What does censorship-resistance look like in code?



@glozow

6B00 2C6E A3F9 1B1B 0DF0 C9BC 8F61 7F12 00A6 D25C



<u>censorship-resistance</u> is one of Bitcoin's most valuable properties



anyone can pay anyone, anywhere in the world



no central authority that can confiscate/freeze your money



even in the presence of adversaries

... but what does that look like in code?



send_payment()

send_payment_uncensorable()

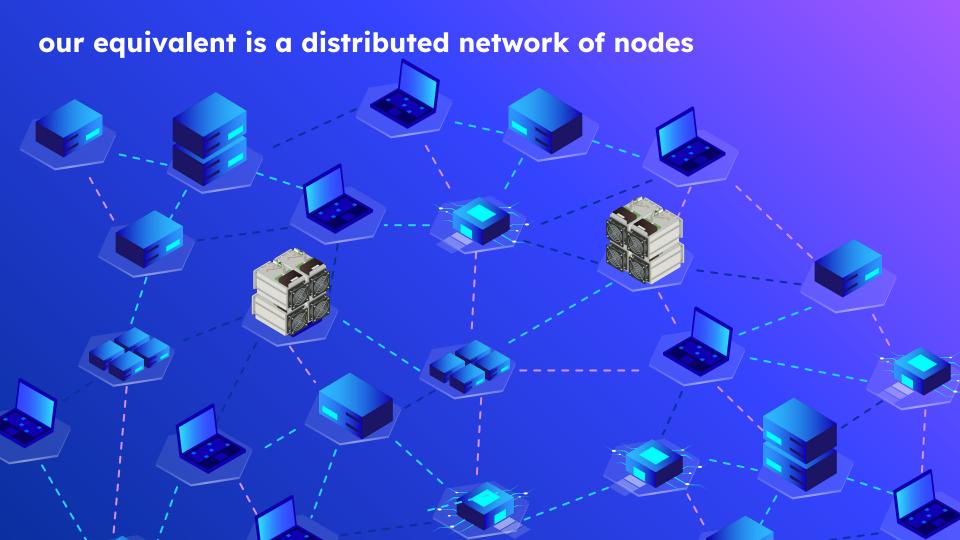
I won't display code, but am linking code for slides:

tinyurl.com/seoul-code-links









... but what does that look like in code?



woah then nobody can shut it down

We want many independent users to run nodes. From a software engineering perspective, this necessitates:



the software must be supported on many platforms



the code is open source and free to use



node operators have control of what is running



running a node must be cheap i.e. minimal operating costs, even as usage scales



mining nodes accept transactions to collect fees

non-mining nodes don't get paid



The Costs

every node has limited:

- memory
- bandwidth
- computation



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tx relay requires:

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tx relay requires:

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block validation requires:

- memory
- bandwidth
- computation



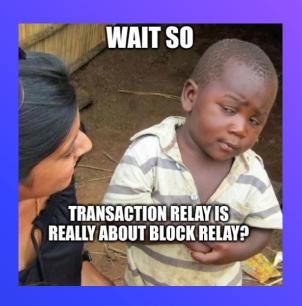
The Benefits

Transaction relay amortizes block validation:

- loading UTXOs from disk
- downloading block data
- verifying signatures and scripts

This means:

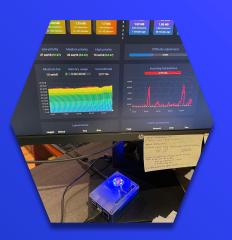
- faster block relay (steadily instead of in spikes)
 - when network-wide = fewer reorgs
- more accurate fee estimation





The challenge = write <u>useful</u> code that can run with constrained resources (e.g. a raspberry pi or a laptop app)







transaction relay vs **memory**



Use memory wisely

Several caches exist:

- UTXO ("coins") cache
- signature cache
- script verification result cache
- memory pool of validated transactions
- orphan pool of transactions for which we are missing UTXOs
- ring buffer of rejected and replaced transactions



A mempool is just a cache

= a <u>memory pool</u> of transactions that may end up in a block

- memory bounds: limit maximum usage (300MB)
- hit rate: what ultimately matters (compact block reconstruction)
- define a <u>utility metric</u> and <u>eviction policy</u>



Pop Quiz:

What happens when there are so many transactions that your mempool exceeds 300MB?



utility metric = fees

a tx is useful if it ends up being confirmed

- => a tx is mined if it would be profitable
 - => incentive compatibility* is a good utility metric

Eviction policy:

- remove transactions > 2 weeks old
- skim away transactions with the worst fees*
- allow replacements with better fees*





"free relay" and min relay fee

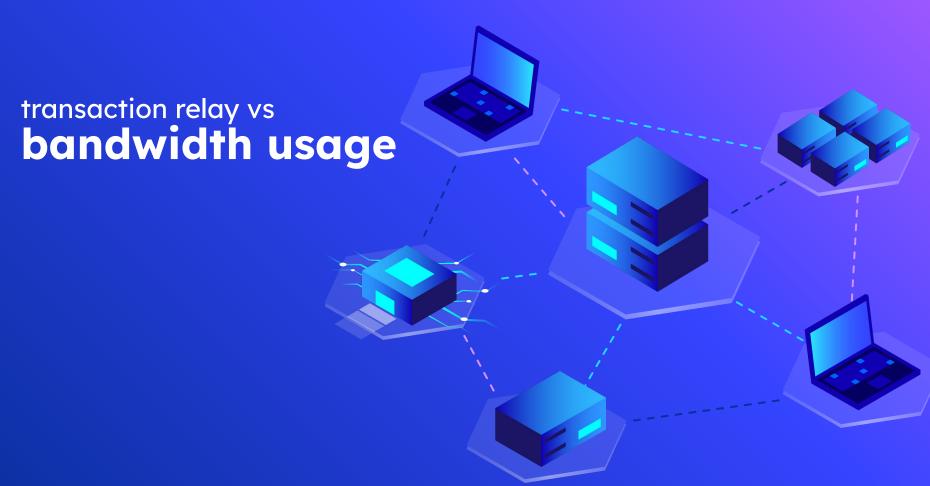
- fees are not paid to the node operator
- however, fees represent a cost to tx creator



<u>free relay</u> = when total cost < memory (+bandwidth) used at 1sat/vB

- -minrelaytxfee: all txns must pay 1sat/vB, even if mempool is empty
- RBF rule #4: replacement must pay additional fees that cover the new transaction's size at 1sat/vB
- the tx also needs to be valid (no ultimate cost if it never confirms!)





everybody gossip to everybody:

- very fast propagation!

- O(n²) tx messages

- lots of wasted bandwidth

- they might even be invalid





instead:

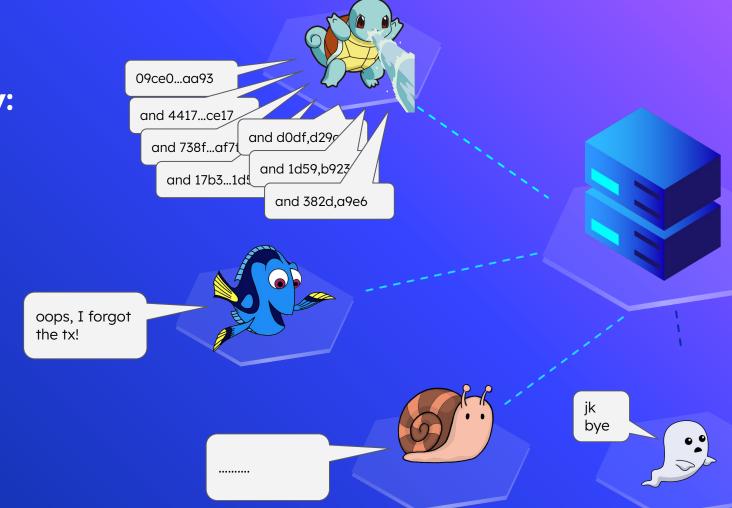
- gossip inventory by tx hash
- request what is missing from 1 peer
 - O(n²) 32b invs
 - O(n) tx messages, usually
- don't re-download rejected transactions





What if they:

- flood
- forget
- stall
- drop off





We need to track all announcers

- double up requests after some timeout
- have preferences (e.g. outbounds)
- load-balance between peers





How else can somebody sabotage a tx?

- malleation of tx data
- pinning through attaching children or replacements







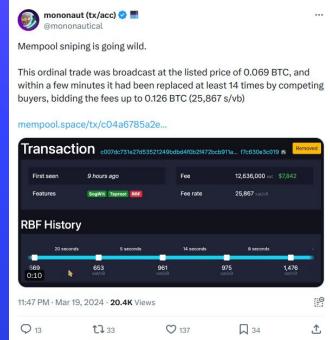




001

"mempool sniping"

- (1) seller creates a malleable transaction, signed SIGHASH_SINGLE | SIGHASH_ANYONECANPAY, exchanging the NFT for the sell price
- (2) buyer signs PSBT providing funds and sending NFT to themselves
- (3) anyone can malleate the tx after broadcast, providing new funds and changing NFT destination (original buyer does not lose their money)



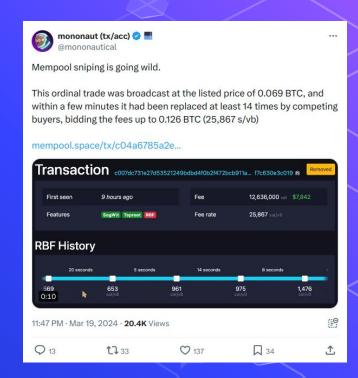


"mempool sniping"

Lesson: if you sign a transaction with SIGHASH_ANYONECANPAY...

... anyone can change its inputs

- on't design stuff like this
- we can't fix this in transaction relay



We can test censorship-resistance

- Add 1 honest peer + n "evil" peers programmed with possible actions:
 - forgetting
 - stalling 🀠
 - flooding 🦠
- - malleating txs
 - disconnecting 🥹
- Simulate random combinations of mesages and assert that we always receive the tx







transaction relay vs **computation**



computation has to be limited

our operating environment = anonymous entities on the internet connect to you and send you data to process

- blocks are naturally DoS-resistant (must have PoW!)
- transactions can be computationally expensive to handle:
 - signature verification
 - mempool package updates
 - etc

signature verification

quadratic sighashing and the 1MB "Megatransaction"

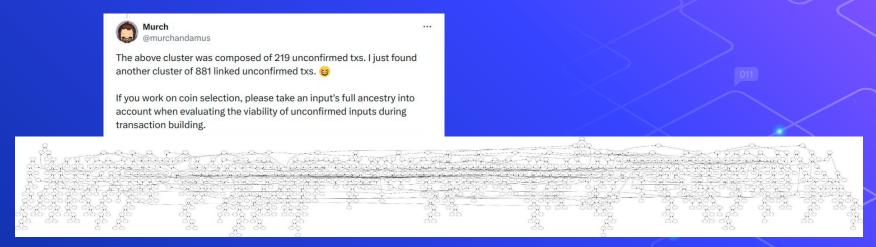


only validate transactions of limited size and sigops, only allow up to 15-of-15 legacy multisig

mempool package updates

insertion and removal require updating ancestors and descendants

fimit the number and size of ancestor and descendant packages





- ... limit the cluster size too

Transaction Relay / Mempool Policy

Bitcoin Core rejects some consensus-valid transactions

- decentralization: minimize memory, bandwidth, computation
- **security: prevent DoS, specific bugs not fixed in consensus
- upgrade hooks: don't use future things before they are defined (avoid collisions and confiscation)
- best practices aka paternalism: discouraging network resource-intensive practices



default policy is *like* a transaction relay interface, but it's optional and varies



Wait, but isn't this CENSORSHIP?



They're all tradeoffs

- accuracy of assessing incentive compatibility / computational complexity
- allowing complex topologies / fixing pinning (censorship) problems
- relaying all consensus-valid transactions / safely updating consensus later

- 👉 to design for censorship-resistance (and other goals), we have to reject some
- the point is to agree on an interface that is both useful for users and acceptable for node operators, not stop users from broadcasting transactions



Practice question from twitter

- In the past, some rules were added to discourage practices proactively
 - uneconomical ("dust") outputs
 - large (nonstandard) scriptPubKeys

Q: Why can't Bitcoin Core just filter {inscriptions, ordinals, brc-20s} too? #FixTheFilters

"Just make Bitcoin Core filter these transactions"



miners are not obligated to run software that throws away fee revenue



this doesn't work on transactions that already bypass policy rules



if you discard transactions that are likely to be mined, your cache is useless

Summary

- Design goals like "censorship resistance" and "decentralization" aren't easy to achieve
- They do create interesting technical puzzles
- All the code is open source that's important
- X YOU can see for yourself how it works

Go deeper!



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Thanks!

Any questions?

@glozow

