

Other Complexity Classes

Computational Complexity

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Introduction

In computational complexity theory, a complexity class is a set of problems of related resource-based complexity.

A typical complexity class has a definition of the form:

The set of problems that can be solved by an abstract machine **M** using **$O(f(n))$** of resource **R**, where **n** is the size of the input.

Time-Bounded Complexity Classes

These are some of the most important complexity classes defined by bounding the time of the algorithm:

Class	Computational Model	Resource Constraint
DTIME	DTM	Time $f(n)$
P	DTM	Time $poly(n)$
EXPTIME	DTM	Time $2^{poly(n)}$
NTIME	NDTM	Time $f(n)$
NP	NDTM	Time $poly(n)$
NEXPTIME	NDTM	Time $2^{poly(n)}$

Space-Bounded Complexity Classes

Some examples of classes defined by bounding the space of the algorithm:

Class	Computational Model	Resource Constraint
DSPACE	DTM	Space $f(n)$
L	DTM	Space $O(\log n)$
L^2	DTM	Space $O(\log^2 n)$
PSPACE	DTM	Space $poly(n)$
EXPSPACE	DTM	Space $2^{poly(n)}$
NSPACE	NDTM	Space $f(n)$
NL	NDTM	Space $O(\log n)$
NPSPACE	NDTM	Space $poly(n)$
NEXPSPACE	NDTM	Space $2^{poly(n)}$

Relations Between Time- and Space-Bounded Classes

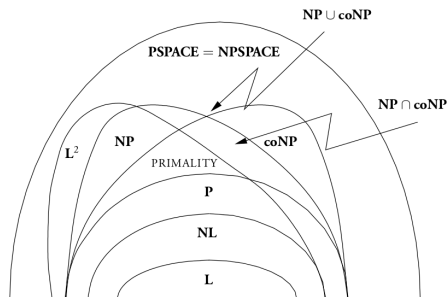


Figure: The relationships among complexity classes

$L \subseteq NL \subseteq P \subseteq NP \subseteq PSPACE \subseteq EXPTIME \subseteq NEXPTIME$

Where

$PSPACE = NPSPACE$ and $P \subset EXPTIME$

Other classes

Other important complexity classes that use different computational models include:

- BPP, ZPP and RP, which are defined using probabilistic Turing machines.
- AC and NC, which are defined using boolean circuits.
- BQP and QMA, which are defined using quantum Turing machines.

Complementary Classes

Complexity classes have a variety of closure properties; for example, decision classes may be closed under negation, disjunction, conjunction, or even under all Boolean operations.

Definition

Each **class X** that is **not closed under negation** has a **complement class co-Y**, which consists of the complements of the languages contained in X.

Example

For the **NP** class there's a **coNP** class which, for every decision problem P in **NP**, contains the complementary decision problem, denoted coP (i.e., the decision problem in which the “Yes” instances are “No” instances of **P** and vice versa).

If **C** is a deterministic time or space complexity class, then **coC** = **C**.

References



John E. Savage.

Models of Computing: Exploring the Power of Computing.
Brown University, 2008.



Wikipedia: Complexity class.