

Chapter 4

The Duality Theorem

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4.1 Primal-Dual Relationship

Consider the following LOP.

Problem 4.1.1

$$\begin{aligned} \text{Max. } z &= 22x_1 + 31x_2 + 29x_3 \\ \text{s.t. } & \quad x_1 + 4x_2 + 6x_3 \leq 73 \\ & \quad 5x_1 - 2x_2 + 3x_3 \leq 68 \\ & \quad x_1, x_2, x_3 \geq 0 \end{aligned}$$

As you are invited to confirm, the simplex algorithm produces the following final tableau.

Tableau 4.1.2

$$\left[\begin{array}{ccc|ccc|c} 0 & 22 & 27 & 5 & -1 & 0 & 297 \\ 22 & 0 & 24 & 2 & 4 & 0 & 418 \\ \hline 0 & 0 & 727 & 199 & 57 & 22 & 18403 \end{array} \right]$$

Tableau 4.1.2 shows the optimum solution $\mathbf{x}^* = (418, 297, 0 \mid 0, 0)^T / 22$ with corresponding optimum value $z^* = 18403/22$. Likewise, consider the dual LOP below.

Problem 4.1.3

$$\begin{aligned} \text{Min. } w &= 73y_1 + 68y_2 \\ \text{s.t. } & \quad y_1 + 5y_2 \geq 22 \\ & \quad 4y_1 - 2y_2 \geq 31 \\ & \quad 6y_1 + 3y_2 \geq 29 \\ & \quad y_1, y_2 \geq 0 \end{aligned}$$

Problem 4.1.3 has the following final tableau.

Tableau 4.1.4

0	22	-4	1	0	0	57
22	0	-2	-5	0	0	199
0	0	-24	-27	22	0	727
0	0	418	297	0	22	-18403

Tableau 4.1.4 shows the optimum solution $y^* = (199, 57 \mid 0, 0, 727)/22$ with corresponding optimum value $w^* = 18403/22$.

Curiously, this data shows certain repetitions of values. It looks like the values of \mathbf{x}^* show up in the final dual objective row, but switched around a little. Likewise, the values of \mathbf{y}^* appear in the final primal objective row, with a similar swap of some sort. To be more precise, the pattern seems to be that the *problem* values of \mathbf{x}^* are the final coefficients of the dual *slack* variables, while the *slack* values of \mathbf{x}^* are the final coefficients of the dual *problem* variables.

Problem 4.1.5

$$\begin{aligned}
 \text{Max. } z &= -22x_1 - 18x_2 - 27x_3 - 23x_4 + 16x_5 - 12x_6 \\
 \text{s.t.} \quad &4x_1 + x_2 - 3x_3 + 2x_5 + 7x_6 \leq 211 \\
 &6x_1 + 2x_3 + 5x_4 - x_5 + 8x_6 \leq 189 \\
 &-5x_1 + 4x_2 - 2x_3 - x_5 - 7x_6 \leq -106 \\
 &3x_1 + 9x_2 - 2x_4 + x_5 + 4x_6 \leq 175 \\
 &\& \quad x_1, x_2, x_3, x_4, x_5, x_6 \geq 0
 \end{aligned}$$

Workout 4.1.6 Consider Problem 4.1.5.

- Use the Simplex algorithm to solve it.
- Without solving the dual linear problem, use the final primal tableau to find the optimal dual variable values \mathbf{y}^* (including slacks).
- Verify that \mathbf{y}^* is dual feasible and optimal.

It will help to articulate this perceived pattern notationally. We return to the general descriptions of primals and duals below.

Problem 4.1.7

$$\text{Max. } z = \sum_{j=1}^n c_j x_j \quad (1)$$

$$\text{s.t.} \quad \sum_{j=1}^n a_{ij} x_j \leq b_i \quad (1 \leq i \leq m) \quad (2)$$

$$\& \quad x_j \geq 0 \quad (1 \leq j \leq n) \quad (3)$$

Problem 4.1.8

$$\text{Min. } w = \sum_{i=1}^m b_i y_i \quad (4)$$

$$\text{s.t.} \quad \sum_{i=1}^m a_{ij} y_i \geq c_j \quad (1 \leq j \leq n) \quad (5)$$

$$\& \quad y_i \geq 0 \quad (1 \leq i \leq m) \quad (6)$$

In order to describe the pattern we will need to look at the optimal primal objective row. Just as with all optimal values, let's use c_k^* for the final coefficient of x_k , as shown.

$$\left[\begin{array}{cccc|cccc|c} c_1^* & \cdots & c_j^* & \cdots & c_n^* & c_{n+1}^* & \cdots & c_{n+i}^* & \cdots & c_{n+m}^* & 1 & z^* \end{array} \right]$$

Notice that there is a 1 written instead of the more general d for the coefficient of z . Why did we do that?

It seems as though the pattern we have witnessed is given below.

$$\left[\begin{array}{cccc|cccc|c} y_{m+1}^* & \cdots & y_{m+j}^* & \cdots & y_{m+n}^* & y_1^* & \cdots & y_i^* & \cdots & y_m^* & 1 & z^* \end{array} \right]$$

That is,

$$y_i^* = c_{n+i}^* \quad (1 \leq i \leq m) \quad \text{and} \quad y_{m+j}^* = c_j^* \quad (1 \leq j \leq n). \quad (7)$$

Furthermore, we have also noticed time and again that, when the primal has an optimal solution, then so does its dual — in fact, with the same optimal value. Quite possibly, we could take advantage of the detailed pattern above to verify such a phenomenon in general. From (1) the value in question is

$$z^* = \sum_{j=1}^n c_j x_j^*. \quad (8)$$

From the Weak Duality Theorem (Inequality 1.4.7) we know that if the dual problem is feasible then its optimum is at most this value z^* . In fact we get equality.

Theorem 4.1.9 *If a linear problem P has an optimum z^* then its dual linear problem D has an optimum w^* ; moreover, $z^* = w^*$.*

Strong Duality Theorem

Proof. Because we already know that every feasible z and w satisfy $z \leq w$, we only need to find a feasible w for which $w = z$. For this we can turn to the y_i^* s defined in (7), and show that they satisfy inequalities (5) and (6) as well as the equality

$$z^* = \sum_{i=1}^m b_i y_i^*. \quad (9)$$

Workout 4.1.10 *Show that the y_i^* s defined in (7) satisfy (6).*

One of the things we can do is write out the optimal objective row, solving for z . With the substitutions from (7) we have

$$z = z^* - \sum_{j=1}^n y_{m+j}^* x_j - \sum_{i=1}^m y_i^* x_{n+j}. \quad (10)$$

Workout 4.1.11 *Use equation (1) and the definition of the slack variables x_{n+j} to derive from equation (10) the equality*

$$\sum_{j=1}^n c_j x_j = \left(z^* - \sum_{i=1}^m b_i y_i^* \right) + \sum_{j=1}^n \left(\left(\sum_{i=1}^m a_{ij} y_i^* \right) - y_{m+j}^* \right) x_j. \quad (11)$$

Interestingly, since these are equations that hold for any values of the x_j s, we may experiment with various choices. For example, if each $x_j = 1$ we obtain

$$\sum_{j=1}^n c_j = \left(z^* - \sum_{i=1}^m b_i y_i^* \right) + \sum_{j=1}^n \left(\sum_{i=1}^m a_{ij} y_i^* - y_{m+j}^* \right).$$

Unfortunately, that experiment tells us nothing to help us show that (5) or (9) hold.

Workout 4.1.12 What choice of values for the x_j s, plugged into (11), show immediately that (9) holds?

Now that (9) holds, we see that (11) reduces to

$$\sum_{j=1}^n c_j x_j = \sum_{j=1}^n \left(\sum_{i=1}^m a_{ij} y_i^* - y_{m+j}^* \right) x_j. \quad (12)$$

Of course, we can try similar experiments on equation (12).

Workout 4.1.13 What choice of values for the x_j s, plugged into (12), shows that

$$c_1 = \sum_{i=1}^m a_{i1} y_i^* - y_{m+1}^* ?$$

Workout 4.1.14 Do for any c_k what you did for c_1 and use your results to show that (5) holds.

Now that (5), (6) and (9) have been verified for the y_i^* s, the Duality Theorem has been proved. \diamond

Workout 4.1.15 Suppose P is a linear problem with 4 variables and 7 constraints.

- (a) If x_3^* is in the basis, what does that say about some optimal objective coefficient?
 - [i] In turn, what does that say about some optimal dual variable value?
 - [ii] In particular, what does that say about some optimal dual constraint?
- (b) If x_8^* is in the basis, what does that say about some optimal objective coefficient?
 - [i] In turn, what does that say about some optimal dual variable value?
 - [ii] Also, what does that say about some optimal primal constraint?

4.2 Complementary Slackness Conditions

Let us return to Problem 4.1.5 and its dual Problem 4.2.1.

Problem 4.2.1

$$\begin{array}{llllllll} \text{Min. } w & = & 211y_1 & + & 189y_2 & - & 106y_3 & + & 175y_4 \\ \\ \text{s.t.} & & 4y_1 & + & 6y_2 & - & 5y_3 & + & 3y_4 & \geq & -22 \\ & & y_1 & & & + & 4y_3 & + & 9y_4 & \geq & -18 \\ & & -3y_1 & + & 2y_2 & - & 2y_3 & & & \geq & -27 \\ & & & & 5y_2 & & & - & 2y_4 & \geq & -23 \\ & & 2y_1 & - & y_2 & - & y_3 & + & y_4 & \geq & 16 \\ & & 7y_1 & + & 8y_2 & - & 7y_3 & + & 4y_4 & \geq & -12 \\ \\ \& & y_1 & , & y_2 & , & y_3 & , & y_4 & \geq & 0 \end{array}$$