

Qubit Mapping and Routing tailored to Advanced Quantum ISAs: Not as Costly as You Think

Abstract

Qubit mapping/routing is critical for executing quantum circuits NISQ and fault-tolerant hardware. Existing scalable methods typically impose several times the routing overhead, especially in circuit depth or duration, due to the disconnect between simplified routing models (e.g., three-CX-unrolled SWAP insertion) and hardware-native gate properties. This gap severely limits compiler optimization potential and thus practical circuit execution performance.

Recent advances in quantum hardware have enabled high-precision implementations of diverse instruction set architectures (ISAs) beyond standard CX-based gates. Advanced ISAs such as \sqrt{i} SWAP and ZZ(θ) offer superior circuit synthesis capabilities while retaining hardware-native noise resilience. However, there are neither systematic compiler optimization strategies tailored to these advanced ISAs nor comprehensive cross-ISA evaluation to unlock their potential.

To address this, we propose CANOPUS, a unified qubit mapping/routing framework tailored to diverse quantum ISAs. With the canonical representation of two-qubit gates, CANOPUS revolves around qubit routing to perform deep co-optimization in an ISA-aware approach. CANOPUS utilizes the canonical forms of arbitrary basis gates and monodromy polytope theory to model the synthesis cost for better SWAP search during routing. Commutation relations between two-qubit gates can be formalized through the canonical representation in our findings, providing a generalized approach to capturing commutative optimizations. Experiments show that CANOPUS consistently results in 30%-40% reduction of routing overhead compared to state-of-the-art methods across versatile ISAs and topologies. Our work also presents a coherent method for co-exploration of program patterns, quantum ISAs, and hardware topologies. Our framework can be extended to integrate more fine-grain hardware information such as qubit-specific basis gate fidelities. We have for the first time demonstrated that advanced quantum ISAs can be efficiently utilized within a unified routing framework, paving the way for more effective co-design of quantum software and hardware.

1 Introduction

Advanced ISAs—Can [4], \sqrt{i} SWAP [13]

CANOPUS (**C**anonical-**O**ptimized **P**lacement **U**tility **S**uite) is a qubit mapping and routing framework that is tailored to advanced quantum ISAs, such as Can [4] and \sqrt{i} SWAP [13], which are adaptive to versatile hardware architectures. CANOPUS is designed to optimize the placement of qubits and the

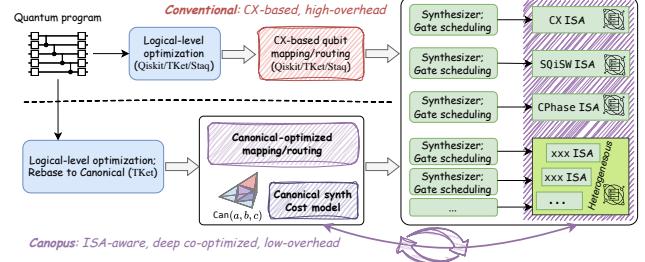


Figure 1. Compilation workflows by means of conventional approaches (top) and CANOPUS (bottom) targeting diverse quantum ISAs. CANOPUS integrates the synthesis cost model (monodromy polytopes within the Weyl chamber) by taking backend ISAs’ synthesis properties into account. CANOPUS routing operates in the 2Q canonical representation while the specific synthesis is completed by backend synthesizer.

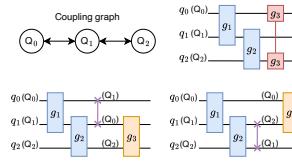


Figure 2. Mapping/routing to resolve physical-qubit topology constraints via SWAP insertion.

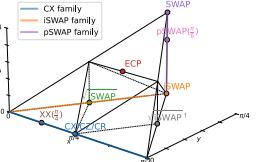


Figure 3. Geometric illustration of canonical gates confined to the Weyl chamber.

routing of quantum gates, taking into account the specific requirements of these advanced ISAs.

Our work addresses the “Babel Tower dilemma” in quantum compilation by establishing a canonical language for diverse two-qubit gates, enabling unified optimization across heterogeneous quantum ISAs. Our key contributions are summarized as follows:

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2 Background

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¹For convenient visualization

2.1 Qubit mapping/routing

2.2 Quantum gates realization in diverse ISAs

Definition 1 (Canonical gate). Any $2Q$ gate $U \in \text{SU}(4)$ can be expressed by the composition of its unique canonical form

$$\text{Can}(a, b, c) := e^{-i\frac{\pi}{2}(aXX+bYY+cZZ)}, \quad \frac{1}{2} \geq a \geq b \geq |c| \quad (1)$$

sandwiched by local $1Q$ gates such that we call U is locally equivalent to the canonical form $\text{Can}(a, b, c)$.

Canonical representation is ubiquitous as an effective ...

[ZY: It is ubiquitously used in many quantum computing task ...]

Although there are other conventions This definition aligns with the TK2 operation definition in TKET, ...

Figure 3

$$\text{pSWAP}(\theta) \sim \text{Can}\left(\frac{1}{2}, \frac{1}{2}, \frac{1}{2} - \frac{\theta}{\pi}\right) \quad (2)$$

$$XX(\theta) = \text{Can}\left(\frac{\theta}{\pi}, 0, 0\right) \sim YY(\theta) \sim ZZ(\theta) \quad (3)$$

3 Motivation

Two-fold motivations:

1. The scalable qubit routing effects (2x-4x) is still a critical challenge in practical quantum computing systems
2. How to utilize the emerging advanced ISAs (hardware breakthroughs); across all phases of compilation, routing is the bottleneck and is the most easily handled for co-optimization

[ZY: Use a “optimal routing benchmark” to illustrate the OVERHEAD of existing methods]

[ZY: There should be many takeaways]

- Previous routing overhead is not precise for hardware execution
- Previous routing is costly and also not precise for hardware execution
- SWAP can be implemented in low overhead (gate duration) with the recent breakthrough gate schemes for advanced ISAs
- How to efficiently capture the rich commutation relations when performing co-optimization during qubit routing and gate scheduling

4 CANOPUS framework

4.1 Overview

4.2 2Q synthesis cost modeling

4.3 Routing in canonical form

In contrast to the regular heuristic cost function used in SABRE:

$$H = \frac{1}{|F|} \sum_{(i,j) \in F} \text{dist}[i,j] + \frac{k_E}{|E|} \sum_{(i,j) \in E} \text{dist}[i,j] \\ = \text{Avg}\{\text{dist}[i,j]\}_F + k_E \text{Avg}\{\text{dist}[i,j]\}_E \quad (4)$$

which involves the basic (left term) and lookahead (right term) components. In practice, there is a w_{decay} decay factor applied to H , which is not shown as it does not affect the composition of H .

The heuristic cost function in CANOPUS is defined as:

$$H = w_d \Delta_{\text{depth}} + c_0 (\Delta_{\text{Avg}\{\text{dist}[i,j]\}_F} + k_E \Delta_{\text{Avg}\{\text{dist}[i,j]\}_E}) \quad (5)$$

- Unified and highly-effective qubit routing approach in canonical form, with properties of quantum ISAs tailored to the routing process

4.4 Enhanced optimization via commutation

- Capture optimization opportunities exposed by gate commutation; while commutation relations can be uniformly described in canonical form

Theorem 1 (Canonical gate commutation). Let $\text{Can}(a, b, c)_{q_0, q_1}$ and $\text{Can}(a', b', c')_{q_1, q_2}$ denote canonical gates acting on qubits (q_0, q_1) and (q_1, q_2) respectively, with an overlapping qubit q_1 . They are commutative if and only if

$$b = b' = c = c' = 0, \quad (6)$$

that is, when both consist solely of XX rotations.

[ZY: Proposition?? Theorem?]

5 Implementation

5.1 Core functionalities

5.2 Extensions

5.3 Scalability

6 Case Studies

6.1 QFT kernel

6.2 Co-exploration of routing and ISA selection

6.3 Stabilizer circuit

6.4 Mapping on FTQC architecture

7 Evaluation

We further holistically evaluate CANOPUS compared to other scalable SOTA methods, with respect to representative quantum ISAs and hardware topologies. The evaluation not only provides cross-compiler but also cross-ISA comparisons under the coherent basis gate cost and routing overhead settings.

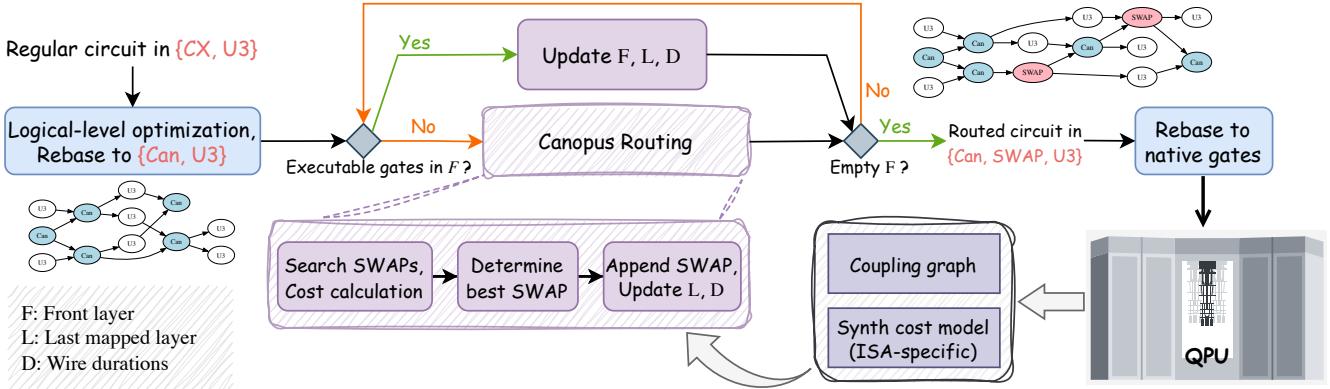


Figure 4. Overview of the CANOPUS framework. ...

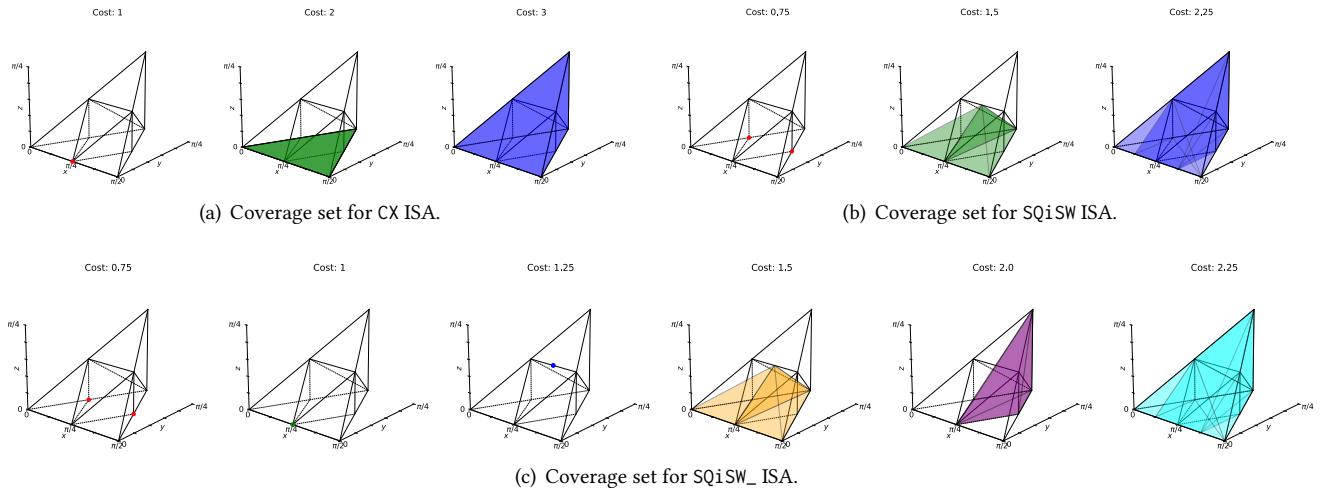


Figure 5. Coverage set examples. CX ISA: $\{\text{CX}, \text{U3}\}$ gate set; SQiSW ISA: $\{\sqrt{i\text{SWAP}}, i\text{SWAP}, \text{U3}\}$ gate set; SQiSW_ ISA: $\{\sqrt{i\text{SWAP}}, i\text{SWAP}, \text{ECP}, \text{CX}, \text{U3}\}$ gate set. Costs of Basis 2Q gates are set as CX ~ 1 , $\sqrt{i\text{SWAP}} \sim 0.75$, iSWAP ~ 1.5 , ECP ~ 1.25 .

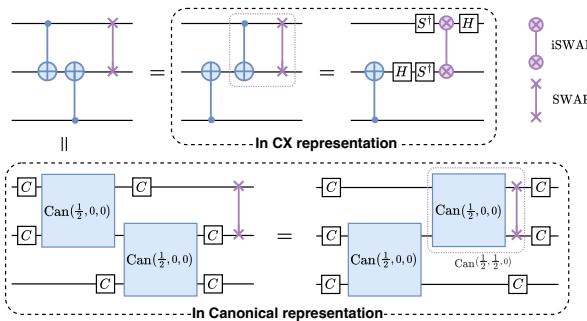


Figure 6. Canonical gate representation enables easily capturing commutative relations within real-world circuits.

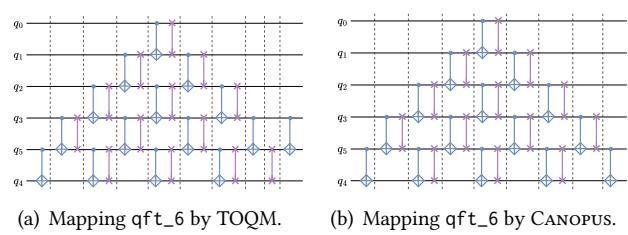


Figure 7. [ZY: Re-draw this figure] Mapping/routing comparison for the QFT kernel. For convenient visualization, only CPhase and SWAP gates are shown. (a) TOQM generates a sub-optimal mapping scheme, with 2Q depth of 10. (b) CANOPUS generates the optimal scheme in a perfect butterfly structure, with 2Q depth of 9.

Algorithm 1: Update L when adding a new 2Q gate

Input : G' (Routed DAG), π (current logic-to-physical mapping), L (last mapped layer), D (wire durations for each qubit), C (commutative pairs within L)

Output: Updated G' , L , D , C

```

/* g: resolved logical gate; g': routed gate */
1  $g' \leftarrow G'.PUSHBACK(g, \pi[g.q_0], \pi[g.q_1]); // g'.q_i = \pi[g.q_i]$ 
2  $d \leftarrow \max(D[g'.q_0], D[g'.q_1]) + \text{SYNTHCOST}(g);$ 
3  $D[g'.q_0] \leftarrow d; D[g'.q_1] \leftarrow d;$ 
4 for pred  $\in G'.\text{PREDECESSORS}(g')$  do
5   if IS2QGATE(pred) then
6     if ISCOMMUTATIVECANONICALPAIR(g', pred) then
7        $C[(\text{pred}.q_0, \text{pred}.q_1)] \leftarrow (g'.q_0, g'.q_1);$ 
8     else
9        $L.\text{POP}((\text{pred}.q_0, \text{pred}.q_1), \text{NONE});$ 
10       $C.\text{POP}((\text{pred}.q_0, \text{pred}.q_1), \text{NONE});$ 
11    else
12      /* pred_pred must be None or a 2Q gate */
13      pred_pred  $\leftarrow \text{NEXT}(G'.\text{PREDECESSORS}(\text{pred}))$ ;
14      if pred_pred  $\neq \text{NONE}$  then
15         $L.\text{POP}((\text{pred\_pred}.q_0, \text{pred\_pred}.q_1), \text{NONE});$ 
16         $C.\text{POP}((\text{pred\_pred}.q_0, \text{pred\_pred}.q_1), \text{NONE});$ 
16  $L[(g'.q_0, g'.q_1)] \leftarrow g';$ 

```

Algorithm 2: Update D when adding a SWAP gate

Input : swap (encountered SWAP gate), can (canonical gate within L on the same qubits as swap), D , C

Output: Updated D

```

1 if (swap.q0, swap.q1)  $\in C$  then
2    $q'_0, q'_1 \leftarrow C[(\text{swap}.q_0, \text{swap}.q_1)];$ 
3   /* Adjust D by finding matched qubits
4      $q_i \in \{\text{swap}.q_0, \text{swap}.q_1\}$  and  $q'_j \in \{q'_0, q'_1\}$  */
5    $D[q_i] \leftarrow D[q'_j] + \text{SYNTHCOST}(\text{can});$ 
6    $D[\text{the other swap qubit}] \leftarrow D[q_i];$ 
5  $d \leftarrow \max(D[\text{swap}.q_0], D[\text{swap}.q_1]) +$ 
6    $\text{SYNTHCOST}(\text{can.MIRROR}()) - \text{SYNTHCOST}(\text{can});$ 
6  $D[\text{swap}.q_0] \leftarrow d; D[\text{swap}.q_1] \leftarrow d;$ 

```

Table 1. Qubit routing comparison for the QFT kernel.

QFT kernel		qft_6		qft_12	
Topology	Method	#Can	Depth2Q	#Can	Depth2Q
1D Chain	Optimal	15	9	66	21
	TOQM	16	10	67	22
	CANOPUS	15	9	67	21
2D Square	TOQM	21	13	100	39
	CANOPUS	15	9	80 ($\pm 10\%$)	30 ($\pm 10\%$)

7.1 Experimental settings

7.1.1 ISAs and basis gate costs. We consider six different ISAs (including the conventional CX ISA) listed in Table 2.

Table 2. Selected quantum ISAs.

ISA	2Q basis gates	Description
CX	{CX}	Conventional CX gate
ZZPhase	$\{\text{ZZ}_{\frac{\pi}{6}}, \text{ZZ}_{\frac{\pi}{4}}, \text{ZZ}_{\frac{\pi}{2}}\}$	Discrete CX-family gates, i.e., $\{\sqrt[3]{\text{CX}}, \sqrt{\text{CX}}, \text{CX}\}$ [24]
SQiSW	$\{\sqrt{i\text{SWAP}}, i\text{SWAP}\}$	Half evolution of iSWAP and iSWAP [13]
ZZPhase __	ZZPhase + $\{\text{pSWAP}_{\frac{\pi}{6}, \frac{\pi}{4}, \frac{\pi}{2}}\}$	ZZPhase ISA with the mirror gates
SQiSW __	SQiSW + {ECP, CX}	SQiSW ISA with the mirror gates [23]
Het	ZZPhase + SQiSW	Heterogeneous CX-family and iSWAP-family gates

These mainly cover a wide range of powerful basis gates from CX-family and iSWAP-family gates. Particularly, SQiSW [13] proves to a more powerful ISA option and has been adopted by recent software projects [11, 23]. ZZPhase ISA containing three fractional $\text{ZZ}(\theta)$ rotation gates is adopted by QISKit’s latest synthesis functionalities [15, 24]. Mirorr [ZY: TODO: find the initial paper about mirror gate]

[23]

We also involve the Het ISA that is the composition of ZZPhase and SQiSW.

The unit costs for the involved basis gates are set as:

$$\left\{ \begin{array}{l} \text{CX : } 1, \text{ZZ}(\pi/t) : 2/t, \sqrt{i\text{SWAP}} : 0.75, \\ \text{iSWAP : } 1.5, \text{ECP : } 1.25, \text{pSWAP}(\pi/t) : 2 - 1/t \end{array} \right\} \quad (7)$$

[ZY: Plot a weyl chamber to illustrate the cost settings]

7.1.2 Benchmarks. We select a set of medium-size benchmarks from QASMBench [20] and MQTBench [26] spanning various categories of quantum programs. These benchmarks first go through logical-level optimization by TKET [27] and are rebased to {Can, U3} as the input of qubit routing compilers. Information of benchmarks after logical-level optimization are summarized in Table 3, where “Circuit cost” denotes the circuit duration assuming each canonical gate will be finally rebased to CX ISA and the duration (cost) of each CX is set to 1.

7.1.3 Baselines.

7.2 Circuit duration reduction

7.3 Co-exploration of routing and ISA selection

7.4 Breakdown analysis

7.4.1 Commutative optimization. In this section, we analyze individual factors in the improvement brought by CANOPUS, mainly about the commutative optimization mechanism and the heuristic depth-cost weight factor.

7.4.2 Depth weight factor. Note that the ...

[ZY: about comm_opt effect]

7.5 Real-system experiments

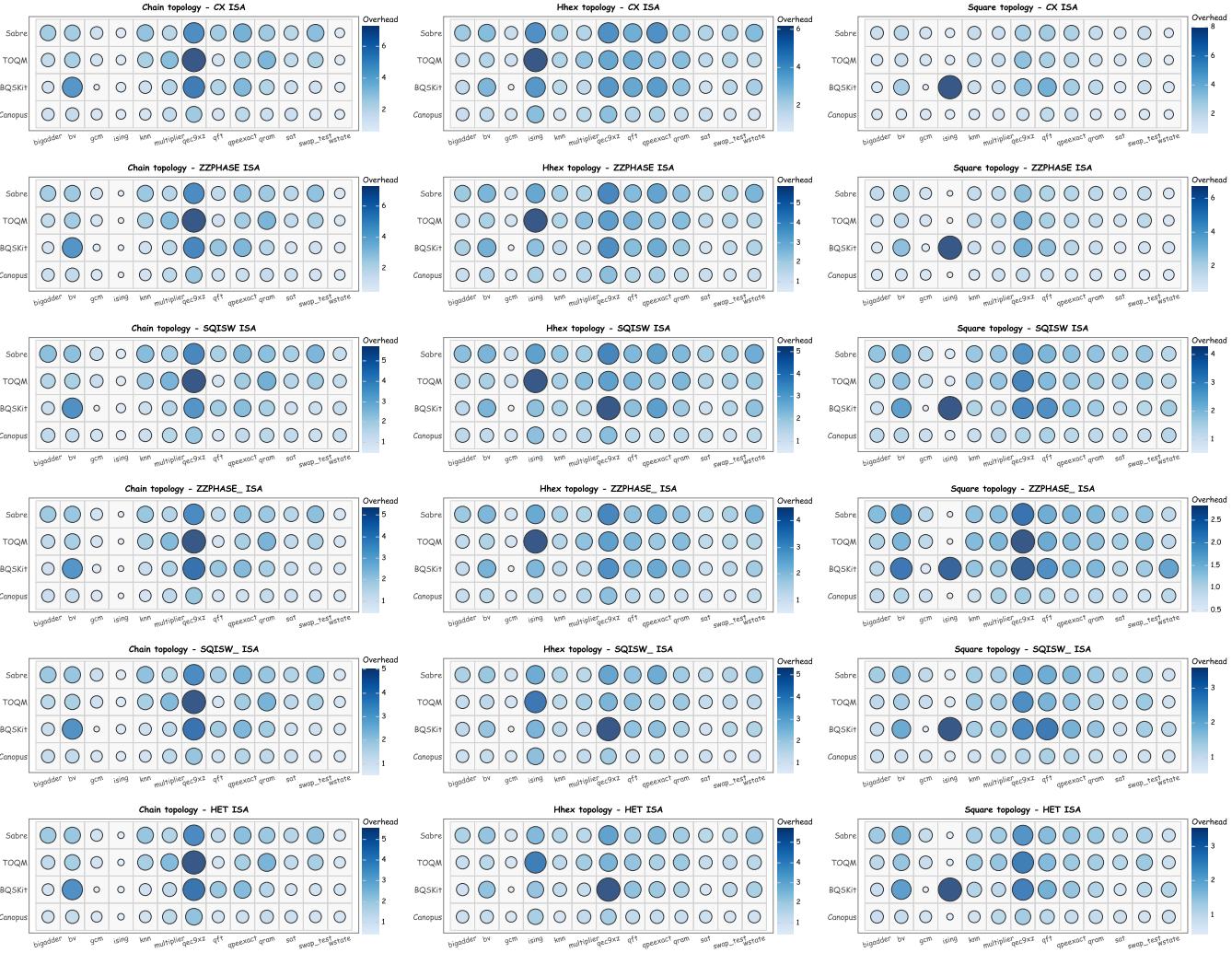


Figure 8. Detailed benchmarking results for different compilers (SABRE, TOQM, BQSKit, CANOPUS) across diverse topologies (columns from left to right: 1D Chain, 2D HHex, 2D Square) and quantum ISAs (rows from top to bottom: CX, ZZPhase, SQiSW, ZZPhase_{_}, SQiSW_{_}, Het). Sizes and colors of the bubbles represent the values for routing overhead (multiples of the routed and rebased circuits compared to that of prior-routing ones).

[ZY: Fractional gate IBM]
[ZY: QFT or QV test?]

7.6 For specific gate scheme

[ZY: For AshN gate scheme]

7.7 Runtime analysis

In our field tests for the set of benchmarks above, CANOPUS exhibits around 1x-2x compilation latency than SABRE, both of which are implemented by QISKIT framework. This result aligns with the complexity analysis in Section 5.3. Herein we specifically demonstrate the end-to-end compilation latency for larger-scale quantum circuits. We use random quantum

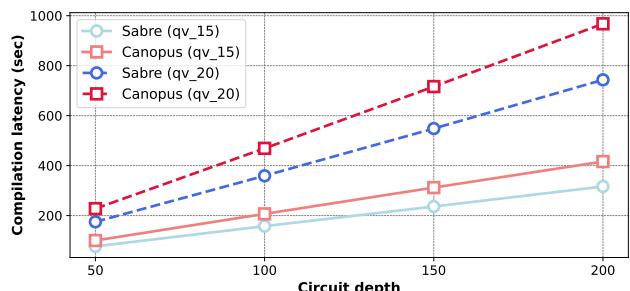


Figure 9. Compilation latency comparison.

Table 3. Benchmarks information. These metrics are collected from the circuits after logical-level optimization by TKET, thus including only Can and U3 gates. Circuit cost (Duration T) is calculated in CX ISA.

Program	#Qubit	#Can	Depth2Q	Circuit cost
bigadder [20]	18	114	79	88.0
bv [20]	19	18	18	18.0
gcm [20]	13	377	376	510.0
ising [20]	26	25	2	4.0
knn [20]	25	72	50	62.0
multiplier [20]	15	198	122	133.0
qec9xz [20]	17	32	12	12.0
qft [26]	18	153	33	66.0
qpeexact [26]	16	127	43	86.0
qram [20]	20	110	70	78.0
sat [20]	11	210	182	204.0
swap_test [20]	25	72	50	62.0
wstate [20]	27	52	28	28.0

volume [8] circuits generated by QISKit for scalability benchmarking, which represents the end of the spectrum with respect to canonical-form circuits. Each canonical gate within the quantum volume circuit contains unique canonical parameters as each 2Q unitary is randomly generated, thus there is no cached synthesis cost calculation for performance improvement in one pass. In the cases of circuit width (number of qubits) set to 15 and 20, we vary the depth of benchmarked circuits from 50 to 200. Quantum volume circuits consists of dense 2Q gates and the largest size for benchmarking is up to thousands of 2Q gates. Figure 9 illustrates the end-to-end compilation latencies, where each data point is tested with the same trial setting (`max_iterations` is 5, both `trials` and `layout_trials` are 10). For each benchmarked circuit, CANOPUS leads to on average 1.31x ($\pm 1\%$) latency than SABRE. Both compilers’ latency scales linearly with circuit depth and width. If we compares the curve slopes, CANOPUS leads to 1.32x (1.30x) latency scaling than SABRE in terms of circuit depth for `qv_10` (`qv_20`) circuits. Overall, although CANOPUS involves sophisticated data structures and calculation mechanisms, its practical compilation scalability is comparable to the industrial-level SABRE algorithm.

8 Related Works

Qubit mapping/routing is one the the most well-explored topic of quantum compiler research, as it shares the similar methodologies with instruction scheduling [6, 12] and register allocation [3, 25] in classical computing. Conventional methods focus on the simplified routing model, that is, #SWAP-minimal insertion, three-CX-unrolled SWAP gate, and CX-based latency metric. That brings a gap between quantum hardware performance and its ceiling, which is particularly evident with the progress of underlying instruction models for modern quantum hardware.

Zulehner et al. [38] introduces an A*-based algorithm to minimize SWAP gate overhead for concurrent CNOT gate layers. The approach partitions the circuit into layers and solves the mapping problem subsequently. Li et al. [21] also utilizes the circuit DAG layering thought to tackle the qubit mapping problem and proposes the bidirectional routing procedure to acquire a better initial mapping desired to result in #SWAP inserted minimization as expected. It also briefly discusses the trade-off between the inserted SWAP count and the circuit depth but does not prioritize optimizing circuit depth. Some other works leverage algorithmic procedures similar to SABRE to improve parallelism among inserted SWAPs and other 2Q gates [1, 19, 37], or attempt to minimize circuit depth via graph matching [5]. Zhang et al. [34] systematically investigates the time (circuit depth) optimality of qubit mapping and proposed an A*-based method TOQM that results in better results than the SOTA solver-based depth-driven algorithm [28]. However, the optimality of qubit routing is a complex task. There are rarely theoretical studies that claims the holistic optimality of some SWAP insertion schemes provided the quantum ISAs, device topologies, and synthesis cost models. In our field tests, TOQM does not lead to time-optimal results compared to our heuristic CANOPUS, and the optimal mapping scheme for specific patterns such as QFT kernel analyzed in [34] are not indeed optimal, according to our case study in Section 6.1.

With the recent development of advanced quantum ISAs such as superconducting fractional gates [14], ion-trapped partial entangling gates [16, 31], and the AshN scheme [4], some works began exploring how to efficiently utilize these ISAs to make compiler optimizations closer to hardware characteristics. McKinney et al. [23] investigates the practical performance of SQiSW ISA proposed by Huang et al. [13] and the synthesis capability when incorporating the basis gates’ mirrors into the ISA. The modified SABRE algorithm in [23] provides an attempt of the collaborative gate decomposition and qubit routing approach, while the optimization opportunities considered therein are limited and the algorithmic techniques are not sophisticated. BQSKit [32] and the series of works behind [9, 18, 30, 33] provides a toolkit to rebase arbitrary 2Q unitaries to specific ISAs through approximate synthesis (structural search and numerical optimization) that is not computational efficient. Approximate synthesis by BQSKit does not ensure an optimal schemes for two-qubit and multi-qubit synthesis cases. In addition, due to the lack of native compilation strategies and rational synthesis cost model, Kalloor et al. [17] claims that alternative ISAs are hard to be comparable to CX when evaluating quantum hardware roofline by BQSKit. As for applicability of expanded ISAs to QEC, Google’s latest theoretical [22] and experimental works [10] demonstrate the CX-i SWAP combination ISA could benefits suppressing fault-tolerant threshold. Zhou et al. [36] proposes a routing-based method enhanced

Table 4. Summarized results for the routing overhead with average (geometric-mean) values emphasized.

Topology	Routing overhead (CX)				Routing overhead (ZZPhase)			
	SABRE	TOQM	BQSKIT	CANOPUS	SABRE	TOQM	BQSKIT	CANOPUS
Chain	2.93x	2.67x	2.34x	1.73x	2.63x	2.51x	2.17x	1.61x
HHex	2.90x	2.56x	2.66x	1.81x	2.62x	2.37x	2.26x	1.66x
Square	2.14x	2.04x	2.47x	1.42x	1.87x	1.82x	1.94x	1.19x
Routing overhead (SQiSW)				Routing overhead (ZZPhase_)				
Topology	SABRE	TOQM	BQSKIT	CANOPUS	SABRE	TOQM	BQSKIT	CANOPUS
Chain	2.57x	2.27x	2.00x	1.55x	2.08x	1.94x	1.87x	1.26x
HHex	2.59x	2.33x	2.28x	1.64x	2.08x	1.90x	1.96x	1.33x
Square	2.05x	1.94x	2.02x	1.39x	1.54x	1.49x	1.65x	1.00x
Routing overhead (SQiSW_)				Routing overhead (Het)				
Topology	SABRE	TOQM	BQSKIT	CANOPUS	SABRE	TOQM	BQSKIT	CANOPUS
Chain	2.17x	1.93x	1.78x	1.38x	2.13x	2.00x	1.74x	1.29x
HHex	2.17x	1.93x	2.10x	1.42x	2.13x	1.94x	1.98x	1.37x
Square	1.67x	1.58x	1.83x	1.15x	1.58x	1.52x	1.56x	1.02x

by CX-iSWAP for overcoming ancilla defects among surface code blocks while preserving encoded logical information.

9 Conclusion

It is promising to explore novel Clifford circuit optimization techniques drawing on the canonical gate representation.

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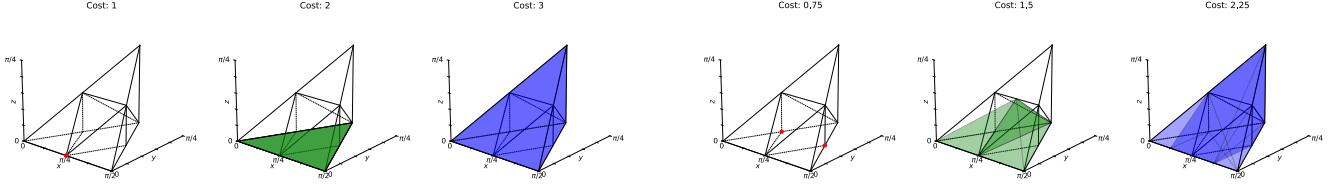


Figure 10. Coverage set for CX ISA.

Figure 11. Coverage set for SQiSW ISA.

A Canonical gate and 2Q circuit synthesis

In this section we show the basic mathematical properties of its canonical form of 2Q unitary and then discuss the synthesis capability of some 2Q basis gates.

A.1 Canonical decomposition

$SU(N)$ is a real manifold with dimension $N^2 - 1$, within which any element is a *special unitary* matrix with determinant equal to 1. Since the global phase does not affect quantum computation processes, it is sufficient to focus on the mathematical properties of special unitaries in the area of circuit synthesis. A generic 2Q gate, despite having 15 real parameters, can have its nonlocal behavior fully characterized by only 3 real parameters. This method, known as *Canonical decomposition* or *KAK decomposition* from Lie algebra theory, is widely adopted in quantum computing [2, 29, 35, 39]. Specifically, for any $U \in SU(4)$, there exists a unique $\vec{\eta} = (x, y, z) \in W \subseteq \mathbb{R}^3$, along with $V_1, V_2, V_3, V_4 \in SU(2)$ and a global phase, such that

$$U = g \cdot (V_1 \otimes V_2) e^{-i\vec{\eta} \cdot \vec{\Sigma}} (V_3 \otimes V_4), g \in \{1, i\} \quad (8)$$

where $\vec{\Sigma} \equiv (XX, YY, ZZ)$ [29]. The set

$$W := \left\{ (x, y, z) \in \mathbb{R}^3 \mid \frac{\pi}{4} \geq x \geq y \geq |z|, z \geq 0 \text{ if } x = \frac{\pi}{4} \right\} \quad (9)$$

is known as the *Weyl chamber* [35], and $\vec{\eta} \in W$ is known as the *Weyl coordinate* of U . We also refer to a gate of the form

$$\text{Can}(a, b, c) := e^{-i\frac{\pi}{2}(aXX + bYY + cZZ)} = \begin{pmatrix} e^{-i\frac{c\pi}{2}} \cos \frac{(a-b)\pi}{2} & 0 & 0 & -ie^{-i\frac{c\pi}{2}} \sin \frac{(a-b)\pi}{2} \\ 0 & e^{i\frac{c\pi}{2}} \cos \frac{(a+b)\pi}{2} & -ie^{i\frac{c\pi}{2}} \sin \frac{(a+b)\pi}{2} & 0 \\ 0 & -ie^{i\frac{c\pi}{2}} \sin \frac{(a+b)\pi}{2} & e^{i\frac{c\pi}{2}} \cos \frac{(a+b)\pi}{2} & 0 \\ -ie^{-i\frac{c\pi}{2}} \sin \frac{(a-b)\pi}{2} & 0 & 0 & e^{-i\frac{c\pi}{2}} \cos \frac{(a-b)\pi}{2} \end{pmatrix} \quad (10)$$

as a *canonical* gate. Two 2Q gates U and V are considered *locally equivalent* if they differ only by 1Q gates, meaning their canonical coefficients can be transformed into one another via the equivalence rules [7]:

1. $(a, b, c) \sim (b, a, c)$ or $(a, b, c) \sim (c, b, a)$, i.e., any permutation of the coefficients;
2. $(a, b, c) \sim (-a, -b, c)$;
3. $(a, b, c) \sim (a - 1, b, c)$;
4. $(1/2, b, c) \sim (1/2, b, -c)$.

Note that we align the conventional that canonical coefficient (a, b, c) differs from Weyl coordinate (x, y, z) by a $\frac{\pi}{2}$ factor. Unless otherwise specified, the canonical coefficients of gates in quantum ISAs and circuits are confined to $\frac{1}{2} \geq a \geq b \geq |c|$. While for the Weyl chamber visualization by means of `weylchamber`, we assume the Weyl coordinates are confined to $\{\frac{\pi}{4} \geq x \geq y \geq z \geq 0\} \cup \{\frac{\pi}{4} \geq \frac{\pi}{2} - x \geq y \geq z \geq 0\}$, as illustrated by Figure 3. Conversion of Weyl coordinates for different conventions is not simple according to the equivalence rules above.

A.2 Quantum ISA and the synthesis capability

A.3 2Q gate mirroring

Mirroring formula:

$$\text{SWAP} \cdot \text{Can}(a, b, c) \sim \left(x + \frac{1}{2}, y + \frac{1}{2}, z + \frac{1}{2} \right) \sim \left(x + \frac{1}{2} - 1, y + \frac{1}{2} - 1, z + \frac{1}{2} - 1 \right) \sim \begin{cases} \left(\frac{1}{2} - z, \frac{1}{2} - y, x - \frac{1}{2} \right), & \text{if } z \geq 0 \\ \left(\frac{1}{2} + z, \frac{1}{2} - y, \frac{1}{2} - x \right), & \text{if } z < 0 \end{cases} \quad (11)$$

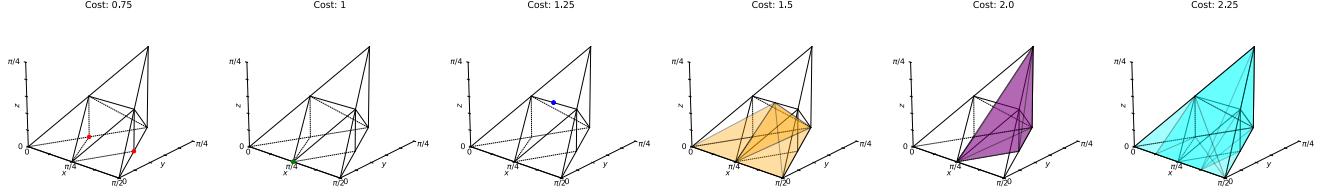


Figure 12. Coverage set for SQiSW_ISA.

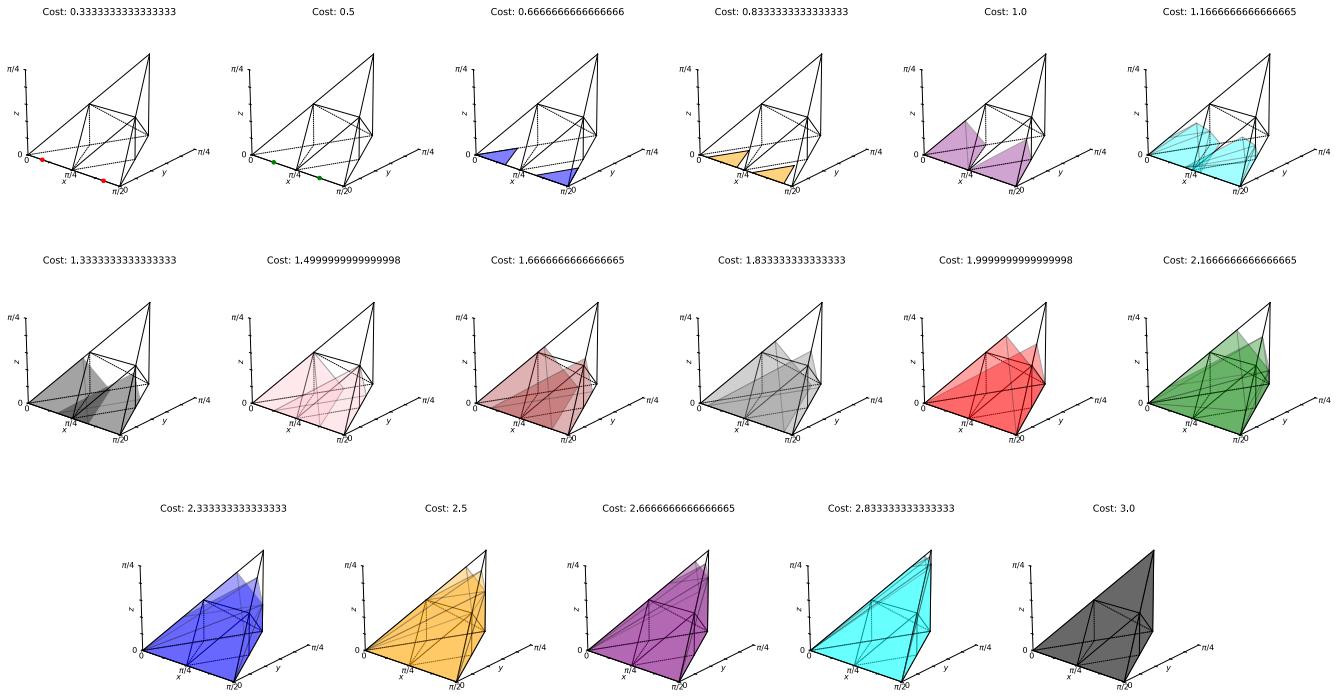


Figure 13. Coverage set for ZZPhase ISA.

A.4 Coverage sets for the selected ISAs used in ...

B Commutative relation of canonical gates

Herein we present detailed proof for Theorem 1. The *if* direction is trivial, and hence we justify the *only if* direction, relying on the following two lemmas.

Lemma 1. Let A, B be two Hermitian matrices with eigenvalues in the range $[-2, 2)$. If $[e^{-i\frac{\pi}{2}A}, e^{-i\frac{\pi}{2}B}] = 0$ then $[A, B] = 0$.

Proof. This follows from the fact that compatible observables (commuting operators) can be simultaneously diagonalized. In this case, the respective unitary matrix $e^{-i\frac{\pi}{2}A}$ commutes with $e^{-i\frac{\pi}{2}B}$. Denote by A_λ the eigenspace corresponding to the eigenvalue λ of $e^{-i\frac{\pi}{2}A}$, i.e. $e^{-i\frac{\pi}{2}A} = \bigoplus_\lambda \lambda A_\lambda$. Then we have

$$\forall \vec{v} \in A_\lambda, e^{-i\frac{\pi}{2}B} e^{-i\frac{\pi}{2}A} \vec{v} = e^{-i\frac{\pi}{2}B} \lambda \vec{v} = \lambda e^{-i\frac{\pi}{2}B} \vec{v} = e^{-i\frac{\pi}{2}A} e^{-i\frac{\pi}{2}B} \vec{v}, \quad (12)$$

and thus $e^{-i\frac{\pi}{2}B} \vec{v} \in A_\lambda$. Thus A_λ is $e^{-i\frac{\pi}{2}B}$ -invariant and the restriction $e^{-i\frac{\pi}{2}B}|_{A_\lambda}$ of $e^{-i\frac{\pi}{2}B}$ to A_λ is still unitary since it preserves inner products. Hence it is diagonalizable and we can find an orthonormal basis $w_{\lambda_1}, w_{\lambda_2}, \dots, w_{\lambda_k}$ consisting of eigenvectors of $e^{-i\frac{\pi}{2}B}|_{A_\lambda}$. Note that these are also eigenvectors of $e^{-i\frac{\pi}{2}A}$ (with eigenvalue λ). Following the same token as above, for each eigenspace E_{λ_i} of $e^{-i\frac{\pi}{2}A}$, we can construct an orthonormal basis β_i for it consisting of eigenvectors of $e^{-i\frac{\pi}{2}B}$. Finally since the eigenspaces of different eigenvalues of $e^{-i\frac{\pi}{2}A}$ are orthogonal to each other, $\beta = \bigcup_i \beta_i$ forms an orthonormal basis of the entire Hilbert space \mathcal{H}_n consisting of the coeigenvectors of both $e^{-i\frac{\pi}{2}A}$ and $e^{-i\frac{\pi}{2}B}$.

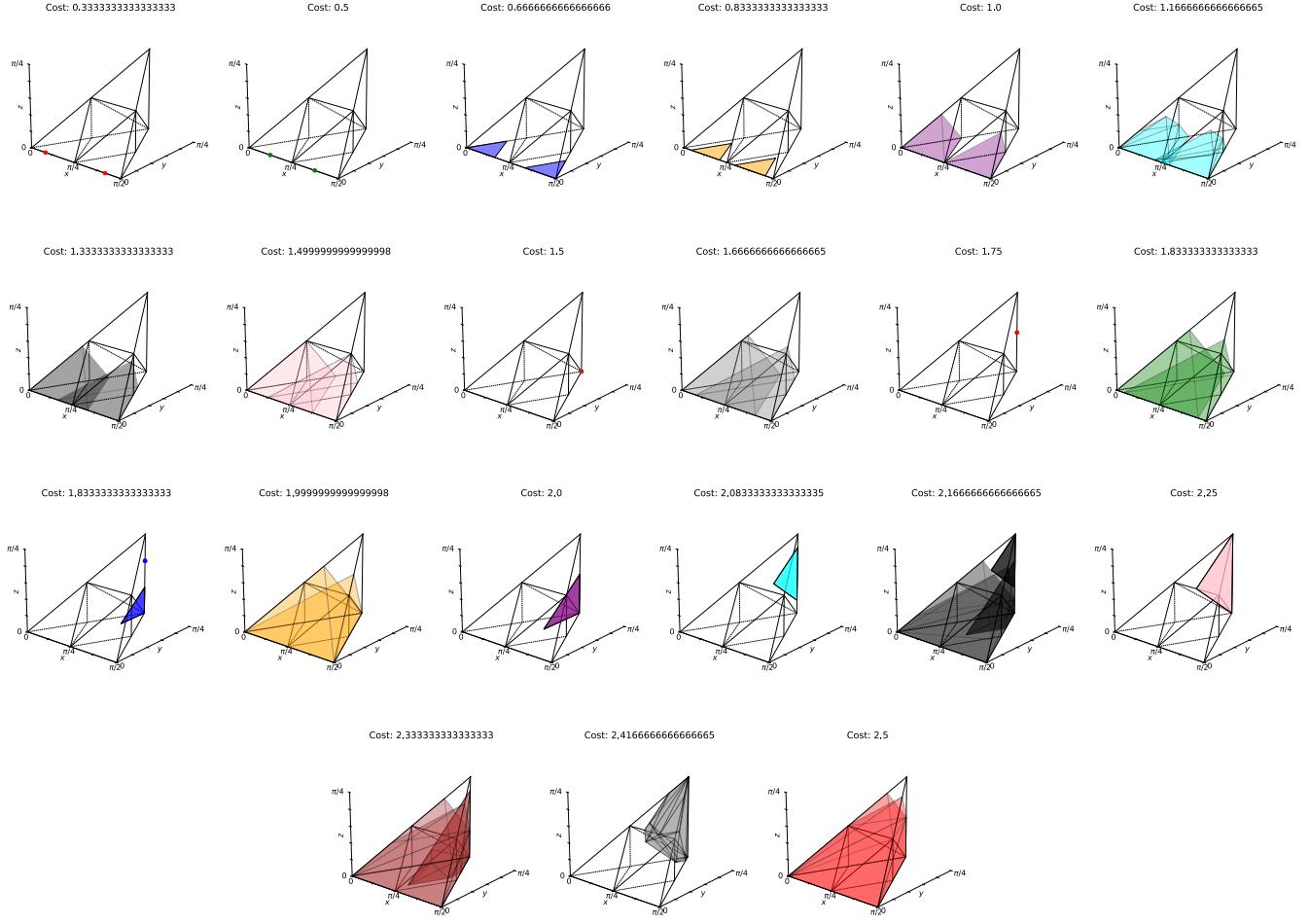


Figure 14. Coverage set for ZZPhase_ISA.

Now let U be a unitary matrix with the vectors in β being its columns, then

$$\begin{aligned} U^\dagger e^{-i\frac{\pi}{2}A}U &= D_A \\ U^\dagger e^{-i\frac{\pi}{2}B}U &= D_B \end{aligned} \tag{13}$$

In general, an eigenvector of $e^{-i\frac{\pi}{2}A}$ need *not* be that of A . However, since A has its eigenvalues in the range $[-2, 2]$, the map

$$f : [-2, 2] \rightarrow U(1), a \rightarrow e^{-i\frac{\pi}{2}a} \tag{14}$$

is injective. Consequently different eigenvalues of A correspond to different eigenvalues of $e^{-i\frac{\pi}{2}A}$, and hence the eigenspaces of $e^{-i\frac{\pi}{2}A}$ and A coincide. Therefore, we have that

$$\begin{aligned} U^\dagger A U &= \Sigma_A \\ U^\dagger B U &= \Sigma_B \end{aligned} \tag{15}$$

and since $[\Sigma_A, \Sigma_B] = 0$ as they are diagonal, $[A, B] = 0$. We obtain the desired result. \square

Lemma 2. Let $P_1 = (a_1 X_1 X_2 + b_1 Y_1 Y_2 + c_1 Z_1 Z_2) I_3$, $P_2 = I_1 (a_2 X_2 X_3 + b_2 Y_2 Y_3 + c_2 Z_2 Z_3)$ with $|c_1| \leq b_1 \leq a_1 \leq \frac{1}{2}$, $|c_2| \leq b_2 \leq a_2 \leq \frac{1}{2}$. If $[P_1, P_2] = 0$ and $P_1, P_2 \neq 0$, then $b_1 = b_2 = c_1 = c_2 = 0$.

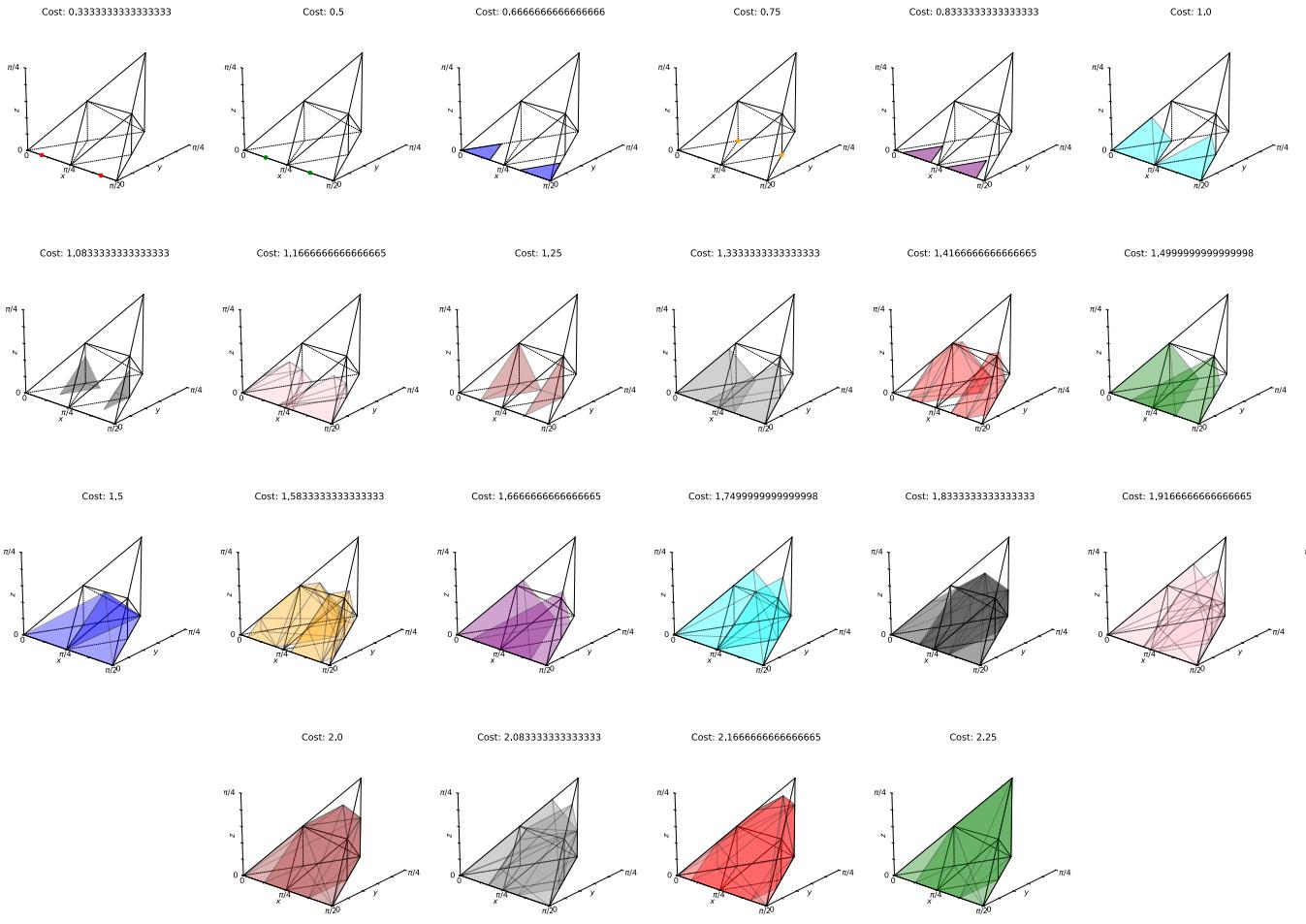


Figure 15. Coverage set for Het ISA.

Proof. Consider the product P_1P_2 . We assume for the sake of contradiction that $b_1 \neq 0$. Using $[X, Y] = 2iZ$, $[Y, Z] = 2iX$, $[Z, X] = 2iY$, we expand

$$[P_1, P_2] = 2i(a_1b_2X_1Z_2Y_3 - b_1a_2Y_1Z_2X_3 + b_1c_2Y_1X_2Z_3) - 2i(a_1c_2X_1Y_2Z_3 + c_1a_2Z_1Y_2X_3 + c_1b_2Z_1X_2Y_3).$$

Since the each Pauli string is linearly independent in the 8×8 operator basis, e.g. term $Y_1Z_2X_3$ cannot be canceled out by any other terms, contradictory to the fact that $[P_1, P_2] = 0$. Hence, vanishing of $[P_1, P_2]$ requires

$$a_1b_2 = a_1c_2 = b_1c_2 = b_1a_2 = c_1a_2 = c_1b_2 = 0.$$

Since $P_1, P_2 \neq 0$, at least a_1, a_2 is nonzero, leading to $b_1 = b_2 = c_1 = c_2 = 0$. \square

Using Lemma 1 and Lemma 2 above, it is straightforward to prove Theorem 1. We see that $\|P_1\| \leq \|a_1X_1X_2I_3\| + \|b_1Y_1Y_2I_3\| + \|c_1Z_1Z_2I_3\| \leq |a_1| + |b_1| + |c_1| \leq \frac{3}{2}$, where $\|\cdot\|$ is the operator norm. Hence, eigenvalues of P_1 are in range of $[-2, 2]$. Same as the eigenvalues of P_2 . Now if $[e^{-i\frac{\pi}{2}P_1}, e^{-i\frac{\pi}{2}P_2}] = 0$, then we have that $[P_1, P_2] = 0$ according to Lemma 1, and thus $b_1 = b_2 = c_1 = c_2 = 0$ according to Lemma 2, which proves the *only if* direction.