Scheme

Guerrilla Section: November 1, 2019 Solutions

1 Scheme

```
What would Scheme do?
scm> (and 0 2 200)
200
scm> (or True (/ 1 0))
True
scm> (and False (/ 1 0))
False
scm> (not 3)
False
What would Scheme display?
scm> (define a (+ 1 2))
scm> a
3
scm> (define b (+ (* 3 3) (* 4 4)))
b
scm> (+ a b)
```

```
2 Scheme
    scm> (= (modulo 10 3) (quotient 5 3))
    #t
    scm> (even? (+ (- (* 5 4) 3) 2))
    #f
    scm> (if (and #t (/ 1 0)) 1 (/ 1 0))
    Error
    scm> (if (> (+ 2 3) 5) (+ 1 2 3 4) (+ 3 4 (* 3 2)))
    13
    scm> ((if (< 9 3) + -) 4 100)
    -96
    scm> (if 0 #t #f)
    #t
    Write two Scheme expressions that are equivalent to the following Python statement
    - one defining a function directly, and the other creating an anonymous lambda that
    is then bound to the name cat:
    cat = lambda meow, purr: meow + purr
    (define cat (lambda (meow purr) (+ meow purr)))
    (define (cat meow purr) (+ meow purr))
1.4 Spot the bug(s). Test out the code and your fixes in the scheme interpreter!
    (https://scheme.cs61a.org/)
    (define (sum-every-other lst)
     (cond ((null? lst) lst)
           (else (+ (cdr lst)
                     (sum-every-other (caar lst)) )))
    1. Missing a paren at the end.
    2. The base case should return 0, not '().
    3. (cdr lst) is a list, so it doesn't make sense to add it to something. Instead, use
```

(car lst), which will give us a number.

4. Using the caar (car of the car) is incorrect because the car is a number and it doesn't make sense to get the car of a number. Instead, we should use cddr (the cdr of the cdr) to skip forward two elements. However, the cdr could be '(), so we need to add a case to our cond to take care of this.

The corrected function:

1.5 Define **sixty-ones**, a funcion that takes in a list and returns the number of times that 1 follows 6 in the list.

1.6 Define **no-elevens**, a function that takes in a number n, and returns a list of all distinct length-n lists of 1s and 6s that do not contain two consecutive 1s.

1.7 Define remember, a function that takes in another zero-argument function f, and returns another function g. When called for the first time, g will call f and pass on its return value. When called subsequent times, g will remember its previous return value and return it directly, without calling f again.

```
(Hint: look up set! in the Scheme spec!)
(define (remember f)
  (define remembered? #f)
  (define remembered nil)
  (lambda ()
    (if remembered?
        remembered
        (begin (set! remembered (f))
               (set! remembered? #t)
               remembered)))
)
scm> (define (f) (print "hello!") 5)
scm> (define g (remember f))
scm>(f)
hello!
5
scm>(g)
hello!
5
scm>(g)
```

Check your understanding

How are call expressions (like (+ 1 2 3)) evaluated? What about special forms, like (or #f #t (/ 1 0))

To evaluate call expressions, Scheme first evaluates the operator, and then evaluates all of the operands from left to right. It then *applies* the operator to the operands (i.e. calls the procedure with the evaluate operands), just like how Python evaluates function calls. In contrast, the first subexpression in a special form is *not* evaluated, but rather detected and treated specially by the interpreter. The remaining subexpressions may or may not be evaluated, depending on the behavior of the special form. For instance, or will short-circuit when it detects a non-false value, so the above example will not error, since or will never reach the divide-by-zero.

• What is the purpose of the quote special form?

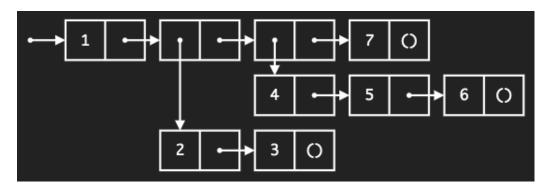
The quote special form is meant to *postpone* the evaluation of an expression. For instance, if we write (1 2 3), Scheme will typically treat it as a call expression, treating 1 as a procedure (which it is not!). Instead, writing

(quote (1 2 3)), or the equivalent shorthand '(1 2 3), will cause the overall expression to evaluate to the second subexpression of the quote special form, allowing us to obtain (1 2 3) after evaluation, as desired.

2 Scheme Lists

2.1 Draw out a box-and-pointer diagram for the following list:

```
scm> (define nested-lst (list 1 (cons 2 (cons 3 'nil)) '(4 5 6) 7)) nested-lst
```



Then, write out what Scheme would display for the following expressions:

```
scm> (cdr nested-lst)
```

 $(1 \ 4 \ 5 \ 6)$

```
((2 3) (4 5 6) 7)
scm> (cdr (cdr nested-lst)))
(3)
scm> (cons (car nested-list) (car (cdr (cdr nested-list))))
```

Extra

2.2 Notice that the builtin append takes in, not a *list* of lists, but an *arbitrary* number of lists as arguments, which it then concatenates together. Implement better-append, which behaves in such a manner, allowing the caller to pass in an arbitrary number of arguments. You may use concat from the previous question.

(Hint: look up "variadic functions" in the Scheme spec!)

```
(define (better-append . args)
  (concat args))

scm> (better-append '(1 2 3))
(1 2 3)
scm> (better-append '(1 2 3) '(2 3 4))
(1 2 3 2 3 4)
scm> (better-append '(1 2 3) '(2 3 4) '(3 4 5))
(1 2 3 2 3 4 3 4 5)
```

Check your understanding

• How can you get the third element of a Scheme list? Draw out a box-and-pointer diagram if you aren't sure.

To get the third element of a Scheme list, we need to get the car of the cdr of the cdr of the list - in other words, the third element of 1st is (car (cdr (cdr 1st))).

• What is the difference between eq? and equal? in the context of Scheme lists? Construct two lists 1st1 and 1st2 such that (equal? 1st1 1st2) is #t but (eq? 1st1 1st2) is #f.

equal? tests *equality*, and behaves like == in Python - in other words, it returns true if all the corresponding elements of two lists are themselves equal. eq?, in contrast, tests *identity*, and returns true only if its two arguments are in fact the same *object*. Thus, one possibility is simply (define lst1 (list 1)) and (define lst2 (list 1)).