

Design and Analysis of Algorithms
CS 575, Spring 2025

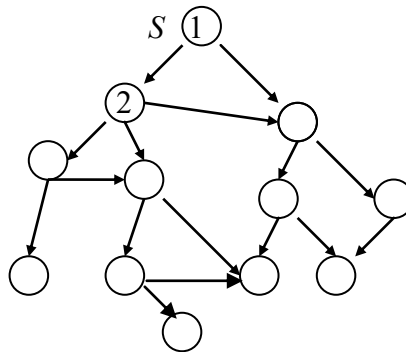
Theory Assignment 3.2

Due on 4/16/25 (Wednesday) at 11:59pm

Remember to include the following statement at the start of your answers with a signature by the side. “I have done this assignment completely on my own. I have not copied it, nor have I given my solution to anyone else. I understand that if I am involved in plagiarism or cheating I will have to sign an official form that I have cheated and that this form will be stored in my official university record. I also understand that I will receive a grade of 0 for the involved assignment for my first offense and that I will receive a grade of “F” for the course for any additional offense.”

Please handwrite or type your answer to each question, scan or save your answers into a pdf file (with a vertical orientation so that we do not need to rotate your file to grade), and upload it to the homework submission site.

1. [20 points] In this problem, you will apply different traversal methods on a given directed graph. The start node is denoted by S . When arbitrary decisions on order must be made assume that child nodes are visited from left to right.
 - a. [10 points] Perform a breadth-first search on the directed graph below (on the left). (i) *Number the nodes* according to the order in which they are visited (become gray). For example, the first node visited is S so S is numbered 1. Then, the left child of S is visited so it is numbered 2. Show the order numbers inside the circles. (ii) Show the distance of each node to S beside each circle. (iii) Show the breadth-first tree.



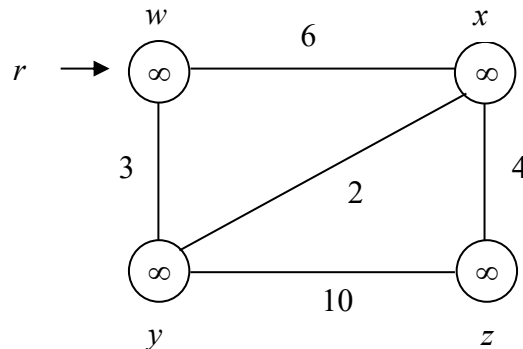
- b. [10 points] Perform a depth-first search on the directed graph below (on the left). (i) Number the nodes according to the order in which they are visited at the first time (when they become gray). Show the order numbers inside the circles. (ii) Show the depth-first tree(s).


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PRIM( $G, w, r$ )
   $Q = \emptyset$ 
  for each  $u \in G.V$ 
     $u.key = \infty$ 
     $u.\pi = \text{NIL}$ 
  INSERT( $Q, u$ )
  DECREASE-KEY( $Q, r, 0$ )      //  $r.key = 0$ 
  while  $Q \neq \emptyset$ 
     $u = \text{EXTRACT-MIN}(Q)$ 
    for each  $v \in G.Adj[u]$ 
      if  $v \in Q$  and  $w(u, v) < v.key$ 
         $v.\pi = u$ 
        DECREASE-KEY( $Q, v, w(u, v)$ )

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Apply the algorithm to the weighted, connected graph below (the initialization part has been done). Show a new intermediate graph after each vertex is processed in the while loop. For each intermediate graph and the final graph, you need to show the vertex being processed, the new key value for each vertex and edges in the current (partial) MST (draw a directed edge from vertex v to u if $v.\pi = u$).



5. [20 points] Apply Kruskal's algorithm to the graph below. Show new intermediate graphs with the shaded edges belong to the forest being grown. The algorithm considers each edge in sorted order by weight. An arrow points to the edge under consideration at each step of the algorithm. If the edges join two distinct trees in the forest, it is added to the forest, thereby merging the two trees.

