

HW #20

1) randomsolution(items, W) - W is max weight of knapsack
 for all item in items:
 int random = random(0 or 1)
 if random == 1:
 if (totalweight + item.weight <= W)
 totalweight += item.weight > Add item to knapsack
 Knapsack[itemindex] = 1
 else: Don't add item → Knapsack[itemindex] = 0
 return knapsack

localsearch(items, W)

RandomKnapsack = randomsolution(items, weight)

currentmax = sum of values of items stored in knapsack

itemindex = 0

while true:

get neighbor { while true:
 flip bit of item in randomknapsack (if = 1, remove, if 0, add)
 if (randomknapsack.totalweight > W)
 revert bit change back to before
 if itemindex == items.size then itemindex = 0, else itemindex++
 else: break

if (randomknapsack.totalvalue > currentmax)

currentmax = randomknapsack.totalvalue

if itemindex == items.size then itemindex = 0, else, itemindex++

else: revert knapsack to prev state then break

return randomknapsack

My algorithm works using a bit array of 0s & 1s representing each item. The random solution loops through the size of items, & uses a random number generator to decide add or not for the initial solution. The weight of the knapsack is checked to ensure a valid solution is made, & returned to the main local search function.

Once the random solution is made, the value of the items with 1 in the bit array is saved as the current max. The main while true loop starts, only stopping when the neighbor solution is $<$ the current max, which indicates a local max has been found for total value, since we are trying to max that.

As for finding neighbors, I chose the processes of flipping 1 bit each iteration of the loop, using an item index iterator to track which bit to change. (Ex. if random solution = 0101 for 4 item knapsack, 1st iteration neighbor = 1101, 2nd iteration = 1001, 3rd 1011, etc) It is not a perfect method to traverse all neighbors, but should be enough to find a local max.

- 2) The whole concept of a local search is greedy in the sense that it takes the first max/min it sees. Graphically speaking, an absolute max/min comes from the set of local max/mins, but a local max/min is not guaranteed to be the absolute max/min. Without trying all possible solutions, there is no way to know if it is an optimal solution. Therefore, there's a chance it is the optimal answer, but we can't guarantee it without testing the rest of the solutions, which goes against the greedy design of this algorithm.