



Algorithmic Game Theory

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Some Formal Definitions

DYNAMIC GAMES II

Dynamic Games II

1. First Mover or Second Mover?
2. Zermelo Theorem
3. Perfect Information/Pure Strategy
4. Imperfect Information/Information Set
5. Information vs Time

First mover advantage

- Is being the first mover always good?
 - **Yes, sometimes**: as in the Cournot Stackelberg model
 - **Not always**, as in the Rock, Paper, Scissors game
 - Sometimes neither being the first nor the second is good

The NIM game

- We have two players
- There are two piles of stones, A and B
- Each player, in turn, decides to delete some stones from whatever pile
- The player that remains with the last stone wins

Let's play the game

The NIM game (2)

- If piles are equal → second mover advantage
- If piles are unequal → first mover advantage
- You'll know who will win the game from the initial setup
- You can solve through backward induction

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The Zermelo Theorem

- Consider a general 2-Player game
- We assume **perfect information**
 - Players know where they are in the game tree and how they got there
- We assume a finite game, i.e. a game-tree with a finite number of nodes
- There can be three or fewer outcomes:
 W_1 (player 1 wins), L_1 (player 2 wins), T (tie)

The Zermelo Theorem

The result (or solution) of this game is:

1. Either player 1 can force a **win (over player 2)**
2. Or player 1 can force a **tie**
3. Or player 2 can force a **loss (on player 1)**

The Zermelo Theorem

- This theorem appears to be trivial:
 - Three possible outcomes
 - Games are subdivided in three categories:
 - Those in which, **whatever player 2 does**, player 1 can win (provided he/she plays well)
 - Those in which player 1 can always force a draw/tie
 - Those in which, player 1 is toast, and can only loose

Examples of games

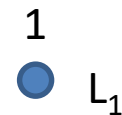
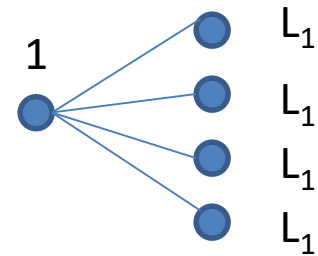
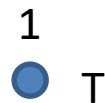
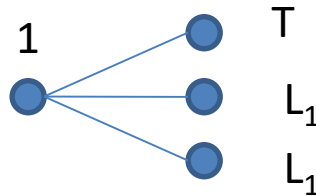
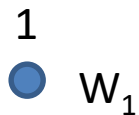
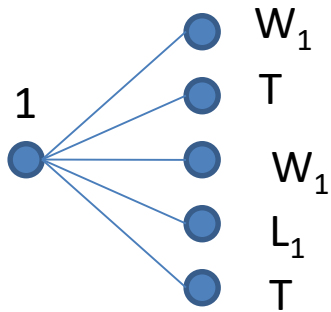
- **NIM**
 - Depends on number of stones in the first stage
- **Tic-tac-toe:**
 - If players play correctly, you can always force a tie
 - If players make wrong moves, they can loose
- **Chess** ➔ has a solution!
- In fact, the theorem **doesn't tell you how to play**, it just tells you there is a solution!

The Zermelo Theorem proof (I)

- We're going to prove the theorem, in a sketchy way, as this is relates to backward induction
- **Proof methodology:**
Induction on maximum length of a game N
 - We'll start with an induction hypothesis
 - And we'll prove this is true for longer games

The Zermelo Theorem proof (2)

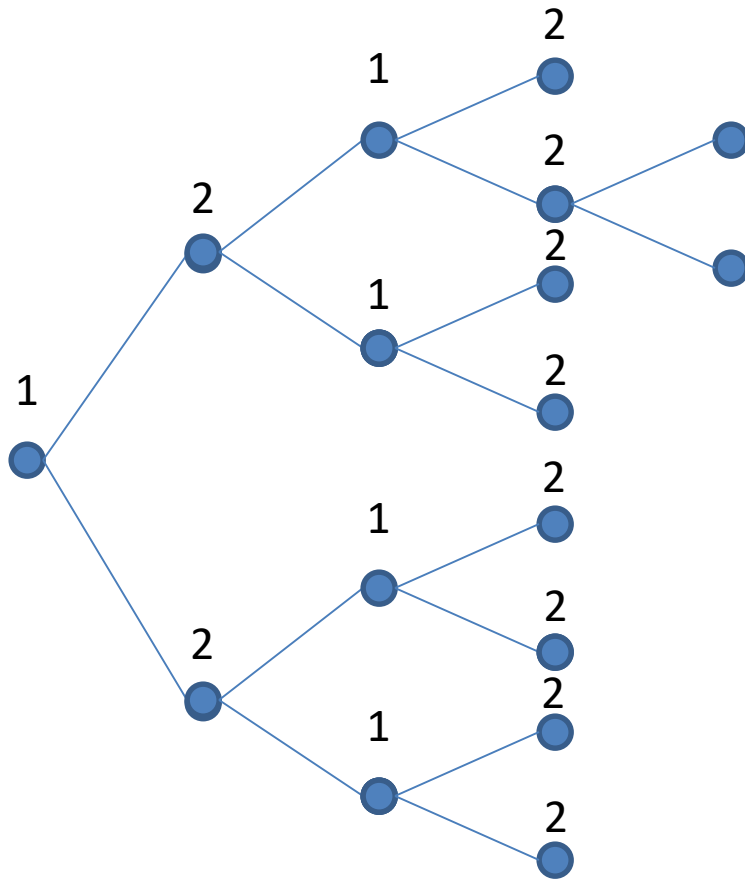
- If $N = 1$



The Zermelo Theorem proof (3)

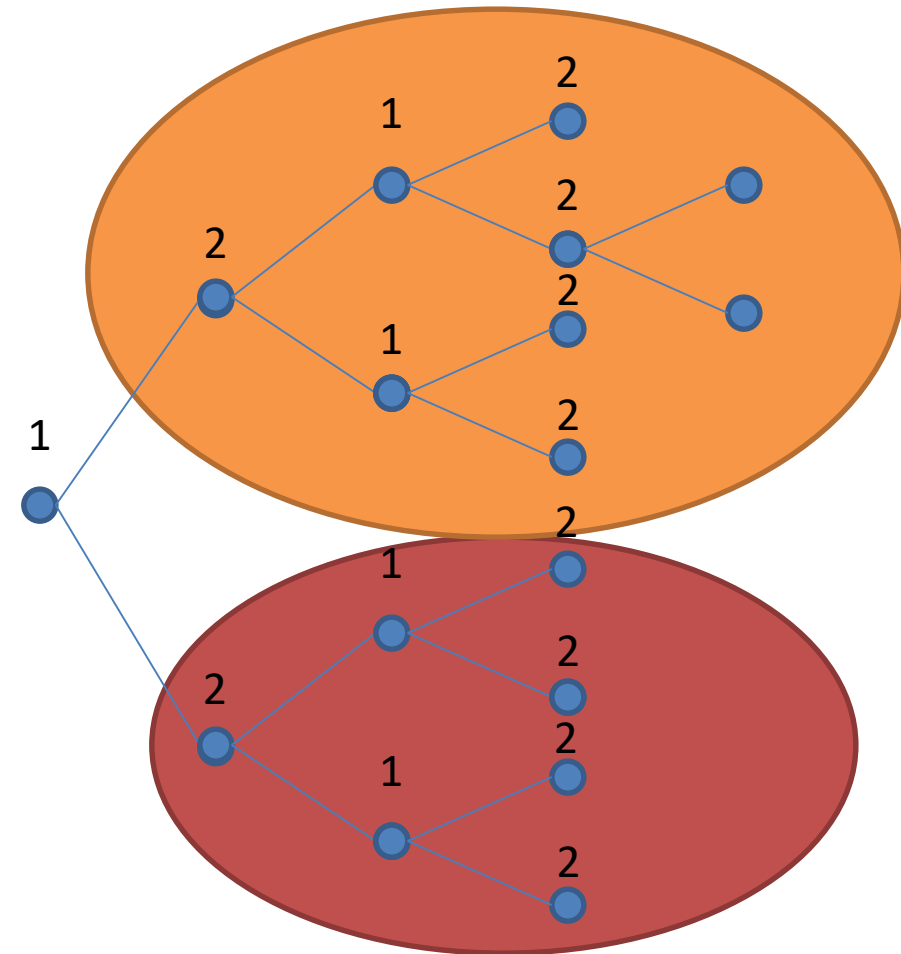
- Induction hypothesis:
Suppose the claim is true for all games of length $\leq N$
- We claim, therefore it will be true for games of length $N+1$
- Let's take an example

The Zermelo Theorem proof (4)



➤ What is the maximum length of the game?

The Zermelo Theorem proof (5)

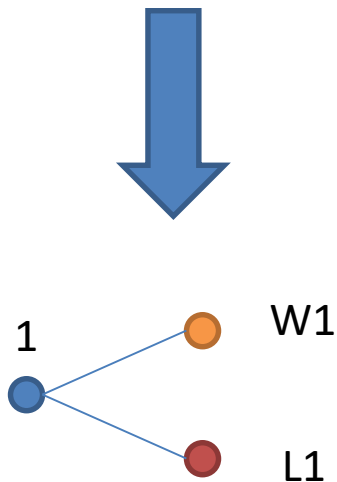


We have two **sub-games**

- The upper **sub-game**: follows “1” and it has length 3
- The lower **sub-game**: follows “1” and has length 2

The Zermelo Theorem proof (6)

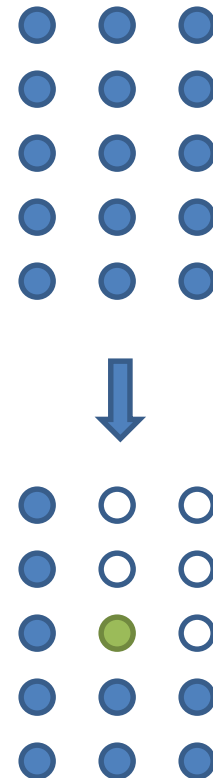
- By induction hypothesis (for $N=3$), upper sub-game has a solution, say “ W_1 ”
- Again, by induction hypothesis ($N=2$), lower sub-game has a solution, say “ L_1 ”



- This game has a solution, it is a game of length 1 we know already!

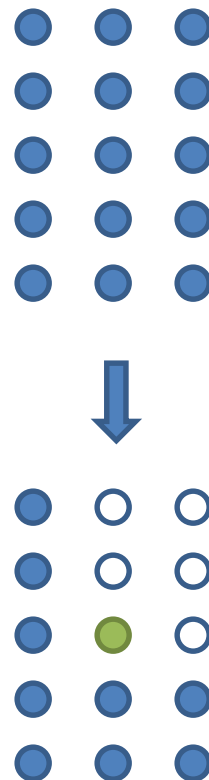
A more Complex Example

- Suppose we have an array of stones, and two players
- Sequential moves, each player can delete some stones
 - Select one, delete all stones that lie above and right
- **The loser is the person who ends up removing the last rock**



A more Complex Example

- According to Zermelo's Theorem, this game has a solution and the advantage depends on $N \times M$, the size of the array
- Think about it!



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FORMAL DEFINITIONS

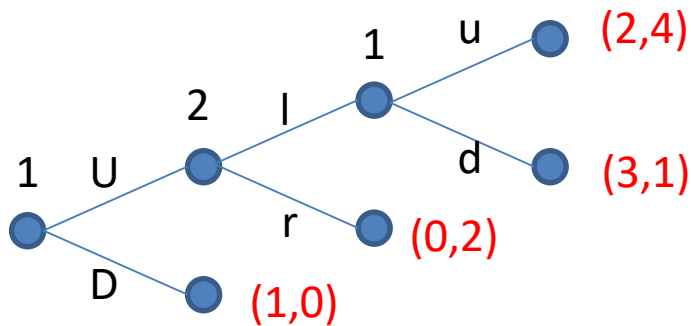
Perfect Information Game

A game of **perfect information** is one in which at each node of the game tree, the player whose turn is to move **knows** which node she is at and **how** she got there

Pure Strategy

A ***pure strategy*** for player i in a game of perfect information is a ***complete plan*** of actions: it specifies which action i will take at each of its decision nodes

Example I



- Strategies

- Player 2:

- [l], [r]

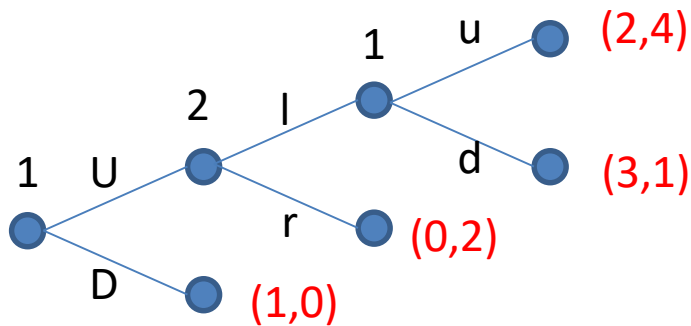
- Player 1:

- [U,u], [U,d]

- [D, u], [D,d]

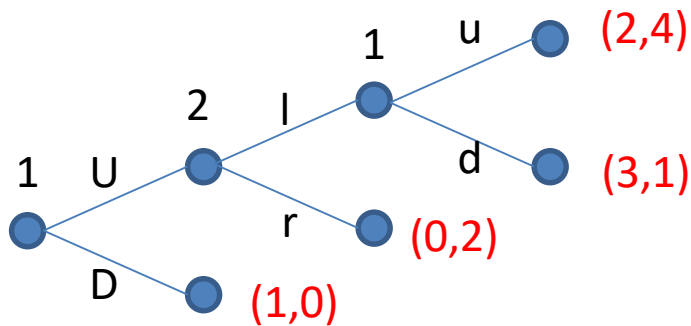
Hey, they look redundant, but we need them!

Example I



- Note:
 - In this game it appears that player 2 may never have the possibility to play her strategies
 - This is also true for player 1!

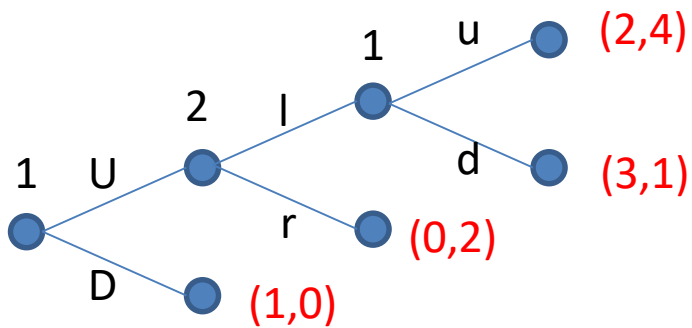
Example I



- Backward Induction
 - Start from the end
 - “d” \rightarrow higher payoff
 - Summarize game
 - “r” \rightarrow higher payoff
 - Summarize game
 - “D” \rightarrow higher payoff

➤ BI :: {[D,d],r}

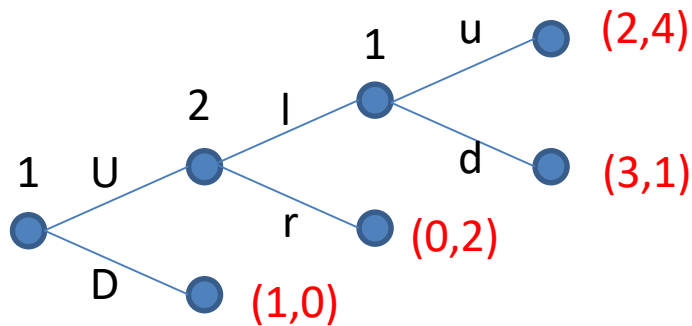
Example I



	l	r
U u	2,4	0,2
U d	3,1	0,2
D u	1,0	1,0
D d	1,0	1,0

From the *extensive form*
To the *normal form*

Example I



	l	r
U u	2,4	0,2
U d	3,1	0,2
D u	1,0	1,0
D d	1,0	1,0

Backward Induction

$\{[D, d], r\}$



Nash Equilibrium

$\{[D, d], r\}$
 $\{[D, u], r\}$

Wait! We will find an answer to this later.

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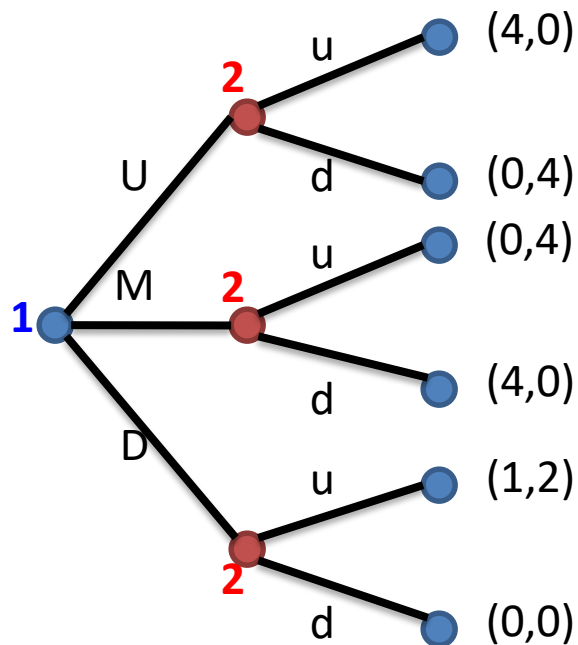
Let's be in the real world!

IMPERFECT INFORMATION

Brief Review

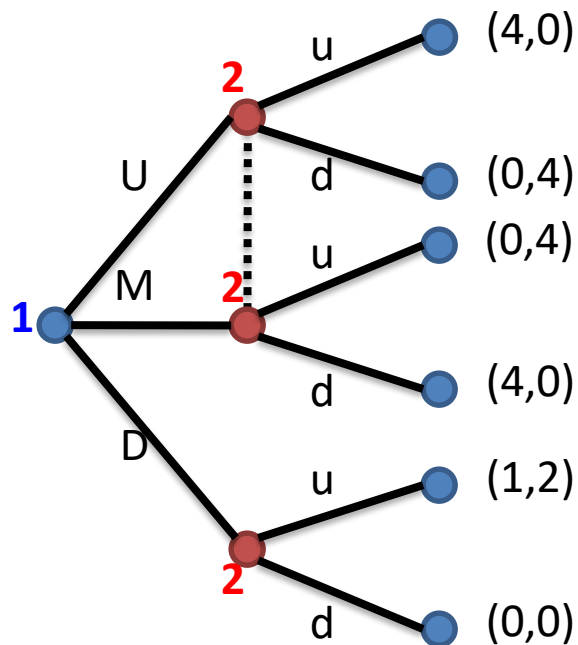
- We have seen **simultaneous** move games, in which players **cannot** observe strategies and have to reason based on the idea of **best response**
- We have seen **sequential** move games, in which **observation** is allowed, and players reason using **backward induction**
- Now, let's study a class of games in which these two approaches **are blended**

A Simple Dynamic Game



- Sequential move game
- Assume for a moment perfect information
- We know how to solve it using **backward induction**
 - Player 1 knows that if he chooses U or M, player 2 can **crush** him
 - Player 2 has a **huge second mover advantage** in the first branches of the tree

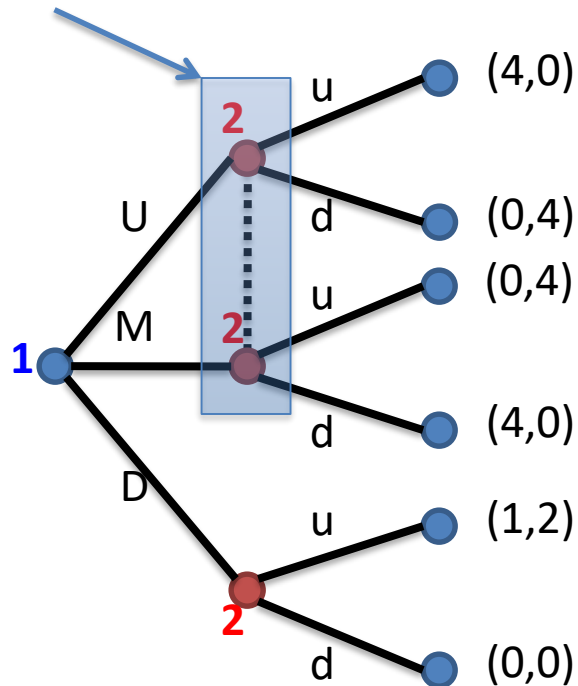
Imperfect Information



- Sequential move game
- **Imperfect information**
 - Player 2 **cannot distinguish** where she is on (some parts of) the tree
- If player 1 chooses D, player 2 can observe it
- If player 1 chooses U or M, player 2 doesn't know which choice was made

Information Set

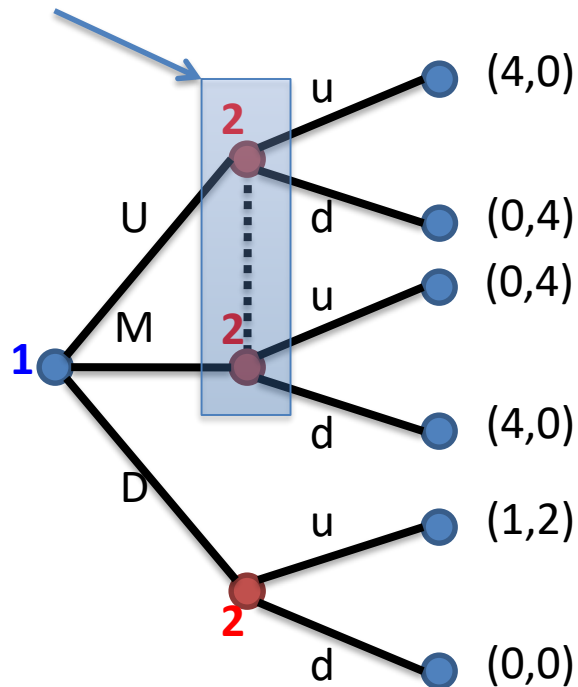
Information set



- The idea is that the two internal nodes are in the same **information set**
 - Player 2 knows that player 1 chose whether U or M, but not which one
- How can we analyze this kind of games?

How to Solve?

Information set

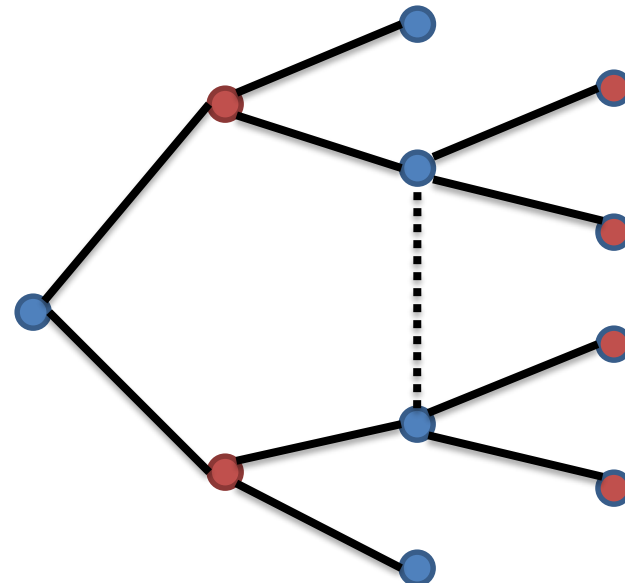
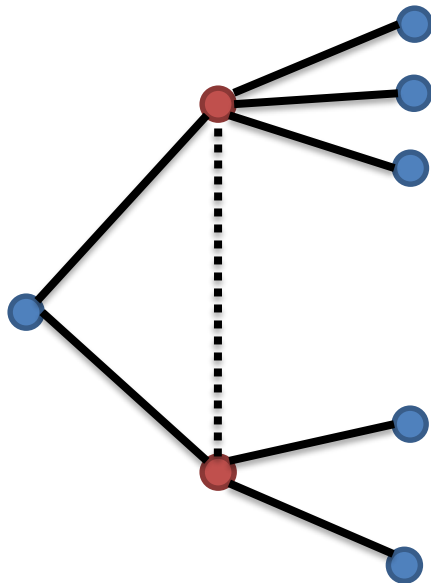


- The simple backward induction argument (player 2 could always crush player 1) does not hold anymore
- Moreover, player 1 **knows** that player 2 cannot distinguish U or M
 - Player 1 **might** decide to randomize over U and M, and hope to get an expected payoff of 2
 - A payoff of 2 is better than what player 1 could ever obtain by choosing D

Information Set

- An **information set** of player i is a collection of player i 's **decision nodes** among which i **cannot** distinguish

Examples: Are these information sets?



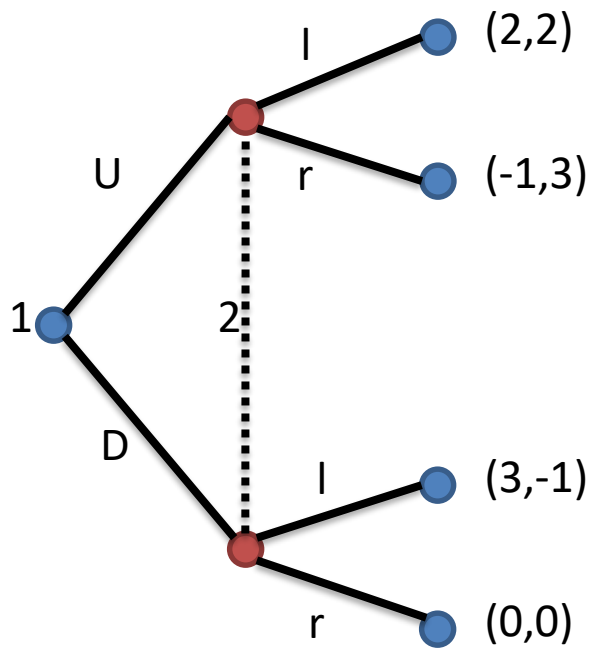
Information Sets: Some Rules

- **Rule 1:** A player must not be able to infer in which node she is by looking at the **number of available strategies** she has
- **Rule 2:** provided a player can recall what she did earlier on in the tree, she shouldn't be able to distinguish where she is
 - This assumption is called **perfect recall**
 - **NOTE:** perfect recall is not always realistic!

Definition: Perfect/Imperfect Information

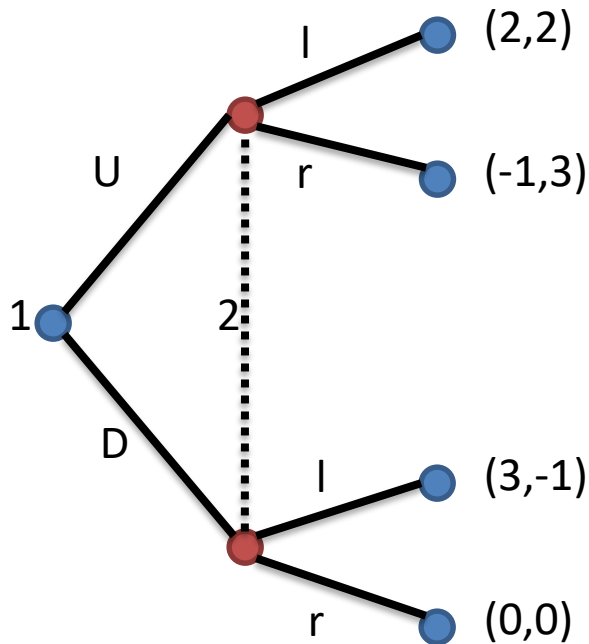
- A game of **perfect information** is a game in which all information sets in the game tree include just one node
- A game of **imperfect information** is not a game of perfect information!

Simple Example



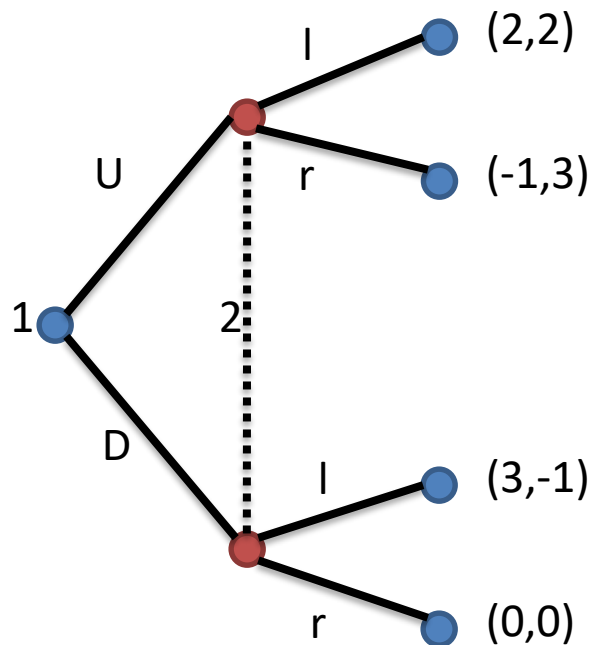
- The **information set** indicates that player 2 cannot observe whether player 1 moved *up* or *down*
 - **Perfect information:** player 2 could have chosen separately, in each node, whether to choose *left* or *right*
 - **Imperfect information:** player 2 has *only the choice of choosing left or right*, for both nodes, since she doesn't know which one she'll be at

Solution



- There's a catch here that makes the game easy:
 - Whatever is the information set, for player 2 choosing *right* is consistently better than choosing *left*
 - This game solves out rather like when using **backward induction**

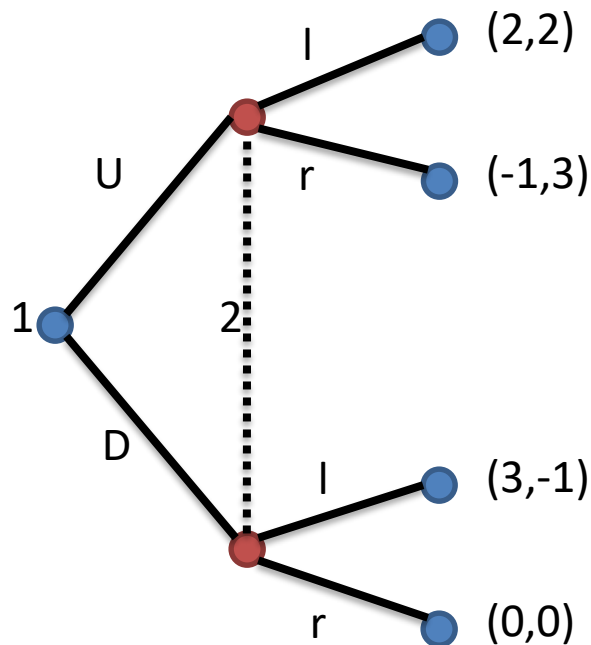
From Dynamic to Static Game



		Player 2	
		l	r
Player 1	U	2,2	-1,3
	D	3,-1	0,0

- **Question:** What game is this?
 - Prisoners Dilemma
- Notice that by using **information sets**, we were able to represent in a tree a simultaneous move game
 - It does not really matter the **time** here, what matters is **information**

From Dynamic to Static Game



		Player 2	
		l	r
Player 1	U	2,2	-1,3
	D	3,-1	0,0

- We don't have **redundant** strategies in the matrix
- We can't have a complete action plan when we don't know where we are in the tree
 - This implies we have to revisit our definition of strategy

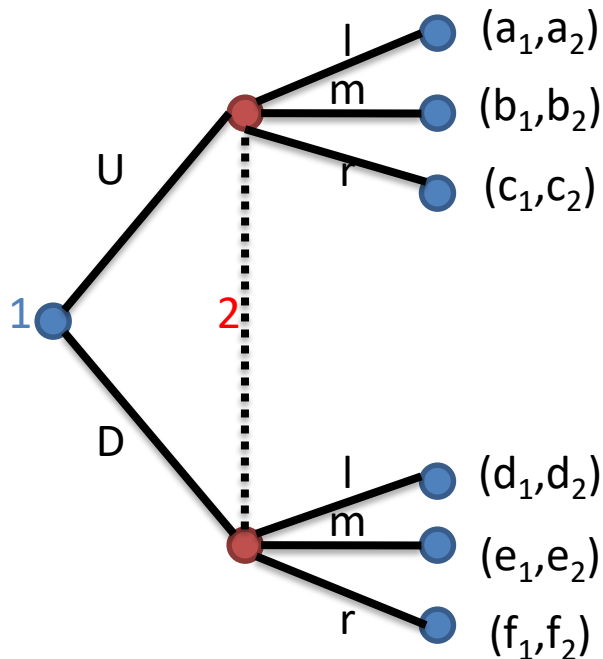
Pure Strategy: A New Definition

- A **pure strategy** of player i is a complete plan of action: it specifies what player i will do at each of its ***information sets***
- It looks like the same definition we saw last time, but this one involves information sets and it is more general
 - The idea remains the same: we want to transform a game tree in a matrix

Dynamic Games II

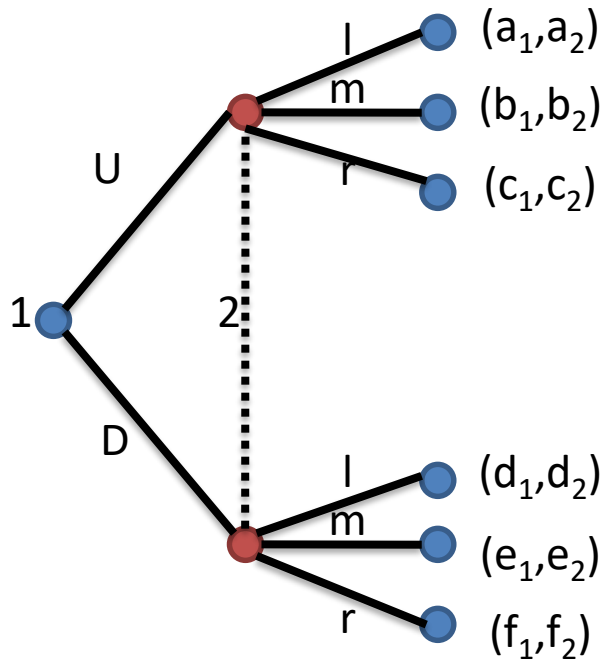
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Information vs Time



- Player 2 does not know if player 1 chooses *up* or *down*
→ Player 2 has just three choices
- Our goal now is to transform the game into a matrix

Information vs Time

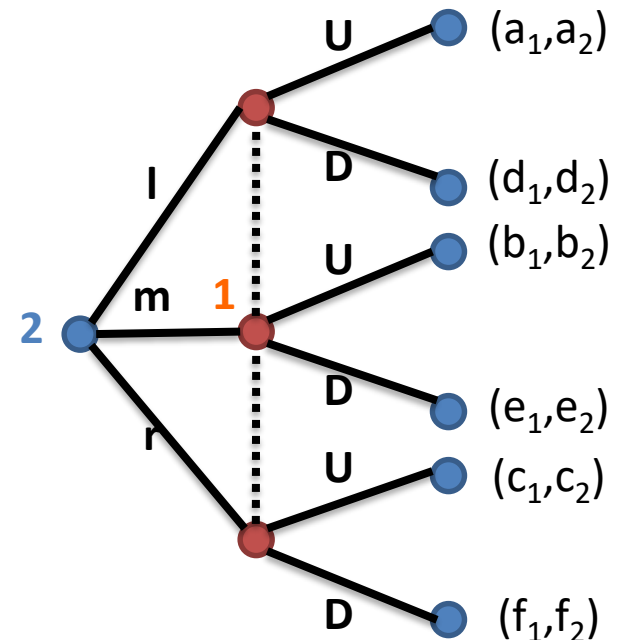


		Player 2		
		l	m	r
Player 1	U	a_1, a_2	b_1, b_2	c_1, c_2
	D	d_1, d_2	e_1, e_2	f_1, f_2

CLAIM: If we look at the matrix above it is **not obvious** that the game tree on the left is the only possible tree that could generate the matrix

Information vs Time

		Player 2		
		l	m	r
Player 1	U	a_1, a_2	b_1, b_2	c_1, c_2
	D	d_1, d_2	e_1, e_2	f_1, f_2



In the game tree to the right,
player 2 moves first, then player 1
moves but she doesn't know which
action player 2 chose

CLAIM: These two games trees are **equivalent**

Observations

- What matters is **not time**, but **information**
- We would like to set-up the machinery to analyze such games and predict what it is going to happen