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# Algorithmic Game Theory

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۱۳۹۶



Some Formal Definitions

# **DYNAMIC GAMES II**

# **Dynamic Games II**

- 1. First Mover or Second Mover?**
- 2. Zermelo Theorem**
- 3. Perfect Information/Pure Strategy**
- 4. Imperfect Information/Information Set**
- 5. Information vs Time**

# First mover advantage

- Is being the first mover always good?
  - Yes, sometimes: as in the Cournot Stackelberg model
  - Not always, as in the Rock, Paper, Scissors game
  - Sometimes neither being the first nor the second is good

# The NIM game

- We have two players
- There are two piles of stones, A and B
- Each player, in turn, decides to delete some stones from whatever pile
- The player that remains with the last stone wins

**Let's play the game**

# The NIM game (2)

- If piles are equal → second mover advantage
- If piles are unequal → first mover advantage
- You'll know who will win the game from the initial setup
- You can solve through backward induction

# **Dynamic Games II**

1. First Mover or Second Mover?
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# The Zermelo Theorem

- Consider a general 2-Player game
- We assume **perfect information**
  - Players know where they are in the game tree and how they got there
- We assume a finite game, i.e. a game-tree with a finite number of nodes
- There can be three or fewer outcomes:  
 $W_1$  (player 1 wins),  $L_1$  (player 2 wins),  $T$  (tie)

# The Zermelo Theorem

The result (or solution) of this game is:

1. Either player 1 can force a **win (over player 2)**
2. Or player 1 can force a **tie**
3. Or player 2 can force a **loss (on player 1)**

# The Zermelo Theorem

- This theorem appears to be trivial:
  - Three possible outcomes
  - Games are subdivided in three categories:
    - Those in which, whatever player 2 does, player 1 can win (provided he/she plays well)
    - Those in which player 1 can always force a draw/tie
    - Those in which, player 1 is toast, and can only loose

# Examples of games

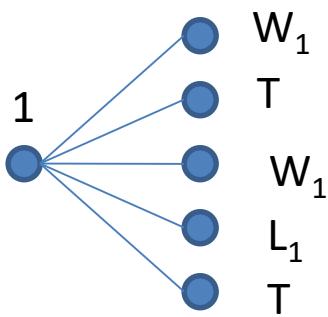
- **NIM**
  - Depends on number of stones in the first stage
- **Tic-tac-toe:**
  - If players play correctly, you can always force a tie
  - If players make wrong moves, they can loose
- **Chess → has a solution!**
- In fact, the theorem **doesn't tell you how to play**, it just tells you there is a solution!

# The Zermelo Theorem proof (I)

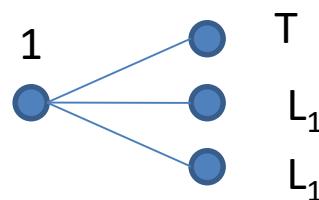
- We're going to prove the theorem, in a sketchy way, as this relates to backward induction
- Proof methodology:  
**Induction on maximum length of a game  $N$** 
  - We'll start with an induction hypothesis
  - And we'll prove this is true for longer games

# The Zermelo Theorem proof (2)

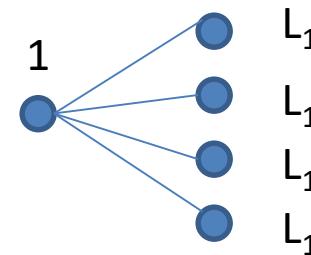
- If  $N = I$



1  
W<sub>1</sub>



1  
T

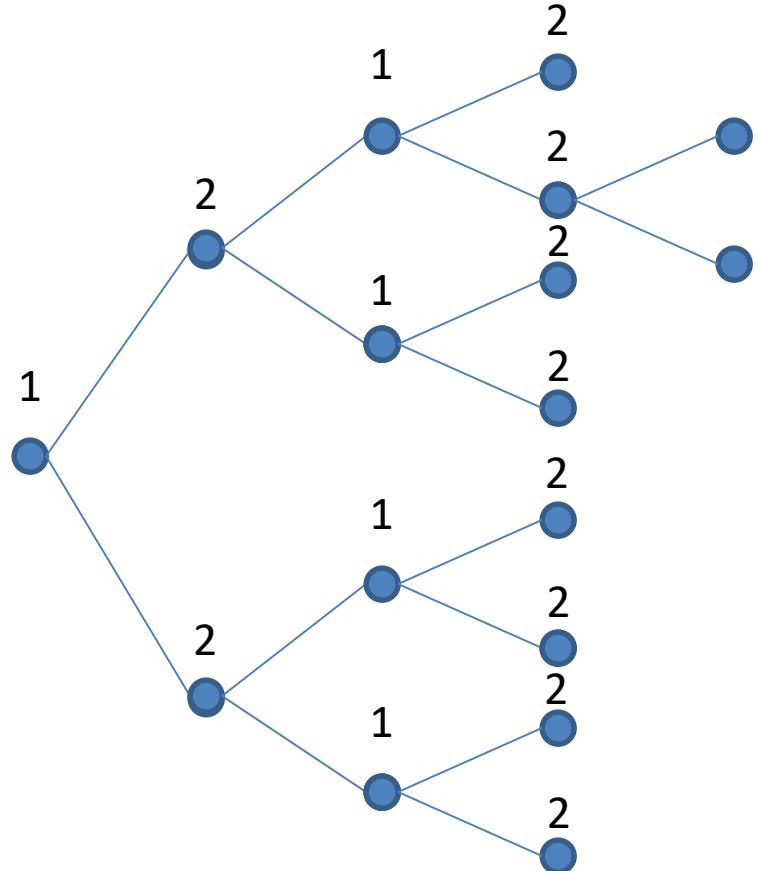


1  
L<sub>1</sub>

# The Zermelo Theorem proof (3)

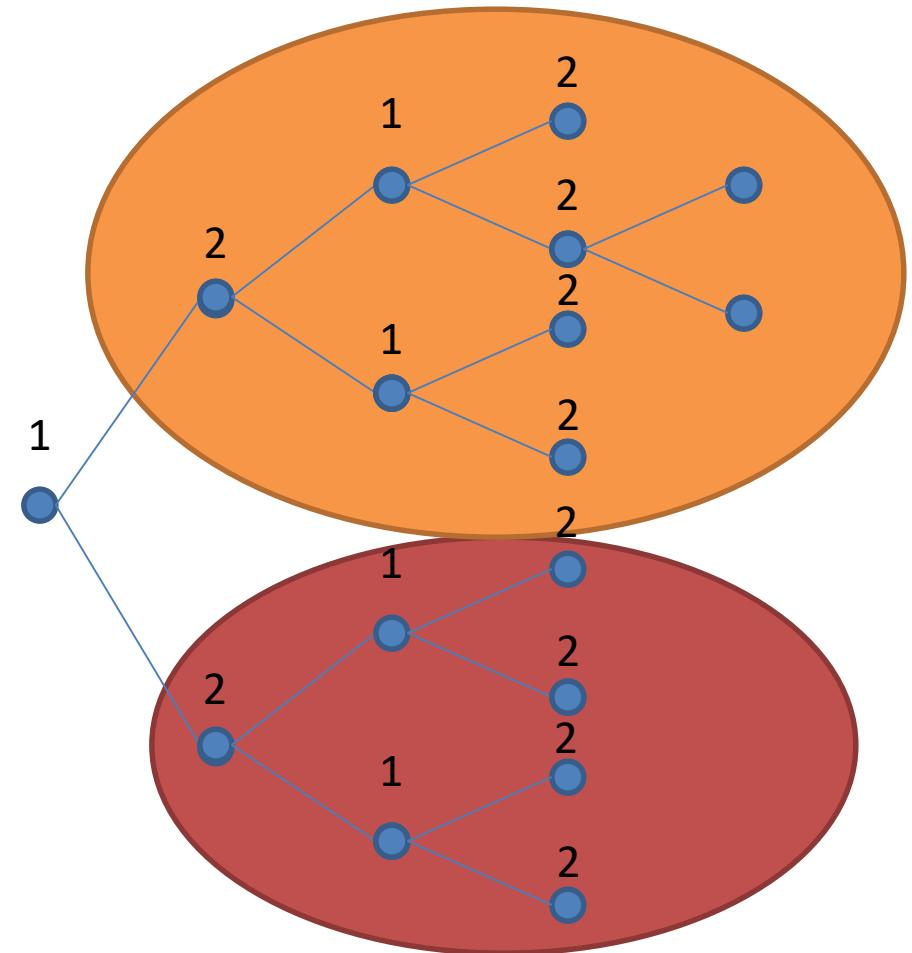
- Induction hypothesis:  
Suppose the claim is true for all games of length  $\leq N$
- We claim, therefore it will be true for games of length  $N+1$
- Let's take an example

# The Zermelo Theorem proof (4)



➤ What is the maximum length of the game?

# The Zermelo Theorem proof (5)

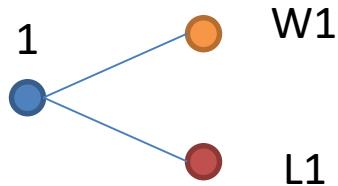


We have two **sub-games**

- The upper **sub-game**: follows “1” and it has length 3
- The lower **sub-game**: follows “1” and has length 2

# The Zermelo Theorem proof (6)

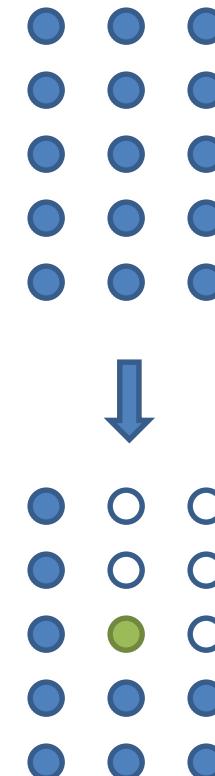
- By induction hypothesis (for  $N=3$ ), upper sub-game has a solution, say “ $W_1$ ”
- Again, by induction hypothesis ( $N=2$ ), lower sub-game has a solution, say “ $L_1$ ”



- This game has a solution, it is a game of length 1 we know already!

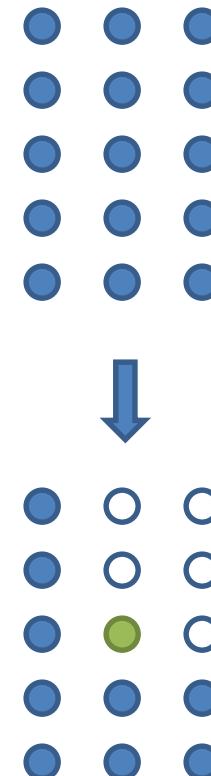
# A more Complex Example

- Suppose we have an array of stones, and two players
- Sequential moves, each player can delete some stones
  - Select one, delete all stones that lie above and right
- **The looser is the person who ends up removing the last rock**



# A more Complex Example

- According to Zermelo's Theorem, this game has a solution and the advantage depends on NxM, the size of the array
- Think about it!



# Dynamic Games II

1. First Mover or Second Mover?
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4. Imperfect Information/Information Set
5. Information vs Time

# **FORMAL DEFINITIONS**

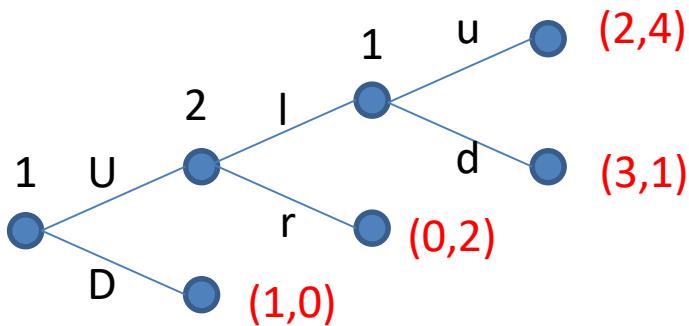
# Perfect Information Game

A game of perfect information is one in which at each node of the game tree, the player whose turn is to move knows which node she is at and how she got there

# Pure Strategy

A **pure strategy** for player  $i$  in a game of perfect information is a **complete plan** of actions: it specifies which action  $i$  will take at each of its decision nodes

# Example I



- Strategies

- Player 2:

- $[l], [r]$

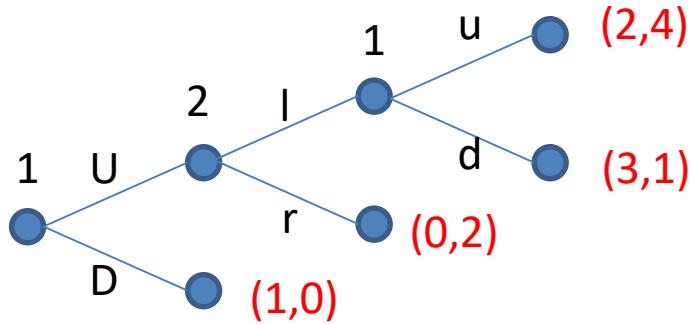
- Player 1:

- $[U,u], [U,d]$

- $[D,u], [D,d]$

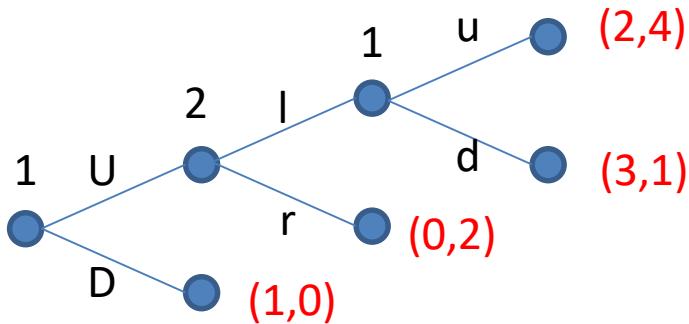
Hey, they look redundant, but we need them!

# Example I



- Note:
  - In this game it appears that player 2 may never have the possibility to play her strategies
  - This is also true for player 1!

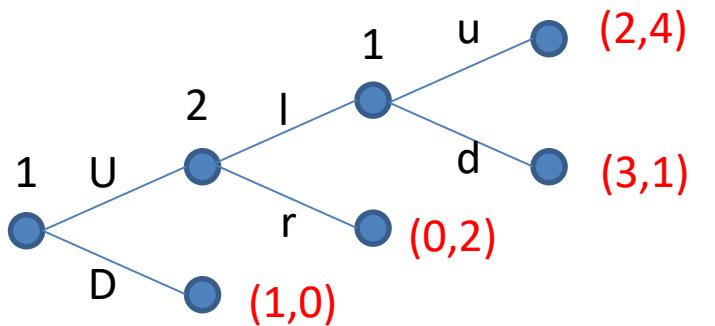
# Example I



- Backward Induction
  - Start from the end
    - “d” → higher payoff
  - Summarize game
    - “r” → higher payoff
  - Summarize game
    - “D” → higher payoff

➤ BI :: {[D,d],r}

# Example I

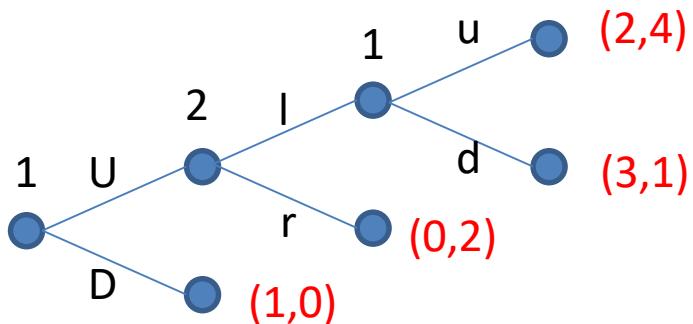


I      r

	I	r
U u	2,4	0,2
U d	3,1	0,2
D u	1,0	1,0
D d	1,0	1,0

From the *extensive form*  
To the *normal form*

# Example I



	I	r
U u	2,4	0,2
U d	3,1	0,2
D u	1,0	1,0
D d	1,0	1,0

Backward Induction

{[D, d], r}



Nash Equilibrium

{[D, d], r}  
 {[D, u], r}

Wait! We will find an answer to this later.

# **Dynamic Games II**

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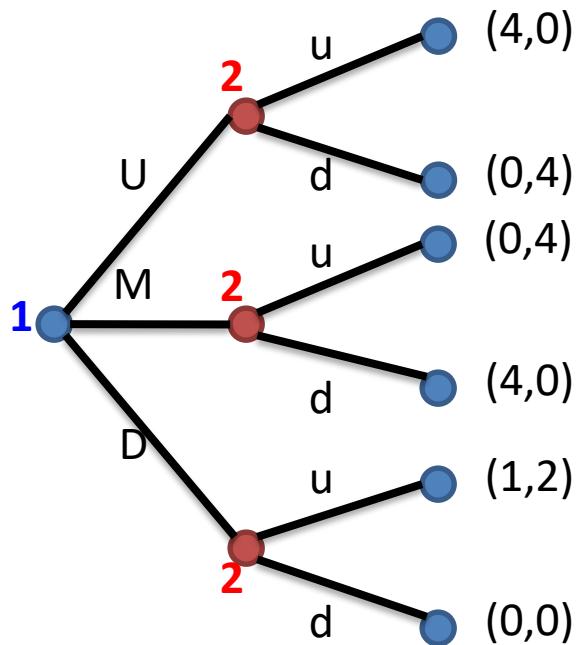
Let's be in the real world!

# **IMPERFECT INFORMATION**

# Brief Review

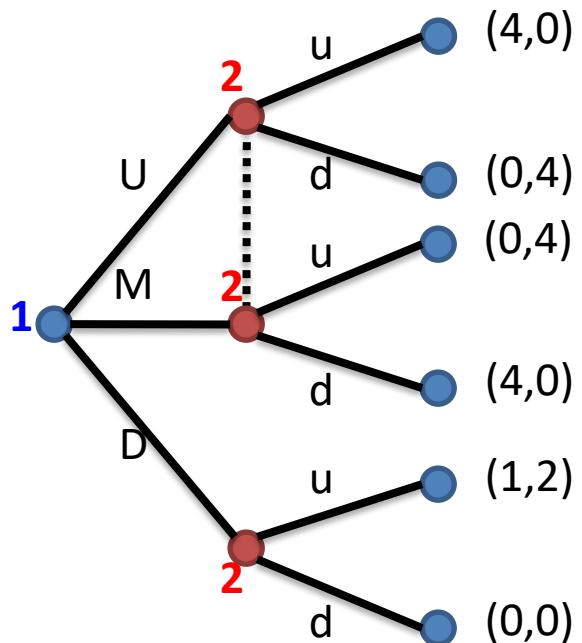
- We have seen **simultaneous** move games, in which players **cannot** observe strategies and have to reason based on the idea of **best response**
- We have seen **sequential** move games, in which **observation** is allowed, and players reason using **backward induction**
- Now, let's study a class of games in which these two approaches **are blended**

# A Simple Dynamic Game



- Sequential move game
- Assume for a moment perfect information
- We know how to solve it using **backward induction**
  - Player 1 knows that if he chooses U or M, player 2 can **crush** him
  - Player 2 has a **huge second mover advantage** in the first branches of the tree

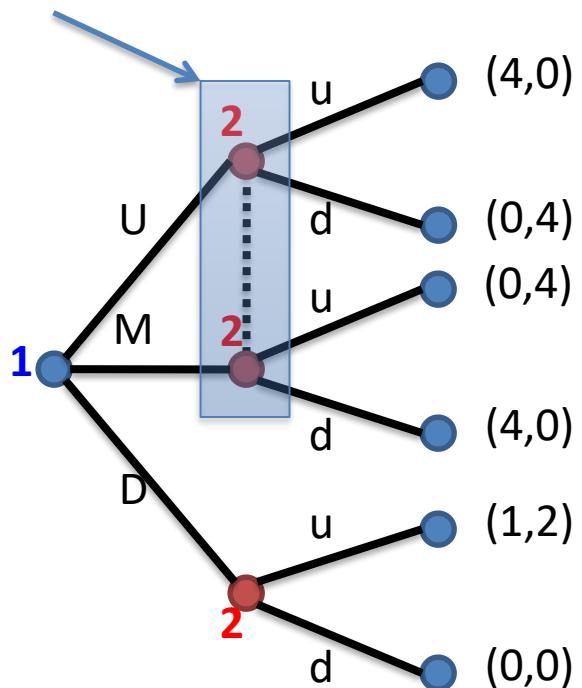
# Imperfect Information



- Sequential move game
- **Imperfect information**
  - Player 2 **cannot distinguish** where she is on (some parts of) the tree
- If player 1 chooses D, player 2 can observe it
- If player 1 chooses U or M, player 2 doesn't know which choice was made

# Information Set

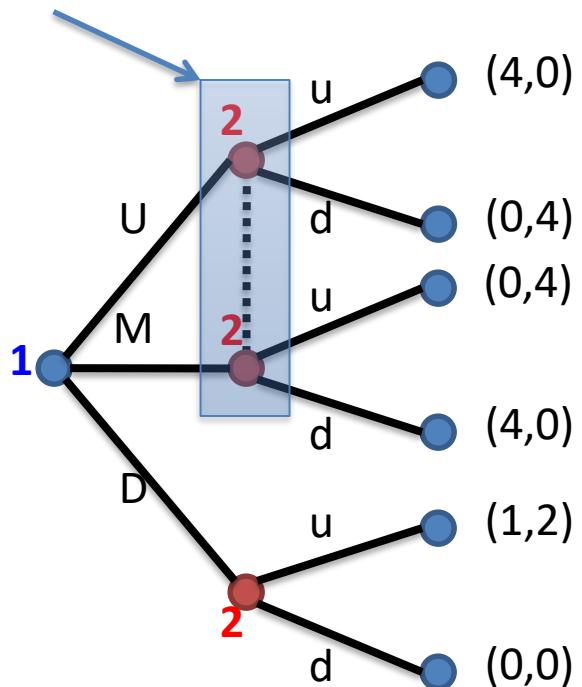
Information set



- The idea is that the two internal nodes are in the same **information set**
  - Player 2 knows that player 1 chose whether U or M, but not which one
- How can we analyze this kind of games?

# How to Solve?

Information set

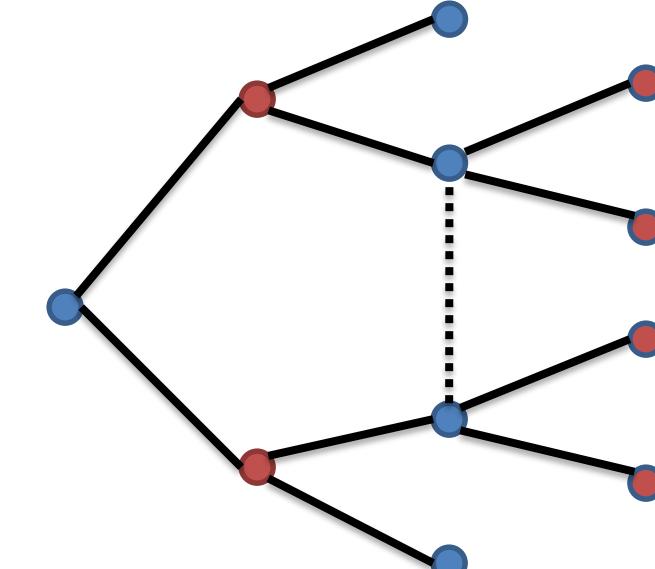
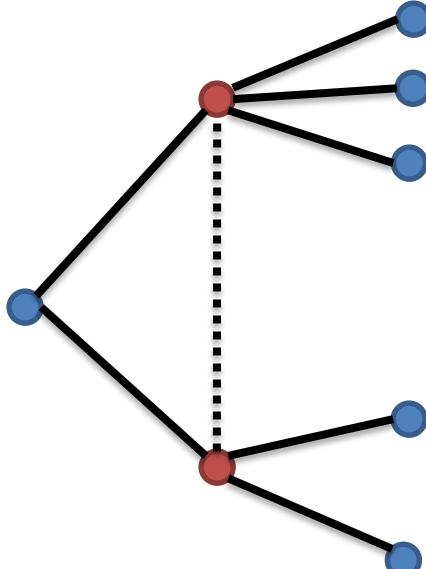


- The simple backward induction argument (player 2 could always crush player 1) does not hold anymore
- Moreover, player 1 **knows** that player 2 cannot distinguish U or M
  - Player 1 **might** decide to randomize over U and M, and hope to get an expected payoff of 2
  - A payoff of 2 is better than what player 1 could ever obtain by choosing D

# Information Set

- An **information set** of player  $i$  is a collection of player  $i$ 's **decision nodes** among which  $i$  **cannot** distinguish

*Examples:* Are these information sets?



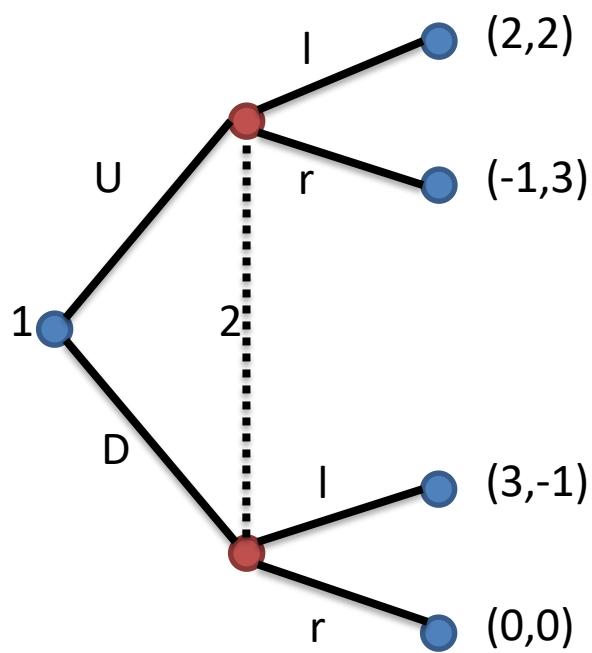
# Information Sets: Some Rules

- **Rule 1:** A player must not be able to infer in which node she is by looking at the **number of available strategies** she has
- **Rule 2:** provided a player can recall what she did earlier on in the tree, she shouldn't be able to distinguish where she is
  - This assumption is called **perfect recall**
  - **NOTE:** perfect recall is not always realistic!

## **Definition: Perfect/Imperfect Information**

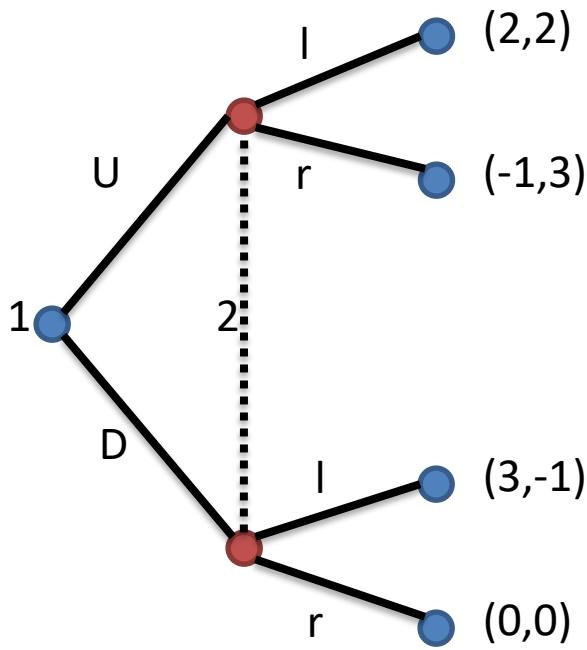
- A game of **perfect information** is a game in which all information sets in the game tree include just one node
- A game of **imperfect information** is not a game of perfect information!

# Simple Example



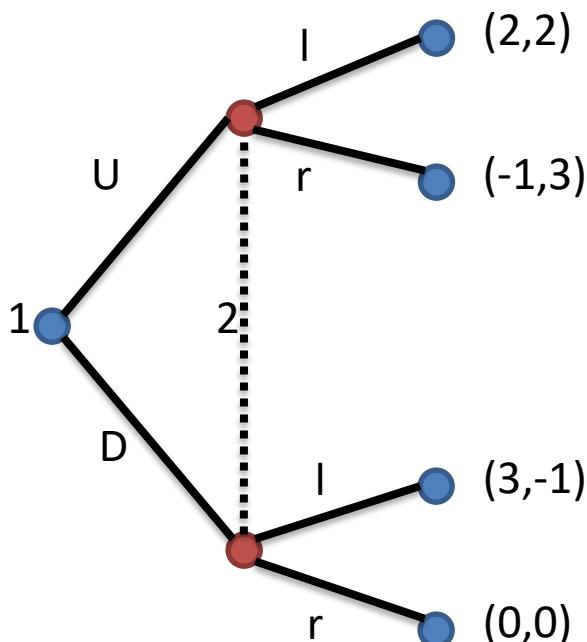
- The **information set** indicates that player 2 cannot observe whether player 1 moved *up* or *down*
  - **Perfect information:** player 2 could have chosen separately, in each node, whether to choose *left* or *right*
  - **Imperfect information:** player 2 has *only the choice of choosing left or right*, for both nodes, since she doesn't know which one she'll be at

# Solution



- There's a catch here that makes the game easy:
  - Whatever is the information set, for player 2 choosing *right* is consistently better than choosing *left*
  - This game solves out rather like when using **backward induction**

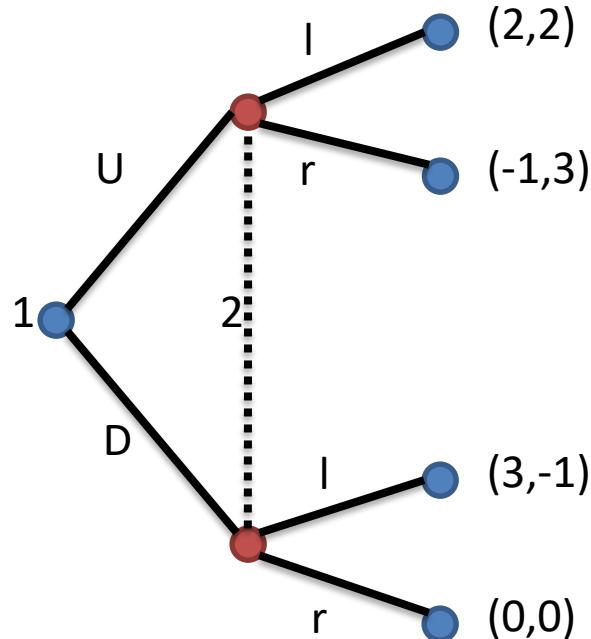
# From Dynamic to Static Game



		Player 2	
		I	r
		2,2	-1,3
Player 1	U	3,-1	0,0
	D		

- **Question:** What game is this?
  - Prisoners Dilemma
- Notice that by using **information sets**, we were able to represent in a tree a simultaneous move game
  - It does not really matter the **time** here, what matters is **information**

# From Dynamic to Static Game



		Player 2	
		I	r
		2,2	-1,3
Player 1	U	2,2	-1,3
	D	3,-1	0,0

- We don't have **redundant** strategies in the matrix
- We can't have a complete action plan when we don't know where we are in the tree
  - This implies we have to revisit our definition of strategy

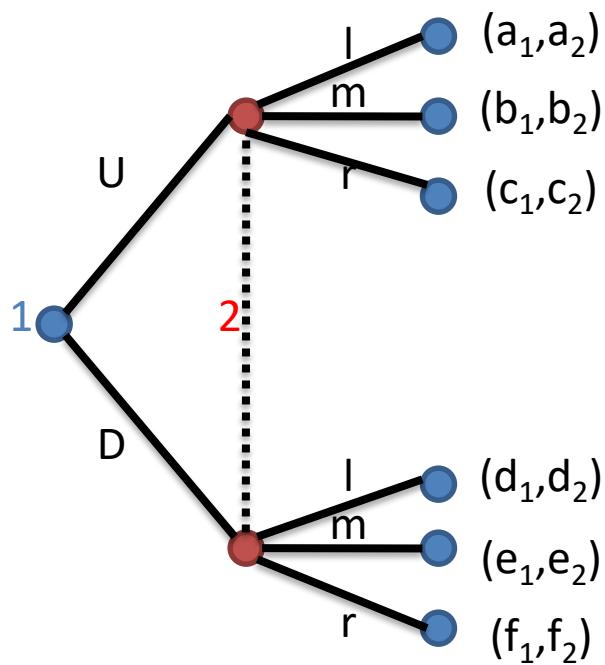
# Pure Strategy: A New Definition

- A **pure strategy** of player  $i$  is a complete plan of action: it specifies what player  $i$  will do at each of its ***information sets***
- It looks like the same definition we saw last time, but this one involves information sets and it is more general
  - The idea remains the same: we want to transform a game tree in a matrix

# **Dynamic Games II**

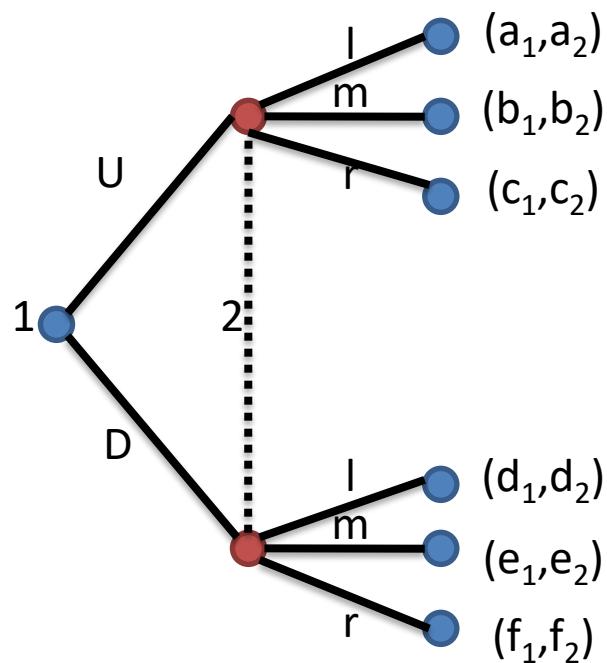
1. First Mover or Second Mover?
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# Information vs Time



- Player 2 does not know if player 1 chooses up or down  
→ Player 2 has just three choices
- Our goal now is to transform the game into a matrix

# Information vs Time

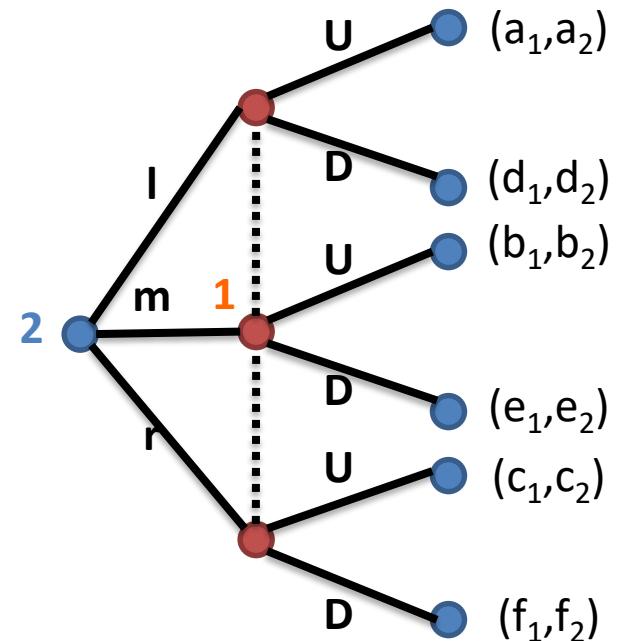


		Player 2		
		l	m	r
		U	$a_1, a_2$	$b_1, b_2$
		D	$d_1, d_2$	$e_1, e_2$
			$f_1, f_2$	

**CLAIM:** If we look at the matrix above **it is not obvious** that the game tree on the left is the only possible tree that could generate the matrix

# Information vs Time

		Player 2		
		I	m	r
Player 1		U	$a_1, a_2$	$b_1, b_2$
		D	$d_1, d_2$	$e_1, e_2$
			$c_1, c_2$	$f_1, f_2$



In the game tree to the right,  
player 2 moves first, then player 1  
moves but she doesn't know which  
action player 2 chose

**CLAIM:** These two games trees are **equivalent**

# Observations

- What matters is **not time**, but **information**
- We would like to set-up the machinery to analyze such games and predict what it is going to happen