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```

Basic

1

1

1.1 Increase Stack Size

```
//stack resize (linux)
#include <sys/resource.h>
 void increase_stack_size() {
   const rlim_t ks = 64*1024*1024;
   struct rlimit rl;
   int res=getrlimit(RLIMIT_STACK, &rl);
   if(res==0){
     if(rl.rlim_cur<ks){</pre>
       rl.rlim_cur=ks;
       res=setrlimit(RLIMIT_STACK, &rl);
1.2 Misc
 編譯參數: -std=c++14 -Wall -Wshadow (-fsanitize=
     undefined)
 //check special cases for example (n==1)
 //check size arrays
 #include <random>
mt19937 gen(0x5EED);
 int randint(int lb, int ub)
 { return uniform_int_distribution<int>(lb, ub)(gen); }
 #define SECs ((double)clock() / CLOCKS_PER_SEC)
 struct KeyHasher {
  size_t operator()(const Key& k) const {
     return k.first + k.second * 100000;
} };
 typedef unordered_map<Key,int,KeyHasher> map_t;
 __builtin_popcountll
                         //換成二進位有幾個1
__builtin_clzll
                         //返回左起第一個1之前0的個數
                         //返回1的個數的奇偶性
 __builtin_parityll
 1.3 check
for ((i=0;;i++))
     echo "$i"
     python3 gen.py > input
     ./ac < input > ac.out
     ./wa < input > wa.out
     diff ac.out wa.out || break
done
 1.4 python-related
 int(eval(num.replace("/","//")))
 from fractions import Fraction
 from decimal import Decimal, getcontext
 getcontext().prec = 250 # set precision
 itwo = Decimal(0.5)
 two = Decimal(2)
 N = 200
def angle(cosT):
    """given cos(theta) in decimal return theta"""
   for i in range(N):
   cosT = ((cosT + 1) / two) ** itwo
sinT = (1 - cosT * cosT) ** itwo
return sinT * (2 ** N)
pi = angle(Decimal(-1))
 2
     flow
 2.1 MinCostFlow
```

```
struct MinCostMaxFlow{
typedef int Tcost;
 static const int MAXV = 20010:
 static const int INFf = 1000000;
 static const Tcost INFc = 1e9;
 struct Edge{
```

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```
int v, cap;
     Tcost w;
     int rev;
     Edge(){}
     Edge(int t2, int t3, Tcost t4, int t5)
     : v(t2), cap(t3), w(t4), rev(t5) {}
  int V, s, t;
  vector<Edge> g[MAXV];
  void init(int n, int _s, int _t){
  V = n; s = _s; t = _t;
  for(int i = 0; i <= V; i++) g[i].clear();</pre>
  void addEdge(int a, int b, int cap, Tcost w){
   g[a].push_back(Edge(b, cap, w, (int)g[b].size()));
   g[b].push_back(Edge(a, 0, -w, (int)g[a].size()-1));
  Tcost d[MAXV];
  int id[MAXV], mom[MAXV];
  bool inqu[MAXV];
  queue<int> q;
  pair<int,Tcost> solve(){
     int mxf = 0; Tcost mnc = 0;
     while(1){
        fill(d, d+1+V, INFc);
       fill(inqu, inqu+1+V, 0);
fill(mom, mom+1+V, -1);
       mom[s] = s;
       d[s] = 0;
       q.push(s); inqu[s] = 1;
       while(q.size()){
          int u = q.front(); q.pop();
          inqu[u] = 0;
for(int i = 0; i < (int) g[u].size(); i++){</pre>
             Edge &e = g[u][i];
             int v = e.v;
             if(e.cap > 0 \& d[v] > d[u]+e.w){
               d[v] = d[u] + e.w;
               mom[v] = u;
               id[v] = i
               if(!inqu[v]) q.push(v), inqu[v] = 1;
        if(mom[t] == -1) break ;
       int df = INFf;
       for(int u = t; u != s; u = mom[u])
  df = min(df, g[mom[u]][id[u]].cap);
for(int u = t; u != s; u = mom[u]){
          Edge &e = g[mom[u]][id[u]];
          g[e.v][e.rev].cap += df;
       mxf += df;
       mnc += df*d[t];
     return {mxf,mnc};
} }flow;
2.2 Dinic
const int MXN = 10000;
struct Dinic{
  struct Edge{ int v,f,re; };
  int n,s,t,level[MXN];
  vector<Edge> E[MXN];
```

```
que.push(it.v);
    } } }
    return level[t] != -1;
  int DFS(int u, int nf){
    if (u == t) return nf;
    int res = 0;
     for (auto &it : E[u]){
       if (it.f > 0 && level[it.v] == level[u]+1){
         int tf = DFS(it.v, min(nf,it.f));
         res += tf; nf -= tf; it.f -= tf;
E[it.v][it.re].f += tf;
         if (nf == 0) return res;
    if (!res) level[u] = -1;
    return res;
  int flow(int res=0){
    while ( BFS() )
      res += DFS(s,2147483647);
     return res;
} }flow;
```

2.3 Kuhn Munkres 最大完美二分匹配

```
struct KM{ // max weight, for min negate the weights
   static const int MXN = 2001; // 1-based
  static const ll INF = 0x3f3f3f3f3f;
  int n, mx[MXN], my[MXN], pa[MXN];
ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
  bool vx[MXN], vy[MXN];
  void init(int _n) {
     n = _n;
     for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);</pre>
  void addEdge(int x, int y, ll w) {g[x][y] = w;}
void augment(int y) {
     for(int x, z; y; y = z)
x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
  void bfs(int st) {
     for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;</pre>
     queue<int> q; q.push(st);
     for(;;) {
       while(q.size()) {
          int x=q.front(); q.pop(); vx[x]=1;
          for(int y=1; y<=n; ++y) if(!vy[y]){</pre>
            li t = lx[x]+ly[y]-g[x][y];
            if(t==0){
              pa[y]=x
               if(!my[y]){augment(y);return;}
               vy[y]=1, q.push(my[y]);
            }else if(sy[y]>t) pa[y]=x,sy[y]=t;
       il cut = INF;
for(int y=1; y<=n; ++y)</pre>
          if(!vy[y]&&cut>sy[y]) cut=sy[y];
       for(int j=1; j<=n; ++j){</pre>
          if(vy[j]) ly[j] += cut;
          else sy[j] -= cut;
       for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){</pre>
          if(!my[y]){augment(y);return;}
         vy[y]=1, q.push(my[y]);
  } } }
  11 solve(){
     fill(mx, mx+n+1, 0); fill(my, my+n+1, 0); fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
     for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)</pre>
       lx[x] = max(lx[x], g[x][y]);
     for(int x=1; x<=n; ++x) bfs(x);</pre>
     for(int y=1; y<=n; ++y) ans += g[my[y]][y];</pre>
     return ans;
} }graph;
```

2.4 Max flow with lower/upper bound

```
// flow use ISAP
// Max flow with lower/upper bound on edges
// source = 1 , sink = n
```

```
int in[N], out[N];
int l[M], r[M], a[M], b[M];//0-base,a下界,b
int solve(){
  flow.init(n); //n為點的數量,m為邊的數量,點是1-
       base
  for( int i = 0 ; i < m ; i ++ ){
    in[ r[ i ] ] += a[ i ];
out[ l[ i ] ] += a[ i ];
flow.addEdge( l[ i ] , r[ i ] , b[ i ] - a[ i ] );
// flow from l[i] to r[i] must in [a[ i ], b[ i ]]
  int nd = 0;
  for( int i = 1 ; i <= n ; i ++ ){
  if( in[ i ] < out[ i ] ){</pre>
       flow.addEdge( i , flow.t , out[ i ] - in[ i ] );
nd += out[ i ] - in[ i ];
     if( out[ i ] < in[ i ] )</pre>
       flow.addEdge( flow.s , i , in[ i ] - out[ i ] );
  // original sink to source
  flow.addEdge( n , 1 , INF );
if( flow.maxflow() != nd )
   // no solution
     return -1;
  int ans = flow.G[ 1 ].back().c; // source to sink
  flow.G[ 1 ].back().c = flow.G[ n ].back().c = \emptyset;
  // take out super source and super sink
  for( size_t i = 0 ; i < flow.G[ flow.s ].size() ; i</pre>
       ++ ){
    flow.G[ flow.s ][ i ].c = 0;
Edge &e = flow.G[ flow.s ][ i ];
     flow.G[e.v][e.r].c = 0;
  for( size_t i = 0 ; i < flow.G[ flow.t ].size() ; i</pre>
     flow.G[ flow.t ][ i ].c = 0;
Edge &e = flow.G[ flow.t ][ i ];
     flow.G[ e.v ][ e.r ].c = 0;
  flow.addEdge( flow.s , 1 , INF
  flow.addEdge( n , flow.t , INF );
  flow.reset();
  return ans + flow.maxflow();
```

2.5 Flow Method

Maximize $c^T x$ subject to $Ax \le b$, $x \ge 0$;

with the corresponding symmetric dual problem,

```
Minimize b^T y subject to A^T y \geq c, y \geq 0.
Maximize c^T x subject to Ax \le b;
with the corresponding asymmetric dual problem,
Minimize b^T y subject to A^T y = c, y \ge 0.
Minimum vertex cover on bipartite graph =
Maximum matching on bipartite graph
找出最小點覆蓋,做完dinic之後,從源點dfs只走還有流量的
邊,紀錄每個點有沒有被走到,左邊沒被走到的點跟右邊被走
到的點就是答案
Maximum density subgraph ( \sum W_e + \sum W_v ) / |V|
Binary search on answer:
For a fixed D, construct a Max flow model as follow:
Let S be Sum of all weight( or inf)
1. from source to each node with cap = S
2. For each (u,v,w) in E, (u->v,cap=w), (v->u,cap=w)
3. For each node v, from v to sink with cap = S + 2 * D - deg[v] - 2 * (W of v)
where deg[v] = \sum weight of edge associated with v
If maxflow < S * IVI, D is an answer.
Requiring subgraph: all vertex can be reached from
    source with
edge whose cap > 0.
```

3 Math

3.1 FFT

```
// const int MAXN = 262144;
// (must be 2^k)
// before any usage, run pre_fft() first
typedef long double ld;
typedef complex<ld> cplx; //real() ,imag()
const ld PI = acosl(-1);
const cplx I(0, 1);
cplx omega[MAXN+1];
void pre_fft(){
  for(int i=0; i<=MAXN; i++)</pre>
    omega[i] = exp(i * 2 * PI / MAXN * I);
// n must be 2^k
void fft(int n, cplx a[], bool inv=false){
  int basic = MAXN / n;
  int theta = basic;
  for (int m = n; m >= 2; m >>= 1) {
    int mh = m >> 1;
for (int i = 0; i < mh; i++) {</pre>
      cplx w = omega[inv ? MAXN-(i*theta%MAXN)
                            : i*theta%MAXN];
      for (int j = i; j < n; j += m) {
        int k = j + mh;
         cplx x = a[j] - a[k];
        a[j] += a[k];
        a[k] = w * x;
    } }
    theta = (theta * 2) % MAXN;
  int i = 0;
  for (int j = 1; j < n - 1; j++) {
    for (int k = n >> 1; k > (i ^= k); k >>= 1); if (j < i) swap(a[i], a[j]);
  if(inv) for (i = 0; i < n; i++) a[i] /= n;
cplx arr[MAXN+1];
inline void mul(int _n,ll a[],int _m,ll b[],ll ans[])
  int n=1,sum=_n+_m-1;
  while(n<sum)</pre>
    n << =1;
  for(int i=0;i<n;i++)</pre>
    double x=(i<_n?a[i]:0), y=(i<_m?b[i]:0);
    arr[i]=complex<double>(x+y,x-y);
  fft(n,arr);
  for(int i=0;i<n;i++)</pre>
    arr[i]=arr[i]*arr[i];
  fft(n,arr,true);
  for(int i=0;i<sum;i++)</pre>
    ans[i]=(long long int)(arr[i].real()/4+0.5);
```

3.2 NTT

```
/ Remember coefficient are mod P
/* p=a*2^n+1
        2^n
                                     root
   n
        65536
                     65537
   16
                                     3 */
                    7340033
        1048576
// (must be 2^k)
template<LL P, LL root, int MAXN>
struct NTT{
  static LL bigmod(LL a, LL b) {
    LL res = 1;
    for (LL bs = a; b; b >>= 1, bs = (bs * bs) % P)
      if(b&1) res=(res*bs)%P;
    return res;
  static LL inv(LL a, LL b) {
    if(a==1)return 1;
    return (((LL)(a-inv(b\%a,a))*b+1)/a)\%b;
 LL omega[MAXN+1];
 NTT() {
    omega[0] = 1;
```

```
LL r = bigmod(root, (P-1)/MAXN);
for (int i=1; i<=MAXN; i++)
  omega[i] = (omega[i-1]*r)%P;</pre>
   // n must be 2^k
   void tran(int n, LL a[], bool inv_ntt=false){
      int basic = MAXN / n , theta = basic;
for (int m = n; m >= 2; m >>= 1) {
          int mh = m >> 1;
for (int i = 0; i < mh; i++) {
  LL w = omega[i*theta%MAXN];</pre>
             for (int j = i; j < n; j += m) {
    int k = j + mh;
    LL x = a[j] - a[k];
    if (x < 0) x += P;
                a[j] += a[k];
                if (a[j] > P) a[j] -= P;
a[k] = (w * x) % P;
          theta = (theta * 2) % MAXN;
       int i = 0;
      for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
          if (j < i) swap(a[i], a[j]);
      if (inv_ntt) {
          LL ni = inv(n,P);
         reverse( a+1 , a+n );
for (i = 0; i < n; i++)
a[i] = (a[i] * ni) % P;
  }
const LL P=2013265921,root=31;
const int MAXN=4194304;
NTT<P, root, MAXN> ntt;
```

3.3 Fast Walsh Transform

```
/* xor convolution:
 * x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
 * x' = ( x0+x1 , x0-x1 ) , y' = ( y0+y1 , y0-y1 )

* z' = ( ( x0+x1 )( y0+y1 ) , ( x0-x1 )( y0-y1 ) )

* z = (1/2) * z''
 * or convolution:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
* x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div */typedef long long LL; const int MAXN = (1<<20)+10;
const LL MOD = 1e9+7;
inline LL inv( LL x ) {
   return mypow( x , MOD-2 );
inline void fwt( LL x[ MAXN ] , int N , bool inv=0 ) {
   for( int d = 1 ; d < N ; d <<= 1 ) {
       int d2 = d << 1;
      for( int s = 0 ; s < N ; s += d2 )
  for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
    LL ta = x[i], tb = x[j];
    x[i] = ta+tb;
    x[j] = ta-tb;
    if( x[i] > MOD > x[i] > MOD > x[i]
              if( x[ i ] >= MOD ) x[ i ] -= MOD;
if( x[ j ] < 0 ) x[ j ] += MOD;</pre>
    if( inv )
       for( int i = 0 ; i < N ; i++ ) {
    x[ i ] *= inv( N );
    x[ i ] %= MOD;</pre>
}
```

3.4 Poly operator

```
struct PolyOp {
#define FOR(i, c) for (int i = 0; i < (c); ++i)
   NTT<P, root, MAXN> ntt;
   static int nxt2k(int x) {
   int i = 1; for (; i < x; i <<= 1); return i;</pre>
```

```
// c[i]=sum{j=0~i}a[j]*b[i-j] -> c[i+j]+=a[i]*b[j](加
      分卷積)
// if c[i-j]+=a[i]*b[j] (減法卷積)
// (轉換成加法捲積) -> reverse(a); c=mul(a,b);
reverse( c );
void Mul(int n, LL a[], int m, LL b[], LL c[]) {
  static LL aa[MAXN], bb[MAXN];
   int N = nxt2k(n+m);
   copy(a, a+n, aa); fill(aa+n, aa+N, 0); copy(b, b+m, bb); fill(bb+m, bb+N, 0);
   ntt.tran(N, aa); ntt.tran(N, bb);
  FOR(i, N) c[i] = aa[i] * bb[i] % P;
ntt.tran(N, c, 1);
void Inv(int n, LL a[], LL b[]) {
   // ab = aa^-1 = 1 mod x^(n/2)
   // (b - a^-1)^2 = 0 mod x^n
   // bb - a^2 + 2 ba^1 = 0
   \frac{1}{1} bba - \frac{1}{1} a^-1 + 2b = 0
   // bba + 2b = a^{-1}
   static LL tmp[MAXN];
   if (n == 1) {b[0] = ntt.inv(a[0], P); return;}
Inv((n+1)/2, a, b);
   int N = nxt2k(n*2);
   copy(a, a+n, tmp);
fill(tmp+n, tmp+N, 0);
   fill(b+n, b+N, 0);
   ntt.tran(N, tmp); ntt.tran(N, b);
   FOR(i, N) {
      LL t1 = (2 - b[i] * tmp[i]) % P;
     if (t1 < 0) t1 += P;
b[i] = b[i] * t1 % P;
   ntt.tran(N, b, 1);
fill(b+n, b+N, 0);
void Div(int n, LL a[], int m, LL b[], LL d[], LL r
      []) {
   // Ra = Rb * Rd mod x^(n-m+1)
   // Rd = Ra * Rb^{-1} mod
   static LL aa[MAXN], bb[MAXN], ta[MAXN], tb[MAXN];
   if (n < m) {copy(a, a+n, r); fill(r+n, r+m, 0);</pre>
         return;}
   // d: n-1 - (m-1) = n-m (n-m+1 terms)
copy(a, a+n, aa); copy(b, b+m, bb);
   reverse(aa, aa+n); reverse(bb, bb+m);
  Inv(n-m+1, bb, tb);
Mul(n-m+1, ta, n-m+1, tb, d);
  fill(d+n-m+1, d+n, 0); reverse(d, d+n-m+1);

// r: m-1 - 1 = m-2 (m-1 terms)

Mul(m, b, n-m+1, d, ta);
   FOR(i, n) \{ r[i] = a[i] - ta[i]; if (r[i] < 0) r[i] \}
          += P; }
void dx(int n, LL a[], LL b[]) { REP(i, 1, n-1) b[i
    -1] = i * a[i] % P; }
void Sx(int n, LL a[], LL b[]) {
   b \lceil 0 \rceil = 0;
   FOR(i, n) b[i+1] = a[i] * ntt.inv(i+1, P) % P;
void Ln(int n, LL a[], LL b[]) {
   // Integral a' a^-1 dx
   static LL a1[MAXN], a2[MAXN], b1[MAXN];
   int N = nxt2k(n*2)
   dx(n, a, a1); Inv(n, a, a2);
   Mul(n-1, a1, n, a2, b1);
Sx(n+n-1-1, b1, b);
   fill(b+n, b+N, 0);
void Exp(int n, LL a[], LL b[]) {
   // Newton method to solve g(a(x)) = \ln b(x) - a(x)
        = 0
  = 0
// b' = b - g(b(x)) / g'(b(x))
// b' = b (1 - lnb + a)
static LL lnb[MAXN], c[MAXN], tmp[MAXN];
assert(a[0] == 0); // dont know exp(a[0]) mod P
if (n == 1) {b[0] = 1; return;}
   Exp((n+1)/2, a, b);
fill(b+(n+1)/2, b+n, 0);
   Ln(n, b, lnb);
```

```
fill(c, c+n, 0); c[0] = 1;
FOR(i, n) {
    c[i] += a[i] - lnb[i];
    if (c[i] < 0) c[i] += P;
    if (c[i] >= P) c[i] -= P;
}
Mul(n, b, n, c, tmp);
copy(tmp, tmp+n, b);
}
polyop;
```

3.5 O(1)mul

```
LL mul(LL x,LL y,LL mod){
  LL ret=x*y-(LL)((long double)x/mod*y)*mod;
  // LL ret=x*y-(LL)((long double)x*y/mod+0.5)*mod;
  return ret<0?ret+mod:ret;
}</pre>
```

3.6 Linear Recurrence

```
// Usage: linearRec({0, 1}, {1, 1}, k) //k'th fib
typedef vector<ll> Poly;
//S:前i項的值,tr:遞迴系數,k:求第k項
ll linearRec(Poly& S, Poly& tr, ll k) {
  int n = tr.size()
  auto combine = [&](Poly& a, Poly& b) {
  Poly res(n * 2 + 1);
    rep(i,0,n+1) rep(j,0,n+1)
       res[i+j]=(res[i+j] + a[i]*b[j])%mod;
     for(int i = 2*n; i > n; --i) rep(j,0,n)
       res[i-1-j]=(res[i-1-j] + res[i]*tr[j])%mod;
    res.resize(n + 1);
    return res;
 Poly pol(n + 1), e(pol);
pol[0] = e[1] = 1;
for (++k; k; k /= 2) {
    if (k % 2) pol = combine(pol, e);
    e = combine(e, e);
  ll res = 0;
  rep(i,0,n) res=(res + pol[i+1]*S[i])%mod;
  return res;
```

3.7 Miller Rabin

```
// n < 4,759,123,141
                              3: 2, 7, 61
// n < 1,122,004,669,633
                                   2, 13, 23, 1662803
// n < 3,474,749,660,383
                                         pirmes <= 13
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
LL magic[]={}
bool witness(LL a, LL n, LL u, int t){
  if(!a) return 0;
  LL x=mypow(a,u,n);
  for(int i=0;i<t;i++) {</pre>
    LL nx=mul(x,x,n);
    if(nx==1&&x!=1&&x!=n-1) return 1;
    x=nx;
  return x!=1;
bool miller_rabin(LL n) {
  int s=(magic number size)
  // iterate s times of witness on n
  if(n<2) return 0;</pre>
  if(!(n\&1)) return n == 2;
  ll u=n-1; int t=0;
// n-1 = u*2^t
  while(!(u&1)) u>>=1, t++;
  while(s--){
    LL a=magic[s]%n;
    if(witness(a,n,u,t)) return 0;
  return 1;
```

3.8 Faulhaber $(\sum_{i=1}^{n} i^{p})$

```
/* faulhaber's formula - 
 * cal power sum formula of all p=1\simk in O(k^2) */
#define MAXK 2500
const int mod = 1000000007;
int b[MAXK]; // bernoulli number
int inv[MAXK+1]; // inverse
int cm[MAXK+1][MAXK+1]; // combinactories
int co[MAXK][MAXK+2]; // coeeficient of x^j when p=i
inline int getinv(int x) {
  int a=x,b=mod,a0=1,a1=0,b0=0,b1=1;
  while(b) {
     int q,t;
     q=a/b; t=b; b=a-b*q; a=t;
t=b0; b0=a0-b0*q; a0=t;
     t=b1; b1=a1-b1*q; a1=t;
  return a0<0?a0+mod:a0;
inline void pre() {
   /* combinational */
  for(int i=0;i<=MAXK;i++) {</pre>
     cm[i][0]=cm[i][i]=1;
     for(int j=1;j<i;j++)
  cm[i][j]=add(cm[i-1][j-1],cm[i-1][j]);</pre>
   /* inverse */
  for(int i=1;i<=MAXK;i++) inv[i]=getinv(i);</pre>
   /* bernoulli */
  b[0]=1; b[1]=getinv(2); // with b[1] = 1/2
  for(int i=2;i<MAXK;i++) {</pre>
     if(i&1) { b[i]=0; continue; }
     b[i]=1;
     for(int j=0;j<i;j++)</pre>
       b[i]=sub(b[i],
                  mul(cm[i][j],mul(b[j], inv[i-j+1])));
  /* faulhaber */
  // sigma_x=1~n \{x^p\} =
// 1/(p+1) * sigma_j=0~p \{C(p+1,j)*Bj*n^(p-j+1)\}
  for(int i=1;i<MAXK;i++) {</pre>
     co[i][0]=0;
     for(int_j=0;j<=i;j++)</pre>
       co[i][i-j+1]=mul(inv[i+1], mul(cm[i+1][j], b[j]))
  }
}
/* sample usage: return f(n,p) = sigma_x=1\sim (x^p) */
inline int solve(int n,int p) {
  int sol=0,m=n;
  for(int i=1;i<=p+1;i++) {</pre>
     sol=add(sol,mul(co[p][i],m));
     m = mul(m, n);
  return sol;
```

3.9 Chinese Remainder

3.10 Pollard Rho

```
// does not work when n is prime
LL f(LL x, LL mod){ return add(mul(x,x,mod),1,mod); }
LL pollard_rho(LL n) {
   if(!(n&1)) return 2;
   while(true){
      LL y=2, x=rand()%(n-1)+1, res=1;
      for(int sz=2; res==1; sz*=2) {
        for(int i=0; i<sz && res<=1; i++) {
            x = f(x, n);
            res = __gcd(abs(x-y), n);
        }
        y = x;
    }
   if (res!=0 && res!=n) return res;
} }</pre>
```

3.11 Josephus Problem

```
int josephus(int n, int m){ //n人每m次
  int ans = 0;
  for (int i=1; i<=n; ++i)
      ans = (ans + m) % i;
  return ans;
}</pre>
```

3.12 Gaussian Elimination

```
const int GAUSS_MOD = 100000007LL;
struct GAUSS{
     int n;
     vector<vector<int>> v
     int ppow(int a , int k){
         if(k == 0) return 1;
         if(k % 2 == 0) return ppow(a * a % GAUSS_MOD ,
             k >> 1);
         if(k % 2 == 1) return ppow(a * a % GAUSS_MOD ,
             k \gg 1) * a % GAUSS_MOD;
     vector<int> solve(){
         vector<int> ans(n);
         REP(now , 0 , n){
    REP(i , now , n) if(v[now][now] == 0 && v[i ][now] != 0)

             swap(v[i] , v[now]); // det = -det;
if(v[now][now] == 0) return ans;
             int inv = ppow(v[now][now] , GAUSS_MOD - 2)
             REP(i, 0, n) if(i != now){
                  int tmp = v[i][now] * inv % GAUSS_MOD;
                  GAUSS_MOD) %= GAUSS_MOD;
             }
             i , 0 , n) ans[i] = v[i][n + 1] * ppow(v[i
][i] , GAUSS_MOD - 2) % GAUSS_MOD;
         return ans;
     // gs.v.clear() , gs.v.resize(n , vector<int>(n + 1
          , 0));
} gs;
```

3.13 ax+by=gcd

```
PII gcd(int a, int b){
   if(b == 0) return {1, 0};
   PII q = gcd(b, a % b);
   return {q.second, q.first - q.second * (a / b)};
}
```

3.14 Discrete sqrt

```
void calcH(LL &t, LL &h, const LL p) {
   LL tmp=p-1; for(t=0;(tmp&1)==0;tmp/=2) t++; h=tmp;
}
// solve equation x^2 mod p = a
bool solve(LL a, LL p, LL &x, LL &y) {
   if(p == 2) { x = y = 1; return true; }
   int p2 = p / 2, tmp = mypow(a, p2, p);
   if (tmp == p - 1) return false;
```

```
if ((p + 1) % 4 == 0) {
    x=mypow(a,(p+1)/4,p); y=p-x; return true;
} else {
    LL t, h, b, pb; calcH(t, h, p);
    if (t >= 2) {
        do {b = rand() % (p - 2) + 2;
        } while (mypow(b, p / 2, p) != p - 1);
        pb = mypow(b, h, p);
} int s = mypow(a, h / 2, p);
for (int step = 2; step <= t; step++) {
        int ss = (((LL)(s * s) % p) * a) % p;
        for(int i=0;i<t-step;i++) ss=mul(ss,ss,p);
        if (ss + 1 == p) s = (s * pb) % p;
        pb = ((LL)pb * pb) % p;
        pb = ((LL)s * a) % p; y = p - x;
} return true;
}</pre>
```

3.15 Prefix Inverse

```
void solve( int m ){
  inv[ 1 ] = 1;
  for( int i = 2 ; i < m ; i ++ )
    inv[ i ] = ((LL)(m - m / i) * inv[m % i]) % m;
}</pre>
```

3.16 Roots of Polynomial 找多項式的根

```
const double eps = 1e-12;
const double inf = 1e+12;
double a[ 10 ], x[ 10 ]; // a[0..n](coef) must be
     filled
int n; // degree of polynomial must be filled
int sign( double x ){return (x < -eps)?(-1):(x>eps);}
double f(double a\square, int n, double x){
  double tmp=1,sum=0;
  for(int i=0;i<=n;i++)</pre>
  { sum=sum+a[i]*tmp; tmp=tmp*x; }
  return sum;
double binary(double 1, double r, double a∏, int n){
  int sl=sign(f(a,n,l)),sr=sign(f(a,n,r));
  if(sl==0) return l; if(sr==0) return r;
if(sl*sr>0) return inf;
  while(r-l>eps){
    double mid=(1+r)/2;
    int ss=sign(f(a,n,mid));
    if(ss==0) return mid;
    if(ss*sl>0) l=mid; else r=mid;
  }
  return 1;
}
void solve(int n,double a[],double x[],int &nx){
  if(n==1){x[1]=-a[0]/a[1]; nx=1; return; }
  double da[10], dx[10]; int ndx;
  for(int i=n;i>=1;i--) da[i-1]=a[i]*i;
  solve(n-1,da,dx,ndx);
  nx=0:
  if(ndx==0){
    double tmp=binary(-inf,inf,a,n);
    if (tmp<inf) x[++nx]=tmp;</pre>
    return;
  double tmp;
  tmp=binary(-inf,dx[1],a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
  for(int i=1;i<=ndx-1;i++){</pre>
    tmp=binary(dx[i],dx[i+1],a,n);
    if(tmp<inf) x[++nx]=tmp;</pre>
  tmp=binary(dx[ndx],inf,a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
} // roots are stored in x[1..nx]
```

3.17 Primes

```
/* 12721, 13331, 14341, 75577, 123457, 222557, 556679
* 999983, 1097774749, 1076767633, 100102021, 999997771
* 1001010013, 1000512343, 987654361, 999991231
* 999888733, 98789101, 987777733, 999991921, 1010101333
* 1010102101, 1000000000039, 10000000000037
* 2305843009213693951, 4611686018427387847
```

```
* 9223372036854775783, 18446744073709551557 */
int mu[N], p_tbl[N];
vector<int> primes;
void sieve() {
   mu[1] = p_tbl[1] = 1;
   for( int i = 2 ; i < N ; i ++ ){
  if( !p_tbl[ i ] ){
    p_tbl[ i ] = i;
}</pre>
         primes.push_back( i );
         mu[i] = -1;
      for( int p : primes ){
         int x = i * p;
         if( x >= M ) break;
         p_{tbl}[x] = p;
         mu[x] = -mu[i];
if(i%p==0){
mu[x]=0;
} } } }
vector<int> factor( int x ){
   vector<int> fac{ 1 };
  while( x > 1 ){
  int fn = SZ(fac), p = p_tbl[ x ], pos = 0;
      while( x \% p == 0 ){
         x /= p;
for( int i = 0 ; i < fn ; i ++ )
fac.PB( fac[ pos ++ ] * p );
   } }
   return fac;
}
3.18 Result
    • Lucas' Theorem :
      For n,m\in\mathbb{Z}^* and prime P, C(m,n)\mod P=\Pi(C(m_i,n_i)) where m_i is the i\text{-th} digit of m in base P.
   • Stirling approximation :
      n! \approx \sqrt{2\pi n} \left(\frac{n}{e}\right)^n e^{\frac{1}{12n}}
   • Stirling Numbers(permutation |P|=n with k cycles):
      S(n,k) = \text{coefficient of } x^k \text{ in } \prod_{i=0}^{n-1} (x+i)
   • Stirling Numbers(Partition n elements into k non-empty set):
      S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n
   • Pick's Theorem : A=i+b/2-1 其面積 A 和內部格點數目 i 、邊上格點數目 b 的關係
   • Catalan number : C_n = {2n \choose n}/(n+1)
      C_n^{n+m} - C_{n+1}^{n+m} = (m+n)! \frac{n-m+1}{n+1} for n \ge m
      C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{(n+1)!n!}
      \begin{array}{lll} C_0 = 1 & and & C_{n+1} = 2(\frac{2n+1}{n+2})C_n \\ C_0 = 1 & and & C_{n+1} = \sum_{i=0}^n C_i C_{n-i} & for & n \geq 0 \end{array}
    • Euler Characteristic:
      planar graph: V-E+F-C=1 convex polyhedron: V-E+F=2
      V,E,F,C: number of vertices, edges, faces(regions), and compo-
      nents
   • Kirchhoff's theorem :
      A_{ii} = deg(i), A_{ij} = (i, j) \in E ? -1 : 0, Deleting any one row, one
      column, and cal the det(A)
   • Polya' theorem (c 為方法數,m 為總數): (\sum_{i=1}^m c^{\gcd(i,m)})/m
   • 錯排公式: (n 個人中,每個人皆不再原來位置的組合數):
      dp[0] = 1; dp[1] = 0;
       dp[i] = (i-1) * (dp[i-1] + dp[i-2]);
   • Bell 數 (有 n 個人, 把他們拆組的方法總數):
      B_0 = 1
      B_n = \sum_{k=0}^{n} s(n, k) \quad (second - stirling)
      B_{n+1} = \sum_{k=0}^{n} \binom{n}{k} B_k
   • Wilson's theorem :
      (p-1)! \equiv -1 \pmod{p}
```

• Fermat's little theorem : $a^p \equiv a (mod \ p)$

Geometry

4.1 definition

```
typedef long double ld;
const ld eps = 1e-8;
int dcmp(ld x) {
  if(abs(x) < eps) return 0;</pre>
  else return x < 0 ? -1 : 1;
}
struct Pt {
  ld x, y;
Pt(ld _x=0, ld _y=0):x(_x), y(_y) {}
  Pt operator+(const Pt &a) const {
    return Pt(x+a.x, y+a.y);
  Pt operator-(const Pt &a) const {
    return Pt(x-a.x, y-a.y);
  Pt operator*(const ld &a) const {
    return Pt(x*a, y*a);
  Pt operator/(const ld &a) const {
    return Pt(x/a, y/a);
  ld operator*(const Pt &a) const {
    return x*a.x + y*a.y;
  ld operator^(const Pt &a) const {
    return x*a.y - y*a.x;
  bool operator<(const Pt &a) const {</pre>
    return x < a.x | | (x == a.x \&\& y < a.y);
    //return dcmp(x-a.x) < 0 \mid \mid (dcmp(x-a.x) == 0 \&\&
         dcmp(y-a.y) < 0);
  bool operator==(const Pt &a) const {
    return dcmp(x-a.x) == 0 \&\& dcmp(y-a.y) == 0;
ld norm2(const Pt &a) {
  return a*a;
ld norm(const Pt &a) {
  return sqrt(norm2(a));
Pt perp(const Pt &a) {
  return Pt(-a.y, a.x);
Pt rotate(const Pt &a, ld ang) {
   return Pt(a.x*cos(ang)-a.y*sin(ang), a.x*sin(ang)+a.y
       *cos(ang));
struct Line {
  Pt s, e, v; // start, end, end-start
  ld ang;
  Line(Pt_s=Pt(0, 0), Pt_e=Pt(0, 0)):s(_s), e(_e) { v
        = e-s; ang = atan2(v.y, v.x); }
  bool operator<(const Line &L) const {</pre>
    return ang < L.ang;</pre>
  }
};
struct Circle {
  Pt o; ld r;
  Circle(Pt _o=Pt(0, 0), ld _r=0):o(_o), r(_r) {}
```

4.2 Intersection of 2 lines

```
Pt LLIntersect(Line a, Line b) {
  Pt p1 = a.s, p2 = a.e, q1 = b.s, q2 = b.e;
  ld f1 = (p2-p1)^{(q1-p1)}, f2 = (p2-p1)^{(p1-q2)}, f;
  if(dcmp(f=f1+f2) == 0)
    return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
  return q1*(f2/f) + q2*(f1/f);
}
```

4.3 halfPlaneIntersection

```
// for point or line solution, change > to >=
```

```
bool onleft(Line L, Pt p) {
  return dcmp(L.v^(p-L.s)) > 0;
// assume that Lines intersect
vector<Pt> HPI(vector<Line>& L) {
  sort(L.begin(), L.end()); // sort by angle
int n = L.size(), fir, las;
  Pt *p = new Pt[n];
  Line *q = new Line[n];
  a[fir=las=0] = L[0];
for(int i = 1; i < n; i++) {
  while(fir < las && !onleft(L[i], p[las-1])) las--;
}</pre>
    while(fir < las && !onleft(L[i], p[fir])) fir++;</pre>
    q[++las] = L[i];
    if(dcmp(q[las].v^q[las-1].v) == 0) {
       las-
       if(onleft(q[las], L[i].s)) q[las] = L[i];
    if(fir < las) p[las-1] = LLIntersect(q[las-1], q[</pre>
         las]);
  while(fir < las && !onleft(q[fir], p[las-1])) las--;</pre>
  if(las-fir <= 1) return {};</pre>
  p[las] = LLIntersect(q[las], q[fir]);
  int m = 0;
  vector<Pt> ans(las-fir+1);
  for(int i = fir ; i <= las ; i++) ans[m++] = p[i];
  return ans;
```

4.4 Convex Hull

```
double cross(Pt o, Pt a, Pt b){
 return (a-o) ^ (b-o);
vector<Pt> convex_hull(vector<Pt> pt){
  sort(pt.begin(),pt.end());
  int top=0;
 vector<Pt> stk(2*pt.size());
 for (int i=0; i<(int)pt.size(); i++){</pre>
   while (top >= 2 && cross(stk[top-2],stk[top-1],pt[i
        ]) <= 0)
      top--;
    stk[top++] = pt[i];
 for (int i=pt.size()-2, t=top+1; i>=0; i--){
   while (top >= t && cross(stk[top-2],stk[top-1],pt[i
        ]) <= 0)
      ton--:
   stk[top++] = pt[i];
 stk.resize(top-1);
 return stk;
```

4.5 Intersection of 2 segments

4.6 Intersection of circle and segment

```
| bool Inter( const Pt& p1 , const Pt& p2 , Circle& cc ){
    Pt dp = p2 - p1;
    double a = dp * dp;
    double b = 2 * ( dp * ( p1 - cc.0 ) );
```

4.7 Intersection of 2 circles

4.8 Circle cover

```
#define N 1021
#define D double
struct CircleCover{
  int C; Circ c[N]; //填入C(圓數量),c(圓陣列) bool g[N][N], overlap[N][N];
   // Area[i] : area covered by at least i circles
  D Area[ N ];
void init( int _C ){ C = _C; }
bool CCinter( Circ& a , Circ& b , Pt& p1 , Pt& p2 ){
     Pt o1 = a.0 , o2 = b.0;
     D r1 = a.R , r2 = b.R;
     if( norm( o1 - o2 ) > r1 + r2 ) return {};
     if( norm( o1 - o2 ) < max(r1, r2) - min(r1, r2) )
     return {};
D d2 = ( o1 - o2 ) * ( o1 - o2 );
     D d = sqrt(d2);
if( d > r1 + r2 ) return false;
     Pt u=(o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
     D A=sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2+d));
     Pt v=Pt( o1.Y-o2.Y , -o1.X + o2.X ) * A / (2*d2);
p1 = u + v; p2 = u - v;
     return true;
   struct Teve {
     Pt p; D ang; int add;
     Teve() {}
     Teve(Pt _a, D _b, int _c):p(_a), ang(_b), add(_c){}
     bool operator<(const Teve &a)const
      {return ang < a.ang;}
   }eve[ N * 2 ];
   \frac{1}{x} strict: x = 0, otherwise x = -1
   bool disjuct( Circ& a, Circ &b, int x )
  {return sign( norm( a.O - b.O ) - a.R - b.R ) > x;}
bool contain( Circ& a, Circ &b, int x )
{return sign( a.R - b.R - norm( a.O - b.O ) ) > x;}
   bool contain(int i, int j){
     /* c[j] is non-strictly in c[i]. */
     return (sign(c[i].R - c[j].R) > 0 ||
(sign(c[i].R - c[j].R) == 0 && i < j) ) &&
                      contain(c[i], c[j], -1);
   void solve(){
     for( int i = 0 ; i \leftarrow C + 1 ; i + + )
        Area[ i ] = 0;
     for( int i = 0; i < C; i ++ )
for( int j = 0; j < C; j ++ )
overlap[i][j] = contain(i, j);
     for( int i = 0; i < C; i ++ )
for( int j = 0; j < C; j ++ )
g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                           disjuct(c[i], c[j], -1));
     for( int i = 0 ; i < C ; i ++ ){
  int E = 0, cnt = 1;</pre>
        for( int j = 0 ; j < C ; j ++ )
  if( j != i && overlap[j][i] )</pre>
             cnt ++;
        for( int j = 0; j < C; j ++)
           if( i != j && g[i][j] ){
Pt aa, bb;
             CCinter(c[i], c[j], aa, bb);

D A=atan2(aa.Y - c[i].0.Y, aa.X - c[i].0.X);

D B=atan2(bb.Y - c[i].0.Y, bb.X - c[i].0.X);
             eve[E ++] = Teve(bb, B, 1);
eve[E ++] = Teve(aa, A, -1);
              if(B > A) cnt ++;
        if( E == 0 ) Area[ cnt ] += pi * c[i].R * c[i].R;
        else{
           sort( eve , eve + E );
           eve[E] = eve[0];
           for( int j = 0 ; j < E ; j ++ ){
              cnt += eve[j].add;
```

4.9 Convex Hull trick

```
/* Given a convexhull, answer querys in O(\lg N)
CH should not contain identical points, the area should
be > 0, min pair(x, y) should be listed first */
double det( const Pt& p1 , const Pt& p2 )
{ return p1.X * p2.Y - p1.Y * p2.X; }
struct Conv{
  int n;
  vector<Pt> a;
  vector<Pt> upper, lower;
  Conv(vector < Pt > \_a) : a(\_a){}
     n = a.size();
     int ptr = 0;
     for(int i=1; i<n; ++i) if (a[ptr] < a[i]) ptr = i;</pre>
     for(int i=0; i<=ptr; ++i) lower.push_back(a[i]);
for(int i=ptr; i<n; ++i) upper.push_back(a[i]);</pre>
     upper.push_back(a[0]);
  int sign( LL x ){ // fixed when changed to double
  return x < 0 ? -1 : x > 0; }
  pair<LL,int> get_tang(vector<Pt> &conv, Pt vec){
     int l = 0, r = (int)conv.size() - 2;
for(; l + 1 < r; ){</pre>
        int mid = (l + r) / 2;
        if(sign(det(conv[mid+1]-conv[mid],vec))>0)r=mid;
        else l = mid:
     return max(make_pair(det(vec, conv[r]), r)
                   make_pair(det(vec, conv[0]), 0));
  void upd_tang(const Pt &p, int id, int &i0, int &i1){
     if(det(a[i0] - p, a[id] - p) > 0) i0 = id;
if(det(a[i1] - p, a[id] - p) < 0) i1 = id;</pre>
   void bi_search(int l, int r, Pt p, int &i0, int &i1){
     if(l == r) return;
     upd_tang(p, 1 % n, i0, i1);
     int sl=sign(det(a[l % n] - p, a[(l + 1) % n] - p));
     for( ; l + 1 < r; ) {</pre>
       int mid = (l + r) / 2;
        int smid=sign(det(a[mid%n]-p, a[(mid+1)%n]-p));
       if (smid == sl) l = mid;
       else r = mid;
     upd_tang(p, r % n, i0, i1);
   int bi_search(Pt u, Pt v, int l, int r) {
     int sl = sign(det(v - u, a[l % n] - u));
     for(; l + \bar{1} < r; )
        int mid = (l + r) / 2;
        int smid = sign(det(v - u, a[mid % n] - u));
       if (smid == s\bar{l}) l = mid;
       else r = mid;
     return 1 % n;
   // 1. whether a given point is inside the CH
  bool contain(Pt p) {
     if (p.X < lower[0].X | l p.X > lower.back().X)
           return 0;
     int id = lower_bound(lower.begin(), lower.end(), Pt
          (p.X, -INF)) - lower.begin();
     if (lower[id].X == p.X) {
     if (lower[id].Y > p.Y) return 0;
}else if(det(lower[id-1]-p,lower[id]-p)<0)return 0;</pre>
     id = lower_bound(upper.begin(), upper.end(), Pt(p.X
            INF), greater<Pt>()) - upper.begin();
     if (upper[id].X == p.X) {
       if (upper[id].Y < p.Y) return 0;</pre>
     }else if(det(upper[id-1]-p,upper[id]-p)<0)return 0;</pre>
     return 1:
   // 2. Find 2 tang pts on CH of a given outside point
  // return true with i0, i1 as index of tangent points
```

```
′return false if inside CH
  bool get_tang(Pt p, int &i0, int &i1) {
    if (contain(p)) return false;
    i0 = i1 = 0;
    int id = lower_bound(lower.begin(), lower.end(), p)
          lower.begin();
    bi_search(0, id, p, i0, i1);
bi_search(id, (int)lower.size(), p, i0, i1);
    id = lower_bound(upper.begin(), upper.end(), p,
         greater<Pt>()) - upper.begin();
    bi_search((int)lower.size() - 1, (int)lower.size()
         -1 + id, p, i0, i1);
    bi_search((int)lower.size() - 1 + id, (int)lower.
         size() - 1 + (int)upper.size(), p, i0, i1);
    return true:
  // 3. Find tangent points of a given vector
  // ret the idx of vertex has max cross value with vec
  int get_tang(Pt vec){
    pair<LL, int> ret = get_tang(upper, vec);
    ret.second = (ret.second+(int)lower.size()-1)%n;
    ret = max(ret, get_tang(lower, vec));
    return ret.second;
  // 4. Find intersection point of a given line
  // return 1 and intersection is on edge (i, next(i))
  // return 0 if no strictly intersection
  bool get_intersection(Pt u, Pt v, int &i0, int &i1){
   int p0 = get_tang(u - v), p1 = get_tang(v - u);
if(sign(det(v-u,a[p0]-u))*sign(det(v-u,a[p1]-u))<0){</pre>
     if (p0 > p1) swap(p0, p1);
     i0 = bi\_search(u, v, p0, p1);
     i1 = bi_search(u, v, p1, p0 + n);
     return 1;
   return 0;
} };
4.10 Tangent line of two circles
vector<Line> go( const Cir& c1 , const Cir& c2 , int
    sign1 ){
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret;
  double d_sq = norm2(c1.0 - c2.0);
  if( d_sq < eps ) return ret;</pre>
  double d = sqrt( d_sq );
  Pt v = ( c2.0 - c1.0 ) / d;
double c = ( c1.R - sign1 * c2.R ) / d;
  if( c * c > 1 ) return ret;
  double h = sqrt( max( 0.0 , 1.0 - c * c ) );
for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
  Pt n = { v.X * c - sign2 * h * v.Y ,
              v.Y * c + sign2 * h * v.X };
    Pt p1 = c1.0 + n * c1.R;
    Pt p2 = c2.0 + n * (c2.R * sign1);
    if( fabs( p1.X - p2.X ) < eps and
         fabs(p1.Y - p2.Y) < eps )
      p2 = p1 + perp(c2.0 - c1.0);
    ret.push_back( { p1 , p2 } );
  return ret;
4.11 KD Tree
const int MXN=100005;
const int MXK=10;
struct KDTree{
  struct Nd{
    LL x[MXK],mn[MXK],mx[MXK];
    int id,f;
    Nd *1.*r
  }tree[MXN],*root;
  LL dis(LL a,LL b){return (a-b)*(a-b);}
  LL dis(LL a[MXK],LL b[MXK]){
    LL ret=0;
    for(int i=0;i<k;i++) ret+=dis(a[i],b[i]);</pre>
```

return ret;

void init(vector<vector<LL>> &ip,int _n,int _k){

```
n=_n,k=_k;
for(int i=0;i<n;i++){</pre>
      tree[i].id=i;
      copy(ip[i].begin(),ip[i].end(),tree[i].x);
    root=build(0,n-1,0);
  Nd* build(int l,int r,int d){
    if(l>r) return NULL;
    if(d==k) d=0;
    int m=(l+r)>>1;
    nth_element(tree+l,tree+m,tree+r+1,[&](const Nd &a,
        const Nd &b){return a.x[d]<b.x[d];});</pre>
    tree[m].f=d;
    copy(tree[m].x,tree[m].x+k,tree[m].mn);
    copy(tree[m].x,tree[m].x+k,tree[m].mx);
    tree[m].l=build(l,m-1,d+1);
    if(tree[m].l){
      for(int i=0;i<k;i++){</pre>
        tree[m].mn[i]=min(tree[m].mn[i],tree[m].l->mn[i
        tree[m].mx[i]=max(tree[m].mx[i],tree[m].l->mx[i
             ١);
      }
    tree[m].r=build(m+1,r,d+1);
    if(tree[m].r){
      for(int i=0;i<k;i++){</pre>
        tree[m].mn[i]=min(tree[m].mn[i],tree[m].r->mn[i
        tree[m].mx[i]=max(tree[m].mx[i],tree[m].r->mx[i
      }
    }
    return tree+m;
  LL pt[MXK],md;
  int mID;
  bool touch(Nd *r){
    LL d=0;
    for(int i=0;i<k;i++){</pre>
      if(pt[i]<=r->mn[i]) d+=dis(pt[i],r->mn[i]);
        else if(pt[i]>=r->mx[i]) d+=dis(pt[i],r->mx[i])
    return d<md;
  void nearest(Nd *r){
    if(!rll!touch(r)) return;
    LL td=dis(r->x,pt);
    if(td<md) md=td,mID=r->id;
    nearest(pt[r->f]< r->x[r->f]?r->l:r->r);
    nearest(pt[r->f]< r->x[r->f]? r->r:r->1);
  pair<LL,int> query(vector<LL> &_pt,LL _md=1LL<<57){</pre>
    mID=-1, md=\_md;
    copy(_pt.begin(),_pt.end(),pt);
    nearest(root);
    return {md,mID};
}tree;
```

4.12 Lower Concave Hull

```
const ll is_query = -(1LL<<62);</pre>
struct Line {
  11 m, b;
  mutable function<const Line*()> succ;
  bool operator<(const Line& rhs) const {</pre>
    if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line* s = succ();
    return s ? b - s->b < (s->m - m) * rhs.m : 0;
}; // maintain upper hull for maximum
struct HullDynamic : public multiset<Line> {
 bool bad(iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
      return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
```

4.13 Min Enclosing Circle

```
struct Mec{
   // return pair of center and r
   static const int N = 101010;
   int n;
   Pt p[N], cen;
   double r2
   void init( int _n , Pt _p[] ){
     n = _n;
     memcpy( p , _p , sizeof(Pt) * n );
   double sqr(double a){ return a*a; }
   Pt center(Pt p0, Pt p1, Pt p2) {
     Pt a = p1-p0;
     Pt b = p2-p0;
     double c1=norm2( a ) * 0.5;
     double c2=norm2( b ) * 0.5;
     double d = a \wedge b;
     double x = p0.X + (c1 * b.Y - c2 * a.Y) / d;
     double y = p0.Y + (a.X * c2 - b.X * c1) / d;
     return Pt(x,y);
   pair<Pt,double> solve(){
     random_shuffle(p,p+n);
      for (int i=0; i<n; i++){</pre>
        if (norm2(cen-p[i]) <= r2) continue;</pre>
        cen = p[i];
        r2 = 0:
        for (int j=0; j<i; j++){
  if (norm2(cen-p[j]) <= r2) continue;
  cen=Pt((p[i].X+p[j].X)/2,(p[i].Y+p[j].Y)/2);</pre>
          r2 = norm2(cen-p[j]);
          for (int k=0; k<j; k++){
  if (norm2(cen-p[k]) <= r2) continue;
  cen = center(p[i],p[j],p[k]);</pre>
             r2 = norm2(cen-p[k]);
       }
     return {cen,sqrt(r2)};
} mec;
```

4.14 Min Enclosing Ball

```
L[0]=(sol[0]*m[1][1]-sol[1]*m[0][1])/det;
L[1]=(sol[1]*m[0][0]-sol[0]*m[1][0])/det;
res=outer[0]+q[0]*L[0]+q[1]*L[1];
       radius=norm2(res, outer[0]);
       break;
     case 4:
       for (i=0; i<3; ++i) q[i]=outer[i+1]-outer[0], sol
   [i]=(q[i] * q[i]);</pre>
       for (i=0;i<3;++i) for(j=0;j<3;++j) m[i][j]=(q[i]
       * q[j])*2;
det= m[0][0]*m[1][1]*m[2][2]
          + m[0][1]*m[1][2]*m[2][0]
+ m[0][2]*m[2][1]*m[1][0]
          - m[0][2]*m[1][1]*m[2][0]
- m[0][1]*m[1][0]*m[2][2]
          - m[0][0]*m[1][2]*m[2][1];
        if ( fabs(det)<eps ) return;</pre>
       for (j=0; j<3; ++j) {
  for (i=0; i<3; ++i) m[i][j]=sol[i];</pre>
          L[j]=( m[0][0]*m[1][1]*m[2][2]
+ m[0][1]*m[1][2]*m[2][0]
                   + m[0][2]*m[2][1]*m[1][0]
                   - m[0][2]*m[1][1]*m[2][0]
- m[0][1]*m[1][0]*m[2][2]
                     m[0][0]*m[1][2]*m[2][1]
                ) / det;
          for (i=0; i<3; ++i) m[i][j]=(q[i] * q[j])*2;
       } res=outer[0];
       for (i=0; i<3; ++i ) res = res + q[i] * L[i];
       radius=norm2(res, outer[0]);
void minball(int n){ ball();
  if( nouter < 4 ) for( int i = 0 ; i < n ; i ++ )
  if( norm2(res, pt[i]) - radius > eps ){
       outer[ nouter ++ ] = pt[ i ]; minball(i); --
             nouter
       if(i>0){ Pt Tt = pt[i]
          memmove(&pt[1], &pt[0], sizeof(Pt)*i); pt[0]=Tt
}}}
double solve(){
  // n points in pt
  random_shuffle(pt, pt+n); radius=-1;
  for(int i=0;i<n;i++) if(norm2(res,pt[i])-radius>eps)
    nouter=1, outer[0]=pt[i], minball(i);
  return sqrt(radius);
```

4.15 Min dist on Cuboid

```
typedef LL T;
Tr:
if(j>=0 && j< 2) turn(i, j+1, x, y0+W+z, y0+W-y,
x0, y0+W, L, H, W);
  if(i<=0 && i>-2) turn(i-1, j, x0-z, y, x-x0, x0-H, y0, H, W, L);
  if(j<=0 && j>-2) turn(i, j-1, x, y0-z, y-y0, x0, y0-H, L, H, W);
T solve(T L, T_W, T H,
        T x1, T y1, T z1, T x2, T y2, T z2){
  if( z1!=0 && z1!=H ){
    if( y1==0 || y1==W )
      swap(y1,z1), swap(y2,z2), swap(W,H);
  else swap(x1,z1), swap(x2,z2), swap(L,H);
  if (z1==H) z1=0, z2=H-z2;
  r=INF; turn(0,0,x2-x1,y2-y1,z2,-x1,-y1,L,W,H);
  return r;
```

4.16 Heart of Triangle

```
Pt inCenter( Pt &A, Pt &B, Pt &C) { // 内心
  double a = norm(B-C), b = norm(C-A), c = norm(A-B);
  return (A * a + B * b + C * c) / (a + b + c);
```

```
Pt circumCenter( Pt &a, Pt &b, Pt &c) { // 外心
   Pt bb = b - a, cc = c - a;
   double db=norm2(bb), dc=norm2(cc), d=2*(bb ^ cc);
   return a-Pt(bb.Y*dc-cc.Y*db, cc.X*db-bb.X*dc) / d;
Pt othroCenter( Pt &a, Pt &b, Pt &c) { // 垂心 Pt ba = b - a, ca = c - a, bc = b - c; double Y = ba.Y * ca.Y * bc.Y,
     A = ca.X * ba.Y - ba.X * ca.Y
     x0= (Y+ca.X*ba.Y*b.X-ba.X*ca.Y*c.X) / A,
     y0 = -ba.X * (x0 - c.X) / ba.Y + ca.Y;
   return Pt(x0, y0);
}
```

5 Graph

5.1 DominatorTree

```
const int MAXN = 100010;
struct DominatorTree{
#define REP(i,s,e) for(int i=(s);i<=(e);i++)</pre>
#define REPD(i,s,e) for(int i=(s);i>=(e);i--)
  int n , m , s;
vector< int > g[ MAXN ] , pred[ MAXN ];
vector< int > cov[ MAXN ];
   int dfn[ MAXN ] , nfd[ MAXN ] , ts;
   int par[ MAXN ]; //idom[u] s到u的最後一個必經點int sdom[ MAXN ] , idom[ MAXN ];
   int mom[ MAXN ] , mn[ MAXN ];
inline bool cmp( int u , int v )
{ return dfn[ u ] < dfn[ v ]; }
int eval( int u ){</pre>
      if( mom[ u ] == u ) return u;
int res = eval( mom[ u ] );
if(cmp( sdom[ mn[ mom[ u ] ] ] , sdom[ mn[ u ] ] ))
      mn[ u ] = mn[ mom[ u ] ];
return mom[ u ] = res;
   void init( int _n , int _m , int _s ){
      ts = 0; n = _n; m = _m; s = _s;
REP( i, 1, n ) g[ i ].clear(), pred[ i ].clear();
   void addEdge( int u , int v ){
  g[ u ].push_back( v );
  pred[ v ].push_back( u );
   void dfs( int u ){
      ts++;
      dfn[ u ] = ts;
nfd[ ts ] = u;
      for( int v : g[ u ] ) if( dfn[ v ] == 0 ){
         par[ v ] = u;
dfs( v );
      }
   void build(){
      REP( i , 1 , n ){
  dfn[ i ] = nfd[ i ] = 0;
  cov[ i ].clear();
  mom[ i ] = mn[ i ] = sdom[ i ] = i;
      dfs( s );
      REPD( i , n , 2 ){
         int u = nfd[ i ];
         if( u == 0 ) continue ;
for( int v : pred[ u ] ) if( dfn[ v ] ){
            eval( v );
            if( cmp( sdom[ mn[ v ] ] , sdom[ u ] ) )
  sdom[ u ] = sdom[ mn[ v ] ];
         cov[ sdom[ u ] ].push_back( u );
mom[ u ] = par[ u ];
         for( int w : cov[ par[ u ] ] ){
            eval( w );
            if( cmp( sdom[ mn[ w ] ] , par[ u ] ) )
            idom[w] = mn[w];
else idom[w] = par[u];
         cov[ par[ u ] ].clear();
      REP( i , 2 , n ){
  int u = nfd[ i ];
```

```
if( u == 0 ) continue;
if( idom[ u ] != sdom[ u ] )
      idom[ u ] = idom[ idom[ u ] ];
}
}
domT;
```

5.2 MaximumClique 最大團

```
#define N 111
struct MaxClique{ // 0-base
  typedef bitset<N> Int;
  Int linkto[N] , v[N];
  int n;
  void init(int _n){
    n = _n;
     for(int i = 0; i < n; i ++){
       linkto[i].reset(); v[i].reset();
  }
  void addEdge(int a , int b)
  \{ v[a][b] = v[b][a] = 1; \}
  int popcount(const Int& val)
  { return val.count(); } int lowbit(const Int& val)
  { return val._Find_first(); }
  int ans , stk[N];
int id[N] , di[N] , deg[N];
  Int cans;
  void maxclique(int elem_num, Int candi){
     if(elem_num > ans){
       ans = elem_num; cans.reset();
for(int i = 0 ; i < elem_num ; i ++)
   cans[id[stk[i]]] = 1;</pre>
     int potential = elem_num + popcount(candi);
     if(potential <= ans) return;</pre>
     int pivot = lowbit(candi);
     Int smaller_candi = candi & (~linkto[pivot]);
     while(smaller_candi.count() && potential > ans){
       int next = lowbit(smaller_candi);
       candi[next] = !candi[next]
       smaller_candi[next] = !smaller_candi[next];
       potential --:
       if(next == pivot || (smaller_candi & linkto[next
            ]).count()){
         stk[elem_num] = next;
         maxclique(elem_num + 1, candi & linkto[next]);
    }
  int solve(){
    for(int i = 0; i < n; i ++){
       id[i] = i; deg[i] = v[i].count();
    sort(id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; });
    for(int i = 0; i < n; i ++) di[id[i]] = i;
for(int i = 0; i < n; i ++)
for(int j = 0; j < n; j ++)</pre>
          if(v[i][j]) linkto[di[i]][di[j]] = 1;
    Int cand; cand.reset();
for(int i = 0 ; i < n ; i ++) cand[i] = 1;</pre>
     ans = 1;
     cans.reset(); cans[0] = 1;
    maxclique(0, cand);
     return ans;
} solver;
```

5.3 MaximalClique 極大團

```
#define N 80
struct MaxClique{ // 0-base
  typedef bitset<N> Int;
  Int lnk[N] , v[N];
  int n;
  void init(int _n){
    n = _n;
    for(int i = 0 ; i < n ; i ++){
       lnk[i].reset();
    }</pre>
```

```
void addEdge(int a , int b)
  \{ v[a][b] = v[b][a] = 1; \}
  int ans , stk[N], id[N] , di[N] , deg[N];
  Int cans:
  void dfs(int elem_num, Int candi, Int ex){
    if(candi.none()&ex.none()){
       cans.reset();
       for(int i = 0; i < elem_num; i ++)
       cans[id[stk[i]]] = 1;
ans = elem_num; // cans is a maximal clique
       return:
     int pivot = (candilex)._Find_first();
    Int smaller_candi = candi & (~lnk[pivot]);
     while(smaller_candi.count()){
       int nxt = smaller_candi._Find_first();
       candi[nxt] = smaller_candi[nxt] = 0;
       ex[nxt] = 1;
       stk[elem_num] = nxt;
       dfs(elem_num+1,candi&lnk[nxt],ex&lnk[nxt]);
  int solve(){
    for(int i = 0; i < n; i ++){
       id[i] = i; deg[i] = v[i].count();
    sort(id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; });
for(int i = 0; i < n; i ++) di[id[i]] = i;
for(int i = 0; i < n; i ++)
       for(int j = 0; j < n; j ++)
    if(v[i][j]) lnk[di[i]][di[j]] = 1;
ans = 1; cans.reset(); cans[0] = 1;
dfs(0, Int(string(n,'1')), 0);
    return ans;
} solver;
```

5.4 Strongly Connected Component

```
struct Scc{
  int n, nScc, vst[MXN], bln[MXN];
vector<int> E[MXN], rE[MXN], vec;
  void init(int _n){
    n = _n;
for (int i=0; i<MXN; i++)</pre>
       E[i].clear(), rE[i].clear();
  void addEdge(int u, int v){
    E[u].PB(v); rE[v].PB(u);
  void DFS(int u){
    vst[u]=1;
    for (auto v : E[u]) if (!vst[v]) DFS(v);
    vec.PB(u);
  void rDFS(int u){
    vst[u] = 1; bln[u] = nScc;
    for (auto v : rE[u]) if (!vst[v]) rDFS(v);
  void solve(){
    nScc = 0:
    vec.clear();
    FZ(vst);
    for (int i=0; i<n; i++)
       if (!vst[i]) DFS(i);
    reverse(vec.begin(),vec.end());
    FZ(vst);
     for (auto v : vec)
      if (!vst[v]){
         rDFS(v); nScc++;
};
```

5.5 Dynamic MST

```
/* Dynamic MST O( Q lg^2 Q )
  (qx[i], qy[i])->chg weight of edge No.qx[i] to qy[i]
  delete an edge: (i, \infty)
  add an edge: change from \infty to specific value */
```

```
const int SZ=M+3*MXQ;
int a[N],*tz;
int find(int xx){
  int root=xx; while(a[root]) root=a[root];
  int next; while((next=a[xx])){a[xx]=root; xx=next; }
  return root;
bool cmp(int aa,int bb){ return tz[aa]<tz[bb]; }</pre>
int kx[N],ky[N],kt, vd[N],id[M], app[M];
bool extra[M];
void solve(int *qx,int *qy,int Q,int n,int *x,int *y,
     int *z,int m1,long long ans){
  if(Q==1){
     for(int i=1;i<=n;i++) a[i]=0;</pre>
    z[qx[0]]=qy[0]; tz = z
    for(int i=0;i<m1;i++) id[i]=i;</pre>
    sort(id,id+m1,cmp); int ri,rj;
for(int i=0;i<m1;i++){</pre>
       ri=find(x[id[i]]); rj=find(y[id[i]]);
       if(ri!=rj){ ans+=z[id[i]]; a[ri]=rj; }
    printf("%lld\n",ans);
    return;
  int ri,rj;
  //contract
  kt=0;
  for(int i=1;i<=n;i++) a[i]=0;</pre>
  for(int i=0;i<Q;i++){</pre>
    ri=find(x[qx[i]]); rj=find(y[qx[i]]); if(ri!=rj) a[
         ri]=rj;
  int tm=0;
  for(int i=0;i<m1;i++) extra[i]=true;</pre>
  for(int i=0;i<Q;i++) extra[ qx[i] ]=false;</pre>
  for(int i=0;i<m1;i++) if(extra[i]) id[tm++]=i;</pre>
  tz=z; sort(id,id+tm,cmp);
  for(int i=0;i<tm;i++){</pre>
    ri=find(x[id[i]]); rj=find(y[id[i]]);
    if(ri!=rj){
       a[ri]=rj; ans += z[id[i]];
       k\bar{x}[k\bar{t}]=x[id[i]]; k\bar{y}[k\bar{t}]=y[id[i]]; kt++;
  for(int i=1;i<=n;i++) a[i]=0;</pre>
  for(int i=0;i<kt;i++) a[ find(kx[i]) ]=find(ky[i]);</pre>
  for(int i=1;i<=n;i++) if(a[i]==0)</pre>
  vd[i]=++n2;
  for(int i=1;i<=n;i++) if(a[i])</pre>
  vd[i]=vd[find(i)];
int m2=0, *Nx=x+m1, *Ny=y+m1, *Nz=z+m1;
  for(int i=0;i<m1;i++) app[i]=-1;</pre>
  for(int i=0;i<Q;i++) if(app[qx[i]]==-1){</pre>
    Nx[m2]=vd[ x[ qx[i] ] ]; Ny[m2]=vd[ y[ qx[i] ] ];
Nz[m2]=z[ qx[i] ];
app[qx[i]]=m2; m2++;
  for(int i=0;i<Q;i++){ z[ qx[i] ]=qy[i]; qx[i]=app[qx[</pre>
       i]]; }
  for(int i=1;i<=n2;i++) a[i]=0;</pre>
  for(int i=0;i<tm;i++){</pre>
     ri=find(vd[ x[id[i]] ]);    rj=find(vd[ y[id[i]] ]);
    if(ri!=rj){
       a[ri]=rj; Nx[m2]=vd[ x[id[i]] ]
       Ny[m2]=vd[y[id[i]]]; Nz[m2]=z[id[i]]; m2++;
  int mid=Q/2;
  solve(qx,qy,mid,n2,Nx,Ny,Nz,m2,ans);
  solve(qx+mid,qy+mid,Q-mid,n2,Nx,Ny,Nz,m2,ans);
int x[SZ],y[SZ],z[SZ],qx[MXQ],qy[MXQ],n,m,Q;
void init(){
  scanf("%d%d",&n,&m);
  for(int i=0;i<m;i++) scanf("%d%d%d",x+i,y+i,z+i);</pre>
  scanf("%d",&Q);
  for(int i=0;i<Q;i++){ scanf("%d%d",qx+i,qy+i); qx[i</pre>
void work(){ if(Q) solve(qx,qy,Q,n,x,y,z,m,0); }
```

5.6 Maximum General graph Matching

```
const int N = 514, E = (2e5) * 2;
struct Graph{
  int to[E],bro[E],head[N],e;
  int lnk[N],vis[N],stp,n;
  void init( int _n ){
    stp = 0; e = 1; n = _n;
    for( int i = 1; i <= n; i ++ )
lnk[i] = vis[i] = 0;
  void add_edge(int u,int v){
    to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u,bro[e]=head[v],head[v]=e++;
  bool dfs(int x){
    vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(!lnk[v]){
        lnk[x]=v, lnk[v]=x;
         return true;
      }else if(vis[lnk[v]]<stp){</pre>
        int w=lnk[v]
         lnk[x]=v, lnk[v]=x, lnk[w]=0;
         if(dfs(w)){
          return true:
         lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    return false;
  int solve(){
    int ans = 0;
    for(int i=1;i<=n;i++)
      if(!lnk[i]){
        stp++; ans += dfs(i);
    return ans;
} graph;
```

5.7 Minimum General Weighted Matching

```
struct Graph {
  // Minimum General Weighted Matching (Perfect Match)
  static const int MXN = 105;
  int n, edge[MXN][MXN];
  int match[MXN],dis[MXN],onstk[MXN];
  vector<int> stk;
  void init(int _n) {
    n = _n;
    for( int i = 0; i < n; i ++)
      for( int j = 0 ; j < n ; j ++ )
  edge[ i ][ j ] = 0;</pre>
  void add_edge(int u, int v, int w)
  \{ edge[u][v] = edge[v][u] = w; \}
  bool SPFA(int u){
    if (onstk[u]) return true;
    stk.PB(u);
    onstk[u] = 1;
    for (int v=0; v<n; v++){
  if (u != v && match[u] != v && !onstk[v]){</pre>
         int m = match[v];
         if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
           dis[m] = dis[u] - edge[v][m] + edge[u][v];
           onstk[v] = 1;
           stk.PB(v)
           if (SPFA(m)) return true;
           stk.pop_back();
           onstk[v] = 0;
    } } }
    onstk[u] = 0;
    stk.pop_back();
    return false;
  int solve() {
    // find a match
    for (int i=0; i<n; i+=2){</pre>
      match[i] = i+1;
      match[i+1] = i;
```

```
while (true){
  int found = 0;
        for( int i = 0 ; i < n_; i ++ )
        onstk[ i ] = dis[ i ] = 0;
for (int i=0; i<n; i++){
          stk.clear()
          if (!onstk[i] && SPFA(i)){
             found = 1
             while (SZ(stk)>=2){
                int u = stk.back(); stk.pop_back();
int v = stk.back(); stk.pop_back();
                match[u] = v;
               match[v] = u;
        } } }
        if (!found) break;
     }
     int ret = 0;
for (int i=0; i<n; i++)
       ret += edge[i][match[i]];
     ret /= 2;
     return ret;
}graph;
```

5.8 Minimum Steiner Tree

```
// Minimum Steiner Tree 重要點的mst
// 0(V 3^T + V^2 2^T)
struct SteinerTree{
#define V 33
#define T 8
#define INF 1023456789
   int n , dst[V][V] , dp[1 << T][V] , tdst[V];
void init( int _n ){</pre>
      n = _n;
      for( int i = 0 ; i < n ; i ++ ){
  for( int j = 0 ; j < n ; j ++ )
    dst[ i ][ j ] = INF;
  dst[ i ][ i ] = 0;</pre>
   void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
   void shortest_path(){
      for( int k = 0 ; k < n ; k ++ )
for( int i = 0 ; i < n ; i ++ )
            int solve( const vector<int>& ter ){
  int t = (int)ter.size();
      for( int i = 0 ; i < (1 << t ) ; i ++ )
  for( int j = 0 ; j < n ; j ++ )
    dp[ i ][ j ] = INF;
for( int i = 0 ; i < n ; i ++ )</pre>
         dp[0][i] = 0;
      for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){</pre>
         if( msk == ( msk & (-msk) ) ){
            int who = __lg( msk );
for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
            continue;
         for( int i = 0 ; i < n ; i ++ )
            for( int submsk = ( msk - 1 ) & msk ; submsk ; submsk = ( submsk - 1 ) & msk ) dp[ msk ][ i ] = min( dp[ msk ][ i ],
                                             dp[ submsk ][ i ] +
                                             dp[ msk ^ submsk ][ i ] );
         for( int i = 0; i < n; i ++){
            tdst[i] = INF;
            for( int j = 0 ; j < n ; j ++ )
  tdst[ i ] = min( tdst[ i ],</pre>
                                 dp[ msk ][ j ] + dst[ j ][ i ] );
         for( int i = 0 ; i < n ; i ++
  dp[ msk ][ i ] = tdst[ i ];</pre>
      int ans = INF;
      for( int i = 0; i < n; i ++)
```

```
ans = min(ans, dp[(1 << t) - 1][i]);
      return ans;
} }solver;
 5.9 BCC based on vertex
struct BccVertex {
   int n,nScc,step,dfn[MXN],low[MXN];
vector<int> E[MXN],sccv[MXN];
   int top,stk[MXN];
   void init(int _n) {
  n = _n; nScc = step = 0;
      for (int i=0; i<n; i++) E[i].clear();</pre>
   void addEdge(int u, int v)
   { E[u].PB(v); E[v].PB(u); } void DFS(int u, int f) { dfn[u] = low[u] = step++;
      stk[top++] = u;
      for (auto v:E[u]) {
        if (v == f) continue;
        if (dfn[v] == -1) {
          DFS(v,u);
low[u] = min(low[u], low[v]);
           if (low[v] >= dfn[u]) {
             int z
             sccv[nScc].clear();
             do {
               z = stk[--top]
               sccv[nScc].PB(z);
             } while (z != v);
             sccv[nScc++].PB(u);
        }else
          low[u] = min(low[u],dfn[v]);
   } }
   vector<vector<int>> solve() {
      vector<vector<int>> res;
      for (int i=0; i<n; i++)</pre>
      dfn[i] = low[i] = -1;
for (int i=0; i<n; i++)
        if (dfn[i] == -1) {
          top = 0;
          DFS(i,i);
      REP(i,nScc) res.PB(sccv[i]);
     return res;
}graph;
 5.10 Min Mean Cycle
/* minimum mean cycle O(VE) */
struct MMC{
#define E 101010
#define V 1021
#define inf 1e9
#define eps 1e-6
   struct Edge { int v,u; double c; };
   int n, m, prv[V][V], prve[V][V], vst[V];
   Edge e[E];
   vector<int> edgeID, cycle, rho;
   double d[V][V];
   void init( int _n )
   { n = _n; m = 0; }
// WARNING: TYPE matters
   void addEdge( int vi , int ui , double ci )
{ e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
     for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
        fill(d[i+1], d[i+1]+n, inf);
for(int j=0; j<m; j++) {
          int v = e[jj.v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
             d[i+1][u] = d[i][v]+e[j].c;
             prv[i+1][u] = v;
             prve[i+1][u] = j;
```

1111

double solve(){

double mmc=inf;

// returns inf if no cycle, mmc otherwise

```
int st = -1
    bellman_ford();
    for(int i=0; i<n; i++) {</pre>
       double avg=-inf;
      for(int k=0; k<n; k++) {</pre>
         if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i</pre>
             ])/(n-k));
        else avg=max(avg,inf);
      if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
    fill(vst,0); edgeID.clear(); cycle.clear(); rho.
         clear();
    for (int i=n; !vst[st]; st=prv[i--][st]) {
      vst[st]++;
      edgeID.PB(prve[i][st]);
      rho.PB(st);
    while (vst[st] != 2) {
      if(rho.empty()) return inf;
      int v = rho.back(); rho.pop_back();
      cycle.PB(v);
      vst[v]++;
    reverse(ALL(edgeID));
    edgeID.resize(SZ(cycle));
    return mmc;
} }mmc;
```

5.11 Directed Graph Min Cost Cycle

```
// works in O(N M)
#define INF 1000000000000000LL
#define N 5010
#define M 200010
struct edge{
  int to; LL w;
  edge(int a=0, LL b=0): to(a), w(b){}
struct node{
 LL d; int u, next;
  node(LL a=0, int b=0, int c=0): d(a), u(b), next(c){}
}b[M];
struct DirectedGraphMinCycle{
  vector<edge> g[N], grev[N];
  LL dp[N][N], p[N], d[N], mu;
  bool inq[N];
  int n, bn, bsz, hd[N];
  void b_insert(LL d, int u){
    int i = d/mu;
    if(i >= bn) return;
    b[++bsz] = node(d, u, hd[i]);
    hd[i] = bsz;
  void init( int _n ){
    n = _n;
for( int i = 1 ; i <= n ; i ++ )
g[ i ].clear();
  void addEdge( int ai , int bi , LL ci )
  { g[ai].push_back(edge(bi,ci)); }
  LL solve(){
    fill(dp[0], dp[0]+n+1, 0);
    for(int i=1; i<=n; i++){</pre>
      fill(dp[i]+1, dp[i]+n+1, INF);
for(int j=1; j<=n; j++) if(dp[i-1][j] < INF){
  for(int k=0; k<(int)g[j].size(); k++)</pre>
           dp[i][g[j][k].to] =min(dp[i][g[j][k].to]
                                       dp[i-1][j]+g[j][k].w);
    mu=INF; LL bunbo=1;
    for(int i=1; i<=n; i++) if(dp[n][i] < INF){</pre>
       LL a=-INF, b=1;
      for(int j=0; j<=n-1; j++) if(dp[j][i] < INF){
  if(a*(n-j) < b*(dp[n][i]-dp[j][i])){</pre>
           a = dp[n][i]-dp[j][i];
           b = n-j;
      if(mu*b > bunbo*a)
         mu = a, bunbo = b;
    if(mu < 0) return -1; // negative cycle</pre>
```

```
if(mu == INF) return INF; // no cycle
     if(mu == 0) return 0;
     for(int i=1; i<=n; i++)</pre>
        for(int j=0; j<(int)g[i].size(); j++)
g[i][j].w *= bunbo;</pre>
     memset(p, 0, sizeof(p));
     queue<int> q;
     for(int i=1; i<=n; i++){
        q.push(i);
        inq[i] = true;
     while(!q.empty()){
        int i=q.front(); q.pop(); inq[i]=false;
for(int j=0; j<(int)g[i].size(); j++){
  if(p[g[i][j].to] > p[i]+g[i][j].w-mu){
             p[g[i][j].to] = p[i]+g[i][j].w-mu;
             if(!inq[g[i][j].to]){
   q.push(g[i][j].to);
                inq[g[i][j].to] = true;
     } } } }
     for(int i=1; i<=n; i++) grev[i].clear();</pre>
     for(int i=1; i<=n; i++)</pre>
        for(int j=0; j<(int)g[i].size(); j++){
  g[i][j].w += p[i]-p[g[i][j].to];</pre>
          grev[g[i][j].to].push_back(edge(i, g[i][j].w));
       }
     LL mldc = n*mu;
     for(int i=1; i<=n; i++){</pre>
        bn=mldc/mu, bsz=0;
       memset(hd, 0, sizeof(hd));
fill(d+i+1, d+n+1, INF);
b_insert(d[i]=0, i);
        for(int j=0; j<=bn-1; j++) for(int k=hd[j]; k; k=
   b[k].next){</pre>
           int u = b[k].u;
           LL du = b[k].d;
           if(du > d[u]) continue;
           for(int l=0; l<(int)g[u].size(); l++) if(g[u][l</pre>
                ].to > i)
             if(d[g[u][i].to] > du + g[u][l].w){
               d[g[u][1].to] = du + g[u][1].w;
                b_insert(d[g[u][l].to], g[u][l].to);
        } } }
        for(int j=0; j<(int)grev[i].size(); j++) if(grev[</pre>
             i][j].to > i)
           mldc=min(mldc,d[grev[i][j].to] + grev[i][j].w);
     return mldc / bunbo;
} }araph;
```

5.12 K-th Shortest Path

```
// time: O(|E| \setminus |E| + |V| \setminus |g| |V| + |K|)
// memory: O(|E| \lg |E| + |V|)
struct KSP{ // 1-base
   struct nd{
      int u, v; ll d;
     nd(int ui = 0, int vi = 0, ll di = INF)
{ u = ui; v = vi; d = di; }
   struct heap{
     nd* edge; int dep; heap* chd[4];
   static int cmp(heap* a,heap* b)
   { return a->edge->d > b->edge->d; }
   struct node{
      int v; ll d; heap* H; nd* E;
      node(){}
     node(ll _d, int _v, nd* _E)
{ d =_d; v = _v; E = _E; }
node(heap* _H, ll _d)
      \{ H = _H; d = _d; \}
      friend bool operator<(node a, node b)
      { return a.d > b.d; }
  int n, k, s, t, dst[N];
nd *nxt[N];
  vector<nd*> g[N], rg[N
heap *nullNd, *head[N];
                               rg[ N ];
  void init( int _n , int _k , int _s , int _t ){
  n = _n;  k = _k;  s = _s;  t = _t;
  for( int i = 1 ; i <= n ; i ++ ){</pre>
```

if(dst[s] == -1) return;

```
g[ i ].clear(); rg[ i ].clear();
nxt[ i ] = NULL; head[ i ] = NULL;
dst[ i ] = -1;
                                                                       ans.push_back( dst[ s ] );
                                                                       if( head[s] != nullNd )
                                                                         Q.push(node(head[s], dst[s]+head[s]->edge->d));
                                                                       for( int _ = 1 ; _ < k and not Q.empty() ; _ ++ ){
  node p = Q.top(), q; Q.pop();
  ans.push_back( p.d );
  } }
  void addEdge( int ui , int vi , ll di ){
    nd* e = new nd(ui, vi, di);
g[ui].push_back( e );
                                                                         if(head[ p.H->edge->v ] != nullNd){
    rg[ vi ].push_back( e );
                                                                            q.H = head[ p.H->edge->v ];
                                                                            q.d = p.d + q.H->edge->d;
  queue<int> dfsQ;
                                                                            Q.push(q);
  void dijkstra(){
                                                                         for( int i = 0 ; i < 4 ; i ++ )
  if( p.H->chd[ i ] != nullNd ){
    q.H = p.H->chd[ i ];
}
    while(dfsQ.size()) dfsQ.pop();
    priority_queue<node> Q;
    Q.push(node(0, t, NULL));
    while (!Q.empty()){
                                                                              q.d = p.d - p.H->edge->d + p.H->chd[i]->
      node p = Q.top(); Q.pop();
                                                                                   edge->d;
       if(dst[p.v] != -1) continue;
                                                                              Q.push( q );
      dst[p.v] = p.d;
nxt[p.v] = p.E;
                                                                     } }
                                                                     void solve(){ // ans[i] stores the i-th shortest path
      dfsQ.push( p.v_);
                                                                       dijkstra();
       for(auto e: rg[ p.v ])
                                                                       build();
         Q.push(node(p.d + e->d, e->u, e));
                                                                       first_K(); // ans.size() might less than k
                                                                  } }solver;
  heap* merge(heap* curNd, heap* newNd){
    if(curNd == nullNd) return newNd;
                                                                   5.13 SPFA
    heap* root = new heap;
memcpy(root, curNd, sizeof(heap));
                                                                  bool spfa(){
    if(newNd->edge->d < curNd->edge->d){
                                                                     deque<int> dq;
       root->edge = newNd->edge;
                                                                     dis[0]=0;
      root->chd[2] = newNd->chd[2];
root->chd[3] = newNd->chd[3];
                                                                     dq.push_back(0);
                                                                     inq[0]=1;
                                                                     while(!dq.empty()){
      newNd->edge = curNd->edge;
      newNd->chd[2] = curNd->chd[2];
                                                                       int u=dq.front()
      newNd - > chd[3] = curNd - > chd[3];
                                                                           dq.pop_front();
                                                                       inq[u]=0;
                                                                       for(auto i:edge[u]){
    if(root->chd[0]->dep < root->chd[1]->dep)
      root->chd[0] = merge(root->chd[0],newNd);
                                                                         if(dis[i.first]>i.second+dis[u]){
                                                                            dis[i.first]=i.second+dis[u];
                                                                            len[i.first]=len[u]+1;
      root->chd[1] = merge(root->chd[1],newNd);
                                                                            if(len[i.first]>n) return 1;
     root->dep = max(root->chd[0]->dep, root->chd[1]->
                                                                            if(inq[i.first]) continue;
         dep) + 1;
                                                                            if(!dq.empty()&&dis[dq.front()]>dis[i.first])
    return root;
                                                                              dq.push_front(i.first);
  vector<heap*> V;
                                                                              dq.push_back(i.first);
  void build(){
    nullNd = new heap;
                                                                            inq[i.first]=1;
    nullNd->dep = 0;
                                                                     } } }
    nullNd->edge = new nd;
                                                                     return 0;
    fill(nullNd->chd, nullNd->chd+4, nullNd);
    while(not dfsQ.empty()){
                                                                  5.14 差分約束
       int u = dfsQ.front(); dfsQ.pop();
      if(!nxt[ u ]) head[ u ] = nullNd;
else head[ u ] = head[nxt[ u ]->v];
                                                                     約束條件 V_j - V_i \leq W 建邊 V_i - > V_j 權重為 W-> bellman-ford or spfa
                                                                        String
      V.clear();
       for( auto&& e : g[ u ] ){
                                                                  6.1 PalTree
         int v = e->v;
         if( dst[ v ] == -1 ) continue;
                                                                 |// len[s]是對應的回文長度
         e->d += dst[ v ] - dst[ u ];
if( nxt[ u ] != e ){
                                                                  // num[s]是有幾個回文後綴
           heap* p = new heap
                                                                  // cnt[s]是這個回文子字串在整個字串中的出現次數
           fill(p->chd, p->chd+4, nullNd);
                                                                  // fail[s]是他長度次長的回文後綴,aba的fail是a
           p->dep = 1;
                                                                  const int MXN = 1000010;
           p->edge = é:
                                                                  struct PalT{
           V.push_back(p);
                                                                     int nxt[MXN][26],fail[MXN],len[MXN];
                                                                     int tot,lst,n,state[MXN],cnt[MXN],num[MXN];
int diff[MXN],sfail[MXN],fac[MXN],dp[MXN];
      if(V.empty()) continue;
      make_heap(V.begin(), V.end(), cmp);
                                                                     char s[MXN] = \{-1\};
#define L(X) ((X<<1)+1)
#define R(X) ((X<<1)+2)
                                                                     int newNode(int 1,int f){
  len[tot]=1,fail[tot]=f,cnt[tot]=num[tot]=0;
                                                                       memset(nxt[tot],0,sizeof(nxt[tot]));
diff[tot]=(l>0?l-len[f]:0);
       for( size_t i = 0 ; i < V.size() ; i ++ ){</pre>
         if(L(i) < V.size()) V[i]->chd[2] = V[L(i)];
else V[i]->chd[2]=nullNd;
                                                                       sfail[tot]=(l>0&&diff[tot]==diff[f]?sfail[f]:f);
         if(R(i) < V.size()) V[i]->chd[3] = V[R(i)];
                                                                       return tot++;
         else V[i]->chd[3]=nullNd;
                                                                     int getfail(int x){
      head[u] = merge(head[u], V.front());
                                                                       while(s[n-len[x]-1]!=s[n]) x=fail[x];
  } }
                                                                       return x;
  vector<ll> ans
  void first_K(){
                                                                     int getmin(int v){
    ans.clear();
                                                                       dp[v]=fac[n-len[sfail[v]]-diff[v]];
    priority_queue<node> Q;
                                                                       if(diff[v]==diff[fail[v]])
```

dp[v]=min(dp[v],dp[fail[v]]);

```
return dp[v]+1;
  int push(){
    int c=s[n]-'a',np=getfail(lst);
    if(!(lst=nxt[np][c])){
      lst=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
      nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
    for(int v=lst;len[v]>0;v=sfail[v])
         fac[n]=min(fac[n],getmin(v));
    return ++cnt[lst],lst;
  void init(const char *_s){
    tot=lst=n=0:
    newNode(0,1), newNode(-1,1);
    for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}palt;
```

6.2 KMP

```
len-failure[k]:
在k結尾的情況下,這個子字串可以由開頭
長度為(len-failure[k])的部分重複出現來表達
failure[k]:
failure[k]為次長相同前綴後綴
如果我們不只想求最多,而且以0-base做為考量
,那可能的長度由大到小會是
failuer[k] \ failure[failuer[k]-1]
^ failure[failure[failuer[k]-1]-1]..
直到有值為0為止
int failure[MXN];
void KMP(string& t, string& p)
    if (p.size() > t.size()) return;
   for (int i=1, j=failure[0]=-1; i<p.size(); ++i)</pre>
    {
       while (j \ge 0 \&\& p[j+1] != p[i])
           j = failure[j];
       if (p[j+1] == p[i]) j++;
       failure[i] = j;
   for (int i=0, j=-1; i<t.size(); ++i)</pre>
       while (j >= 0 \&\& p[j+1] != t[i])
           j = failure[j];
       if (p[j+1] == t[i]) j++;
       if (j == p.size()-1)
           cout << i - p.size() + 1<<" ";
           j = failure[j];
  }
}
```

6.3 SAIS

```
const int N = 300010;
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )</pre>
#define REP1(i,a,b) for ( int i=(a); i<=int(b); i++ )
  bool _t[N*2];
  int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
  hei[N], r[N];
int operator [] (int i){ return _sa[i]; }
void build(int *s, int n, int m){
    memcpy(_s, s, sizeof(int) * n);
    sais(_s, _sa, _p, _q, _t, _c, n, m);
mkhei(n);
  void mkhei(int n){
    REP(i,n) r[\_sa[i]] = i;
    hei[0] = 0;
    REP(i,n) if(r[i]) {
  int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;
      while(_s[i+ans] == _s[_sa[r[i]-1]+ans]) ans++;
      hei[r[i]] = ans;
```

```
void sais(int *s, int *sa, int *p, int *q, bool *t,
       int *c, int n, int z){
    bool uniq = t[n-1] = true, neq;
    int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
         lst = -1;
#define MSO(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
    memcpy(x, c, sizeof(int) * z); \
    XD;
    memcpy(x + 1, c, sizeof(int) * (z - 1)); \
REP(i,n) if(sa[i] && !t[sa[i]-1]) sa[x[s[sa[i
         ]-1]]++] = sa[i]-1; \setminus
    MSO(c, z);
    REP(i,n) uniq \&= ++c[s[i]] < 2;
    REP(i,z-1) c[i+1] += c[i];
     if (uniq) { \overline{REP(i,n)} \overline{sa[--c[s[i]]]} = i; return; }
    for(int i = n - 2; i >= 0; i--) t[i] = (s[i]==s[i +1] ? t[i+1] : s[i]<s[i+1]);
    MAGIC(REP1(i,1,n-1) if(t[i] && !t[i-1]) sa[--x[s[i]]]=p[q[i]=nn++]=i);
    REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
       neq=lst<0|lmemcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa[i])|
            [i])*sizeof(int));
       ns[q[lst=sa[i]]]=nmxz+=neq;
    sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
    MAGIC(for(int i = nn - 1; i \ge 0; i--) sa[--x[s[p[
         nsa[i]]]] = p[nsa[i]];
  }
}sa;
int H[ N ], SA[ N ];
void suffix_array(int* ip, int len) {
  // should padding a zero in the back
  // ip is int array, len is array length
// ip[0..n-1] != 0, and ip[len] = 0
  ip[len++] = 0;
  sa.build(ip, len, 128);
for (int i=0; i<len; i++) {</pre>
    H[i] = sa.hei[i + 1];
    SA[i] = sa.\_sa[i + 1];
  // resulting height, sa array \in [0,len)
}
```

6.4 SuffixAutomata

```
// any path start from root forms a substring of S // occurrence of P : iff SAM can run on input word P \,
// number of different substring : ds[1]-1
// total length of all different substring : dsl[1]
// max/min length of state i : mx[i]/mx[mom[i]]+1
// assume a run on input word P end at state i:
// number of occurrences of P : cnt[i]
// first occurrence position of P : fp[i]-IPI+1
// all position of P : fp of "dfs from i through rmom"
const int MXM = 1000010;
struct SAM{
  int tot, root, lst, mom[MXM], mx[MXM]; //ind[MXM]
  int nxt[MXM][33]; //cnt[MXM],ds[MXM],dsl[MXM],fp[MXM]
  // bool v[MXM]
  int newNode(){
    int res = ++tot;
    fill(nxt[res], nxt[res]+33, 0);
mom[res] = mx[res] = 0; //cnt=ds=dsl=fp=v=0
    return res;
  void init(){
    tot = 0;
    root = newNode();
    lst = root;
  void push(int c){
    int p = lst;
    int np = newNode(); //cnt[np]=1
    mx[np] = mx[p]+1; //fp[np]=mx[np]-1
    for(; p && nxt[p][c] == 0; p = mom[p])
      nxt[p][c] = np;
```

```
if(p == 0) mom[np] = root;
    else{
      int q = nxt[p][c];
      if(mx[p]+1 == mx[q]) mom[np] = q;
      else{
        int nq = newNode(); //fp[nq]=fp[q]
        mx[nq] = mx[p]+1;
        for(int i = 0; i < 33; i++)
          nxt[nq][i] = nxt[q][i];
        mom[nq] = mom[q];
        mom[q] = nq;
        mom[np] = nq;
        for(; p && nxt[p][c] == q; p = mom[p])
          nxt[p][c] = nq;
    lst = np;
  }
  void calc(){
    calc(root);
    iota(ind,ind+tot,1)
    sort(ind,ind+tot,[&](int i,int j){return mx[i]<mx[j</pre>
        ];});
    for(int i=tot-1;i>=0;i--)
    cnt[mom[ind[i]]]+=cnt[ind[i]];
  void calc(int x){
    v[x]=ds[x]=1;dsl[x]=0; //rmom[mom[x]].push_back(x);
    for(int i=1; i<=26; i++){</pre>
      if(nxt[x][i]){
        if(!v[nxt[x][i]]) calc(nxt[x][i]);
ds[x]+=ds[nxt[x][i]];
        dsl[x]+=ds[nxt[x][i]]+dsl[nxt[x][i]];
  } } }
  void push(const string& str){
    for(int i = 0; i < str.size() ; i++)</pre>
      push(str[i]-'a'+1);
} sam;
6.5 Aho-Corasick
struct ACautomata{
```

```
struct Node{
  int cnt,i
  Node *go[26], *fail, *dic;
  Node (){
cnt = 0; fail = 0; dic=0;
    memset(go,0,sizeof(go));
}pool [1048576],*root;
int nMem,n_pattern;
Node* new_Node(){
  pool[nMem] = Node()
  return &pool[nMem++];
void init() {nMem=0;root=new_Node();n_pattern=0;}
void add(const string &str) { insert(root, str, 0); }
void insert(Node *cur, const string &str, int pos){
  for(int i=pos;i<str.size();i++){</pre>
    if(!cur->go[str[i]-'a'])
  cur->go[str[i]-'a'] = new_Node();
    cur=cur->go[str[i]-'a'];
  }
  cur->cnt++; cur->i=n_pattern++;
void make_fail(){
  queue<Node*> que;
  que.push(root);
  while (!que.empty()){
    Node* fr=que.front(); que.pop();
    for (int i=0; i<26; i++){
      if (fr->go[i]){
        Node *ptr = fr->fail;
        while (ptr && !ptr->go[i]) ptr = ptr->fail;
        fr->go[i]->fail=ptr=(ptr?ptr->go[i]:root);
        fr->go[i]->dic=(ptr->cnt?ptr:ptr->dic);
        que.push(fr->go[i]);
} } } }
void query(string s){
    Node *cur=root;
    for(int i=0;i<(int)s.size();i++){</pre>
        while(cur&&!cur->go[s[i]-'a']) cur=cur->fail;
```

```
cur=(cur?cur->go[s[i]-'a']:root);
if(cur->i>=0) ans[cur->i]++;
             for(Node *tmp=cur->dic;tmp;tmp=tmp->dic)
                 ans[tmp->i]++;
    }// ans[i] : number of occurrence of pattern i
}AC;
```

6.6 Z Value

```
char s[MAXN]
int len,z[MAXN];
void Z_value() { //z[i] = lcp(s[1...],s[i...])
  int i,j,left,right;
  left=right=0; z[0]=len;
  for(i=1;i<len;i++) {</pre>
    j=max(min(z[i-left],right-i),0);
    for(;i+j<len&&s[i+j]==s[j];j++);</pre>
    z[i]=i
    if(i+z[i]>right) {
      right=i+z[i];
      left=i;
```

6.7 ZValue Palindrome

```
void z_value_pal(char *s,int len,int *z){
  len=(len<<1)+1
  for(int i=len-1;i>=0;i--)
    s[i]=i&1?s[i>>1]:'@';
  z[0]=1;
  for(int i=1,l=0,r=0;i<len;i++){</pre>
    z[i]=i<r?min(z[l+l-i],r-i):1;
    while(i-z[i] \ge 0\&\&i+z[i] < len&\&s[i-z[i]] == s[i+z[i]])
         ++z[i];
    if(i+z[i]>r) l=i,r=i+z[i];
} }
```

Smallest Rotation 6.8

```
//rotate(begin(s),begin(s)+minRotation(s),end(s))
int minRotation(string s) {
  int a = 0, N = s.size(); s += s;
  rep(b,0,N) rep(k,0,N) {
     if(a+k == b | | s[a+k] < s[b+k])
     {b += max(0, k-1); break;}
if(s[a+k] > s[b+k]) {a = b; break;}
  } return a;
}
```

6.9 Cyclic LCS

pred[i][j]=L;

```
#define L 0
#define LU 1
#define U 2
const int mov[3][2]={0,-1, -1,-1, -1,0};
int al,bl;
char a[MAXL*2],b[MAXL*2]; // 0-indexed
int dp[MAXL*2][MAXL]
char pred[MAXL*2][MAXL];
inline int lcs_length(int r) {
  int i=r+al,j=bl,l=0;
  while(i>r) {
    char dir=pred[i][j];
    if(dir==LU) l++;
    i+=mov[dir][0];
    j+=mov[dir][1];
  return 1;
inline void reroot(int r) { // r = new base row
  int i=r, j=1;
  while(j<=bl&&pred[i][j]!=LU) j++;</pre>
  if(j>bl) return;
  pred[i][j]=L;
while(i<2*al&&j<=bl) {</pre>
    if(pred[i+1][j]==U) {
      pred[i][j]=L;
    } else if(j<bl&&pred[i+1][j+1]==LU) {</pre>
      i++;
      j++
```

```
} else {
       j++;
} } }
int cyclic_lcs() {
  // a, b, al, bl should be properly filled
  // note: a WILL be altered in process
               -- concatenated after itself
  char tmp[MAXL];
  if(al>bl) {
    swap(al,bl)
    strcpy(tmp,a);
    strcpy(a,b);
    strcpy(b,tmp);
  strcpy(tmp,a);
  strcat(a,tmp);
  // basic lcs
  for(int i=0;i<=2*al;i++) {
  dp[i][0]=0;</pre>
    pred[i][0]=U;
  for(int j=0;j<=bl;j++) {
  dp[0][j]=0;</pre>
    pred[0][j]=L;
  for(int i=1;i<=2*al;i++) {
    for(int j=1; j<=bl; j++)</pre>
       if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;
       else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
if(dp[i][j-1]==dp[i][j]) pred[i][j]=L;
else if(a[i-1]==b[j-1]) pred[i][j]=LU;
       else pred[i][j]=U;
 } }
// do cyclic lcs
  int clcs=0;
  for(int i=0;i<al;i++) {</pre>
    clcs=max(clcs,lcs_length(i));
    reroot(i+1);
  // recover a
  a[al]='\0':
  return clcs;
```

7 Data Structure

7.1 Segment tree

struct seg_tree{

```
static const int MXN=1e5+5,NO_TAG=0; // to be set
ll a[MXN], val[MXN*4], tag[MXN*4], v;
int n,ql,qr;
void push(int i,int l,int r){
  if(tag[i]!=NO_TAG){
    val[i]+=tag[i]; // update by tag
    if(l!=r)
      tag[cl(i)]+=tag[i]; // push
      tag[cr(i)]+=tag[i]; // push
    tag[i]=NO_TAG;
} }
void pull(int i,int l,int r){
  int mid=(l+r)>>1;
  push(cl(i),l,mid);push(cr(i),mid+1,r);
  val[i]=max(val[cl(i)],val[cr(i)]); // pull
void build(int i,int l,int r){
  if(l==r){
    val[i]=a[l]; // set value
    return:
  int mid=(l+r)>>1;
  build(cl(i),1,mid);build(cr(i),mid+1,r);
  pull(i,l,r);
void update(int i,int l,int r){
  push(i,l,r);
  if(ql \le l\&r \le qr){
    tag[i]+=v; // update tag
    return;
  int mid=(l+r)>>1;
```

```
19
     if(ql<=mid) update(cl(i),l,mid);</pre>
    if(qr>mid) update(cr(i),mid+1,r);
    pull(i,l,r);
  void query(int i,int l,int r){
    push(i,1,r);
     if(ql <= l\&r <= qr){
      v=max(v,val[i]); // update answer
      return:
     int mid=(l+r)>>1;
    if(ql<=mid) query(cl(i),l,mid);</pre>
    if(qr>mid) query(cr(i),mid+1,r);
} }tree;
7.2 Treap
struct Treap{
  int sz , val , pri , tag;
Treap *l , *r;
  Treap( int _val ){
    val = _val; sz = 1;
    pri = rand(); l = r = NULL; tag = 0;
};
void push( Treap * a ){
  if( a->tag ){
    Treap *swp = a -> 1; a -> 1 = a -> r; a -> r = swp;
    int swp2;
    if( a->l ) a->l->tag ^= 1;
if( a->r ) a->r->tag ^= 1;
    a \rightarrow tag = 0;
} }
inline int Size( Treap * a ){ return a ? a->sz : 0; }
void pull( Treap * a ){
  a -> sz = Size(a -> l) + Size(a -> r) + 1;
Treap* merge( Treap *a , Treap *b ){
  if( !a | | | b ) return a ? a : b;
  if( a->pri > b->pri ){
    push( a );
    a \rightarrow r = merge(a \rightarrow r, b);
    pull( a );
    return a:
  }else{
    push( b );
    b->l = merge(a, b->l);
    pull( b );
    return b;
void split_kth( Treap *t , int k, Treap*&a, Treap*&b ){
  if( !t ){ a = b = NULL; return; }
  push( t )
  if( Size( t -> l ) + 1 <= k ){
    a = t:
    split_kth( t->r , k - Size( t->l ) - 1 , a->r , b )
    pull( a );
  }else{
    b = t;
    split_kth( t->l , k , a , b->l );
    pull( b );
void split_key(Treap *t, int k, Treap*&a, Treap*&b){
  if(!t){ a = b = NULL; return; }
  push(t);
  if(k \le t - val)
    b = t;
split_key(t->l,k,a,b->l);
    pull(b);
  }
  else{
    split_key(t->r,k,a->r,b);
    pull(a);
} }
7.3 Disjoint Set
```

struct DisjointSet{

// save() is like recursive
// undo() is like return

```
int n, fa[ N ], sz[ N ];
vector< pair<int*,int> > h;
vector<int> sp;
void init( int tn ){
  n=tn;
  for( int i = 0 ; i < n ; i ++ ){</pre>
    fa[ i ]=i;
sz[ i ]=1;
  sp.clear(); h.clear();
void assign( int *k, int v ){
  h.PB( {k, *k} );
  *k = v;
}
void save(){ sp.PB(SZ(h)); }
void undo(){
  assert(!sp.empty());
  int last=sp.back(); sp.pop_back();
  while( SZ(h)!=last ){
     auto x=h.back(); h.pop_back();
     *x.first = x.second;
void uni( int x , int y ){
  x = f(x); y = f(y);

if(x == y) return;

if(sz[x] < sz[y]) swap(x, y);
  assign(\&sz[x], sz[x] + sz[y]);
  assign( &fa[ y ] , x);
} }djs;
```

7.4 Black Magic

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
typedef tree<int,null_type,less<int>,rb_tree_tag,
    tree_order_statistics_node_update> set_t;
#include <ext/pb_ds/assoc_container.hpp>
typedef cc_hash_table<int,int> umap_t;
typedef priority_queue<int> heap;
#include<ext/rope>
using namespace __gnu_cxx;
int main(){
 // Insert some entries into s.
  set_t s; s.insert(12); s.insert(505);
  // The order of the keys should be: 12, 505.
 assert(*s.find_by_order(0) == 12);
 assert(*s.find_by_order(3) == 505);
  // The order of the keys should be: 12, 505.
 assert(s.order_of_key(12) == 0);
 assert(s.order_of_key(505) == 1);
  // Erase an entry.
 s.erase(12);
 // The order of the keys should be: 505.
 assert(*s.find_by_order(0) == 505);
 // The order of the keys should be: 505.
 assert(s.order_of_key(505) == 0);
 heap h1 , h2; h1.join( h2 );
 rope<char> r[ 2 ];
r[ 1 ] = r[ 0 ]; // persistenet
string t = "abc";
 r[1].insert(0, t.c_str());
 r[1].erase(1,1);
cout << r[1].substr(0,2);
```

8 Others

8.1 Find max tangent(x,y is increasing)

```
const int MAXN = 100010;
Pt sum[MAXN], pnt[MAXN], ans, calc;
inline bool cross(Pt a, Pt b, Pt c){
   return (c.y-a.y)*(c.x-b.x) > (c.x-a.x)*(c.y-b.y);
}//pt[0]=(0,0);pt[i]=(i,pt[i-1].y+dy[i-1]),i=1~n;dx>=l
double find_max_tan(int n,int l,LL dy[]){
   int np, st, ed, now;
   sum[0].x = sum[0].y = np = st = ed = 0;
   for (int i = 1, v; i <= n; i++)
      sum[i].x=i,sum[i].y=sum[i-1].y+dy[i-1];</pre>
```

```
ans.x = now = 1,ans.y = -1;
for (int i = 0; i <= n - 1; i++){
    while(np>1&&cross(pnt[np-2],pnt[np-1],sum[i]))
        np--;
    if (np < now && np != 0) now = np;
    pnt[np++] = sum[i];
    while(now<np&&!cross(pnt[now-1],pnt[now],sum[i+l]))
        now++;
    calc = sum[i + l] - pnt[now - 1];
    if (ans.y * calc.x < ans.x * calc.y)
        ans = calc,st = pnt[now - 1].x,ed = i + l;
}
return (double)(sum[ed].y-sum[st].y)/(sum[ed].x-sum[st].x);
}</pre>
```

8.2 Exact Cover Set

```
// given n*m 0-1 matrix
// find a set of rows s.t.
\ensuremath{//} for each column, there's exactly one 1
#define N 1024 //row
#define M 1024 //column
#define NM ((N+2)*(M+2))
char A[N][M]; //n*m 0-1 matrix
int used[N]; //answer: the row used
int id[N][M];
int L[NM],R[NM],D[NM],U[NM],C[NM],S[NM],ROW[NM];
void remove(int c){
  L[R[c]]=L[c]; R[L[c]]=R[c];
  for( int i=D[c]; i!=c; i=D[i] )
  for( int j=R[i]; j!=i; j=R[j] )
       U[D[j]]=U[j]; D[U[j]]=D[j]; S[C[j]]--;
void resume(int c){
  for( int i=D[c]; i!=c; i=D[i] )
  for( int j=L[i]; j!=i; j=L[j] ){
    U[D[j]]=D[U[j]]=j; S[C[j]]++;
  L[R[c]]=R[L[c]]=c;
int dfs(){
  if(R[0]==0) return 1;
  int md=100000000,c;
  for( int i=R[0]; i!=0; i=R[i] )
  if(S[i]<md){ md=S[i]; c=i; }</pre>
  if(md==0) return 0;
  remove(c);
  for( int i=D[c]; i!=c; i=D[i] ){
     used[ROW[i]]=1;
     for( int j=R[i]; j!=i; j=R[j] ) remove(C[j]);
     if(dfs()) return 1
     for( int j=L[i]; j!=i; j=L[j] ) resume(C[j]);
     used[ROW[i]]=0;
  resume(c);
  return 0;
int exact_cover(int n,int m){
  for( int i=0; i<=m; i++ ){
    R[i]=i+1; L[i]=i-1; U[i]=D[i]=i;</pre>
     S[i]=0; C[i]=i;
  R[m]=0; L[0]=m;
  int t=m+1;
  for( int i=0; i<n; i++ ){
     int k=-1;
     for( int j=0; j<m; j++ ){
   if(!A[i][j]) continue;</pre>
        if(k==-1) L[t]=R[t]=t;
        else{ L[t]=k; R[t]=R[k]; }
k=t; D[t]=j+1; U[t]=U[j+1];
        L[R[t]]=R[L[t]]=U[D[t]]=D[U[t]]=t;
        C[t]=j+1; S[C[t]]++; ROW[t]=i; id[i][j]=t++;
   for( int i=0; i<n; i++ ) used[i]=0;</pre>
  return dfs();
```