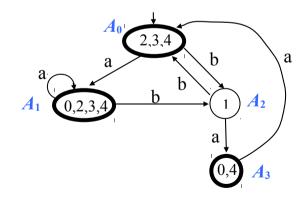
$$A_0 \rightarrow \mathbf{a} \ A_0 \mid \mathbf{b} \ A_1 \mid \mathbf{c} \ A_2 \mid \mathbf{\epsilon}$$
  
 $A_1 \rightarrow \mathbf{b} \ A_1 \mid \mathbf{c} \ A_2$   
 $A_2 \rightarrow \mathbf{\epsilon}$ 



- Generate the equivalent DFA;
- Specify the BNF of the regular language of the DFA.



- Generate the equivalent DFA;
- Specify the BNF of the regular language of the DFA.



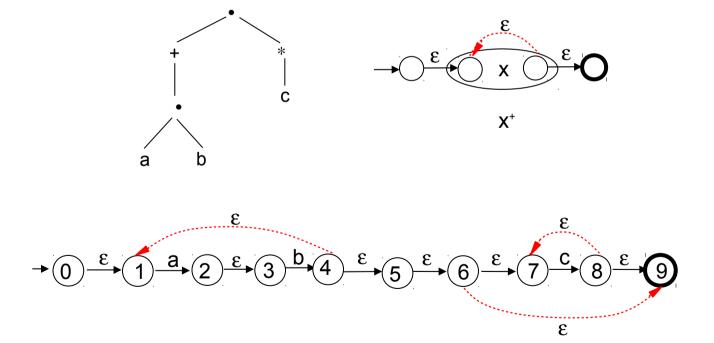
$$A_0 \rightarrow \mathbf{a} \ A_1 \mid \mathbf{b} \ A_2 \mid \mathbf{\epsilon}$$
 $A_1 \rightarrow \mathbf{a} \ A_1 \mid \mathbf{b} \ A_2 \mid \mathbf{\epsilon}$ 
 $A_2 \rightarrow \mathbf{a} \ A_3 \mid \mathbf{b} \ A_0$ 
 $A_3 \rightarrow \mathbf{a} \ A_0 \mid \mathbf{\epsilon}$ 

Given the regular expression  $\mathbf{r} = (\mathbf{a} \mathbf{b})^{+} \mathbf{c}^{*}$ , we ask to:

- Outline the tree of **r**;
- Define the (extended) construction rule of Thompson for the operator + (repetition one or more times);
- Based on the previous points, outline, by means of the construction of Thompson, the NFA which recognizes the regular language of **r**;
- Specify the BNF of the regular language of the NFA.

Given the regular expression  $\mathbf{r} = (\mathbf{a} \mathbf{b})^{\dagger} \mathbf{c}^{*}$ , we ask to:

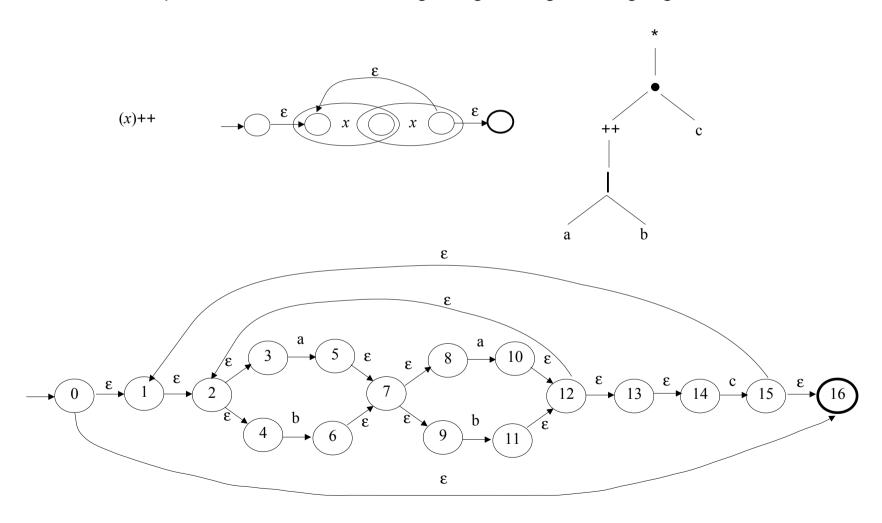
- Outline the tree of **r**;
- Define the (extended) construction rule of Thompson for the operator + (repetition one or more times);
- Based on the previous points, outline, by means of the construction of Thompson, the NFA which recognizes the regular language of **r**;
- Specify the BNF of the regular language of the NFA.

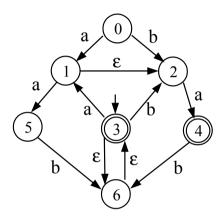


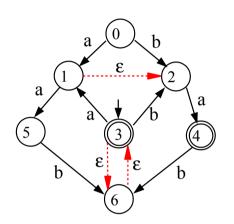
$A_0 \rightarrow A_1$
$A_1 \rightarrow \mathbf{a} A_2$
$A_2 \rightarrow A_3$
$A_3 \rightarrow \mathbf{b} A_4$
$A_4 \rightarrow A_1 \mid A_5$
$A_5 \rightarrow A_6$
$A_6 \rightarrow A_7 \mid A_9$
$A_7 \rightarrow \mathbf{c} A_8$
$A_8 \to A_7 \mid A_9$
$A_9 \rightarrow \varepsilon$

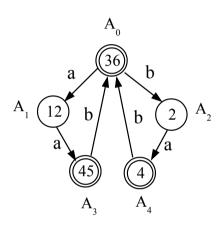
Specify the extended construction rule of Thompson for the operator ++, defined as repetition 2n times,  $n \in [1, 2, ...]$ . Then, draw the tree of the regular expression  $\mathbf{r} = ((\mathbf{a} \mid \mathbf{b})^{++} \mathbf{c})^*$ . Finally, based on the construction of Thompson, outline the NFA recognizing the regular language of  $\mathbf{r}$ .

Specify the extended construction rule of Thompson for the operator ++, defined as repetition 2n times,  $n \in [1, 2, ...]$ . Then, draw the tree of the regular expression  $\mathbf{r} = ((\mathbf{a} \mid \mathbf{b})^{++} \mathbf{c})^*$ . Finally, based on the construction of Thompson, outline the NFA recognizing the regular language of  $\mathbf{r}$ .

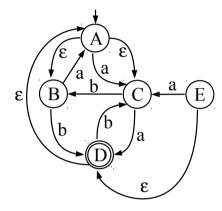


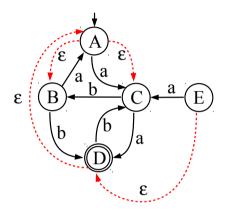


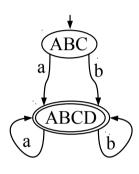




$$A_0 \rightarrow \mathbf{a} A_1 \mid \mathbf{b} A_2 \mid \mathbf{\epsilon}$$
 $A_1 \rightarrow \mathbf{a} A_3$ 
 $A_2 \rightarrow \mathbf{a} A_4$ 
 $A_3 \rightarrow \mathbf{b} A_0 \mid \mathbf{\epsilon}$ 
 $A_4 \rightarrow \mathbf{b} A_0 \mid \mathbf{\epsilon}$ 







$$A_0 \rightarrow \mathbf{a} A_1 \mid \mathbf{b} A_1$$
  
 $A_1 \rightarrow \mathbf{a} A_1 \mid \mathbf{b} A_1 \mid \mathbf{\epsilon}$