

## HW12 - S.26 + S.28

S.26) For set of variables  $x_1, \dots, x_n$   $n$  & input  $m$  of constraints of form " $x_i = x_j$ " or " $x_i \neq x_j$ ", give efficient algorithm to decide if all constraints can be satisfied.

check constraints ( $m, n$ )  $m$  = variables,  $n$  = constraints

make set for all variables to  $x_n$  in  $n \rightarrow \{x_1\}, \{x_2\}, \dots, \{x_n\}$  disjoint set

for all constraints,  $c$  in  $m$  \* Iterate Through each constraint, assuming pre work

left = left side variable of  $c$  to make sure all equalities core first has been done \*

right = right side variable of

if  $c$  is equality

Union(left, right) \* Add two variables into same disjoint set to create equivalence classes

if  $c$  is disequality

if find(left, right) == True \* If two variables that are supposed to not be equal

are in the same disjoint set, the constraint is violated: return false

return True

S.28) Alice's party -  $n$  people to choose from, list of pairs who know each other. Invitee must 1) have at least  $S$  people they know & 2)  $S$  people they don't know  
choose invitees ( $n$ , pairs)

inviteeset = makeset of tuple objects of  $\{person, count=0\}$  for all people in  $n$

while person removed is True Person removed = False

for all pairs in pairs

increment count for person1 & person2 of pair

for all invitees in inviteeset:

if count  $\leq S$  \* does not know

remove invitee from inviteeset

person removed = True remove pair containing person from pairs

else if count  $> n - S$  \* knows too many people

remove invitee from inviteeset

person removed = True

remove pairs containing person from pairs

return inviteeset