### Boyce-Codd normal form

A relation (table) is called Boyce-Codd normal fom if for any  $A_1,A_2,\ldots A_n\to B_1,B_2,\ldots B_m$   $A_1,A_2,\ldots A_n$  is a key

That is:  $A_1, A_2, \dots A_n$  determines the whole tuble

Student(S-id, Name, address, U-id, U-Name, U-city, Bac, Rating)

 $\mathsf{Bac} \to \mathsf{Rating}$ 

S-id  $\rightarrow$  Name, adress, Bac

U-id  $\rightarrow U$ -Name, U-city

U-Name, U-city  $\rightarrow$  U-id

### Multivalued dependency

Relation 
$$R(A, B, C)$$
  
 $A \rightarrow B$  if

for all 
$$L_1, L_2$$
 in  $R$  with  $R_1[A] = R_2[A]$ , then

there exists  $L_3$  in R such that

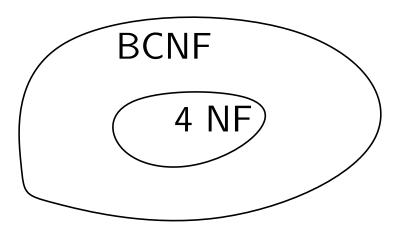
$$R_3[A] = R_1[A]$$
 and

$$R_3[B] = R_1[B]$$
 and

$$R_3[C] = R_2[C]$$

#### 4th Normal form

A realation is in 4th normal form, if for any  $A \twoheadrightarrow B$ , we have that A is a key



### Bad points of BCNF and 4 NF

- 1. Over-decomposition
- 2. Query workload

Consider a relation R(A,B,C,D). For which of the following sets of FDs is R in Boyce-Codd Normal Form

$$A o B$$
 and  $B o C$  and  $C o D$  and  $D o A$   $C o B$  and  $D o A$  and  $C o D$  and  $A o C$ 

# Example

Can apply to each college once for one major Universities have non-overlapping application dates

Apply(sID, U-Name, date, discipline)

**FDs** 

Keys

**BCNF** 

Good design?

# Example 2

every student might have several Bac or English scores

```
Student(sID, s-name, Bac score, English score) 
FDs and Keys sID \rightarrow sname 
MVDs sID , sname \rightarrow Bac score 
4NF? No
```

Good design?

## Example 3

BCNF?

A student might apply to several universities. Result is determined by Bac score

```
Student(sID, U-name, Bac score,Result)  FDs \quad sID \rightarrow Bac \quad Bac \rightarrow Result \quad sID \rightarrow Result  Keys
```

Good design?