```
f = O(g) \rightarrow f \leq g

f = \Omega(g) / g = O(f) \neq (f \geq g)
f=0(g)/+=(x(g) ( += 2(g) = (f=g)
                         CS 312 - HW1
              1) a) f(n)=n-100 a(n) g(n)·n-200 a(n) -
                    \lim_{n\to\infty}\frac{n-100}{n-200}=\frac{\infty}{60}\xrightarrow{R}\lim_{n\to\infty}\frac{1}{1}=1 \in \mathbb{R}^+\to \Big\{f(n)=\Theta(g(n))\Big\}
               b) f(n) = n^{1/2} g(n) = n^{1/3}
                     \lim_{n\to\infty} \frac{n^{1/2}}{n^{2/5}} = \lim_{n\to\infty} \frac{-1/6}{n^{1/6}} = 0 \Rightarrow f(n) = O(g(n))
               c) f(n)=100n + logn 19(n) n + (logn)
                     1/m 100x+logn = hm 100+n = lm 100+n = lm 100 = 100 ER+

NOO 1+(logn) = 100 1+2logn(+) 1100 1+2logn = 100 ER+

11(0)=6(a)
               d) flatenlogn g(n)=10n logion
                    In alogn = 1 lm logn = 1 lm = 15(1)= 15 ER= Hulodge
               e) f(n)=log 2n g(n)-log3n
                   10 10920 = lin (m) = lin 1 = 1 ER -> P(n) = O(g(n))
              f) f(n)=10 logn g(n)=nlog2(n)
                   1 m 10 log n = 1 m 0(log n) + 10(f) = 1 m 16 m = 0 = 0 = 0
                   (log (2)) 210g(x). 1 (1)
                                                                 (f(n) = O(g(n))
```

```
HW1 (Cont)
    3 a) fab. L(n), Exponential Fqb
if n-<2 then
          return 1
          return fabl(n-1) + fabl(n-2) * fabl(n-3)
   For in levels of function calls, we get 3"
          leaves worth of function calls in our tree, so
          we can say fabl(n) is an exponential tanction.
          fab2(n) Linear Fab
           if n <2 then
           neturn 1
           creeke any foo[0 to n]
           f[0] f[1] f[2] = 1
             fab[i] = fab[n-1] + (fab[n-2] * fab[n-3]) | performs 1 add & i multiply
           for 1=3 to 1
           return fab [n]
           fab2(n) will perform (n-3) adds & (n-3) multiplications
          because the for loop runs (n-3) times. This
          -can also be written as 2(n-3) operations total
          during runtime.
```