4. Project 2. Manage Multiple Access ¶

Important

Please read the following instructions carefully:

- This project is to be completed by each group individually.
- This project contains 7 tasks. 2 of them are optional (12 points + 6 points).
- The task score does not reflect the implementation difficulty but rather the recommended level.
- Suggested workload is 6~8 FULL days. Manage your time.
- Submit your code through Blackboard. The submission due date is Nov. 19, 2023.
- Each group needs to submit the code once and only once. Immediately after TAs' checking. The submission is performed by one of the group members.
- Tasks with "Optional" tag are optional tasks. Their due date is the end of the semester.
- Unless otherwise mentioned, the instructor is responsible for grading optional tasks. Contact the instructor to check if you have finished one or more of them.
- Task 5 and 6 are graded by TAs.
- Tasks are graded according to their hierarchy. The hierarchy of this project is Task 1 <
 Task 2 < Task 3 < Task 6, and Task 1 < Task 2 < Task 3 < Task 4. A full score of one
 task automatically guarantees the full score of non-overdue preceding tasks (left side of
 "<").
- Free to use any tools for debugging, but only the provided toolkit is permitted in performance assessment.

4.1. Overview

This project aims to enable multiple devices to share the audio medium for relatively reliable data transmission, building upon the results of Project 1. To simplify the physical layer implementation, network devices will no longer be connected through sound but via audio cables. These cables carry analog electrical signals driving the loudspeakers, acting as an 'electromagnetic mirror' of the actual sound. It is worth noting that such wired connections are similar to Ethernet, making the CSMA/CD mechanism an good reference for this project.

However, unlike Ethernet, there is a significant latency (~20 ms) when using the sound card for carrier sensing. This latency arises because the sound card is not designed for this purpose, and we have to process carrier information in the digital domain. This delay reduces the efficiency of CSMA and, more importantly, prevents the implementation of an effective collision detection mechanism. Therefore, we will use ACK to handle error frames. In this sense, the access mechanism is more akin to what is adopted in wireless situations like Wi–Fi. Figure 4.1 illustrates the program architecture of this project.

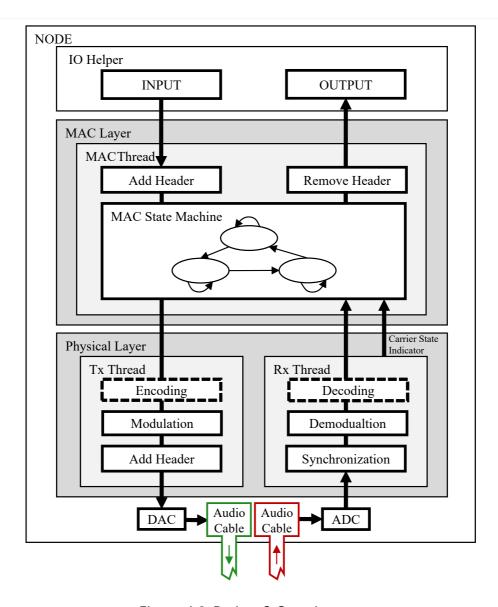


Figure 4.1. Project 2 Overview.

4.2. Task 0: (0 point) Audio Toolkit

Computers connect to audio peripherals through the audio connectors of the sound card. The typical use case involves inserting the 3.5 mm plug (M) of the microphone and/or speaker into the 3.5 mm jack (F) of the sound card. The project will use 3.5 mm audio cables to connect computers starting from this task. Due to the confusing structures of 3.5 mm <u>audio connectors</u> and diverse specifications of sound cards, we will provide the following toolkit. Please borrow it from TAs.

<u>USB Sound Card</u> ×4

This device includes 1 audio input connector (Red F), 1 audio output connector (Green F), and a USB type–A connector. When connected to the computer via USB, it will be recognized as a sound card. Its audio input/output can be read/generated, e.g., through ASIO. Note that while its connectors are all stereo, it only supports mono recording.

<u>Mixer</u> ×1

This device contains 4 input connectors (F) and 1 output connector (F). Its function is to combine the signals from 4 inputs and output them to the out-

put connector. The strength of each input/output signal can be adjusted by knobs. All gain knobs must be turned fully maximized during performance assessment.

• Audio Cable ×6

(M to M) audio cables, with lengths of 1.0 m or 0.5 m.

1-to-2 Cable Splitter ×3

This splitter has three internally-connected connectors (M, F, F). The input signal from any connector is split into the other two connectors, and the input signals from any two connectors are combined into a signal output to the remaining.

The above toolkit will be used to complete all the remaining projects. Additional cables can be requested for specific tasks, and replacements can be requested in case of suspected hardware issues.

4.3. Task 1: (4 points) Cable Connection

The ADC and DAC of the USB sound card are directly connected to the input and output connectors, so the signals flowing between connected computers are analog electrical signals (observable via an oscilloscope). The signal intensity corresponds to the intensity of the sound waveform. Compared to actual sound, electrical signals are quiet, quick, and clean. They undergo minimal attenuation, deformation, multipath, ring effects, etc., factors that can lead to noise in the ideal waveform during cable transmission (sampling frequency offset may still exist). This means a faster and more stable physical layer should be achievable.

This task revisits <u>Project1.Task3</u>. While there are equivalent connection topology, please prioritize using the one in <u>Figure 4.2</u>, as it can smoothly transition to other tasks. The mixer functions like an <u>Ethernet hub</u>, allowing all connected devices to hear each other. Connecting through hubs provides better scalability and convenience. The provided audio mixer can connect 4 computers together.

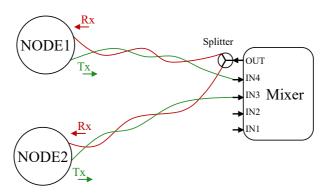


Figure 4.2. Connecting Two Nodes.

Tip

Project 1 uses carrier waves to convey information over the air, but it is no longer necessary for cable connections. You might want to simplify and speed up the physical layer with line coding. A raw rate of at least 12 Kbit/s is suggested to complete the projects with ease.

- You can design a shorter preamble to reduce the header overhead.
- It is recommended to unplug the power adapters of your computers when connecting them. The reason is that different AC-to-DC adapters, especially those without grounding, may cause non-zero voltages in the device's "ground". In some cases, this issue can lead to biased DC levels in the received signals. Occasionally, it may even cause permanent damages. Be very careful if you have to connect two devices with different "ground", such as a charging laptop and a desktop computer.

Performance Assessment

The group provides two devices: <u>NODE1</u> and <u>NODE2</u>, and connects them with the toolkit according to <u>Figure 4.2</u> (the mixer's connecter IDs are interchangeable).

• Objective (4 points). TAs provide a binary file INPUT.bin, which contains 6,250 Bytes. NODE1 sends the bits from the file to NODE2. NODE2 stores the received bits in OUTPUT.bin.

Transmission must be completed within 20 seconds:

Completion Time	Percentage Earned
<10 s	100%
<15 s	75%
<20 s	50%
>20 s	0%

TAs use diff tool to compare INPUT.bin and OUTPUT.bin:

Similarity	Percentage Earned
<80%	0%
<95%	80%
>95%	100%

4.4. Task 2: (5 points) Acknowledgement

The cable-based physical layer can dramatically reduce transmission errors, but errors cannot be completely avoided. For example, non-real-time scheduling of the ASIO process by the operating system may occasionally lead to audio samples missing. More importantly, when multiple nodes share the same cable, decoding errors caused by collisions are almost inevitable. This task will utilize the unreliable data link posed by the physical layer to achieve a much higher level of transmission reliability by using ACKs.

The logic of ACK is straightforward: the sender keeps re-/transmitting until it receives the receiver's acknowledgement of correct reception. The state machine of a simple ACK protocol is shown in <u>Figure 4.3</u>, where each circle represents the working state of the MAC thread. The arrows between states highlight events triggering state transitions (uppercase events are from the physical layer) and the associating operations (>>). In the Rx Frame and Tx Frame states, the MAC thread invokes the data transmission and reception functions provided by the physical layer.

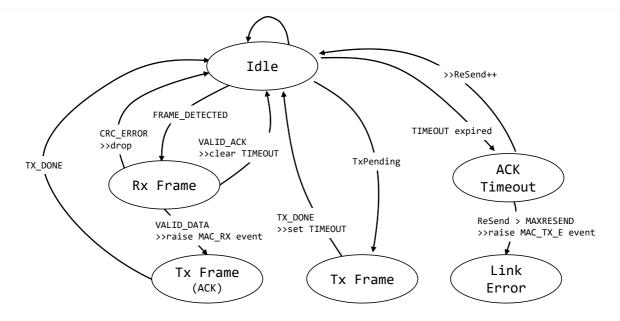


Figure 4.3. Example: State Machine of ACK Protocol.

Specifically,

• The upper-layer application appends a DATA frame to the transmission buffer and raise the TxPending flag. When the MAC thread is idle, it sends the DATA frame and sets a timer with a duration of TIMEOUT.

If no ACK frame is received before TIMEOUT expires, the sender retransmits the frame.

If the maximum number of retransmissions is reached, the MAC terminates and raises an error.

The MAC thread continues to receive/detect frames while idle.

If the received frame is corrupted, it is discarded.

If the received frame is error-free, then

If it is an <u>ACK</u> frame, indicating the previously sent <u>DATA</u> frame was received correctly, the <u>TIMEOUT</u> timer is cleared to prevent retransmission.

If it is a DATA frame, the received content is appended to the reception buffer. The MAC thread reports to the upper-layer, and sends an ACK frame immediately.

In the above protocol, it is important for the receiver to be able to distinguish between DATA frames and ACK frames. This is because, although ACK frames are sent by the receiver itself, due to driver delays (approximately 30~40 ms round-trip time), the sender might receive its own ACK while it is in the idle state. Therefore, the protocol requires adding auxiliary information to the frames, i.e., the MAC layer header. For example, in Figure 4.4, the Type field can be used to incorporate a flag to differentiate between ACK and DATA. Note that the aforementioned issue can also be resolved by specifying the receiver ID, i.e., the MAC address, in the frames. MAC addresses provide other essential functionalities, which will become apparent in the subsequent tasks.

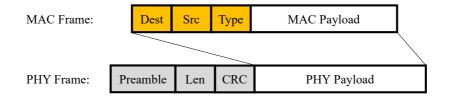


Figure 4.4. Example: MAC Layer Frame Structure.

Tip

• The length of the MAC address depends on the number of devices using the protocol. It can be shortened to fit the actual scale and reduce protocol overhead.

Performance Assessment

The group provides two devices: NODE1 and NODE2, and connects them with the toolkit according to Figure 4.2 (the mixer's connecter IDs are interchangeable).

Objective 1 (4 points). TAs provide a binary file <u>INPUT.bin</u>, which contains 6,250 Bytes.
 NODE1 sends the bits from the file to <u>NODE2</u>. <u>NODE2</u> stores the received bits in <u>OUTPUT.bin</u>.

This part's assessment criteria and procedures are similar to Task 1, with modifications made to the completion time and accuracy considering the protocol overhead and effectiveness of the ACK mechanism.

Transmission must be completed within 40 seconds:

Completion Time	Percentage Earned
<20 s	100%
<30 s	75%
<40 s	50%
>40 s	0%

TAs use diff tool to compare INPUT.bin and OUTPUT.bin:

Similarity	Percentage Earned
<100%	0%
100%	100%

• Objective 2 (1 point). Repeat the above file transfer with the same implementation. The TAs will unplug one of the cables and start timing until the transmitter raises the "link error" warning. This delay is recorded as the detection delay.

Detection Delay	Percentage Earned
<5 s	100%
<8 s	50%
>8 s	0%

4.5. Task 3: (2 points) Carrier Sense Multiple Access

If both nodes in Task 2 want to send data to each other, the ACK protocol alone will no longer be sufficient. This is because both nodes will directly send packets at will, leading to transmission collisions. To allow multiple nodes to share the transmission medium, MAC (Medium Access Control) protocols can be used. This task will implement the CSMA (Carrier Sense Multiple Access) protocol, which has been adopted in early Ethernet and many current wireless networks. CSMA can be described as listen before transmit (carrier sensing) and wait for a random period (backoff) after a collision occurs before attempting to retransmit.

<u>CSMA</u> implementations of different level of aggressiveness differ in the backoff strategy. One simple strategy is shown in <u>Figure 4.5</u>: the node performs medium sensing before each transmission attempt. If the medium is found busy or if a transmission error occurs, the node waits for a random period before attempting to transmit again.

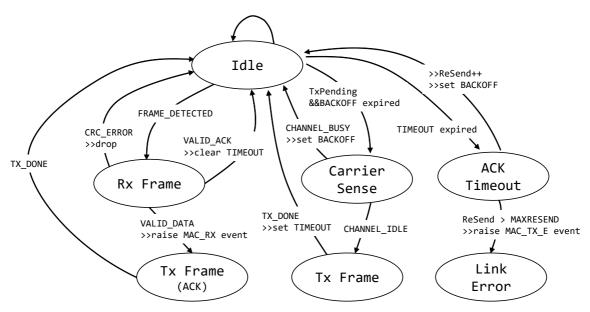


Figure 4.5. Example: State Machine of CSMA Protocol.

Specifically,

- After sending DATA and not receiving an ACK before the TIMEOUT, set the BACKOFF timer to a random value.
- When the BACKOFF timer reaches zero, sense the medium (generally, at least longer than the frame interval between a DATA and the corresponding ACK).

If the medium is idle, proceed with immediate transmission.

If there are other signals on the medium, set the backoff timer to a random value and attempt transmission again.

The <u>WARP CSMAMAC</u> provides a detailed and clear description of a CSMA protocol (with slight differences from the above) from the implementation's perspective, which should be considered an important reference for this task.

Performance Assessment

The group provides two devices: <u>NODE1</u> and <u>NODE2</u>, and connects them with the toolkit according to <u>Figure 4.2</u> (the mixer's connecter IDs are interchangeable).

• Objective (2 points). TAs provide a binary file INPUT1to2.bin, which contains 6,250 Bytes. NODE1 sends the bits from the file to NODE2 stores the received bits in OUTPUT1to2.bin. TAs provide a binary file INPUT2to1.bin, which contains 6,250 Bytes. NODE2 sends the bits from the file to NODE1. NODE1 stores the received bits in OUTPUT2to1.bin. NODE1 and NODE2 start transmission simultaneously. TAs record the completion times for NODE1 and NODE2 stransmissions as T1 and T2.

Transmission must be completed within 80 seconds:

MAX(T1,T2)	Percentage Earned
<40 s	100%
<60 s	75%
<80 s	50%
>80 s	0%

Their completion times should be close enough:

ABS(T1-T2)	Percentage Earned
<10 s	100%
>10 s	0%

TAs use diff tool to compare INPUT*.bin and OUTPUT*.bin:

Similarity	Percentage Earned
<100%	0%
100%	100%

4.6. Task 4: (1 point) CSMA with Interference

This task will apply CSMA in a more realistic environment:

- There might be more than two nodes participating in the network transmission.
- Hidden terminals could exist in the network.
- Outdated, selfish, and poorly-implemented nodes might also be present in the network.

To understand the performance of CSMA in these scenarios, a jamming source, denoted as the Jammer, is connected to the network, as shown in <u>Figure 4.6</u>. The Jammer does not adhere to CSMA and emits interfering signals randomly. However, the <u>Jammer</u> is not always active, allowing data transmission during its silent periods.

The Jammer's behavior can be described as follows: its jamming duration follows a uniform distribution between [50, 100) ms. During the jamming period, the Jammer emits a white noise signal. After each jamming period, it switches to silent before initiating the

next jamming, with each silent period following a uniform distribution between [100, 200) ms.

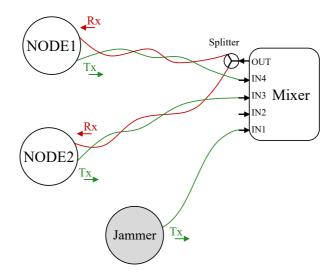


Figure 4.6. Two Nodes and a Jammer.

Performance Assessment

The group provides two devices: NODE1 and NODE2 and TAs provide a Jammer. They are connected according to Figure 4.6. The Jammer uses the USB sound card from the toolkit to emit jamming signals. The script for generating the jamming signal is uploaded to Blackboard. The output gains/volumes of the Jammer and NODE* operating systems are set to similar levels (± 10%).

Transmission must be completed within 120 seconds, other assessment criteria and procedures are the same as in Task 3:

MAX(T1,T2)	Percentage Earned
<80 s	100%
<100 s	75%
<120 s	50%
>120 s	0%

Any group intending to finish Task 5 must inform the TAs explicitly during this task. TAs will repeat the transmission task five times and record the mean of MAX(T1, T2) to rank the performance.

4.7. Task 5: (Optional, 3 points) Performance Rank

A primary design objective of network systems is to enhance data transmission performance, but achieving high performance is often challenging. This task is to reward the best-performing implementation in this project.

Performance Assessment

This task is automatically graded according to the completion time of Task 4.

Rank	Percentage Earned
1 st	100%
2nd	33%
3rd	0%

4.8. Task 6: (Optional, 3 points) X

Let's try with more nodes.

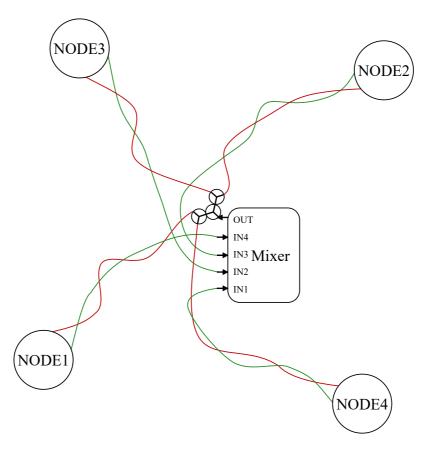


Figure 4.7. Connecting Four Nodes.

Tip

- The provided mixer is a <u>passive mixer</u>. Its output power depends on the number of connected input and output connectors. When connecting more or fewer devices, such as monitoring headphones, to it, the signal strength at the receiver should be reevaluated.
- You may need additional cables.

Performance Assessment

The group provides **four** devices: NODE1, NODE2, NODE3, and NODE4. They are connected according to Figure 4.7.

TAs provide a binary file INPUT1to2.bin, which contains 6,250 Bytes. NODE1 sends the bits from the file to NODE2. NODE2 stores the received bits in OUTPUT1to2.bin. TAs provide a binary file INPUT2to1.bin, which contains 6,250 Bytes. NODE2 sends the bits from the file to NODE1. NODE1 stores the received bits in OUTPUT2to1.bin.

TAs provide a binary file INPUT3to4.bin, which contains 6,250 Bytes. NODE3 sends the bits from the file to NODE4. NODE4 stores the received bits in OUTPUT3to4.bin. TAs provide a binary file INPUT4to3.bin, which contains 6,250 Bytes. NODE4 sends the bits from the file to NODE3. NODE3 stores the received bits in OUTPUT4to3.bin.

 $\underline{\text{NODE1}}$, $\underline{\text{NODE2}}$, $\underline{\text{NODE3}}$, and $\underline{\text{NODE4}}$ start transmission simultaneously. TAs record their completion times as T1, T2, T3, and T4.

Transmission must be completed within 160 seconds:

MAX(Ti)	Percentage Earned
<120 s	100%
<140 s	75%
<160 s	50%
>160 s	0%

Their completion times should be close enough:

MAX(Ti)- MIN(Ti)	Percentage Earned
<15 s	100%
>15 s	0%

TAs use diff tool to compare INPUT*.bin and OUTPUT*.bin:

Similarity	Percentage Earned
<100%	0%
100%	100%