

# Collegio Carlo Alberto

## Game Theory

### Problem Set 11

1. A seller owns an object that a buyer wants to buy. The value of the object to the seller is  $c$ . The value of the object to the buyer is private information. The buyer's valuation  $v$  is a random variable distributed over the interval  $[0, V]$  according to the (continuous) c.d.f.  $F$ . Assume that  $[1 - F(v)] / f(v)$  is a decreasing function of  $v$ . The von Neumann-Morgenstern utility of a type  $v$  from getting a unit at price  $p$  is  $v - p$  and the utility of no purchase in 0.
  - (i) Suppose the seller is constrained to charge just one price. Show that the profit maximizing price satisfies  $p = c + [1 - F(p)] / f(p)$ .
  - (ii) Suppose that the seller can commit to a menu of offers  $[q(v), p(v)]$ , where  $q(v)$  is the probability with which a consumer who chooses offer  $v$  will get a unit, and  $p(v)$  is the price she will pay in the event that she gets a unit. Prove that the menu that maximizes the seller's profit consists of a single price, which is the one found in (i), and that any buyer can get the good at this price with probability 1.
2. Consider the following auction environment. A seller has a single object for sale and can commit to any selling mechanism (the seller's valuation of the object is zero). There are two potential bidders, indexed by  $i = 1, 2$ . The valuation of the object of bidder  $i = 1, 2$  is denoted by  $v_i$  and is distributed uniformly over the unit interval. Valuations are independent between the two bidders. Bidder 1 knows her own valuation  $v_1$ . However, bidder 2 does *not* know  $v_2$ .

The bidders' payoffs are as follows. Suppose bidder  $i = 1, 2$  has type  $v_i$  and pays the amount  $t_i$  to the seller. Her payoff is equal to  $v_i - t_i$  if she gets the object, and equal to  $-t_i$  otherwise.

  - (i) Construct the optimal direct mechanism for the seller (i.e., find the incentive compatible, individually rational mechanism that maximizes the seller's expected revenues). Compute the seller's revenues.
  - (ii) Can you find a simple *indirect* mechanism that gives to the seller the same expected revenues as the optimal direct mechanism?
3. A seller has a unit for sale. Its quality is either high ( $H$ ) or low ( $L$ ). The quality is known to the seller but not to the buyer, whose prior probability that the quality is

high is  $1/2$ . Their valuations of the unit are as follows.

	Quality $H$	Quality $L$
Buyer	$V$	2
Seller	7	0

where  $V > 7$ . Thus, the utility to the buyer of getting the unit at price  $p$  is  $2 - p$  if it is of the low quality, and  $V - p$  if it is of the high quality. Similarly, the utility to the seller is  $p$  and  $p - 7$ , respectively.

- (i) Find the ex-post efficient outcomes.
- (ii) Identify the range of  $V$  (above 7) for which there is, and the range of  $V$  for which there is no incentive compatible, individual rational mechanism that will achieve the ex-post efficient outcome.
- (iii) Describe the best outcome (in the maximizing of the sum of expected utilities) that can be achieved for each  $V$  (above 7) and the mechanism that achieves it.

HINT: A mechanism for this Bayesian bargaining problem consists of a pair of functions  $q : \{L, H\} \rightarrow [0, 1]$  and  $t : \{L, H\} \rightarrow \mathbb{R}$ , where  $q(i)$  is the probability that the object will be sold to the buyer and  $t(i)$  is the expected net payment from the buyer to the seller if  $i = L, H$  is the type reported by the seller to a mediator.

4. A seller owns an object that a buyer wants to buy. The quality of the object is a random variable  $v$ , with support  $[0, 1]$  and distribution function  $F(v) = v^\alpha$ , where  $\alpha > 0$ . The seller knows the quality of the object but the buyer does not. When the quality of the object is  $v$ , the value of the object is  $v$  to the seller and  $zv$  to the buyer, where  $z > 1$ . Thus, if the object of quality  $v$  is traded at price  $p$ , the seller gets  $p - v$  and the buyer gets  $zv - p$ . Both players have utility equal to zero if there is no trade.

Consider the function  $G : (0, \infty) \times (1, \infty) \rightarrow [0, 1]$  defined as follows. For each pair  $(\alpha, z)$  construct the incentive-compatible individually rational mechanism that maximizes the (ex-ante) probability of trade. Denote this probability by  $G(\alpha, z)$ . Derive the function  $G$ .

(N.B. If the probability of trade is  $q(v)$  when the quality is  $v$ , then the (ex-ante) probability of trade is equal to  $\int_0^1 q(v) dF(v)$ .)