## Linear Algebra

Matrices and Vectors. First Introduction, Basic Operations.

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MDS FCS HSE

#### **Matrix**

i Definition

A  ${\bf matrix}$  is an ordered array of numbers arranged in n rows and m columns.

$$A_{n\times m} = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1m} \\ a_{21} & a_{22} & \cdots & a_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nm} \end{bmatrix}$$

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- If n=m, the matrix is called *square*; if  $n \neq m$ , it's called *rectangular*

# Basic operations: transpose

### Definition

**Transpose** of a matrix is an operation where rows and columns are swapped. If  $A \in \mathbb{R}^{n \times m}$ , then  $B = A^T \in \mathbb{R}^{m \times n}$ , where  $b_{ij} = a_{ji}$ 

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \end{bmatrix} \xrightarrow{A^T} \begin{bmatrix} a_{11} & a_{21} \\ a_{12} & a_{22} \\ a_{13} & a_{23} \end{bmatrix}$$

$$A_{2\times 3}$$

$$B_{3\times 2} = A^T$$

### Basic operations: matrix addition

i Definition

**Matrix addition** is only possible for matrices of the same size. The result is obtained by adding corresponding elements. If  $A,B\in\mathbb{R}^{n\times m}$ , then C=A+B, where  $c_{ij}=a_{ij}+b_{ij}$ 

Example:

$$\begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix} + \begin{bmatrix} 5 & 6 \\ 7 & 8 \end{bmatrix} = \begin{bmatrix} 1+5 & 2+6 \\ 3+7 & 4+8 \end{bmatrix} = \begin{bmatrix} 6 & 8 \\ 10 & 12 \end{bmatrix}$$

# Basic operations: scalar multiplication

i Definition

**Scalar multiplication** - each element of the matrix is multiplied by the given number. If  $A \in \mathbb{R}^{n \times m}$  and  $\alpha \in \mathbb{R}$ , then  $C = \alpha A$ , where  $c_{ij} = \alpha \cdot a_{ij}$ 

Example:

$$3 \cdot \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix} = \begin{bmatrix} 3 \cdot 1 & 3 \cdot 2 \\ 3 \cdot 3 & 3 \cdot 4 \end{bmatrix} = \begin{bmatrix} 3 & 6 \\ 9 & 12 \end{bmatrix}$$

Another example - Linear combination:

$$2\begin{bmatrix}1&0\\0&1\end{bmatrix}+3\begin{bmatrix}1&1\\1&1\end{bmatrix}=\begin{bmatrix}2&0\\0&2\end{bmatrix}+\begin{bmatrix}3&3\\3&3\end{bmatrix}=\begin{bmatrix}5&3\\3&5\end{bmatrix}$$

In the simplest representation, we treat a vector as a special case of a matrix:

#### Column vector

 $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \in \mathbb{R}^{n \times 1}$ 

Row vector

$$\mathbf{y}^T = \begin{bmatrix} y_1 & y_2 & \cdots & y_m \end{bmatrix} \in \mathbb{R}^{1 \times m}$$

**Dimension:**  $1 \times m$  (matrix with one row)

**Dimension:**  $n \times 1$  (matrix with one column)

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- $\bullet$  Transpose:  $\mathbf{x}^\top$  converts column to row

### Standard approach

$$\begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1,m-1} & a_{1m} \\ a_{21} & a_{22} & \cdots & a_{2,m-1} & a_{2m} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{n,m-1} & a_{nm} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_{m-1} \\ x_m \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{bmatrix}$$

$$y_1 = a_{11}x_1 + a_{12}x_2 + \dots + a_{1,m-1}x_{m-1} + a_{1m}x_m$$

$$y_j = \sum_{k=1}^m a_{jk} x_k$$

### Computational complexity and a brief look at parallel computing

### Recall the general formula:

$$y_j = \sum_{k=1}^m a_{jk} x_k, \quad j = 1, 2, \dots, n$$

#### Operations analysis:

- For computing one element  $y_j$ : m multiplications (each  $a_{jk} \cdot x_k$ ), m-1 additions (summing m products)
- For the entire vector  $\mathbf y$  (n elements):  $n \cdot (2m-1)$  total operations.

#### Time complexity:

- For square matrix  $n \times n$ :  $\mathcal{O}(2n^2 n) = \mathcal{O}(n^2)$
- For rectangular matrix  $n \times m$ :  $\mathcal{O}(nm)$
- Natural parallelism

Computing each element  $y_i$  is **independent** of other elements!

ullet Row-wise parallelization: each processor computes its own  $y_i$ 

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{bmatrix} = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix}$$

$$A \qquad \qquad \mathbf{y}$$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{bmatrix} = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ y \end{bmatrix}$$

$$A$$

$$\mathbf{x}$$

$$y_1 = a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + a_{14}x_4$$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{bmatrix} = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ y_4 \end{bmatrix}$$

$$\begin{split} y_1 &= a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + a_{14}x_4 \\ y_2 &= a_{21}x_1 + a_{22}x_2 + a_{23}x_3 + a_{24}x_4 \end{split}$$

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$$\mathbf{y} = \begin{array}{c} x_1 \begin{bmatrix} a_{11} \\ a_{21} \\ a_{31} \end{bmatrix} + x_2 \begin{bmatrix} a_{12} \\ a_{22} \\ a_{32} \end{bmatrix} + x_3 \begin{bmatrix} a_{13} \\ a_{23} \\ a_{33} \end{bmatrix} + \begin{array}{c} x_4 \begin{bmatrix} a_{14} \\ a_{24} \\ a_{34} \end{bmatrix}$$

### Visual comparison

# Matrix-by-vector product

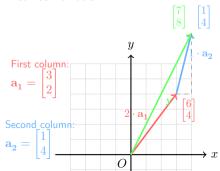
$$\begin{bmatrix} 3 & 1 \\ 2 & 4 \end{bmatrix} \qquad \begin{bmatrix} 2 \\ 1 \\ x \end{bmatrix} = \begin{bmatrix} 7 \\ 8 \\ y \end{bmatrix}$$

# Traditional View: Row-wise

### Calculations:

- $y_1 = 3 \cdot 2 + 1 \cdot 1 = 7$
- $y_2 = 2 \cdot 2 + 4 \cdot 1 = 8$

# Guru View: Column Linear Combination Linear combination:



### Standard approach

$$\begin{array}{c} a_{11} \ a_{12} \ \cdots \ a_{1k} \\ \vdots \ \vdots \ \ddots \ \vdots \\ \\ a_{i1} \ a_{i2} \ \cdots \ a_{ik} \\ \vdots \ \vdots \ \ddots \ \vdots \\ \\ a_{n1} \ a_{n2} \ \cdots \ a_{nk} \end{array}$$

*j*-th column

$$\begin{array}{c} \textbf{\textit{i-th row}} \end{array} \left[ \begin{array}{c} a_{11} \ a_{12} \ \cdots \ a_{1k} \\ \vdots \ \vdots \ \ddots \ \vdots \\ a_{i1} \ a_{i2} \ \cdots \ a_{ik} \\ \vdots \ \vdots \ \ddots \ \vdots \\ a_{n1} \ a_{n2} \ \cdots \ a_{nk} \end{array} \right] \end{array} \right] \quad \left[ \begin{array}{c} b_{11} \ \cdots \ b_{1j} \ \cdots \ b_{1m} \\ b_{2j} \ \cdots \ b_{2m} \\ \vdots \ \ddots \ \vdots \\ b_{k1} \ \cdots \ b_{kj} \ \cdots \ b_{km} \end{array} \right] \ = \ \left[ \begin{array}{c} c_{11} \ \cdots \ c_{1j} \ \cdots \ c_{1m} \\ \vdots \ \ddots \ \vdots \ \ddots \ \vdots \\ c_{i1} \ \cdots \ c_{ij} \ \cdots \ c_{im} \\ \vdots \ \ddots \ \vdots \ \ddots \ \vdots \\ c_{n1} \ \cdots \ c_{nj} \ \cdots \ c_{nm} \end{array} \right]$$

$$\begin{aligned} c_{ij} &= a_{i1} \cdot b_{1j} + a_{i2} \cdot b_{2j} + \dots + a_{ik} \cdot b_{kj} \\ \\ c_{ij} &= \sum_{l=1}^k a_{il} b_{lj} \end{aligned}$$

### Column-wise approach

$$\begin{bmatrix} a_{11} \ a_{12} \ \cdots \ a_{1k} \\ \vdots \ \vdots \ \ddots \ \vdots \\ a_{i1} \ a_{i2} \ \cdots \ a_{ik} \\ \vdots \ \vdots \ \ddots \ \vdots \\ a_{n1} \ a_{n2} \ \cdots \ a_{nk} \end{bmatrix} = \begin{bmatrix} b_{11} \ b_{12} \ b_{13} \ \cdots \ b_{1m} \\ b_{21} \ b_{22} \ b_{23} \ \cdots \ b_{2m} \\ \vdots \ \vdots \ \vdots \ \ddots \ \vdots \\ b_{k1} \ b_{k2} \ b_{k3} \ \cdots \ b_{km} \end{bmatrix} = \begin{bmatrix} c_{11} \ c_{12} \ c_{13} \ \cdots \ c_{1m} \\ \vdots \ \vdots \ \vdots \ \ddots \ \vdots \\ c_{i1} \ c_{i2} \ c_{i3} \ \cdots \ c_{im} \\ \vdots \ \vdots \ \vdots \ \ddots \ \vdots \\ c_{n1} \ c_{n2} \ c_{n3} \ \cdots \ c_{nm} \end{bmatrix}$$

Each column of  $C = A \times$  corresponding column of B

$$\mathbf{C_j} = A\mathbf{B_j}, \quad j = 1, 2, \dots, m$$

### Properties of matrix multiplication

1. Associativity:

$$(AB)C = A(BC)$$

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3. Scalar multiplication:

$$\alpha(AB) = (\alpha A)B = A(\alpha B)$$

### Important limitations

4. Non-commutativity:

$$AB \neq BA$$
 (in general)

**Example:** For  $2 \times 2$  matrices:

$$\begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} = \begin{bmatrix} 3 & 2 \\ 1 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 2 \\ 1 & 3 \end{bmatrix}$$

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5. Transpose of product:

$$(AB)^T = B^T A^T$$

і Уведомление

Note the reversed order of matrices when transposing!

### Computational complexity and parallelization

#### Recall the general formula:

$$c_{ij} = \sum_{l=1}^{k} a_{il} b_{lj}, \quad i = 1, 2, \dots, n, \quad j = 1, 2, \dots, m$$

## Operations analysis for matrices $A_{n\times k}$ and $B_{k\times m}$ :

- ullet For computing one element  $c_{ij}$ : k multiplications, k-1 additions
- Total number of operations:  $\vec{n} \times m \times (2k-1)$

- For square matrices  $n \times n$ :  $\mathcal{O}(2n^3 n^2) = \mathcal{O}(n^3)$ 
  - For rectangular matrices: O(nmk)
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    - ullet By elements: each  $c_{ij}$  is computed independently

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  - By rows: each processor handles its rows of matrix A
     By columns: each processor handles its columns of matrix B
  - Block approach: dividing matrices into blocks for efficient cache usage

# Example. Simple but important idea about matrix computations.

Suppose you have the following expression

$$b = A_1 A_2 A_3 x,$$

where  $A_1,A_2,A_3\in\mathbb{R}^{3\times 3}$  are random square dense (fully filled with numbers) matrices, and  $x\in\mathbb{R}^3$  is a vector. You need to compute b.

Which approach is best to use?

- 1.  $A_1A_2A_3x$  (left to right)
- 2.  $(A_1(A_2(A_3x)))$  (right to left)
- 3. It doesn't matter
- 4. The results of the first two options will not be the same.

Check the attached .ipynb file in the repository.