



# Modeling and Designing Database

# UML Data Modeling

How to represent data for application

1. Relational model (tables)
2. XML
3. Graphes
  - (a) Entity-Relationship Model (E/R)
  - (b) Unified Modeling Language (UML)

Both can be translated to relations  
automatically (or semi-automatically)

# Unified Modeling Language (UML)

1. Classes
2. Associations
3. Association Classes
4. Subclasses
5. Composition & Aggregation

# Classes

Name, attributes, methods

For data modeling: add primary key, delete methods

# Unified Modeling Language (UML)

1. Classes
2. Associations
3. Association Classes
4. Subclasses
5. Composition & Aggregation

# Associations

Relationships between objects of two classes

# Multiplicity of Associations

Each object of class  $C_1$  is related to at least  $m$  and at most  $n$  objects of class  $C_2$

special

$0 \dots *$

$1 \dots 1$  (default)

Complete



# Unified Modeling Language (UML)

1. Classes
2. Associations
3. Association Classes
4. Subclasses
5. Composition & Aggregation

## UML Data Modeling: Association Classes

Relationships between objects of two classes, with attributes on relationships

# Association Classes: Self-Associations

## UML Data Modeling: Subclasses

Relationships between objects of two classes, with attributes on relationships

# UML Data Modeling: Subclasses

Superclass = Generalization

Subclass = Specialization

Incomplete (Partial) vs. Complete

Disjoint (Exclusive) vs. Overlapping

## UML Data Modeling: Composition & Aggregation

Objects of one class belong to objects of another class

Company & Movie & Reviewer

# Designing Database

The same set of information can be captured by different schemas.



The same set of information can be captured into different schemas.

but some schemas are better than others

## Database for students applying to French Universities

123 Marc lives in Nancy applies to Paris-Sorbonne, U-Lille and Lyon-I  
has bac 15/20

sID	Name	Adress	Bac	University
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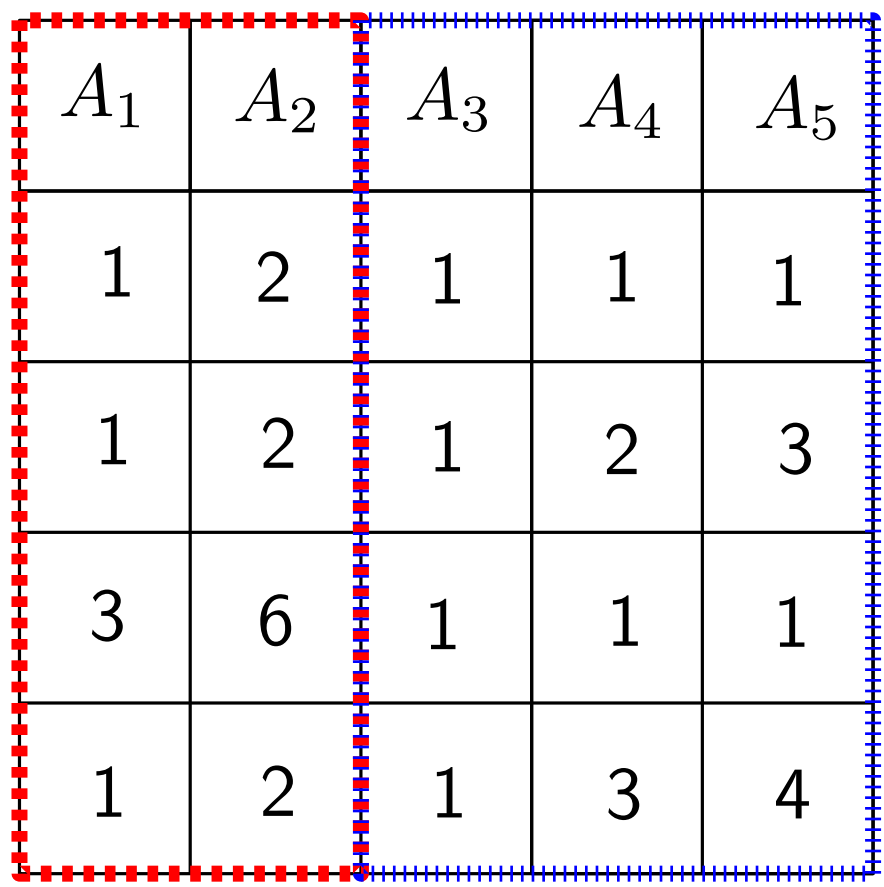
Problem on this schema (anomalies):

1. Redundancy
2. Update anomaly
3. Deletion anomaly

How to solve this problem?

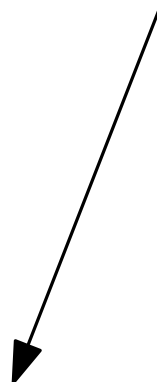
$A_1$	$A_2$	$A_3$	$A_4$	$A_5$
1	2	1	1	1
1	2	1	2	3
3	6	1	1	1
1	2	1	3	4

How to solve this problem?

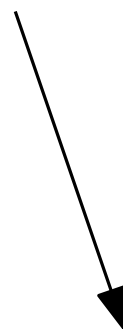


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3	6	1	1	1
1	2	1	3	4



$A_1$	$A_2$
1	2
3	6



$A_3$	$A_4$	$A_5$
1	1	1
1	2	3
1	3	4

## Functional Dependency

$A \rightarrow B$  we read it:  $B$  functionaly depends on  $A$

A	B	C

## Exercise

Consider relation  $R(A, B, C, D, E)$  with functional dependencies:  
 $A, B \rightarrow C$  and  $C \rightarrow D$  and  $B, D \rightarrow E$ ,

Which of the following sets of attributes does not functionally determine E?

1. A,B
2. A,B,C
3. B,C
4. C



## Closure of Attributes

Given relation, FDs, set of attributes  $\overline{A}$

Find all  $B$  such that  $\overline{A} \rightarrow B$

Closure of  $\overline{A}$  is  $\overline{A}^+$

## Boyce-Codd normal form

A relation (table) is called Boyce-Codd normal form if

for any  $A_1, A_2, \dots, A_n \rightarrow B_1, B_2, \dots, B_m$

$A_1, A_2, \dots, A_n$  is a key

That is:  $A_1, A_2, \dots, A_n$  determines the whole table

Student(S-id, Name, address, U-id, U-Name, U-city, Bac, Rating)

Bac  $\rightarrow$  Rating

S-id  $\rightarrow$  Name, address, Bac

U-id  $\rightarrow$  U-Name, U-city

U-Name, U-city  $\rightarrow$  U-id

Student(S-id, Name, address, U-id, U-Name, U-city, Bac, Rating)

Bac  $\rightarrow$  Rating

S-id  $\rightarrow$  Name, address, Bac

U-id  $\rightarrow$  U-Name, U-city

U-Name, U-city  $\rightarrow$  U-id

What is are the keys on this relation?

Student(S-id, Name, address, U-id, U-Name, U-city, Bac, Rating)

Bac  $\rightarrow$  Rating

S-id  $\rightarrow$  Name, address, Bac

U-id  $\rightarrow$  U-Name, U-city

U-Name, U-city  $\rightarrow$  U-id

S-id, U-id  $\rightarrow$  All others



BCNF Violation

What are the keys on this relation?

Student(S-id, Name, address, U-id, U-Name, U-city, Bac, Rating)

Bac  $\rightarrow$  Rating

S-id  $\rightarrow$  Name, address, Bac

U-id  $\rightarrow$  U-Name, U-city

U-Name, U-city  $\rightarrow$  U-id

Is this relation in Boyce-Codd normal form ?

# Exercises

Consider the relation  $R(A, B, C, D, E)$  and suppose we have the functional dependencies  $A, B \rightarrow C$  and  $A, E \rightarrow D$  and  $D \rightarrow B$ . Determine of the keys for  $R$ ?

## BCNF decomposition algorithm

Input: relation  $R$  + FDs for  $R$

Output: decomposition of  $R$  into BCNF relations with lossless join

1. Compute keys for  $R$
2. Repeat until all relations are in BCNF:
  - (a) Pick any  $R$  with  $A \rightarrow B$  that violates BCNF
  - (b) Decompose  $R$  into  $R_1(A, B)$  and  $R_2(A, \text{rest})$
  - (c) Compute FDs and keys for  $R_1$  and  $R_2$



Is BCNF always good?

Apply(S-id,U-id, hoppy)

# Is BCNF always good?

Apply(S-id,U-id, hoppy)

1. Functional dependency ?
2. Keys?
3. BCNT
4. Is it a good design ?

## Multivalued dependency

Relation  $R(A, B, C)$

$A \twoheadrightarrow B$  if

for all  $L_1, L_2$  in  $R$  with  
 $R_1[A] = R_2[A]$ , then

there exists  $L_3$  in  $R$  such that

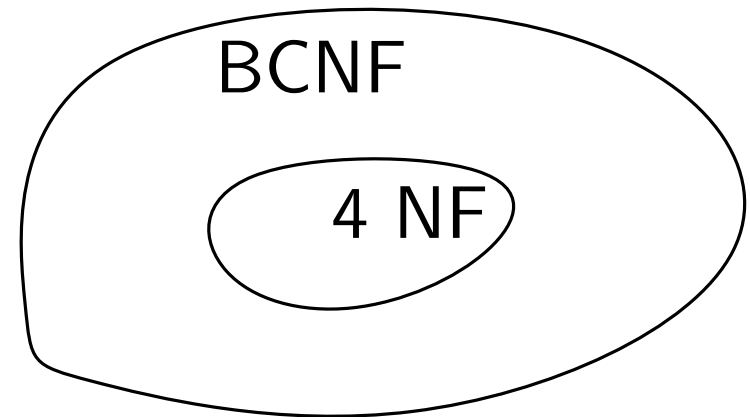
$R_3[A] = R_1[A]$  and

$R_3[B] = R_1[B]$  and

$R_3[C] = R_2[C]$

## 4th Normal form

A relation is in 4th normal form, if for any  $A \twoheadrightarrow B$ , we have that  $A$  is a key



Apply(S-id,U-id, hoppy) + condition

It is BCNF, but not 4 NF

(S-id,U-id)

(S-id, hoppy)

## Bad points of BCNF and 4 NF

1. Over-decomposition
2. Query workload

Consider a relation  $R(A,B,C,D)$ . For which of the following sets of FDs is  $R$  in Boyce-Codd Normal Form

$A \rightarrow B$  and  $B \rightarrow C$  and  $C \rightarrow D$  and  $D \rightarrow A$

$C \rightarrow B$  and  $D \rightarrow A$  and  $C \rightarrow D$  and  $A \rightarrow C$

