



# Algorithmic Game Theory

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1. Formal Definitions/Notations
2. Strict versus Weak Dominance
3. Iterative Deletion of Dominated Strategy
4. Median Voter Theorem
5. Model Simplification for Engineering  
Application: Examples

# Notations

	Notation	Pick a Number Game
<b>Players</b>	$i, j, \dots$	You all
<b>Strategy</b>	$s_i$ : a particular strategy of player $i$  $s_{-i}$ : the strategy of everybody else except player $i$	$S_4=12, s_8=22$
<b>Strategy Set</b>	$S_i$ : the set of possible strategies of player $i$	$\{1, 2, \dots, 100\}$
<b>Strategy Profile</b>	$s$ : a particular play of the game “strategy profile” (vector, or list)	The collection of your pieces of paper
<b>Payoffs</b>	$u_i(s_1, \dots, s_i, \dots, s_N) = u_i(s)$	$u_i(s) = \begin{cases} \$10 - .01 * \Delta & \text{if you win} \\ 0 & \text{otherwise} \end{cases}$

# Assumptions

- We assume all the ingredients of the game to be known
  - Everybody knows the possible strategies everyone else could choose
  - Everybody knows everyone else's payoffs

## Complete Information Game

- This is not very realistic, but we start from this class of games

# Classification of games

Non-cooperative	Cooperative
Static	Dynamic (repeated)
Strategic-form	Extensive-form
Perfect information	Imperfect information
Complete information	Incomplete information

Perfect information: each player can observe the action of each other player.

Complete information: each player knows the identity of other players and, for each of them, the payoff resulting of each strategy.

# Example

		Player 2		
		L	C	R
Player 1	T	5, -1	11, 3	0,0
	B	6, 4	0, 2	2,0

Players	1, 2	
Strategy sets	$S_1 = \{T, B\}$	$S_2 = \{L, C, R\}$
Payoffs	$U_1(T, C) = 11$	$U_2(T, C) = 3$

NOTE: This game is not symmetric

# Game Analysis

- How is the game going to be played?
- Does player 1 have a dominated strategy?
- Does player 2 have a dominated strategy?
- For a strategy to be dominated, we need another strategy for the same player that does always better (in terms of payoffs)

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## Definition: Strict dominance

We say player  $i$ 's strategy  $s_i'$  is strictly dominated by player  $i$ 's strategy  $s_i$  if:

$$u_i(s_i, s_{-i}) > u_i(s_i', s_{-i}) \text{ for all } s_{-i}$$

No matter what other people do, by choosing  $s_i$  instead of  $s_i'$ , player  $i$  will always obtain a higher payoff.

# “Hannibal” game

- An invader is thinking about invading a country, and there are **2 ways through** which he can lead his army.
- You are the defender of this country and you have to decide **which of these ways you choose to defend**: you can only defend one of these routes.
- **One route is a hard pass**: if the invader chooses this route he will **lose one battalion** of his army (over the mountains).
- If the invader **meets your army**, whatever route he chooses, he will **lose a battalion**



# “Hannibal” game

		Attacker	
		e	h
Defender	E	1, 1	1, 1
	H	0, 2	2, 0

## Strategies

1. e, E = Easy Path ;
2. h, H = Hard Path

## Payoffs:

1. **Attacker:** Number of battalions in your country
2. **Defender:** Number of attacker's lost battalions

# “Hannibal” game

- You’re the defender: what would you do?
- Is it true that defending the easy route dominates defending the hard one?
- You’re the attacker: what would you do?
- Now, what the defender should do, if he would put himself in the attacker shoes?

## Definition: Weak dominance

We say player  $i$ 's strategy  $s_i'$  is weakly dominated by player  $i$ 's strategy  $s_i$  if:

$$\begin{aligned} u_i(s_i, s_{-i}) &\geq u_i(s_i', s_{-i}) \text{ for all } s_{-i} \\ u_i(s_i, s_{-i}) &> u_i(s_i', s_{-i}) \text{ for some } s_{-i} \end{aligned}$$

No matter what other people do, by choosing  $s_i$  instead of  $s_i'$ , player  $i$  will always do **at least as well**, and in some cases she does strictly better.

**It turns out that, historically, Hannibal chose H!**

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# Pick a Number Game

*Without showing your neighbor what you're doing, write down an integer number between 1 and 100. I will calculate the average number chosen in the class. The winner in this game is the person whose number is closest to two-thirds ( $\frac{2}{3}$ ) of the average in the class. The winner will win **10** \$ minus the difference in cents between her choice and that two-thirds of the average.*

Example: 3 students

Numbers: 25, 5, 60

Total: 90, Average: 30,  $\frac{2}{3}$ \*average: 20

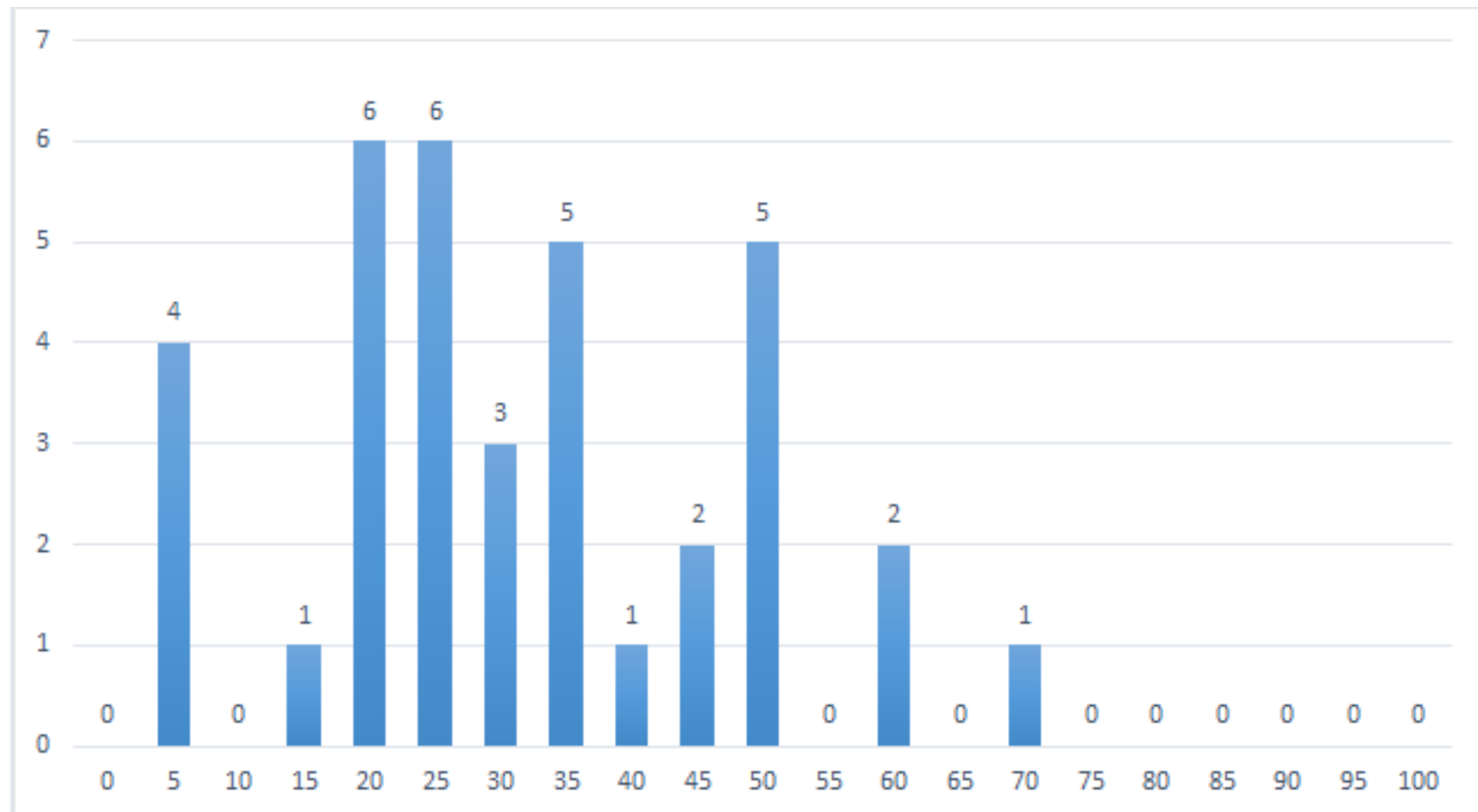
25 wins:  $10 \$ - .01 * 5 = 9.95 \$$

# What did you do?

- What we know:
  - Do not choose a strictly dominated strategy
  - Also, do not choose a weakly dominated strategy
  - You should put yourself in others' shoes, try to figure out what they are going to play, and respond appropriately



# How did you play?



# First Clue!

- A possible assumption:
  - People chose numbers uniformly at random
  - The average is 50
  - $\frac{2}{3} * \text{average} = 33.3$
- What's wrong with this reasoning?

**IUT students are not rand()!**

# Dominated Strategy?

- Let's try to find out whether there are dominated strategies
- If everyone would chose 100, then the winning number would be 67
- ➔ Numbers **bigger than 67** are weakly dominated by 67
- ➔ Rationality tells not to choose numbers **bigger than 67**

# New Game!

- So now we've eliminated dominated strategies, it's like a new game played over the set  $[1, \dots, 67]$
- Once you figured out that nobody is going to choose a number above 67, the conclusion is

Also strategies above 45 are ruled out

- This means:
  1. Rationality
  2. Knowledge that others are rational as well
- Note:

They are weakly dominated, only once we delete 68-100

# Iterative Deletion

- Eventually, we can show that also strategies above 30 are weakly dominated, once we delete previously dominated strategies
- We can go on with this line of reasoning and end up with the conclusion that:
- 1 was the winning strategy!

# Common Knowledge

- **Common knowledge:** you know that others know that others know ... and so on that rationality is underlying all players' choices

# Theory vs. Practice

- **Q:** Why was it that 1 wasn't the winning answer?
- **A:** We need a strong assumption, that is that all players are rational and they know that everybody else's rational as well

# Common Knowledge

Rationality

Rationality and Knowledge of Other's Rationality

Rationality, Knowledge of Other's Rationality, and Knowledge of Knowledge of Rationality (know that you know that I know ....)

Rationality, Knowledge of Other's Rationality, Knowledge of Knowledge of Rationality, and Knowledge of knowledge of knowledge of Rationality

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# Our Game Results: 2015

—Average number was: 29.83

(2015: 24.76, 2014: 35.7, 2013: 35.78, 2012: 26.14)

—Winning number was:  $\frac{2}{3} \times \text{Average}$   
= 19.89

(2015: 16.51, 2014: 23.80, 2013: 23.85, 2012: 17.43)

# Summary

- We've explored a bit the idea of deleting dominated strategies
  - Look at a game
  - Figure out which strategies are dominated
  - Delete them
  - Look at the game again
  - Look at which strategies are dominated now
  - ... and so on ...

# Summary

- **Iterative deletion of dominated strategies** seems a powerful idea, but it's also dangerous if you take it literally
- In some games, iterative deletion converges to a single choice, in others it may not (we'll see shortly an example)
- **HINT**: try to identify all dominated strategies at once before you delete, this may save you some rounds...

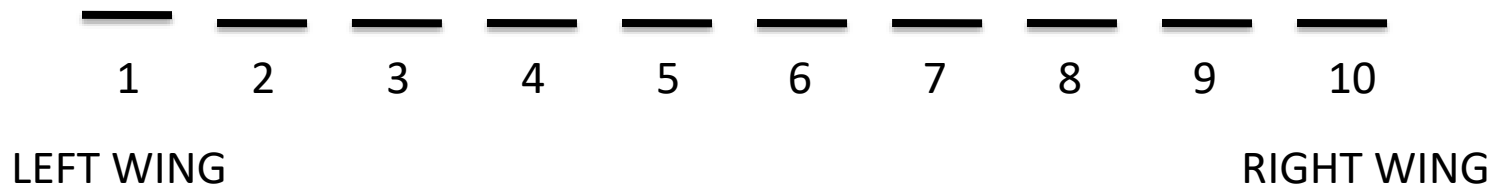
**Let us play  
the same Game, again!**

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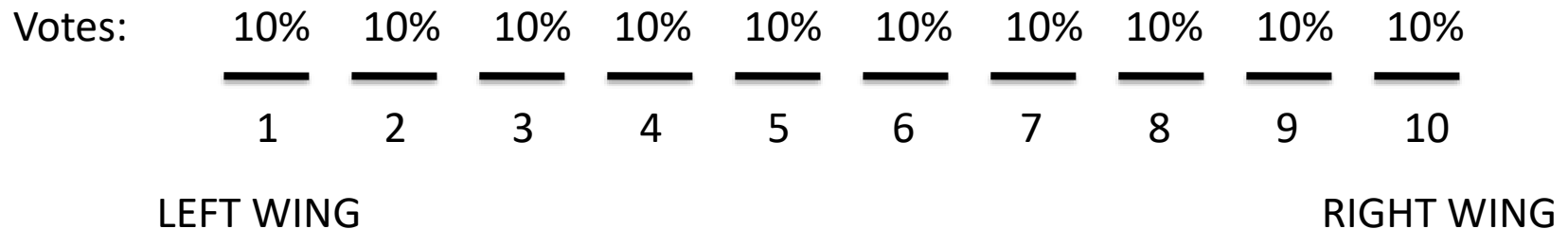
# Election Game Model

- **2 candidates**
- Choosing their political positions on a **spectrum**
- Assume the spectrum has 10 positions



# Election Game Model

- There are 10% of the voters at each of these positions:
  - Voters are uniformly distributed
- Voters will eventually vote for the closest candidate (i.e., for the candidate whose position is closest to their own)
- We break ties by splitting votes equally



# Election Game Strategies

- We assume payoffs follow the idea that the candidates aim to **maximize** their share of vote (**Win the Election**)
- Are there any dominated strategies here?



# Election Game

## Analysis of Dominated Strategy

- Is position 1 dominated? If so, what dominates it?  
Let's test, e.g. how is 1 vs. 2

$s_{-i}$	1's Payoff for $s_i=1$		1's Payoff for $s_i=2$
Vs. 1	$u_1(1,1) = 50 \%$	$<$	$u_1(2,1) = 90\%$
Vs. 2	$u_1(1,2) = 10 \%$	$<$	$u_1(2,2) = 50\%$
Vs. 3	$u_1(1,3) = 15 \%$	$<$	$u_1(2,3) = 20\%$
Vs. 4	$u_1(1,4) = 20 \%$	$<$	$u_1(2,4) = 25\%$
...	...	...	....

- Do you see a pattern coming up here?  
→ We conclude that 2 **strictly dominates** 1
- We're **not** saying that 2 wins over 1

# Election Game

## Analysis of Dominated Strategy

- Using a similar argument, we have that:

**→ 9 strictly dominates 10**

- Is there anything else dominated here?
- What about 2 being dominated by 3?

Vs. 1	$U_1(2,1) = 90\%$	$>$	$U_1(3,1) = 85\%$
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# Election Game

## Analysis of Dominated Strategy

- Even though 2 is not a dominated strategy, if we do the process of iterative deletion and delete dominated strategies (1 and 9)...
- Would 3 dominate 2?

Vs. 2	$u_1(2,2) = 50 \%$	$<$	$u_1(3,2) = 80\%$
Vs. 3	$u_1(2,3) = 20 \%$	$<$	$u_1(3,3) = 50\%$
Vs. 4	$u_1(2,4) = 25 \%$	$<$	$u_1(3,4) = 30\%$
Vs. 5	$U_1(2,5) = 30 \%$	$<$	$u_1(3,5) = 35\%$
...	...	...	....

# Election Game

## Analysis of Dominated Strategy

- Strategies 2 and 8 are **not** dominated
  - ➔ They are dominated once we realize that strategies 1 and 10 won't be chosen
- If we continued the exercise, where would we get?

# Election Game Result

- It turns out that 5 and 6 are not dominated
  - What's the prediction that game theory suggests here?
- Candidates will be squeezed towards the center, they're going to choose positions very close to each other

In political science this is called the  
**Median Voter Theorem**

# Election Game: Similar Examples

- The same model has applications in economics as well (and computer science): **product placement**
- **Example: Placing Gas Stations (abroad) or banks (here!)**
  - Spread themselves evenly out over the town
  - On every road
- As we all know, this doesn't happen:  
All gas stations tend to crowd into the same corners, all the fast foods crowd as well, ... you name it

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# Model Simplification

- We have been using a model of a real-world situation, and tried to predict the outcome using game theory
- What is missing? Is there anything wrong with the model?



# Simplification in Median Voter

- Voters are not evenly distributed
- Some people do not vote
- There may be more than 2 candidates
- There may be higher “dimensions” to the problem

# Model Simplification

- So if we're missing so many things, our model is useless, and in general modeling (as an abstraction effort) is useless!!
- No: first, analyze a problem with simplifying assumptions, then relax them and see what happens
- E.g.: would a different voters distribution change the result?

# **Model Simplification: Engineering Approach**

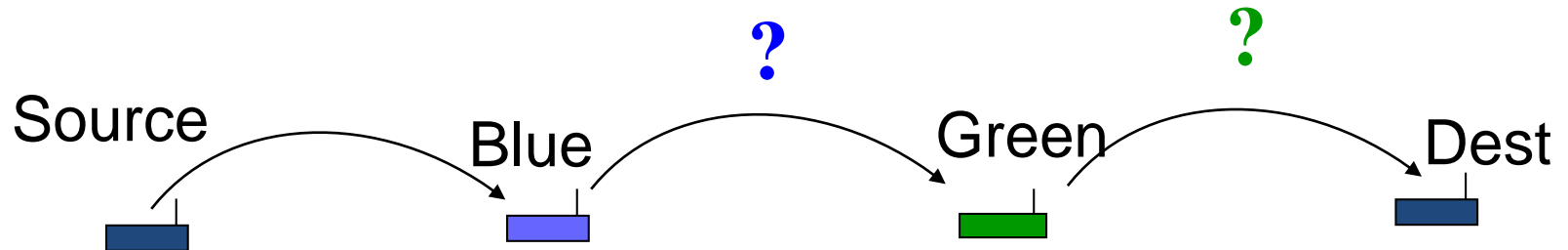
- We basically make lots of abstractions in make game theoretical models for our engineering problems
- Not a bad idea to start with abstraction, but you must be careful about what you design

# George Edward Pelham Box (1919-2013)

- Remember that all models are wrong; the practical question is how wrong do they have to be to not be useful.
- **Essentially, all models are wrong, but some are useful.**



# The Joint Packet Forwarding Game



- Reward for packet reaching the destination: 1
- Cost of packet forwarding:  $c$  ( $0 < c \ll 1$ )

		Green	
		Forward	Drop
Blue	Forward	$(1-c, 1-c)$	$(-c, 0)$
	Drop	$(0, 0)$	$(0, 0)$

**No strictly dominated strategies !**

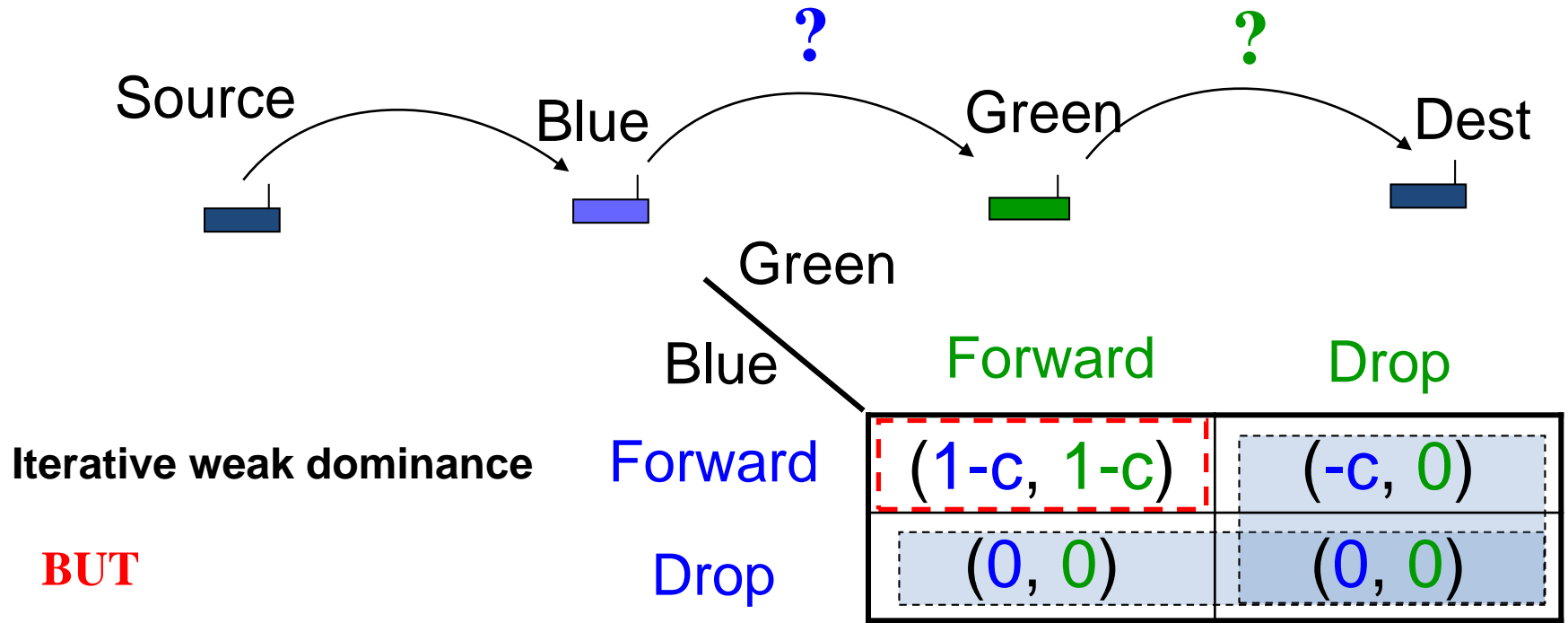
# Weak dominance

**Weak dominance:** strictly better strategy for at least one opponent strategy

Strategy  $s_i$  weakly dominates strategy  $s'_i$  if:

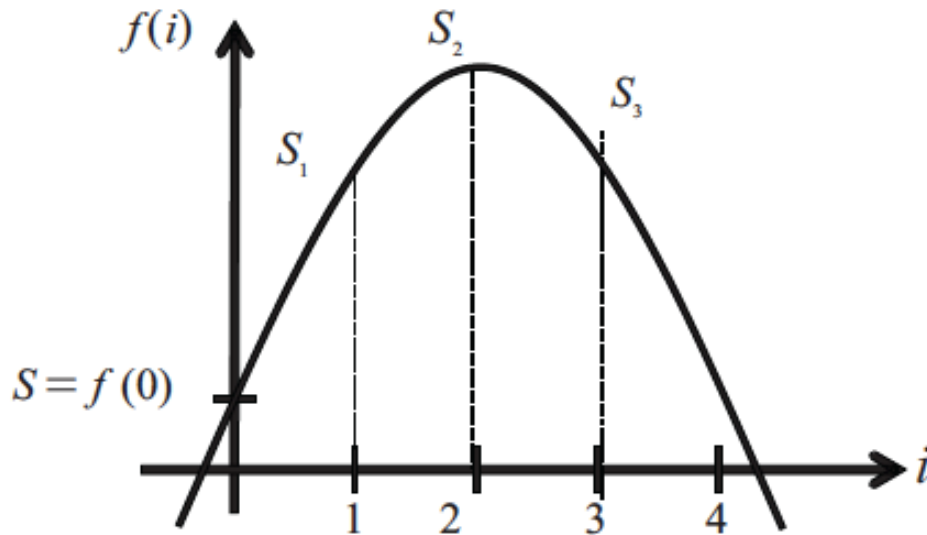
$$u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i}), \forall s_{-i} \in S_{-i}$$

with strict inequality for at least one  $s_{-i}$

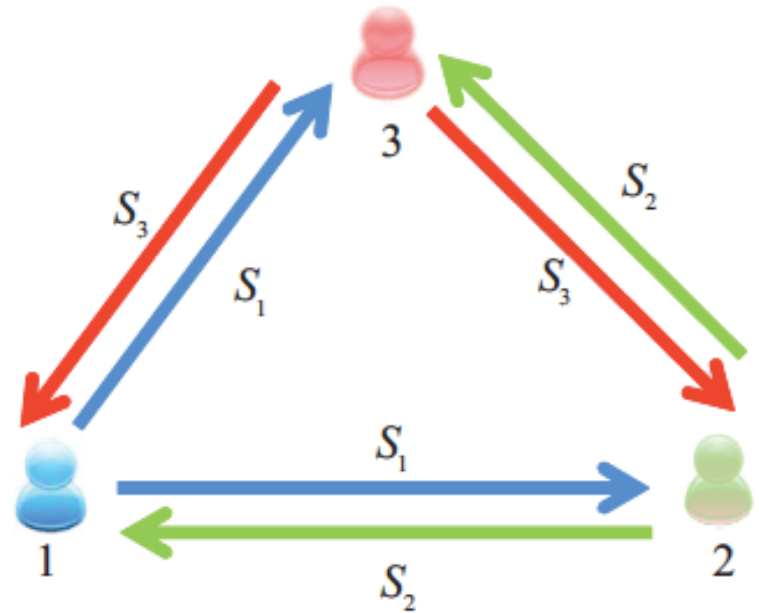


The result of the iterative weak dominance is not unique in general !

# Weak Dominance in Threshold Cryptography



$S = f(0)$  is the secret, and each  $S_i$  is calculated using a polynomial function



Each party should receive the other two secret shares to calculate the secret.

J. Halpern and V. Teague, "Rational secret sharing and multiparty computation"  
In *Proceedings of the 36th Annual ACM Symposium on Theory of Computing*, 2004

# Weak Dominance in Threshold Cryptography

- The parties are rational and that they cooperate if it is in their interest to share a part of the secret (it increases its payoff)
- Given the rationality assumption:  
*“Rational parties will not broadcast their shares”*
- **Not sending the share (Defect)** is a **weakly dominating** strategy in the game between the parties
- Results make sense if we consider that the parties have common knowledge about the running time of the protocol