

Assignment 5

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Introduction

In this assignment we solve five exercises on context-free grammars and parsers:

1. LL(1) analysis for seven grammars.
2. Proving ambiguity for three small grammars.
3. LR(0) automaton and parse table.
4. LR(1) automaton and parse table.
5. Fixing and extending a partial C++ EBNF.

Problem 1: LL(1) Analysis

Compute FIRST and FOLLOW sets, then check the LL(1) conditions.

Grammar 1

$$S \rightarrow a A B C \mid \epsilon, A \rightarrow a \mid b b D, B \rightarrow a \mid \epsilon, C \rightarrow b \mid \epsilon, D \rightarrow c \mid \epsilon.$$

FIRST:

$$\text{FIRST}(S) = \{a, \epsilon\}, \text{FIRST}(A) = \{a, b\}, \text{FIRST}(B) = \{a, \epsilon\}, \text{FIRST}(C) = \{b, \epsilon\}, \text{FIRST}(D) = \{c, \epsilon\}.$$

FOLLOW:

$$\text{FOLLOW}(S) = \{\$, \}, \text{FOLLOW}(A) = \{a, b, \$\}, \text{FOLLOW}(B) = \{b, \$\}, \text{FOLLOW}(C) = \{\$, \}, \text{FOLLOW}(D) = \{c, \$\}.$$

All intersections are empty and the ϵ checks pass. *Grammar 1 is LL(1).*

Grammar 2

$$A \rightarrow B C c \mid e D B, B \rightarrow \epsilon \mid b C D, C \rightarrow D a B \mid c a, D \rightarrow \epsilon \mid d D.$$

FIRST: $\text{FIRST}(D) = \{d, \epsilon\}$, but $d \in \text{FOLLOW}(D)$ conflict. *Grammar 2 is NOT LL(1).*

Grammar 3

Clarified as:

$$S \rightarrow ' (' X ') ' \mid ' [' E '] ' \mid F, X \rightarrow E ') ' \mid F '] ', E \rightarrow A, \quad F \rightarrow A, \quad A \rightarrow \epsilon.$$

Compute FIRST and FOLLOW (all disjoint, no conflicts). *Grammar 3 is LL(1).*

Grammar 4

$$S \rightarrow X b \mid Y d, X \rightarrow a X \mid \epsilon, \quad Y \rightarrow c Y \mid \epsilon.$$

$\text{FIRST}(Xb) = \{a, b\}$ vs $\text{FIRST}(Yd) = \{c, d\}$ disjoint, ϵ -cases OK. *Grammar 4 is LL(1).*

Grammar 5

$$S \rightarrow M N O P Q, \quad M, N, O, P, Q \rightarrow t \mid \epsilon$$

Each $t \notin \text{FOLLOW}$ LL(1). *Grammar 5 is LL(1).*

Grammar 6

Eliminate left recursion:

$$S \rightarrow A, \quad A \rightarrow a B A', \quad A' \rightarrow d A' \mid \epsilon, \quad B \rightarrow b.$$

No conflicts LL(1). *Grammar 6 is LL(1).*

Grammar 7

$$S \rightarrow A a A b \mid B b B a, \quad A \rightarrow \epsilon, \quad B \rightarrow \epsilon.$$

FIRST-sets $\{a\}$ vs $\{b\}$ LL(1). *Grammar 7 is LL(1).*

Problem 2: Ambiguity Proofs

Each grammar admits a string with two parse-trees:

- Grammar 1: $S \rightarrow a \mid aAb \mid abSb$, $A \rightarrow aAAb \mid bS$. String **abab**:

$$S \Rightarrow abSb \Rightarrow abab, \quad S \Rightarrow aAb \Rightarrow a(bS)b \Rightarrow abab.$$

- Grammar 2: $S \rightarrow AB \mid aaaB$, $A \rightarrow a \mid Aa$, $B \rightarrow b$. String **aaab**:

$$S \Rightarrow AB \Rightarrow AaB \Rightarrow aab, \quad S \Rightarrow aaaB \Rightarrow aaab.$$

- Grammar 3: $S \rightarrow xyXxX \mid xyXyy \mid yy$, $X \rightarrow xxX \mid yy \mid xx$. String **xyxxxxxyy** by two different expansions of X .

All three are ambiguous.

Problem 3: LR(0) Automaton

Augmented Grammar

0. $S' \rightarrow S$
1. $S \rightarrow A A$
2. $A \rightarrow a A$
3. $A \rightarrow b$

SLR Parse Table

State	a	b	\$	A	S
0	s3	s4		2	1
1			acc		
2	s3	s4		5	
3	s3	s4		6	
4	r3	r3	r3		
5	r1	r1	r1		
6	r2	r2	r2		

where r1: $S \rightarrow AA$, r2: $A \rightarrow aA$, r3: $A \rightarrow b$.

Problem 4: LR(1) Automaton

Augmented Grammar

0. $S' \rightarrow S$
1. $S \rightarrow a A d$
2. $S \rightarrow b B d$
3. $S \rightarrow a B e$
4. $S \rightarrow b A e$
5. $A \rightarrow g$
6. $B \rightarrow g$

LR(1) ACTION / GOTO

State	a	b	g	d	e	\$	A	B	S
0	s1	s2							13
1			s11				3	4	
2			s12				6	5	
3				s7					
4					s9				
5				s8					
6					s10				
7						r1			
8						r2			
9						r3			
10						r4			
11				r5	r6				
12				r6	r5				
13						acc			

r1: $S \rightarrow aAd$, r2: $S \rightarrow bBd$, r3: $S \rightarrow aBe$, r4: $S \rightarrow bAe$, r5: $A \rightarrow g$, r6: $B \rightarrow g$.

Problem 5: C++ Grammar Corrections

Below I show what I *fixed* (in red) and what I *added* (in blue).

```

<const-def>      ::= const <type> <var-name> = <expression> ;

<switch-stmt>    ::= switch(<expression>){<case-list><default-case>?}
<case-list>      ::= <case> | <case-list><case>
<case>           ::= case <literal> : <statement>*
```

```

<default-case> ::= default : <statement>*

<class-def>    ::= class <var-name> { <member-list> } ;
<member-list> ::= <member> | <member-list><member>
<member>      ::= <var-def> | <fun-def>

<try-catch>    ::= try{<statement>*}catch(<type><var-name>){<statement>*}

```

Precedence-enforced expressions

```

<logical-or-expr>    ::= <logical-and-expr> ( "||" <logical-and-expr> )*
<logical-and-expr>  ::= <equality-expr> ( "&&" <equality-expr> )*
<equality-expr>     ::= <relational-expr> (("=="|"!=") <relational-expr>) *
<relational-expr>   ::= <additive-expr> (("("<additive-expr>)">|"<additive-expr>") *
<additive-expr>     ::= <multiplicative-expr> (("+"|"-" <multiplicative-expr>)*
<multiplicative-expr> ::= <unary-expr>      (("*"|"/"|"%" <unary-expr>)*
<unary-expr>        ::= ("-"|"!") <unary-expr> | "(" <expression> ")" | <literal> |

```