Process and Thread Synchronization

Semaphores and Condition Variables

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The purpose of this chapter

- Present different synchronization mechanisms: semaphores and condition variables
- Present some classical synchronization patterns (problems): producer / consumer

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- Present different synchronization mechanisms: semaphores and condition variables
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Bibliography

- A. Tanenbaum, *Modern Operating Systems*, 2nd Edition, 2001, Chapter 2, Processes, p. 100 – 132
- A. Downey, *The Little Book of Semaphores*, 2nd Edition, 2008, p. 1 – 106

Outline

1. Semaphores
2. The "Producer/Consumer" Problem
3. Condition Variables
4. Conclusions

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1 Semaphores

2 The "Producer/Consumer" Problem

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4 Conclusions

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Characteristics

- **generalization of locks**
  - more threads could be given permission
  - to enter simultaneously in the critical regions protected by the semaphore
- consists of
  - **a value**
  - **two primitives**
- main difference from locks
  - any thread is allowed to increase a semaphore value (no of permissions), not only those that previously got one
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Possible Implementation (uni-processor systems)

```
P() // sema_down()
{
    disable_interrupts();
    while (value == 0) {
        insert(crt_thread, w_queue);
        block_current_thread();;
    }
    value = value - 1;
    enable_interrupts();
}
```

```
V()  // sema_up()
{
    disable_interrupts();
    value = value + 1;
    if (! empty(waiting_queue)) {
        th = remove_first(w_queue);
        insert(ready_queue, th);
    }
    enable_interrupts();
}
```

- **interrupts are disabled**, to provide **code atomicity**
  - after critical code is executed, interrupts are enabled back
- "**sleep / wake-up technique**" is used instead of the "*busy waiting*"
  - "**sending to sleep**" (**blocking**) a thread
    - append it to a waiting queue and
    - take the CPU from it
  - "**waking up**" (**unblocking**) a thread
    - remove it from waiting queue
    - append it to the ready queue
    - eventually being given the CPU again

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Usage Example

```c
#include <pthread.h>

// restrict the number of "persons" in a room
// i.e number of threas in their critical region
sem_t sem;

void three_in_a_room();

main()
{
    // initialize the semaphore!
    // do that before having threads working with it!
    sem_init(&sem, 1, 3);

    // create the competing threads
    for (int i=0; i < 100; i++)
        pthread_create(&t[i], NULL, three_in_a_room, NULL);

    // wait for all threads' termination
    for (int i=0; i < 100; i++)
        pthread_join(t[i], NULL);

    // remove the semaphore
    sem_destroy(&sem);
}
```

Usage Example (cont.)

```
void three_in_a_room()
{
    // ask for a permission, i.e try decrementing the semaphore
    // the thread will be blocked, if no permission is available,
    // i.e. semaphore's value is 0
    sem_wait(&sem);

        // stay in room (critical region) for a while

    // release the obteained permission
    // i.e. increment the semaphore's value
    // this operation NEVER blocks the thread
    sem_post(&sem);
}
```

Practice (1)

- Using semaphores, write the (pseudo)code for the functions in the code below such that to not allow more than 22 players enter simultaneously on the football field. A player is represented by a thread executing the *football_player()* function.

```
void football_player()
{
    enter_football_field();
    play_football();
    exit_the_footbal_field();
}
```

Practice (2)

- Using semaphores, write the (pseudo)code that limits to 5 the number of threads executing simultaneously the critical section in the function below.

```c
#define FREE 1
#define BUSY 0
#define MAX_TH 5

int available_unit[MAX_TH] = {FREE, FREE, FREE, FREE, FREE};
int unit[MAX_TH] = {0, 0, 0, 0, 0};

void limited_area()
{
    int pos = -1;

    for (pos = 0; pos < MAX_TH; pos++)
        if (available_unit[pos] == FREE) {
            available_unit[pos] = BUSY;
            break;
        }

    unit[pos]++;

    available_unit[pos] = FREE;
}
```

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Problem Description

- two types of processes
  - *producers*: produce messages
  - *consumers*: consume messages
- use a bounded buffer
- waiting conditions
  - for producer: “buffer is full”
  - for consumer: “buffer is empty”
- wakeup conditions
  - for producer: “there is space in buffer”
  - for consumer: “there are messages in buffer”

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Basic Solution With Race Conditions

```
// Global variables
const N = 100;          // number of slots in buffer
int count = 0;          // number of messages in the buffer
```

```
void producer(int item)
{
    // wait for available space
    while (count == N) {} // full buffer

    insert_item(item);

    // wakeup a consumer
    count = count + 1;
}
```

```
int consumer(int *item)
{
    // wait for available message
    while (count == 0) {}  // empty buffer

    *item = remove_item();

    // wakeup a producer
    count = count - 1;
}
```

- **race conditions**
  - uncontrolled concurrent access to variable count
  - same message could be consumed more times
  - messages could be overwritten
- **busy waiting**
  - suggests us that **synchronization is needed**

Basic Solution With Race Conditions (cont.)

- **should be avoided** in practice
- **replace it** with a sleep / wakeup **synchronization mechanism**

Producers/Consumers Problem's Implementation With Semaphores

```
// Global variables and synchronization mechanims
const N = 100;                  // number of slots in buffer
Semaphore mutex(1);             // provides mutual exclusion to buffer
Semaphore can_produce_msg(N);   // controls producers' access to buffer
Semaphore can_consume_msg(0);   // controls consumers' access to buffer
```

```
void producer(int item)
{
    // check if buffer is full
    can_produce_msg.P();

    // gets mutual exclusion to buffer
    mutex.P();

    insert_item(item);

    // releases the buffer
    mutex.V();

    // new permission for consumers
    can_consume_msg.V();
}
```

```
int consumer(int *item)
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    // check if buffer is empty
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    *item = remove_item();

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    // new permission for producers
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}
```

Remarks On Producers/Consumers Implementation With Semaphores

- there are two ways of using semaphores
  - as a (generalized) lock to provide limited access (mutual exclusion in the example with `mutex` semaphore)
  - as a condition checking and event counter (semaphores `can_consume_msg` and `can_produce_msg`)
- the order of getting the semaphores is important and it can lead to deadlock
  - see the solution with `mutex` semaphore get first

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Deadlock In The Producers/Consumers Problem’s Implementation With Semaphores

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    *item = insert_item();

    // new permission for producers
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}
```

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  - waited condition reflect the shared resource state
  - lock must be taken before testing / changing the condition
- **should not wait (sleep) keeping the lock**
  - could block the other threads that could change condition to true
- $\Rightarrow$ waiting thread should **release the lock while waiting** (sleeping)
- $\Rightarrow$ another thread can take the lock in the meantime and change condition
- that thread **wakes up the sleeping (waiting) thread**

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Condition Variables: Possible Deadlock Scenario

```
// Thread waiting
// for a condition
// to be fulfilled
thread_1()
{
  mutex.lock();

  while (cond == FALSE) {
    // e.g. (no_msg == 0)
    // for a consumer
    sleep(1);
  }

  mutex.unlock();
}
```

```
// Thread changing
// the waited condition
// to TRUE
thread_2()
{
  mutex.lock();

  cond = TRUE;
  // e.g. no_msg++
  // for a producer



  mutex.unlock();
}
```

Condition Variables: Functionality

- what is it?
  - **a specialized waiting mechanism**
  - transparently release the lock while blocking the lock holder
  - at wakeup transparently reacquire the lock before resuming the waiting thread
- consists of
  - a waiting queue
  - two (three) primitives
    - wait: the calling process is inserted in the waiting queue and suspended
    - signal: one sleeping process is removed from the waiting queue and woken up
    - broadcast: all sleeping processes are woken up

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Condition Variables. Basic Pattern Usage

```
// Thread waiting for
// an event (condition)
// to occur
thread_1()
{
  mutex.lock();

  while (! cond)
    c.wait(mutex);

  mutex.unlock();
}
```

```
// Thread generating
// the event
//
thread_2()
{
  mutex.lock();

  cond = TRUE;
  c.signal();

  mutex.unlock();
}
```

Condition Variables: Implementation

```
wait(Lock mutex)
{
  disable_interrupts();

  mutex.unlock();

  insert(crt_th, w_queue);
  block_current_thread();

  mutex.lock();

  enable_interrupts();
}
```

```
signal()
{
  disable_interrupts();


  if (! empty(waiting_queue)) {
    th = remove_first(w_queue);
    insert(ready_queue, th);
  }


  enable_interrupts();
}
```

Condition Variables. General Pattern Usage (Var 1)

```
// Synchronization mechanims: one condition variable
Lock mutex;
Condition c;
```

```
// Thread 1's Function
thread_1()
{
    // critical region entrace
    mutex.lock();
    while (! cond_1)
        c.wait(mutex);
    mutex.unlock();

    // thread T1 in its critical region
    ....

    // critical region exit
    mutex.lock();
    cond_2 = TRUE;
    c.broadcast();
    mutex.unlock();
}
```

```
// Thread 2's Function
thread_2()
{
    // critical region entrace
    mutex.lock();
    while (! cond_2)
        c.wait(mutex);
    mutex.unlock();

    // thread T2 in its critical region
    ...

    // critical region exit
    mutex.lock();
    cond_1 = TRUE;
    c.broadcast();
    mutex.unlock();
}
```

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Condition Variables. General Pattern Usage (Var 2)

```
// Synchronization mechanims: two condition variables
Lock mutex;
Condition c1, c2;
```

```
// Thread 1's Function
thread_1()
{
    // critical region entrace
    mutex.lock();
    while (! cond_1)
        c1.wait(mutex);
    mutex.unlock();

    // thread T1 in its critical region
    ....

    // critical region exit
    mutex.lock();
    cond_2 = TRUE;
    c2.signal(); // c2.broadcast();
    mutex.unlock();
}
```

```
// Thread 2's Function
thread_2()
{
    // critical region entrace
    mutex.lock();
    while (! cond_2)
        c2.wait(mutex);
    mutex.unlock();

    // thread T2 in its critical region
    ...

    // critical region exit
    mutex.lock();
    cond_1 = TRUE;
    c1.signal();  // c1.broadcast();
    mutex.unlock();
}
```

Condition Variables. Multiple Conditions (Var 1)

```
// Synchronization mechanims: multiple conditions & one condition variable
Lock mutex;
Condition c;
```

```
// Thread 1's Function
thread_1()
{
    // critical region entrace
    mutex.lock();
    while (! cond_1 || ! cond_2)
        c.wait(mutex);
    mutex.unlock();

    // T1's critical region
    ....

    // critical region exit
    mutex.lock();
    .....
    mutex.unlock();
}
```

```
// Thread 2's Function
thread_2()
{
    ....
    ....
    // critical region exit
    mutex.lock();
    switch (x) {
      case 1:
        cond_1 = TRUE;
        c.broadcast(); // c.signal() ?
        break;
      case 2:
        cond_2 = TRUE;
        c.broadcast(); // c.signal() ?
        break;
    }
    mutex.unlock();
}
```

Condition Variables. Multiple Conditions (Var 2 - BAD)

```
// Synchronization mechanims: multiple conditions & multiple condition variables & one lock
Lock mutex;
Condition c1, c2;
```

```
// Thread 1's Function
thread_1()
{
    // critical region entrace
    mutex.lock();
    while (! cond_1)
        c1.wait(mutex);

    while (! cond_2)
        c2.wait(mutex);
    mutex.unlock();

    // T1's critical region
    ....

    // critical region exit
    mutex.lock();
    .....
    mutex.unlock();
}
```

```
// Thread 2's Function
thread_2()
{
    ....
    ....
    // critical region exit
    mutex.lock();
    switch (x) {
      case 1:
        cond_1 = TRUE;
        c1.signal(); // c1.broadcast();
        break;
      case 2:
        cond_2 = TRUE;
        c2.signal(); // c2.broadcast();
        break;
    }
    mutex.unlock();
}
```

Condition Variables. Multiple Conditions (Var 2 - OK)

```
// Synchronization mechanims: multiple conditions & multiple condition variables & multiple
Lock mutex1, mutex2;
Condition c1, c2;
```

```
// Thread 1's Function
thread_1()
{
    // critical region entrace
    mutex1.lock();
    while (! cond_1)
        c1.wait(mutex1);

    mutex2.lock();
    while (! cond_2)
        c2.wait(mutex2);
    mutex2.unlock();
    mutex1.unlock();

    // T1's critical region
    ....

    // critical region exit
    .....
}
```

```
// Thread 2's Function
thread_2()
{
    ....
    // critical region exit

    switch (x) {
    case 1:
        mutex1.lock();
        cond_1 = TRUE;
        c1.signal(); // c1.broadcast();
        mutex1.unlock();
        break;
    case 2:
        mutex2.lock();
        cond_2 = TRUE;
        c2.signal(); // c2.broadcast();
        mutex2.unlock();
        break;
    }
}
```

Condition Variables: Comparison With Semaphores

- *wait* always suspends the calling process
  - while *P* does this only in case the semaphore's value is zero
- a `signal` can be lost (i.e. not seen)
  - while a *V* is not, because it increments the semaphore's value
- safe to use inside a mutual exclusion area
  - while semaphore could lead to deadlock

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Rules to Remember About Using Condition Variables

1. **always use them** (`wait` and `signal`) **inside a mutual exclusion area** (protected by a **lock**)
2. the lock does (normally) not protect the shared resource
   - is would simply provide mutual exclusion, but maybe not needed
   - does not support deadlock-free waiting
3. the lock protects the entrance in and exit from critical region
   - where some condition checks or changes are done
   - in a mutual exclusion manner
4. it **could be needed to recheck the condition** after returning from `wait`
   - use *while* instead of *if*
5. **do not use `wait` without checking a condition**
   - it is not sure the thread must wait (we do not control thread scheduling)
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2. **ears** for hearing
3. **mind** for thinking (understanding, learning, remembering)
4. see `https://www.youtube.com/watch?v=RZxAc3Grkck`
5. review previous slide and make use of the "elements" mentioned above

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Producers/Consumers Problem's Implementation With Locks and Condition Variables

```
// Global variables and synchronization mechanisms
const N = 100;           // number of slots in buffer
int count = 0;           // number of messages in buffer
Lock mutex;              // provides mutual exclusion to buffer
Condition producers;     // controls producers' access to buffer
Condition consumers;     // controls consumers' access to buffer
```

```
void producer(int item)
{
    // gets mutual exclusion to buffer
    mutex.lock();

    // check if buffer is full
    while (count == N)
      producers.wait(mutex);

    insert_item(item);
    count = count + 1;

    // wakes up a consumer
    consumers.signal();

    // releases the buffer
    mutex.unlock();
}
```

```
int consumer(int *item)
{
    // gets mutual exclusion to buffer
    mutex.lock();

    // check if buffer is empty
    while (count == 0)
      consumers.wait(mutex);

    *item = remove_item();
    count = count - 1;

    // wakes up a producer
    producers.signal();

    // releases the buffer
    mutex.unlock();
}
```

Practice (3)

- Using locks and condition variables, write the (pseudo)code for the functions in the code below such that to not allow more than 22 players enter simultaneously on the football field. A player is represented by a thread executing the *football_player()* function.

```
void football_player()
{
    enter_football_field();
    play_football();
    exit_the_footbal_field();
}
```

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Outline

1 Semaphores

2 The "Producer/Consumer" Problem

3 Condition Variables

4 Conclusions

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What we talked about

- semaphores
  - a generalization of locks ⇒ could allow more threads pass the barrier
  - important to initialize the semaphore
  - "`P()`" and "`V()`" primitives
  - could also be used as an event counter, usually initialized with 0
- condition variables
  - `wait`: provide a specialized way to wait for a condition to be fulfilled
  - `signal`: provide a way to signal that a condition is fulfilled
- producer-consumer synchronization pattern

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