Objective: To practice B+-tree and Concurrency Control

## Part 1: B+Tree

- 1. Consider the B+ tree index of order n = 5 shown in Figure 10.1.
  - i. Show the tree that would result from inserting a data entry with key 9 into this tree.
  - ii. Show the B+ tree that would result from inserting a data entry with key 3 into the original tree.
  - iii. Show the B+ tree that would result from deleting the data entry with key 8 from the original tree, assuming that the left sibling is checked for possible redistribution.
  - iv. Show the B+ tree that would result from deleting the data entry with key 8 from the original tree, assuming that the right sibling is checked for possible redistribution.
  - v. Show the B+ tree that would result from starting with the original tree, inserting a data entry with key 46 and then deleting the data entry with key 52.
  - vi. Show the B+ tree that would result from deleting the data entry with key 91 from the original tree.
  - vii. Show the B+ tree that would result from starting with the original tree, inserting a data entry with key 59, and then deleting the data entry with key 91.
  - viii. Show the B+ tree that would result from successively deleting the data entries with keys 32, 39, 41, 45, and 73 from the original tree.

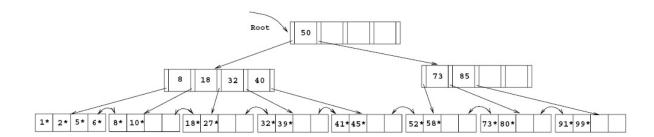


Figure 10.1 Tree for Exercise 10.1

## **Part 2: Concurrency Control**

1. Consider the following two transactions:

```
T1: r1(A); r2(C); r1(B); r2(B); r2(A); r2(A)
```

- For each of the following histories/schedules:
  - a) Is it a valid history?
  - b) Use *serializability graphs* to check whether it is serializable or not, and if it is, what is the equivalent serial history/schedule?

- 2. Consider the following history, with lock and unlock statements added for each transaction:
- a) Does the history follow 2PL protocol?
- b) Did deadlock happen?

T1	T2
rl1(A)	
r1(A)	
rl1(B)	
r1(B)	
	rl2(C)
	r2(C)
	rl2(B)
	r2(B)
wl1(B)	
w1(B)	
	rl2(A)
	r2(A)
	w12(C)
	w2(C)
	w12(A)
	w2(A)
Commit	Commit