

# Combinatorial Game Theory: Surreal Numbers and the Void

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Any number can be written as a tuple of games played by the void with itself.

Denote the void by the empty set  $\emptyset$ . We write:  $\{\emptyset|\emptyset\} = 0$ , with  $|$  as partition. ‘Tuple’ signifies ordering matters, so that  $\{0|\emptyset\} = 1$  and  $\{\emptyset|0\} = -1$ . Then recursively construct the integers:  $\{n|\emptyset\} = n + 1$ . Plug  $\{\emptyset|\emptyset\}$  into  $\{0|\emptyset\}$  to get  $\{\{\emptyset|\emptyset\}|\emptyset\} = 1$ , then this into  $\{1|\emptyset\}$  to get  $\{\{\{\emptyset|\emptyset\}|\emptyset\}|\emptyset\} = 2...$

So if games exist, numbers exist. Or rather: if games exist, numbers don’t have to.

Mixed orderings generate fractions, e.g.  $\{0|1\} = \{\{\emptyset|\emptyset\}|\{\{\emptyset|\emptyset\}|\emptyset\}\} = \frac{1}{2}$ . Games with infinity (written  $\omega$ ) or infinitesimals ( $\epsilon = 1/\omega$ ) permit irrationals, and thus all reals. Further, it is valid to define  $\{\omega|\emptyset\} = \omega + 1$ , etc. Once arithmetic operations are defined, more complex games can define and *use* such quantities as  $\sqrt[3]{\omega}$  and  $\omega^\omega$ .

Therefore: by defining numbers as games, we can construct the **surreal numbers**.<sup>1</sup>

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As well as defining numbers as games, we can treat games like numbers.

$\{\emptyset|\emptyset\}$  can be played as the *zero game*. Simply: player 1 cannot move, and loses. Any game where player 2 has a winning strategy is equivalent to the zero game. Take two games  $G$  and  $H$ , with  $G$  a 2nd player win. The player with a winning strategy in  $H$  can treat the games separately, only moving in  $G$  to respond to the opponent. Player 2 wins  $G$ , but does not affect  $H$ ’s outcome. Conversely, given  $G'$  and  $H'$ , with  $G'$  a 1st player win, player 1 is last to move in  $G'$ , giving player 2 an extra move in  $H'$ , potentially altering its outcome. In terms of outcomes, we say  $G + H = 0 + H \rightarrow G = 0$ .

A game  $-G$  is  $G$  with the roles reversed, as by turning the board around in chess.

$G = H$  if  $G + (-H) = 0$ , i.e. is a player 2 win, and so equivalent to the zero game.

Two more properties are obvious:  $G + (H + K) = (G + H) + K$  (associativity) and  $G + H = H + G$  (commutativity).

We can see that  $G + (-G) = 0$  by a clever example called the Tweedledum & Tweedledee argument. In the game Blue-Red Hackenbush, players are given a drawing composed of separate edges. Each turn, player 1 removes a blue edge, plus any other edges no longer connected to the ground, and player 2 does likewise for red edges. Since in  $G + (-G)$  the number of pieces is the same, player 2 can just copy the moves of player 1 until all pieces are taken. Player 1 will not be able to move, and so will lose. Hence player 2 has a winning strategy.

Thus games are a proper mathematical object, namely: an Abelian group.<sup>10</sup>

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A new notation links all this to surreal numbers. For any set of games  $G^L$  and  $G^R$ , there exists a game  $G = \{G^L|G^R\}$ . Intuitively, the Left player moves in any game in  $G^L$ , and likewise for Right in  $G^R$ . The zero game  $\{\emptyset|\emptyset\} = 0$  is valid, and we may construct the surreals as before. Now we can write the surreals more easily with sets:  $\{1, 2, \dots, n|\} = n + 1$ .

To ‘play’ a surreal number  $\{n|m\}$ , a player chooses a game on their side, which the next player then plays. A player with no choice ( $\emptyset$ ) loses. If a player selects the zero game, the next player is the first to play it, and loses.

Consider a game  $\{1|\emptyset\}$ . If Left moves first, they choose  $1 = \{0|\emptyset\}$  from  $G^L$ ; Right loses. Right moving first is an instant lose. By extension, any game with positive surreal number is a win for Left. The converse holds for negative games  $\{\emptyset|n\}$ . Hence we say:  $G > 0$  if Left has a winning strategy,  $G < 0$  if Right does. Order does not matter.

Yet, recall the class of games  $G = 0$  where player 2 always wins: order matters, Left and Right do not. There is one more case: player 1 always wins. By elimination, such games are not positive, negative, or zero—so are not numbers. We say they are *fuzzy*:  $G \parallel 0$ . One such game is *star*, denoted  $*$  =  $\{0|0\}$ : player 1 chooses the zero game, and thus wins.

The surreal numbers are a subset of games where no member of  $G^L$  has value greater or equal to any member of  $G^R$ . In games violating this rule, players eventually reach  $*$ . Thus there exist more games than numbers—infinitely more. Games such as 0,  $*$ , or  $\{n|n\}$  that are symmetric for both players are called *impartial*. We will turn to these first.

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Combinatorial games have two players who take turns. No information is hidden, ruling out poker. As for probabilities, no dice. A player who can't move loses, and there are no draws (unlike chess). Last, the game ends.

The most famous impartial game is Nim, after the German *nimm*: to take.<sup>11</sup> Given zero or more piles, a player takes one or more (or all) stones from one pile at each turn. When no stones remain, the next player cannot move, and loses. One-pile Nim is trivial: player 1 takes the whole pile to win. Yet, how to win for multiple piles  $\{x, y, z, \dots\}$  is not clear.

Consider the case of two piles with  $n$  stones,  $\{n, n\}$ . Player 2 can win by a Tweedledum & Tweedledee strategy of copying the moves of player 1. Similarly, Player 1 can win  $\{n, m\}$  by taking from  $m$  so that the piles are equal:  $\{n, n\}$ .

For more than two piles, we need a more general tool: the *numbers*. Numbers uniquely express each positive integer as a sum of powers of two multiplied by zero or one, e.g.  $1 = 1 \times 2^0$ , which we write as (1). Likewise,  $2 = 1 \times 2^1 + 0 \times 2^0 = (10)$ ,  $3 = 1 \times 2^1 + 1 \times 2^0 = (11)$ ,  $4 = 1 \times 2^2 + 0 \times 2^1 + 0 \times 2^0 = (100)$ ... Each number counts stones in a nim-pile.

We denote a *nim-sum* by the operator  $\oplus$ , a form of addition with the rule  $1 + 1 = 0$ . This is called modulo arithmetic (mod 2), meaning: after dividing by 2, a nim-sum with remainder 0 is 0, one with remainder 1 is 1.<sup>100</sup> Thus:

$$\begin{array}{rcl} 10 & = & 2 \\ 101 & = & 5 \\ \oplus 1001 & = & 9 \\ \hline 1110 & = & 14 \end{array} \qquad \begin{array}{rcl} 1100 & = & 12 \\ 101 & = & 5 \\ \oplus 1001 & = & 9 \\ \hline 0000 & = & 0 \end{array}$$

A nim-sum usually differs from the sum of the corresponding integers, unless the numbers do not share any 1's in the same position, e.g.  $(010) \oplus (101) = (111) = 7 = 2 + 5$ ; such cases without modulo addition are called disjunctive sums.

A number added to itself is zero. For two piles, the Tweedledum & Tweedledee strategy amounts to keeping the piles' nim-sum at zero. In fact, this generalizes for  $n$  piles. To see why, note that the endpoint, a nim-heap without stones, has nim-sum 0. We call this a **P**-position, since the **P**revious player (the one who just moved to it) wins. Any position that can reach a **P**-position in one move is called an **N**-position, since the **N**ext player can move to the **P**-position.

Further, *any* position with nim-sum zero is a **P**-position, or 'safe position', since no player can move from a **P**-position to a **P**-position, but must move to an **N**-position. For two piles: given a **P**-position  $\{n, n\}$ , the next player can only make the piles unequal, i.e. nonzero nim-sum. For  $n$  piles, **P**-positions' safety is guaranteed by numbers' uniqueness. Consider: if, out of three numbers with nim-sum 0 we know two, we can always figure out the third: there's only one it can be. So if  $n$  piles nim-sum to zero and a player only takes from one pile, the new nim-sum *must* be nonzero.<sup>101</sup>

Since the endpoint is a **P**-position, the player who can first reach a **P**-position wins by moving from **P**-position to **P**-position up to the endpoint. If the initial  $n$  piles have nonzero nim-sum, player 1 wins; if zero, then player 2 wins.<sup>110</sup>

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The **Sprague-Grundy theorem** states that any impartial game is equivalent to a nim-pile of appropriate size. The size is determined recursively by the SG-function, which we may write in two equivalent forms:

$$g(x) = \min\{n \geq 0 : n \neq g(y), \forall y \in F(x)\} \Leftrightarrow g(x) = \text{mex}\{g(y) : y \in F(x)\}$$

That is, from any  $x$  (except an endpoint) we can move to one or more points in the set of followers  $F(x)$ , each with its own SG-value  $g(y)$ ; listing those  $g(y)$ 's in order, the SG-value for the point  $x$  is the smallest number, zero or greater, not in that list. To avoid spelling this out, we define **mex** ('minimal excludant') as the smallest non-negative integer not in the set. If  $x$  is an endpoint,  $F(x) = \emptyset$ , and so the **mex** is just 0. If another  $x$  can (only) reach the endpoint, now called  $y$ , in one move then  $y$  is in  $F(x)$ , and since  $g(y) = 0$ , the new **mex** is 1. And so on. Clearly, the SG-value is unique.



We may generalize  $\uparrow$  and  $\downarrow$ . Define *Tiny* as  $\mathfrak{X}_x = \{0|\{0|-x\}\}$  with  $x \geq 0$ , and is positive; its (negative) opposite is *Miny*,  $\mathfrak{X}_x = \{\{x|0\}|0\}$  for  $x \geq 0$ . Setting  $x = 0$  gives  $\mathfrak{X}_0 = \{0|\{0|0\}\} = \{0|*\} = \uparrow$  and  $\mathfrak{X}_0 = \{\{0|0\}|0\} = \{*\}|0 = \downarrow$ . In *Domineering*, Left (Vera) puts dominos vertically on a grid, Right (Horace) does so horizontally. Games Vera wins have values less than 0, Horace's are greater than 0. In a  $2 \times 4$  game, Vera moving first makes a game where player 2 wins, i.e. 0. In Horace's game, Vera going first makes a zero game, *or* Horace wins and leaves games valued  $-1 + -1 = -2$ . Thus:

$$\begin{array}{|c|c|} \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline \end{array} = \left\{ \begin{array}{|c|c|} \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline \end{array} \mid \begin{array}{|c|c|} \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline & \\ \hline \end{array} \right\} = \{0|\{0|-2\}\} = \mathfrak{X}_2$$

Partizan games have a *temperature*. In a *hot game* players improve their position by moving. *Cold games* worsen players' position each move. A player in a cold position who would prefer not to move at all (but must) is 'in *zugzwang*'. Domineering is hot, numbers are cold. Cold games are expressible as surreal numbers, hot games as pseudo-numbers. In Go jargon, a move that increases a game's temperature has *sente* (先手), one decreasing it has *gote* (後手). Games can be cooled by *taxing* moves (e.g.  $G - t$ ), and *thermographs* show how each move's temperature alters as  $t$  increases.<sup>1001</sup>

Now we may say: the theory of games is *combinatorial* in that it observes the void's play with every combination of itself.

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The philosopher Alain Badiou claims that the surreals let us speak of 'Number' as a unified concept, and open the way to "thinking number as a unified figure of multiple-being" (2008: 107), but he disregards their roots in combinatorial game theory. Badiou's main idea is that "mathematics is ontology" (ibid., 122). Games help show its implicit Platonism: it is considered bad form to include a term in its own definition, but this is just what occurs with the little word 'is'.

Rather, combinatorial games support the idea that numbers do not strictly 'exist'—mathematical fictionalism. They show also that games are primordial to numbers. However, games of the void with itself take the form of (and so presuppose) the concept of set supplemented with that of ordering, so as to form a tuple. The notion of difference is thus required as a catalyst, perhaps with  $\{\emptyset|\emptyset\}$  as a form of 'difference-in-itself'. Difference is play, and games are its syntax.

As combinatorial games show, games need not *exist* to be played. More strictly, players never play a *game*, but play the play of the void. Games' affinity to numbers explains why game theory is so often useful in science. Some ask what makes such models *games*. Answer: there is only one game, and mathematical models are instances of it.

Therefore: if mathematics is ontology, game theory is the science of the void.

Current research on combinatorial games aims to expand their standard definition, which still excludes many ordinary games, such as those with over two players. Moreover, while no bridge from combinatorial games to economic games has yet been made, research on cooperative and incomplete information combinatorial games shows promise.

The standard definition also leads to the *hypergame paradox* (Zwicker, 1987). In hypergame, player 1 chooses any game at all for player 2 to play. The problem is that players can keep choosing hypergame so that the game never terminates. So no formal system of games can include hypergame, which is counterintuitive. In practice, theorists define a class of non-terminating *loopy games* whose graphical representations have loops. These, in turn, open up to *transfinite games*.

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Applications of combinatorial games are few, but do exist. *Lexicodes* are a type of error-correcting code where “code words correspond to winning positions in certain impartial games,” so that if the underlying game is known, it can be solved to reveal the message (Milvang-Jensen, 2000: 91). Several well-known codes (e.g. Hamming codes) are cases of lexicodes. Games have also been used in computer science to describe types of data structures (Burke, 2015).<sup>1010</sup>

By far the most interesting (potential) application is *surreal analysis*. Like real analysis, the aim is to provide foundations for calculus on surreal numbers. Many functions used in physics diverge to infinity, giving nonsensical answers to well-posed questions; hence mathematicians hope surreal analysis will let them rigorously work *with* these infinities. The main obstacle is that no one has yet been able to satisfactorily define integration (Matthews, 1995).

Perhaps someday all of these research programmes may converge to a new foundation for science—a glass bead game.

Nevertheless, the proposition “All the world’s a game” is not well-formed. We know that all combinatorial games can be reduced to tuples of the void, and like surreal numbers, they do not *exist*. If all the world is a game, then there is no ‘world’ to be a game. Conversely, if everything is fully describable in game theoretic terms, the ‘is’ remains undefined.

Thus: Existence is itself played by an *arché-game* prior to difference and ontology. We call this the One-in-One.

## Notes

<sup>1</sup>Practically, we write the surreals as  $\{L|R\}$ , with  $L$  and  $R$  as sets of numbers. Surreal numbers obey the rule: no member of  $L$  is  $\geq$  any member of  $R$ , i.e.  $\ell \not\geq r (\forall \ell \in L)(\forall r \in R)$ . For  $x = \{X^L|X^R\}$  and  $y = \{Y^L|Y^R\}$ , we say  $x \leq y$  iff no  $x^L$  in  $X^L$  is  $\geq y$  and no  $y^R$  in  $Y^R$  is  $\leq x$ . Then  $x = y$  iff  $x \geq y$  and  $x \leq y$ . Define addition as  $x + y = \{X^L + y, x + Y^L | X^R + y, x + Y^R\}$ , e.g.  $\frac{1}{2} + \frac{1}{2} = \{0|1\} + \{0|1\} = \{0 + \frac{1}{2}, \frac{1}{2} + 0 | 1 + \frac{1}{2}, 1 + \frac{1}{2}\} = \{\frac{1}{2}|\frac{3}{2}\} = 1$ . For the last equality, note:  $\frac{1}{2} \not\geq 1$  and  $\frac{3}{2} \not\leq \frac{1}{2} + \frac{1}{2}$ , so  $\{\frac{1}{2}|\frac{3}{2}\} \leq 1$ . Next, writing  $1 = \{0|\}$ , note:  $0 \not\geq \frac{1}{2} + \frac{1}{2}$  and  $\frac{3}{2} \not\leq 1$ , so  $1 \leq \{\frac{1}{2}|\frac{3}{2}\}$ . Thus  $\{\frac{1}{2}|\frac{3}{2}\} = 1$ . For a step-by-step derivation from first principles, see Knuth's quasi-novel *Surreal Numbers* (1974) which coined their name.

<sup>10</sup>This is from Nowakowski (2009: 7-8). A *group* has a set and an operation combining two elements to get a third. For games, the operation is addition. Games count as a group since they have an identity element ( $G + 0 = G$ , so  $G$  is identical), have an inverse ( $-G$ ), are closed (any  $G + H$  is a game), and are associative. Since games are commutative, they are Abelian.

<sup>11</sup>Nim was formalized by Bouton in 1902. The origin of its name is confirmed by Walsh (1953).

<sup>100</sup>As well as nim-sums, one can also do nim-multiplication; thus, in the terminology of group theory, the numbers are a *field* (cf. Ferguson, 2014).

<sup>101</sup>Another way to think about uniqueness is that numbers make a closed system: if  $a \oplus b \oplus c = 0 \rightarrow a \oplus b = c$ . This can be rearranged for any missing number in the system. Systems of  $n$  numbers may be solved similarly.

<sup>110</sup>We've only covered normal play, where a player who can't move loses. An alternate convention is *misère play*: a player who can't move *wins*. In some simple misère games, such as Nim, misère and normal play are largely similar; yet, in the majority of cases misère games are far harder to solve. Notably, the SG-function does not always apply. Thanks goes to Thane Plambeck for confirming that any misère game is writable in surreal form, in keeping with the main argument of this paper.

<sup>111</sup>The two proofs are from Ferguson (2014), followed by Albert, Nowakowski & Wolfe's (2007) notation for numbers.

<sup>1000</sup>This example is from Demaine (2009). The next paragraph on  $\boxtimes_x$  and  $\boxtimes_x$  is from Uiterwijk & Barton (2015: 5).

<sup>1001</sup>Another partizan game is Hex, invented by John Nash. It has a rhombus-shaped board with hexagonal spaces. Each player places a stone on some hexagon, and the goal is to form a diagonal line reaching from one end to the other: player 1 uses white stones to form an off-diagonal line (bottom-left to top-right), player 2 wants a black diagonal line. Curiously, draws are impossible in Hex: this is guaranteed by Brouwer's fixed point theorem, the idea behind Nash equilibrium. Further, Nash showed player 1 has a winning strategy by a 'strategy-stealing' argument: if player 2 has an optimal strategy, player 1 can start with an arbitrary move, then copy all of player 2's moves. So player 1 can 'steal' player 2's optimal strategy, plus have the advantage of an extra stone on the board, and thus can win. Yet, this proof is non-constructive: we know an optimal strategy *exists*, but can't find it. The proof for Nash equilibrium is similarly non-algorithmic.

<sup>1010</sup>A variant of Nim can even be used to derive Chebyshev's inequality (Stolarsky, 1991). As well, algorithmic tools can be used to prove that certain combinatorial games are PSPACE-complete, EXPTIME-complete, or EXPSPACE-complete.

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## TL;DR

- Any number can be written as a tuple of games played by the void ( $\emptyset$ ) with itself.
- We can also treat games like numbers. In fact, numbers are a subclass of games.
- In Nim, players take stones from piles. To find a winning strategy, we invent *numbers*.
- Sprague-Grundy theorem: any impartial (symmetric) game is equivalent to a number.
- Solutions for partizan (asymmetric) games use various other pseudo-numbers.
- Mathematicians want to use surreal numbers for new foundations of mathematics.
- Surreal numbers, numbers, and pseudo-numbers can all be written as tuples of  $\emptyset$ .
- Therefore combinatorial games support the thesis that numbers do not truly 'exist'.
- It is invalid to say "All the world's a game," since games are solely composed of  $\emptyset$ .