

Lecture 7

Repeated Games

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Outline of Lecture: I

- Description and analysis of finitely repeated games.
- Example of a finitely repeated game with a unique equilibrium
- A general theorem on finitely repeated games.

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Outline of Lecture: II

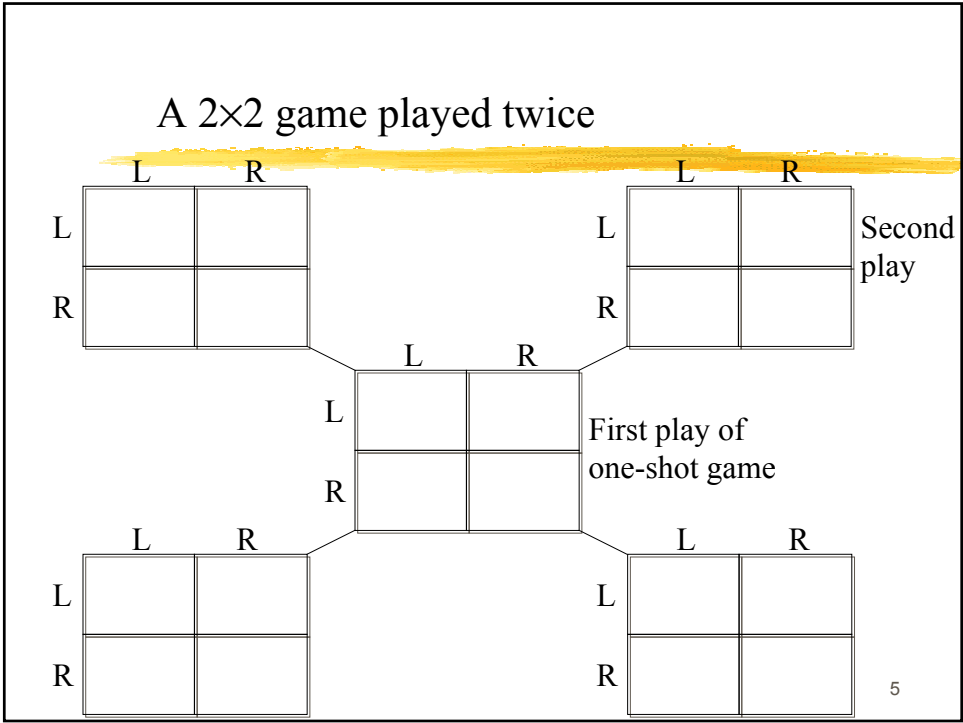
- A formula for computing discounted payoffs in repeated games.
- A description of infinitely repeated games.
- Examples of strategies in infinitely repeated games.
- How to support “better” equilibria in infinitely repeated games.
- Application to pricing games.

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Strategies and payoffs for games played twice

- Finitely repeated games
- Discounted utility
- Complete plans of play for 2×2 games played twice
- Trigger strategies

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Strategies for playing a 2×2 game twice

Strategy No.	Round 1	After (L, L)	After (L, R)	After (R, L)	After (R, R)
1	L	L	L	L	L
2	L	L	L	L	R
3	L	L	L	R	L
4	L	L	L	R	R
...
...
17	R	L	L	L	L
18	R	L	L	L	R
...
...
31	R	R	R	R	L
32	R	R	R	R	R ⁶

- A repeated game is just an extensive form game.
- Selten's theorem on unique subgame perfect equilibria
- Repetition by itself does not solve a credibility problem

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The diagram illustrates a game structure with four identical 2x2 payoff matrices arranged in a square. Each matrix is labeled "Confess" above it. The matrices are connected by lines forming a central square. Each matrix has the following payoffs:

10,10	1,25
25,1	3,3

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Prisoner's Dilemma, backward induction (second play): Pay-off matrix

		Player 2	
		Confess	Deny
Player 1	Confess	10,10	1,25
	Deny	25,1	3,3

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Prisoner's Dilemma, backward induction (second play): Player 1's strategy

		Player 2	
		Confess	Deny
Player 1	Confess	<u>10</u> ,10	<u>1</u> ,25
	Deny	25,1	3,3

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Prisoner's Dilemma, backward induction (second play): Player 2's strategy

A 2x2 payoff matrix for the Prisoner's Dilemma. The rows represent Player 1's strategies (Confess, Deny) and the columns represent Player 2's strategies (Confess, Deny). The payoffs are (Player 1, Player 2). Arrows indicate Player 2's optimal choice for each of Player 1's strategies: an arrow points left from the top row (Confess) to the 'Confess' column, and an arrow points left from the bottom row (Deny) to the 'Confess' column.

	Player 2 Confess	Player 2 Deny
Player 1 Confess	10, <u>10</u>	1, 25
Player 1 Deny	25, <u>1</u>	3, 3

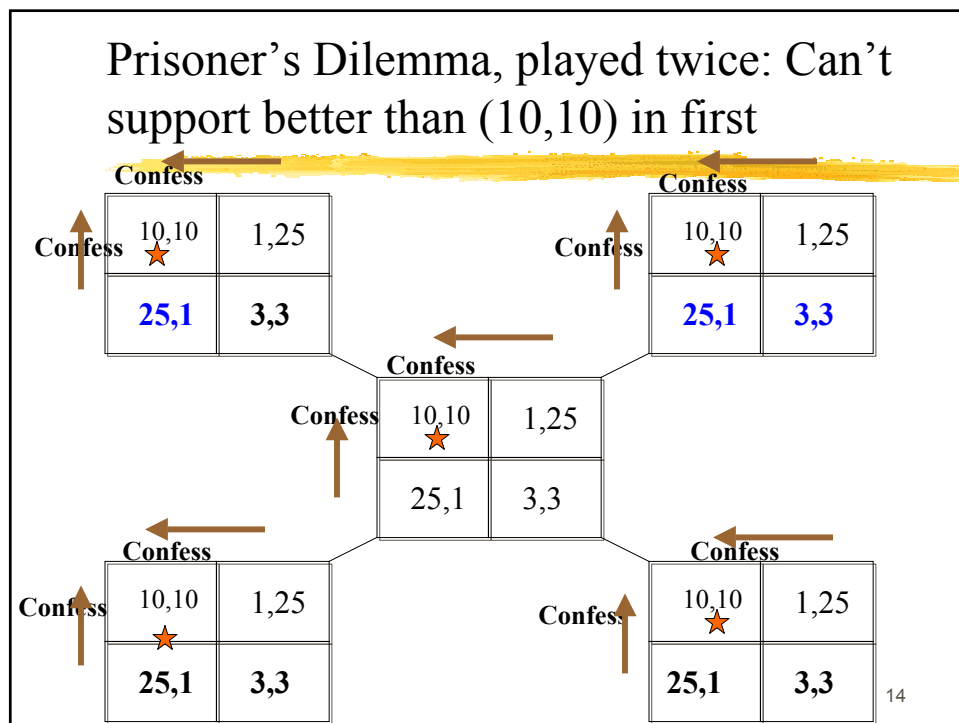
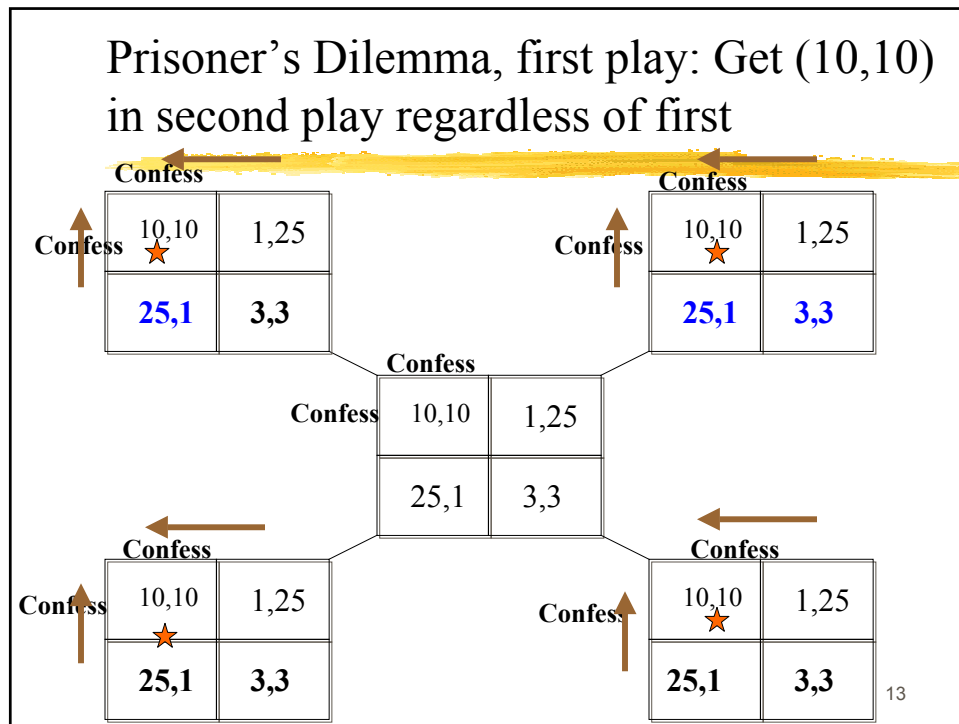
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Prisoner's Dilemma, second play, by backward induction

A 2x2 payoff matrix for the Prisoner's Dilemma, identical to the one above. Arrows indicate the backward induction path: an arrow points left from the top row to the 'Confess' column, an arrow points left from the bottom row to the 'Confess' column, an arrow points up from the 'Confess' column to the top row, and an arrow points up from the 'Deny' column to the bottom row. The top-left cell (10, 10) is marked with a red star, indicating the final outcome of the game.

	Player 2 Confess	Player 2 Deny
Player 1 Confess	10, 10 ★	1, 25
Player 1 Deny	25, 1	3, 3

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Selten's Theorem

If a game with a unique equilibrium is played finitely many times, its solution is that equilibrium played each and every time

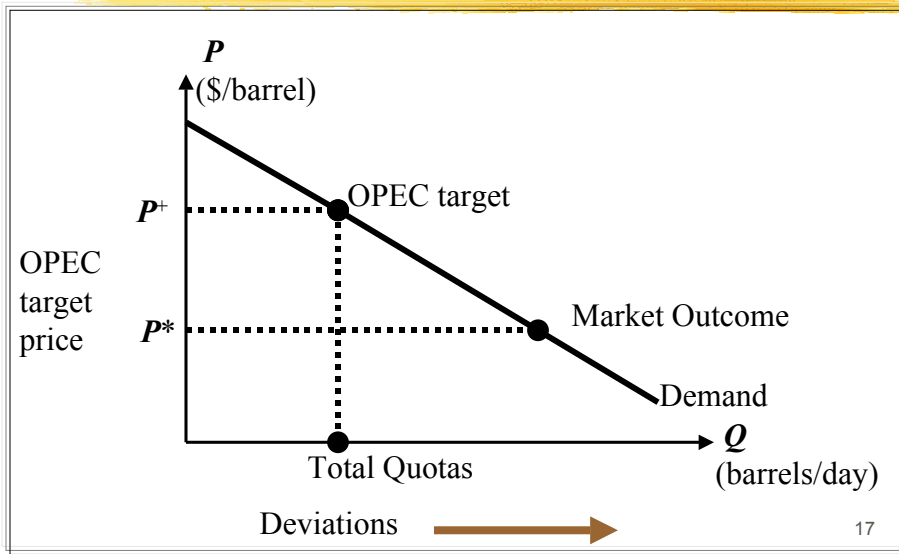
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OPEC drops Quotas

- OPEC's quota system, 1973-1993
- The attempt to improve upon a one-shot Cournot equilibrium
- Finiteness of a resource and finiteness of a game

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OPEC quotas



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Achieving Higher Payoffs in Infinitely Repeated Games.

- Selten's theorem suggests that if the single period game has a unique equilibrium, then the repetition of that equilibrium is all that will occur in repeated games.
- This conclusion is only true for games that are repeated "finitely".

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Infinitely repeated games: strategies and payoffs

- The “as if” interpretation of infinite repetition.
- Complete plans for infinite play.
- Discounting infinite series of payoffs.

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Evaluating payoffs for an infinitely repeated game

Total payoff for player 1, $u_1 = \sum d^t u_1(t)$
 t goes from 0 to ∞ ; d is the discount factor where $d < 1$

Note, if u is profits, $d = 1/(1+r)$ where r is the interest rate or the internal rate of return for the firm.

RULE: When $u_1(t) = 1$ for all t , $u_1 = 1 + d + d^2 + d^3 + \dots$

For $0 < d < 1$, the series sums to

$$u_1 = 1/(1 - d) = 1 + d + d^2 + d^3 + \dots$$

When $u_1(t) = k$ for all t , $u_1 = k/(1 - d)$

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Evaluating payoffs for an infinitely repeated game: Extension

An important modification: Suppose that in every period, there is a probability, $1-p$, that the game ends.

Now, total payoff for player 1, $u_1 = \sum (dp)^t u_1(t)$
 t goes from 0 to ∞ ; d is still the discount factor where $d < 1$ and $p < 1$ so $dp < 1$ as well.

If we call R the “effective rate of return”, then

$$1/(1+R) = dp, \text{ or } R = (1-dp)/dp.$$

See Dixit and Skeath, Chapter 8 and Appendix to Chapter 8.

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Different Types of Strategies for repeated games.

- The repeated “One-shot” strategy:
 - if a profile of strategies form an equilibrium in the one-shot game, then any repetition of these strategies form an equilibrium in the repeated game.
- Thus, in an infinitely repeated Prisoner’s Dilemma, Confess (or Do not Collude) in every period is an equilibrium of the repeated game.

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Different Types of Strategies.

- The "Grim Trigger Strategy":
 - in the prisoner's dilemma game, a promise to Deny forever as long as the rival Denies supported by a threat to Confess forever if the rival Confesses can sometimes be an equilibrium.

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Different Types of Strategies.

- The "Tit for Tat Strategy":
 - in the prisoner's dilemma game, a promise to Deny as long as the rival Denies supported by a threat to Confess for a period if the rival Confesses and then revert to Deny (as long as the deviant rival Denied in the punishment period) can sometimes be an equilibrium.
- Dixit and Skeath focus on this strategy, I look at the Grim Trigger Strategy.

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Duopolists Play the Prisoner's Dilemma

- Consider the game where two firms choose high (cooperative) prices or low (deviant) prices in each period.
- If both choose High, they get profits of 3 each. If one chooses High and the other low, the High gets 0 and the low gets 4.
- If both choose Low, they get 2 each.
- They play this game infinitely and $dp=.25$.

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Prisoner's Dilemma: Duopoly Version

		Firm 2	
		Cooperate	Deviate
Firm 1	Cooperate	3, 3	0, <u>4</u>
	Deviate	4, 0	2, <u>2</u>

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Duopolists Play the Prisoner's Dilemma

- We know that Deviate (Low) forever is an equilibrium. This pays each firm
 - $\Pi_1 = \sum (dp)^t u_1(t)$
 - $= \sum (.25)^t 2 = 2 * (1/(1-.25)) = 8/3$

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Duopolists Play the Prisoner's Dilemma

- If they play the grim trigger strategy,
 - $\Pi_1 = \sum (.25)^t 3 = 3 * (1/(1-.25)) = 4.$
- But won't one Firm deviate?
- If I deviate, I get 4 right away, and then I go to the deviate equilibrium from next period onward which gives me 8/3.
- This strategy would give me
 - $4 + (.25) * (8/3) = 4 + 2/3$

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Duopolists Play the Prisoner's Dilemma

- Since this is bigger than what I get from Cooperate forever, I should deviate.
- But what if $dp = .75$?
- Deviate forever gives me $2 * (1 / (1 - .75)) = 8$
- Cooperate forever gives me $3 * (1 / (1 - .75)) = 12$
- Now if I deviate, I get 4 plus a payoff of 8 from then on, but $4 + .75 * 8 = 10 < 12$.
- I do not want to "Trigger" a price war!

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Infinitely repeated Cournot market games

- A one-shot Cournot equilibrium, repeated infinitely often, is a subgame perfect equilibrium path
- Better paying equilibria than one-shot Cournot
- Monopoly-like equilibria when firms attach enough importance to the future
- A Folk Theorem for infinitely repeated games
- Infinitely repeated Bertrand market games

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When Can Collusion Occur in Infinitely Repeated Games?

- Consider the General Version of the Prisoner's Dilemma Game.

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Prisoner's Dilemma: Duopoly Version

		Firm 2	
		Cooperate	Deviate
Firm 1	Cooperate	C,C	H,L
	Deviate	L,H	D,D

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Collusion In Prisoner's Dilemma Games

- Deviate forever gives $(1/(1-dp))D$.
- Cooperate forever gives $(1/(1-dp))C$.
- But a single period deviation from Cooperate forever gives $H+(dp/(1-dp))D$.
- We need to make sure that a firm does not want to deviate, or that $H+(dp/(1-dp))D < (1/(1-dp))C$.

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Collusion In Prisoner's Dilemma Games

- $H+(dp/(1-dp))D < (1/(1-dp))C$.
- If and only if $(1-dp)*H+dp*D < C$. **
- Presumably $H > C$ and $D < C$ (or else this is not an interesting problem).
- Therefore, (**) is false for $dp=0$ and true for $dp=1$.

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Collusion In Prisoner's Dilemma Games: Conclusions

- If the future matters a lot, (δ is close to one), then collusion is easier to support.
- If interest rate is very high (δ close to zero) or the probability of end game very high (p close to zero), collusion is hard.
- If D is high, collusion is hard to attain.
- If H is high, collusion is hard to attain.
- If C is low, collusion is hard to attain.

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Collusion In Prisoner's Dilemma Games: Conclusions

- Final Remarks: Observe why this suggests that it is easier to have collusion in Bertrand pricing games than in Cournot quantity games.
- In Bertrand pricing games, $D=0$. In Cournot games, $D>0$.
- Ironical conclusion: Collusion is easier to support in repeated games which have MORE competitive stage games!

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