

# Energy aware scheduling

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## 1 Problem formulation

We consider a problem of applications scheduling in a center with limited power and resources.

This data center runs two different applications : active applications (also called web applications) and batch applications.

### 1.1 Active applications

We consider multiple active applications  $\{A_i\}, i \in 1..n$ , which run continuously over a given amount of time intervals 1..24.

Each application  $A_i$  has, at a given time interval  $j$ , an execution modes ; each application's mode produces a given power consumption and profit when applied over an interval, with

$M_{i,j} \in 1..3$  mode of activity  $i$  at time  $j$ .

$E_{i,j} \in 0..1000$  power consumption of activity  $i$  at time  $j$ .

$P_{i,j} \in 0..1000$  profit for running activity  $i$  at time  $j$ .

Those informations can be stored in tables, eg :

	Mode 1	Mode 2	Mode 3
$A_1$	50	60	70
$A_2$	40	45	55

Table 1: Energy consumption for two active applications

	Mode 1	Mode 2	Mode 3
$A_1$	100	110	120
$A_2$	90	100	110

Table 2: Profit for two active applications

- element  $(M_{i,j}, Row_i^{Energy}, E_{i,j})$  %  $E_i = Row_i^{Energy}[M_i]$
- element  $(M_{i,j}, Row_i^{Profit}, P_{i,j})$  %  $P_i = Row_i^{Profit}[M_i]$

## 1.2 Batch applications

We consider multiple batch jobs  $\{B_i\}, i \in 1..m$ , each with its own parameters :

$Duration_i$  is the number of intervals this job must be run before being finished,

$Deadline_i$  deadline is the interval at which the job must be finished,

$Q_i \in 0..1000$  is the profit earnt only if the whole job is finished before or at deadline,

$Power_i$  is the power consumed by the job when it is run over an interval

The batch jobs do not have execution modes.

	On	Off
$B_1$	60	0
$B_2$	80	0

Table 3: Energy consumption for Batch applications

Each batch job  $i$  can be decomposed in as many subjobs  $i, j$  as its duration. We note  $S_{i,j}$  the interval at which the  $j^{th}$  subjob of jow  $i$  is executed. Since each subjob must be executed over a different time interval,

$$0 < S_{i,1} < S_{i,2} \dots < S_{i,Duration_i}.$$

If the job meets its deadline, then also

$$S_{i,Duration_i} \leq Deadline_i$$

Energy consumption of batch job  $i$  during an interval is  $E_{i,j}$ , where  $i \in [1,m]$  and  $j \in [1....Duration_i]$ .

- element  $(S_{i,Duration_i}, [.....Profit_i, Pen_i, ....], Q_i)$

## 1.3 Energy cap

We also introduce the maximum available energy at  $j$  as  $Capacity_j$ , where  $j \in [1,24]$ . Since, in cumulative constraint modeling, the limit of maximum capacity of available energy can not be changed, we introduce fake jobs in each slot to match with the  $Capacity_j$ . So the maximum capacity is defined as Limit, where  $Limit = \max(Capacity_h), h \in [1,24]$ . For scheduling purpose we slice the total time i.e., 24/48 hours to 24/48 slots meaning 1 hour as slot and schedule each slot in advance with known information. So, Total Profit is  $P = \sum_{i=1}^n P_{i,j} + \sum_{i=1}^m Q_i$

- For each active job  $A_i$  starts at  $j$ , has duration of 1 slot and height of power consumption is  $E_{ij}$ , where,  $\forall j \in [1,24]$
- For each batch jobs  $B_i$  the start time is  $S_{ij}$ , duration is 1 slot and and height of power consumption is  $E_i$ , where,  $\forall i \in [1,m]$  and  $\forall j \in [1,duration_i]$
- For every slot  $j \in [1,24]$  we have a fake job  $F_h$  that starts at fixed  $j$ , duration of 1 slot and height of power consumption of  $Limit - Capacity_j$ , where,  $\forall j \in [1,24]$

## 1.4 Memory use

The data center contains  $l$  servers  $\{S_k\}, k \in 1..l$ . Each server  $S_k$  has a memory capacity  $M_k$  which limits the number of applications this server can execute.

Executing an application  $A_i$  or a job  $B_i$  on a server  $S_k$  consumes a fixed amount of memory on the server over the execution interval. Reciprocally, each application must be run on a server at any time and any job executed during an interval must be run on a server.

We note

$HA_{i,j}$  the host of the application  $A_i$  during interval  $j$

$HB_{i,j}$  the host of the job  $B_i$  during interval  $j$

$Load_{k,j}$  the memory load of the server  $k$  during interval time  $j$

$mem(C)$  the memory use of the job or application  $C$ .

then we know that

$$Load_{k,j} = \sum_{A_i} (mem(A_i) if HA_{i,j} = k) + \sum_{B_i} (mem(B_i) if HB_{i,j} = k) \quad (1)$$

## 2 Solution formulation

A solution to such a problem consists in :

- $\forall$  Active application  $A_i$  and time interval  $j$ , the execution mode  $M_{i,j}$  at which to run the application during the interval
- $\forall$  Batch job  $B_i$ , a series time interval  $S_{i,1}..S_{i,Duration_i}$  at which to execute the job.
- $\forall$  time interval  $j$ ,  $\forall$  server  $S_k$ , the set of Applications and jobs hosted on the server during the time interval.

with respect to the previously specified conditions.

## 3 Objective

Our objective function is to maximize  $P$ ,

$$P = \sum_{i=1}^n P_{i,j} + \sum_{i=1}^m Q_i$$