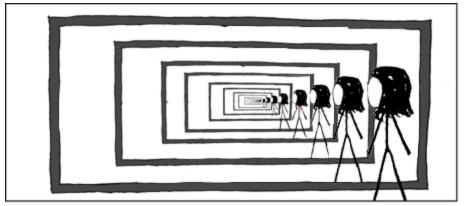
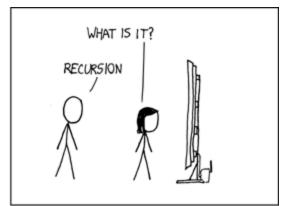
Chapter 14





http://xkcdsw.com/1105

 A recursive function is a function that calls itself

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https://www.google.com/?
gws_rd=ssl#q=recursion

A simple recursive function looks like this

```
void recurse()
{
      cout << "aaaaaaa";
      recurse();
}
cout << "... go to the castle A";
recurse();
cout << "rrrrrrg" << endl;</pre>
```

A simple recursive function looks like this

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void recurse()
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```

 When writing a recursive function, don't forget a base stopping case!

 When writing a recursive function, don't forget a base stopping case!

```
void recurse()
{
     cout << "aaaaaa";
     recurse();
}
How can we make this function stop?
cout << "... go to the castle A";
recurse();
cout << "rrrrrrg" << endl;</pre>
```

There are two forms of recursion

- There are two forms of recursion
 - Direct
 - The form we have seen so far
 - A function that calls itself, ie function B calls function B

- There are two forms of recursion
 - Direct
 - The form we have seen so far
 - A function that calls itself, ie function B calls function B
 - Indirect
 - A function that calls another function
 - Function A calls B calls C calls A

What good are recursive functions?

What good are recursive functions?

 Some algorithms work better, or seem cleaner, written as a recursive function

What good are recursive functions?

- Some algorithms work better, or seem cleaner, written as a recursive function
 - Binary Search
 - O ...

Now, we will look at some recursive examples

A recursive Factorial function

- A recursive Factorial function
 - In mathematics, the factorial of a non-negative integer n, denoted by n!, is the product of all positive integers less than or equal to n. For example, The value of 0! is 1, according to the convention for an empty product.

```
3! =
```

Factorial example

```
3! = 3 * 2!
```

2! =

```
3! = 3 * 2!
2! = 2 * 1!
1! =
```

```
3! = 3 * 2!
2! = 2 * 1!
1! = 1 * 0!
0! =
```

```
3! = 3 * 2!
2! = 2 * 1!
1! = 1 * 0!
0! = 1
```

```
3! = 3 * 2 * 1 * 1
```

Factorial example

How can we write this factorial function?

What's the base case?

- What's the base case?
 - o 0! == 1

```
int factorial(int value)
      if (num == 0) // The algorithm base case
            return 1;
      else
            return num * factorial(value - 1);
```

 Another example of a recursive function is the Greatest Common Divisor

- Another example of a recursive function is the Greatest Common Divisor
- In mathematics, the greatest common divisor (gcd) of two or more integers, when at least o them is not zero, is the largest positive integer divides the numbers without a remainder. For example, the GCD of 8 and 12 is 4.

Euclids Algorithm

```
gcd(x,y) = y If y divides x with no remainder = gcd(y, remainder of x/y) Otherwise
```

GCD Example

gcd(49,28)

GCD Example

$$gcd(49,28) == 7$$

What is the base case?

- What is the base case?
 - \circ x modulo y == 0

- What is the base case?
 - \circ x modulo y == 0

How do we implement GCD recursively?

```
int gcd(int x, int y)
      if (x \% y == 0) // The algorithm base case
            return y;
      else
            return gcd (y, x % y);
```

 The Fibonacci Numbers calculations is a classic example of recursion

- The Fibonacci Numbers
 - A sequence of numbers where
 - $= F_0 = 0$
 - $F_1 = 1$
 - \blacksquare $F_{N}^{\cdot} = F_{N-1} + F_{N-2}$ for all N>= 2

The Fibonacci Numbers

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, ...

What are the base case(s)?

- What are the base case(s)?
 - \circ $F_0 = 0$ \circ $F_1 = 1$

What are the base case(s)?

- \circ $F_0 = 0$ \circ $F_1 = 1$

How would we implement this recursively?

Recursion Examples: Searching

Recursion can also be used in searching

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Recursion can also be used in searching

 How can we use recursion to find an element in a list?

 You might remember from Chapter 9 that the Binary Search is more efficient than a recursive search

 You might remember from Chapter 9 that the Binary Search is more efficient than a recursive search

A reminder of the algorithm...

- Where x is the desired item
- 1. Check the middle item, y
 - 1.1. If x == y, then return
 - 1.2. if x < y, go to step 1 with the smaller half of the array
 - 1.3. if x > y, go to step 1 with the larger half of the array

How can we implement this as a recursive algorithm?

How can we implement this as a recursive algorithm?

What are our base cases?

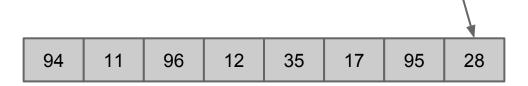
```
int binarySearch(const int* pArray, int first, int last, int value){
      int middle;
      if (first > last)
            return -1;
      middle = (first + last) / 2;
      if (pArray[middle] == value)
            return middle;
      if (pArray[middle] < value)
            return binarySearch(array, middle + 1, last, value);
      else
            return binarySearch(array, first, middle - 1, value);
```

There are also sorting algorithms that can use recursion

- There are also sorting algorithms that can use recursion
 - Quick Sort

- Basic Algorithm
 - o if the number of elements in array is 0 or 1, return
 - pick any element as pivot
 - partition array into two disjoint groups
 - left group are items less than the pivot
 - right group are items greater than the pivot
 - Perform the quick sort on the left and right sides

Choose pivot



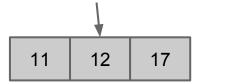
Create two lists

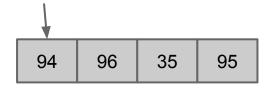
11 | 12 | 17

94 96 35	95
----------	----

Continue. quicksort(left), Quisort(right).

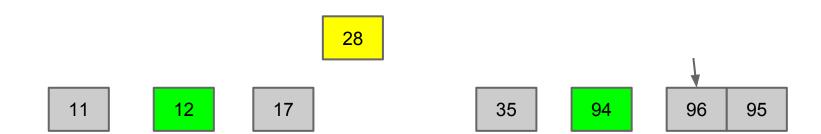
Choose pivot(s)





Create two arrays

Continue. quicksort(left) and quicksort (right). Choose pivot(s)



Break into two arrays

 28

 11
 12

 17
 35

 94

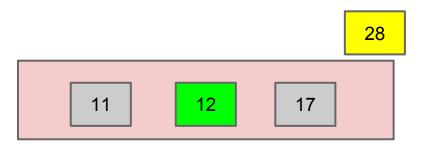
95

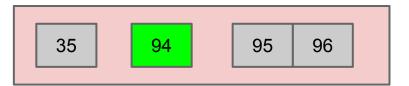
Move back up recursion. Join lists. Right + pivot + left

 11
 12
 17
 35
 94

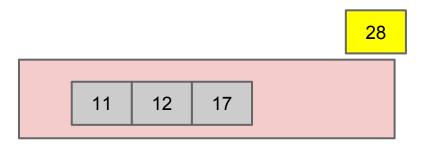
 95
 96

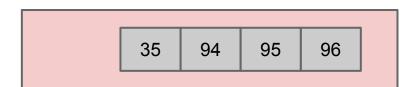
Move back up recursion. Join lists. Right + pivot + left



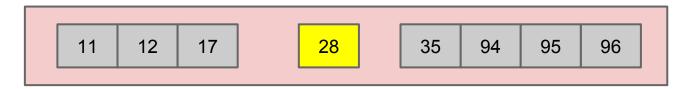


Move back up recursion. Join lists. Right + pivot + left

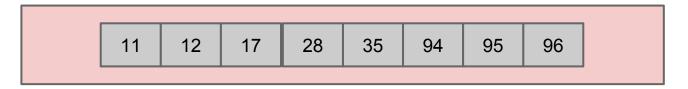




Move back up recursion. Join lists. Right + pivot + left



Move back up recursion. Join lists. Right + pivot + left



Done

11	12	17	28	35	94	95	96
----	----	----	----	----	----	----	----

How can we implement this recursively?

How can we implement this recursively?

What are the base cases?

```
void quickSort(int* pArray, int start, int end){
    if (start < end){
        int p = partition(pArray, start, end);
        quickSort(pArray, start, p - 1);
        quickSort(pArray, p + 1, end);
    }
}</pre>
```

```
int partition(int* pArray, int start, int end){
      int pivotValue = pArray[start];
      int pivotPosition = start;
      for (int pos = start + 1; pos \leq end; pos++){
            if (pArray[pos] < pivotValue){</pre>
                   swap(pArray[pivotPosition + 1], pArray[pos]);
                   swap(pArray[pivotPosition], pArray[pivotPosition + 1]);
                   pivotPosition++;
      return pivotPosition;
```

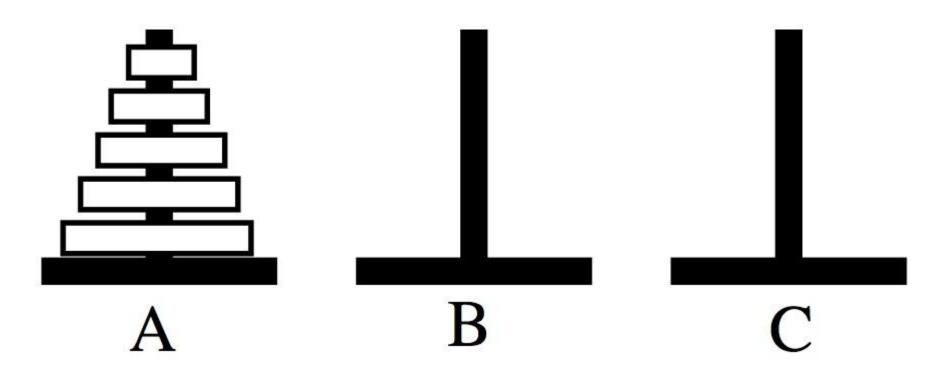
Recursion Examples: Towers of Hanoi

 The Towers of Hanoi is a classic example of a problem that is simple with recursion but otherwise difficult

Recursion Examples: Towers of Hanoi

- In the Towers of Hanoi, the rules are
 - All disks must rest on a peg except while moving
 - Only one disk may move at a time
 - No large disk can be placed on top of a smaller disk

- The game
 - A board exists with three pegs in a row
 - All of the disks are stacked greatest to least on a single peg
 - The pegs must be moved from peg one to peg three



The myth/Legend

There is a story about an Indian temple in Kashi Vishwanath which contains a large room with three time-worn posts in it surrounded by 64 golden disks. Brahmin priests, acting out the command of an ancient prophecy, have been moving these disks, in accordance with the immutable rules of the Brahma, since that time. The puzzle is therefore also known as the Tower of Brahma puzzle. According to the legend, when the last move of the puzzle will be completed, the world will end

http://en.wikipedia.org/wiki/Tower_of_Hanoi

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http://en.wikipedia.org/wiki/Tower_of_Hanoi

- The minimum number of moves is 2^N 1
 - o for N = 64, 2^{64} 1 == 1.8446744e+19

How do we solve this?

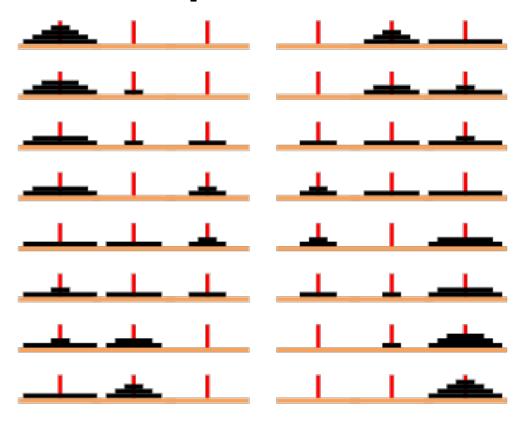
How do we solve this?

To move n disks from peg 1 to peg 3, using peg 2 as a temporary peg if n > 0 Then

Move n - 1 disks from peg 1 to peg 2, using peg 3 as temporary peg

Move a disk from peg 1 to peg 3

Move n - 1 disks from peg 2 to peg 3, using peg 1 as a temporary peg



 How do we implement the solution recursively?

 How do we implement the solution recursively?

What are our base cases?

```
HINT:
void moveDisks(int n, string source, string dest, string temp){
...
}
moveDisks(3, "peg 1", "peg 3", "peg 2");
```

```
void moveDisks(int n, string source, string dest, string temp){
     if (n > 0) {
           moveDisks(n - 1, source, temp, dest);
           cout << "Move a disk from " << source << " to " << dest << endl:
           moveDisks(n - 1, temp, dest, source);
moveDisks(3, "peg 1", "peg 3", "peg 2");
```

Some problems require a different approach

- Enumeration Algorithm
 - An algorithm that generates all possible combinations of items of a certain type

- Enumeration Algorithm
 - An algorithm that generates all possible combinations of items of a certain type
- Exhaustive Algorithm
 - An algorithm that searches through an enumeration set to find the best one

 IE examine all possible combinations and choose the best one

 IE examine all possible combinations and choose the best one

 One great example of this type of algorithm is making change in US currency

What currency do we have?

What currency do we have?

How would you make change for \$1.21?

What currency do we have?

- How would you make change for \$1.21?
 - What way is the best (least number of coins and bills)?

What currency do we have?

- How would you make change for \$1.21?
 - What way is the best (least number of coins and bills)? - Greedy Strategy
 - What way is best (most number of coins)?

 How can we implement this algorithm using recursion?

 How can we implement this algorithm using recursion?

What are the base cases?

 The greedy algorithm doesn't work for all monetary systems

- The greedy algorithm doesn't work for all monetary systems
 - Make change for \$.44 with only coins .01, .20, .25

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 - Make change for \$.44 with only coins .01, .20, .25
 - \circ Greedy says .25, and 19 x .01

- The greedy algorithm doesn't work for all monetary systems
 - Make change for \$.44 with only coins .01, .20, .25
 - Greedy says .25, and 19 x .01
 - o The best is 2 x .20 and 4 x .01

 How would we calculate the number of different ways to make change for a specified amount?

```
const int COIT_SET_SIZE = 6;
const int coinValues[] = {1, 5, 10, 25, 50, 100};
int mkChange(int amount, int largestIndex) {
```

```
const int COIT SET SIZE = 6;
const int coinValues[] = {1, 5, 10, 25, 50, 100};
int mkChange(int amount, int largestIndex) {
        while (coinValue[largestIndex] > amount)
                largestIndex--;
        if (amount == 0 || largestIndex == 0)
                return 1;
        int nWays = 0, nCouns = 0;
        while (nCoins <= amount / coinValues[largestIndex]) {
                int amountLeft = amount - nCouns * coinValues[largestIndex];
                nWays = nWays + mkChange(amountLeft, largestIndex - 1);
                nCoins++:
        return nWays;
```

 How could we modify the example to give us the least number of coins?