



Towards Foundational Verification of Cyber-physical Systems

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SoSCYPS 2016











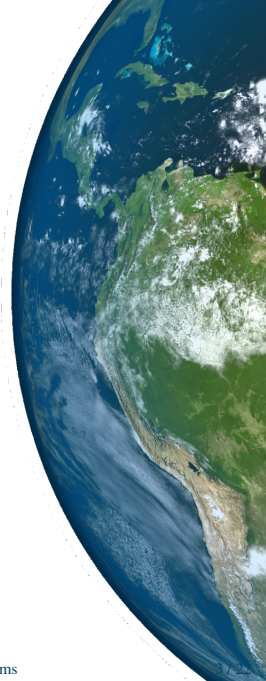
System
(Hw+C)





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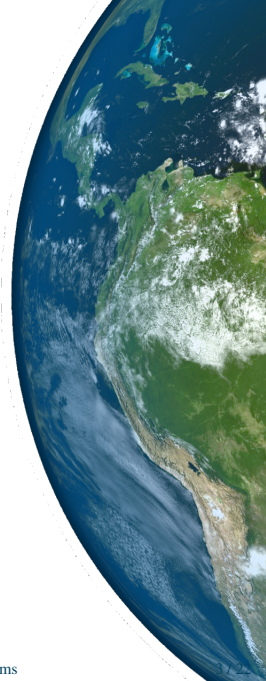
Specification
(Safety)





System
(Hw+C)

Verifier
(KeYmaera)

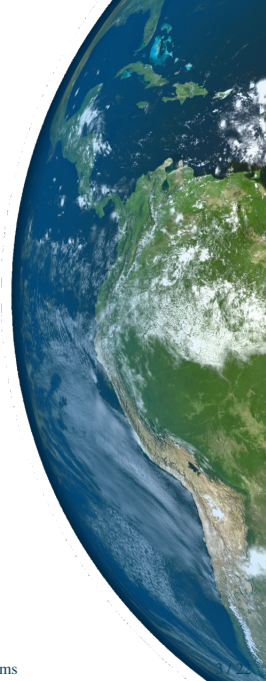




System
(Hw+C)

Model
(∂ DL)

Verifier
(KeYmaera)





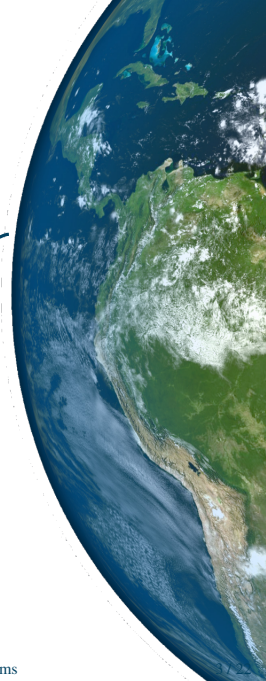
System
(Hw+C)

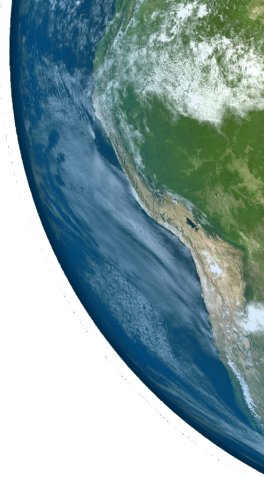


Model
(∂ DL)



Verifier
(KeYmaera)







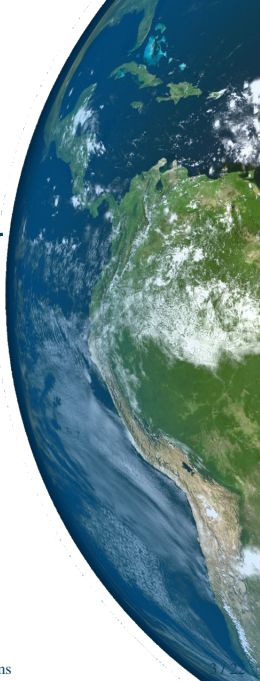
System
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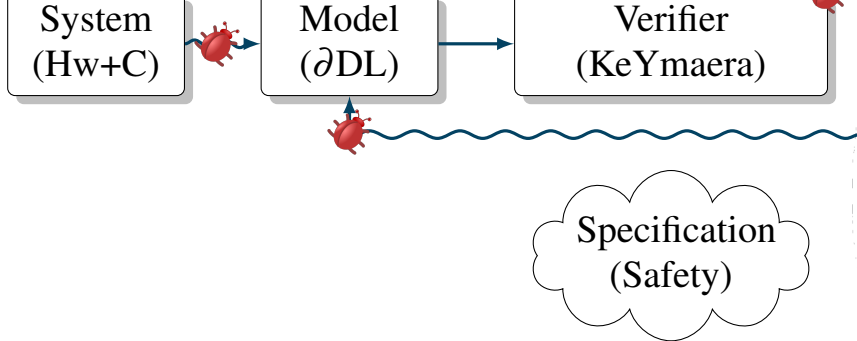


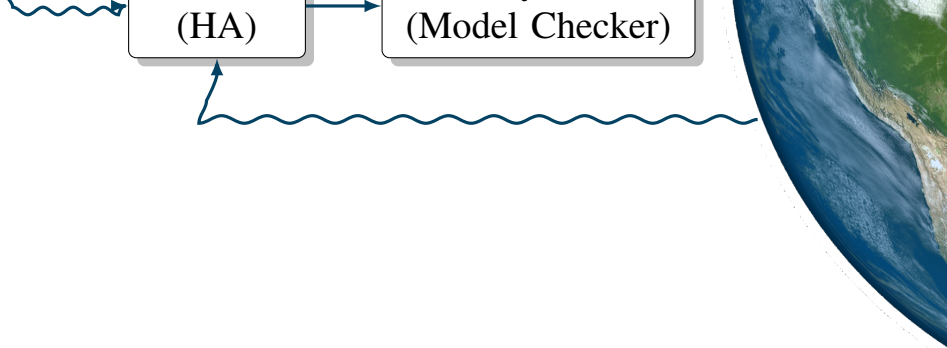
Model
(∂ DL)

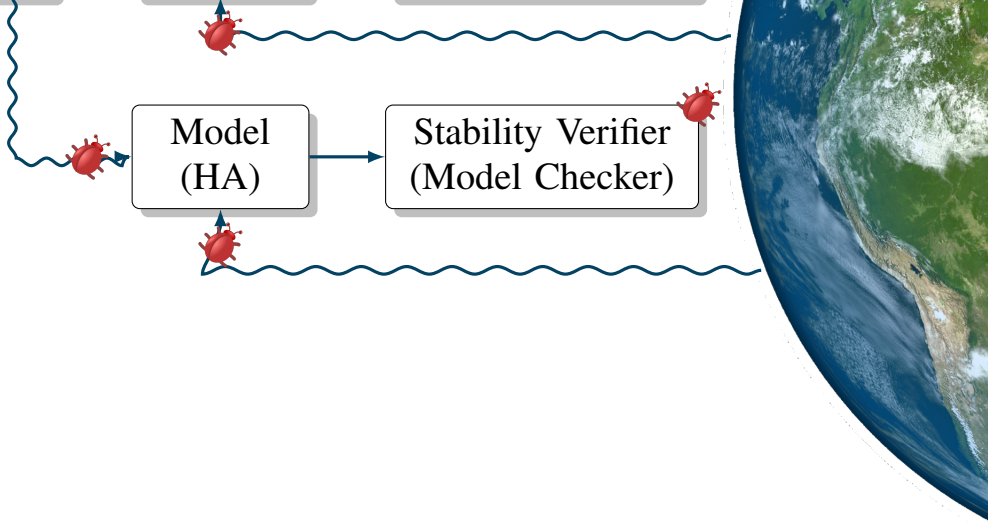


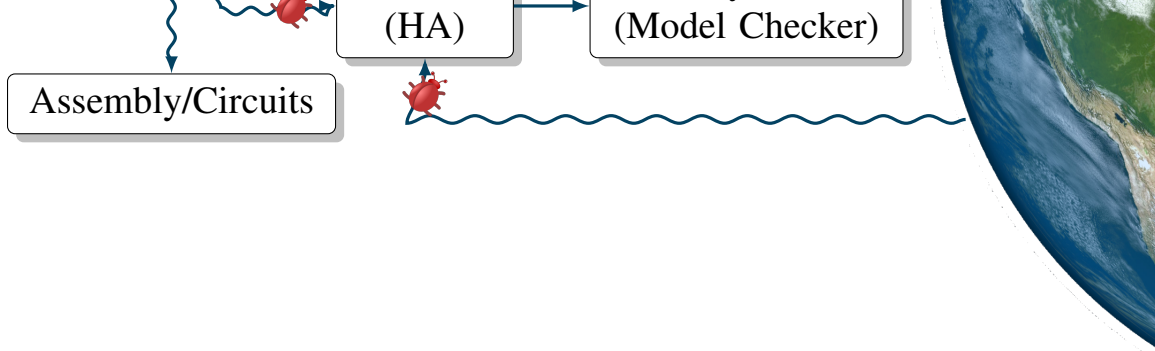
Verifier
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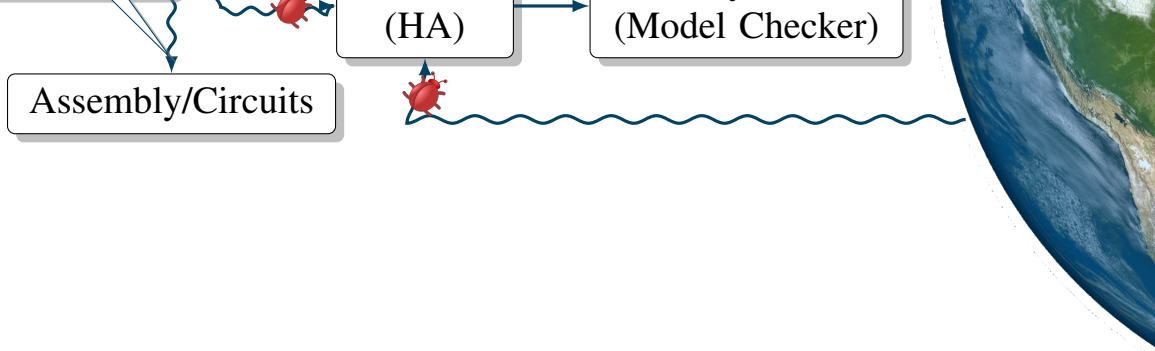


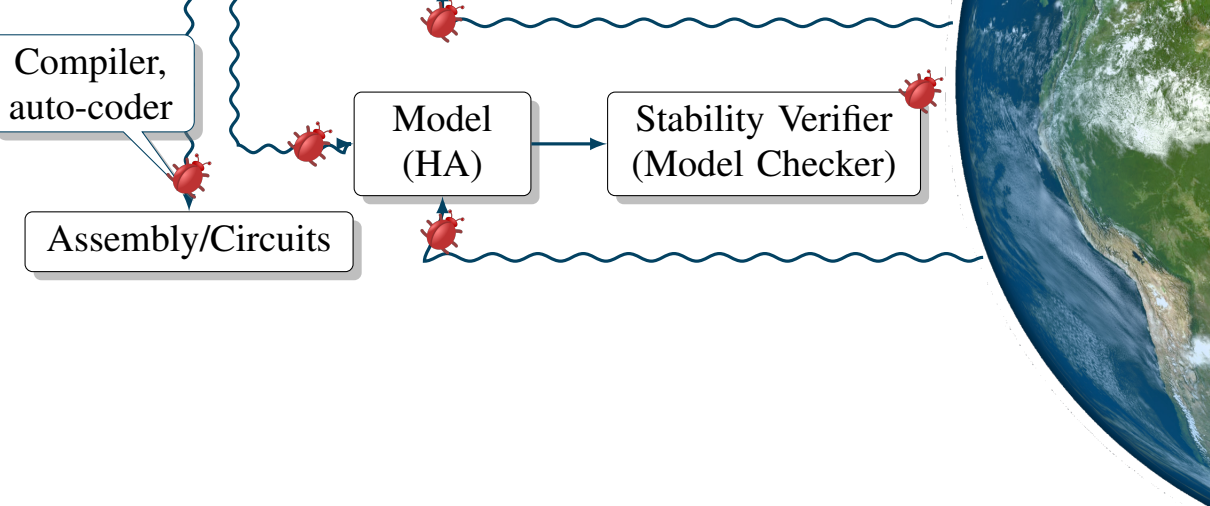


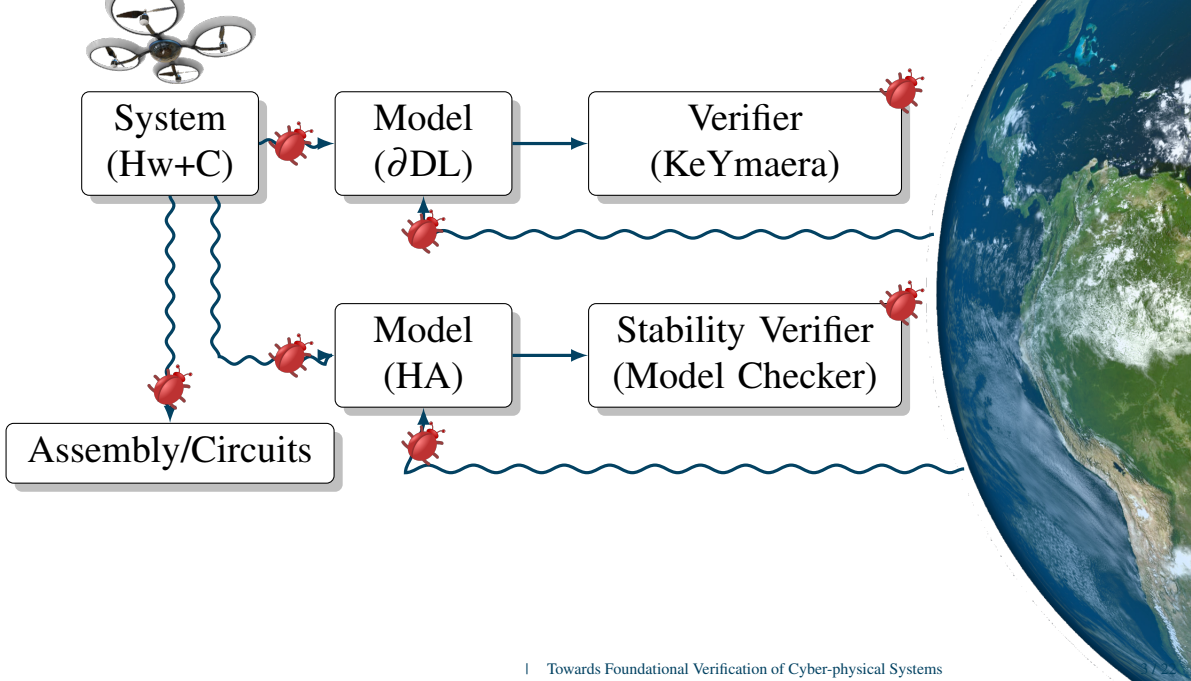


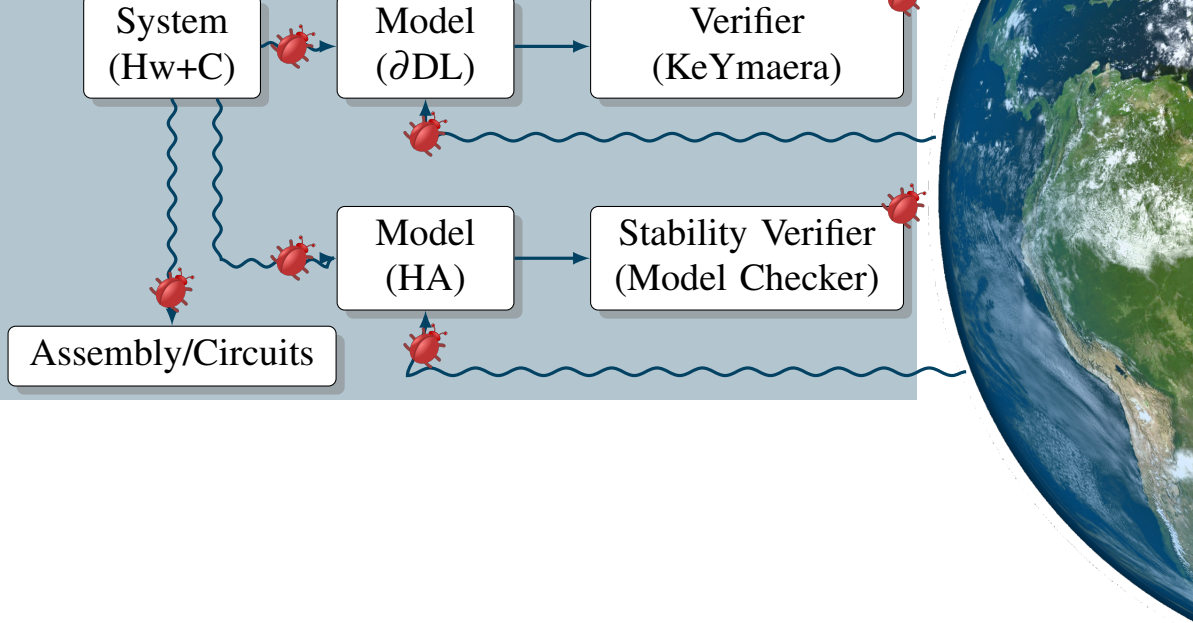












V ERIDRONE



System
(Hw+C)



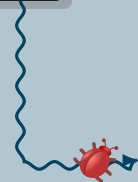
Model
(∂ DL)



Verifier
(KeYmaera)



Assembly/Circuits



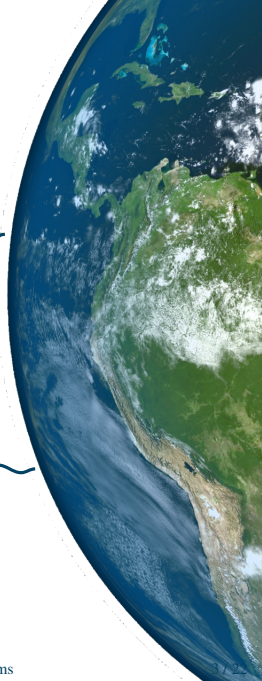
Model
(HA)



Stability Verifier
(Model Checker)



CoQ



Benefits of Foundational Verification [YCER11]

Compiler	Bugs
GCC	122
Clang/LLVM	181
CompCert	

Benefits of Foundational Verification [YCER11]

Compiler	Bugs
GCC	122
Chromium	181
Verified in Coq	0 [†]
CompCert	

[†] In verified code

Benefits of Foundational Verification [YCER11]

Compiler	Bugs
GCC	122
Clang/LLVM	181
CompCert	0 [†]

Verified in Coq

Foundational verification

- Strong guarantees, **and**
- Expressive logic

Verification in Coq



Foundational verification

- Strong guarantees, **and**
- Expressive logic

Verification in Coq

Definitions

```
Def opt (c : c_prog) : c_prog :=  
  ... C ...
```

Foundational verification

- Strong guarantees, **and**
- Expressive logic

Verification in Coq

```
Def opt (c : c_prog) : c_prog :=  
  ... c ...
```

```
Thm opt_sound : forall (c : c_prog),  
  c ~ opt c.
```

Definitions

Theorems



Foundational verification

- Strong guarantees, **and**
- Expressive logic

Verification in Coq

```
Def opt (c : c_prog) : c_prog :=  
  ... C ...
```

```
Thm opt_sound :  $\forall c : c\_prog,$   
   $c \sim \text{opt } c.$ 
```

Proof.

```
induction (* proof for each case *)
```

Definitions

Theorems

Foundational verification

- Strong guarantees, and
- Expressive logic



Verification in Coq

Definitions

```
Def opt (c : c_prog) : c_prog :=  
  ... c ...
```

Theorems

```
Thm opt_sound :  $\forall c : c\_prog,$   
   $c \sim opt\ c.$ 
```

Proof.

```
induction c  
  (* proof for each case *)
```

Qed.

Interactive proof scripts

Foundational verification

- Strong guarantees, **and**
- Expressive logic

(HA)

(Model Checker)

Assembly/Circuits



Coq



(HA)

(Model Checker)

Assembly/Circuits



Coq



Outline

CPS Logic

Coq

Outline

∂ -Ind

CPS Logic

Discrete + Cont

Coq

Case Study: Runtime Monitors [Dan15]

Case Study: Runtime Monitors [Dan15]

Sensors (v, y)

Case Study: Runtime Monitors [Dan15]

Sensors (v, y)

Case Study: VERIDRONE & Runtime Monitors [Dan15]

Unverified Code

Sensors

Controller

World

Sensors (v, y)







Case Study: Runtime Monitors [Dan15]

Case Study: Runtime Monitors [Dan15]

“Sampled-Data System”

Case Study: Runtime Monitors [Dan15]

“Always”

Case Study: Runtime Monitors [Dan15]

Initial condition

$\text{Init} \wedge \Box \text{Sys}_\Delta$ (Monitor , Wend) $\models \Box P$

Case Study: Runtime Monitors [Dan15]

Boundness, stability,
robustness, etc.

A Flavor of Customizable Verification

$$\frac{}{P \wedge \Box \text{Sys}_\Delta(D, W) \vdash \Box P} \text{SYS-IND}$$

Custom proof rules

A Flavor of Customization

' = after the transition

(Discrete)

$$\frac{P \wedge D \wedge 0 \leq \tau' \leq \Delta \vdash P'}{P \wedge \Box \text{Sys}_\Delta(D, W) \vdash \Box P} \text{SYS-IND}$$

A Flavor of Customizable Verification

(Discrete)

$$\frac{P \wedge D \wedge 0 \leq \tau' \leq \Delta \vdash P'}{P \wedge \Box \text{Sys}_{\Delta}(D, W) \vdash \Box P} \text{SYS-IND}$$

A Flavor of Customizable Verification

(Continuous) $P \wedge \text{Cont}(W \wedge \dot{\tau} = -1) \wedge 0 \leq \tau' \vdash P'$

(Discrete)
$$\frac{P \wedge D \wedge 0 \leq \tau' \leq \Delta \vdash P'}{P \wedge \Box \text{Sys}_{\Delta}(D, W) \vdash \Box P} \text{SYS-IND}$$

A Flavor of Customizable Verification

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A Flavor of Customizable Verification

$$\frac{\begin{array}{l} P \wedge \text{Cont}(W \wedge \dot{\tau} = -1) \wedge 0 \leq \tau' \vdash P' \\ P \wedge D \wedge 0 \leq \tau' \leq \Delta \vdash P' \end{array}}{P \wedge \Box \text{Sys}_{\Delta}(D, W) \vdash \Box P} \text{SYS-IND}$$

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A Flavor of Customizable Verification

$$P \wedge \text{Cont}(W \wedge \dot{\tau} = -1) \wedge 0 \leq \tau' \vdash P'$$

$$P \wedge D \wedge 0 \leq \tau' \leq \Delta \vdash P'$$

$$\frac{P \wedge D \wedge 0 \leq \tau' \leq \Delta \vdash P'}{P \wedge \Box \text{Sys}_{\Delta}(D, W) \vdash \Box P} \text{SYS-IND}$$

Thm

Def

Aside: Simplified Quadcopter Dynamics

Linear dynamics

World dynamics

$$W_{QC} \triangleq \left(\begin{array}{lcl} C_{\theta\phi} & \rightarrow & \dot{\mathbf{x}} = \mathbf{v}_x \wedge \dot{\mathbf{y}} = \mathbf{v}_y \wedge \dot{\mathbf{z}} = \mathbf{v}_z \\ & \wedge & \dot{\mathbf{v}}_x = \mathbf{T} \cos \phi \sin \theta \\ & \wedge & \dot{\mathbf{v}}_y = -\mathbf{T} \sin \phi \sin \theta \\ & \wedge & \dot{\mathbf{v}}_z = \mathbf{T} \cos \phi \cos \theta - g \\ & \wedge & \dot{\phi} = 0 \wedge \dot{\theta} = 0 \wedge \dot{\mathbf{T}} = 0 \end{array} \right)$$

Angular thrus

Pitch

Roll

$$C_{\theta\phi} \triangleq |\theta| \leq 30^\circ \wedge |\phi| \leq 30^\circ \wedge 0 \leq \mathbf{T}$$

Simple dynamics if “mostly level” ($< 30^\circ$)

Aside: Simplified Quadcopter Dynamics

$$W_{\text{acc}} \triangleq \left(\begin{array}{l} C_{\theta\phi} \rightarrow \dot{\mathbf{x}} = \mathbf{v}_x \wedge \dot{\mathbf{y}} = \mathbf{v}_y \wedge \dot{\mathbf{z}} = \mathbf{v}_z \\ \wedge \dot{\mathbf{v}}_x = \mathbf{T} \cos \phi \sin \theta \\ \wedge \dot{\mathbf{v}}_y = -\mathbf{T} \sin \phi \\ \wedge \dot{\mathbf{v}}_z = \mathbf{T} \cos \phi \cos \theta \\ \wedge \dot{\phi} = 0 \wedge \dot{\theta} = 0 \wedge \dot{\mathbf{T}} = 0 \end{array} \right)$$

“Small-angle constraint”

Instantaneous change

$$C_{\theta\phi} \triangleq |\theta| \leq 30^\circ \wedge |\phi| \leq 30^\circ \wedge 0 \leq \mathbf{T}$$

[†] Approximation justified by fast angular dynamics and the small angle constraint [GHH⁺11].

Outline

∂ -Ind

CPS Logic

Discrete + Cont

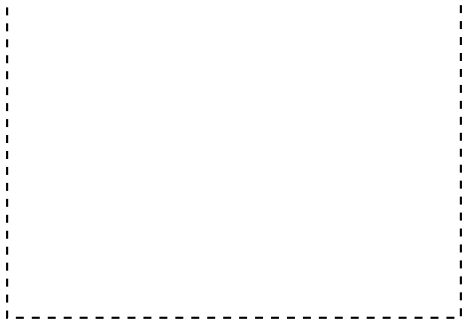
Coq

Outline

CPS Logic

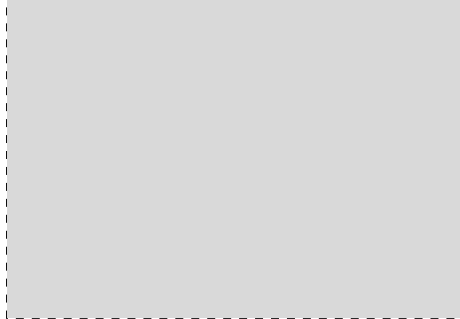
Coq

Composing Monitors [Dan16a]



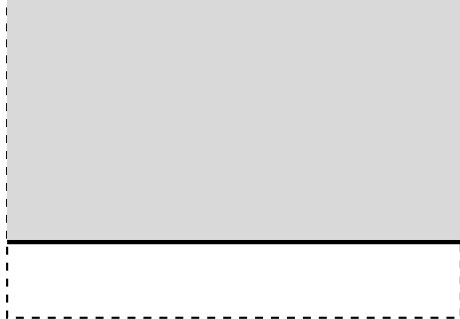
Composing Monitors [Dan16a]

- Spatial transformation



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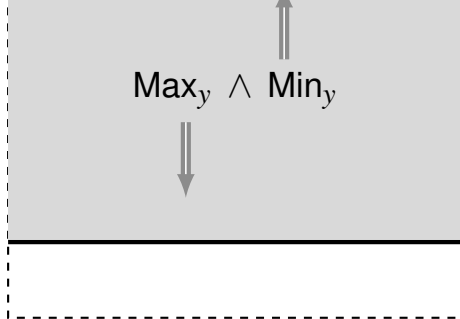
Composing Monitors [Dan16a]

A diagram of a monitor box, consisting of a light gray upper rectangle and a white lower rectangle, both enclosed within a dashed black border. A solid black horizontal line separates the two rectangles. The text $\text{Max}_y \wedge \text{Min}_y$ is centered in the gray upper section.
$$\text{Max}_y \wedge \text{Min}_y$$

- Spatial transformation
- Conjunctive composition

Composing Monitors [Dan16a]

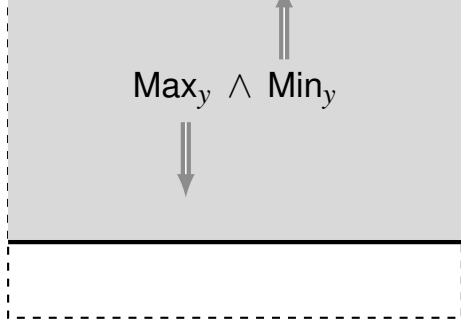
- Spatial transformation
- Conjunctive composition



Composing Monitors [Dan16a]

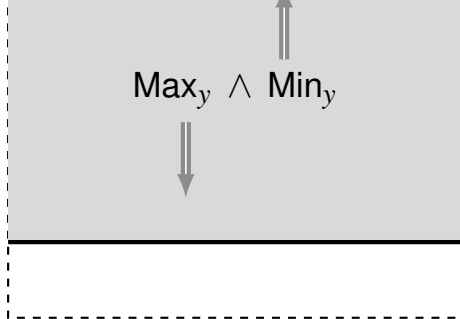
Monitors might not be compatible!

- Spatial transformation
- Conjunctive composition



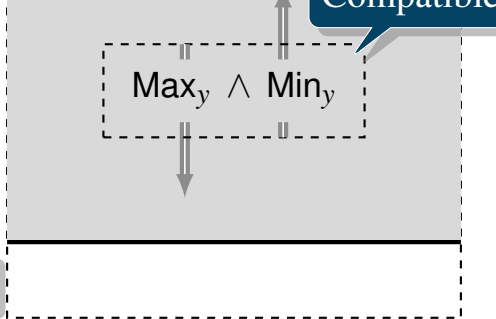
Composing Monitors [Dan16a]

- Spatial transformation
- Conjunctive composition



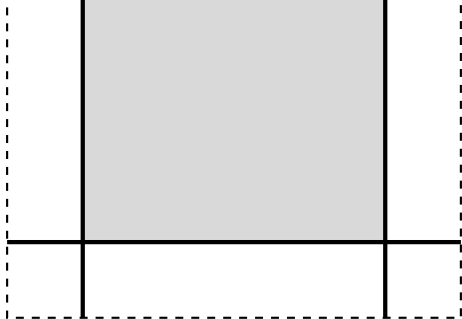
Composing Monitors [Dan16a]

- Spatial **Formalize and build proof rules**
- Conjunctive composition



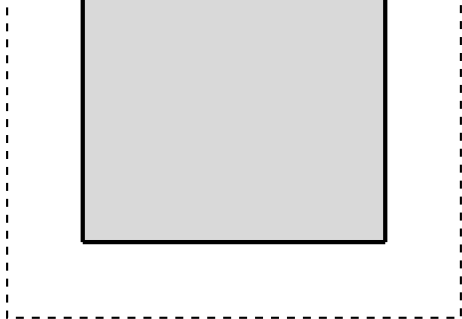
Composing Monitors [Dan16a]

- Spatial transformation
- Conjunctive composition



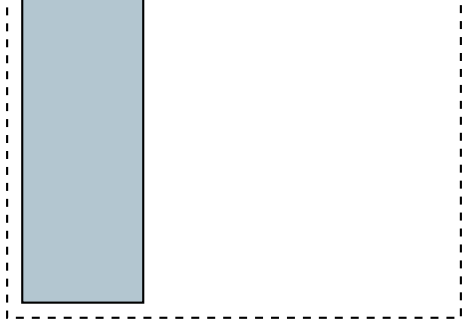
Composing Monitors [Dan16a]

- Spatial transformation
- Conjunctive composition



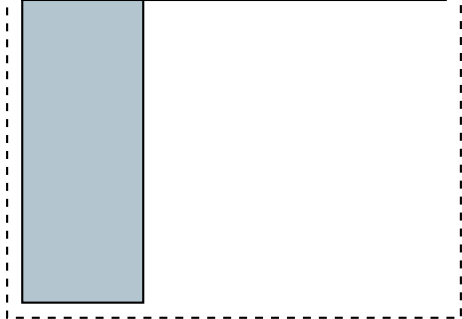
Composing Monitors [Dan16a]

- Spatial transformation
- Conjunctive composition
- Disjunctive composition



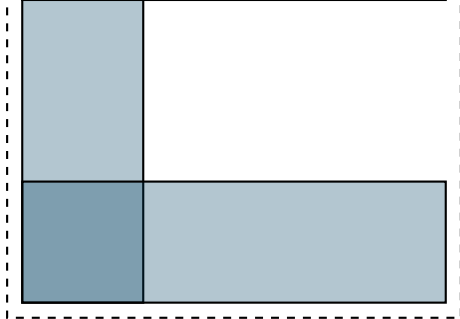
Composing Monitors [Dan16a]

- Spatial transformation
- Conjunctive composition
- Disjunctive composition



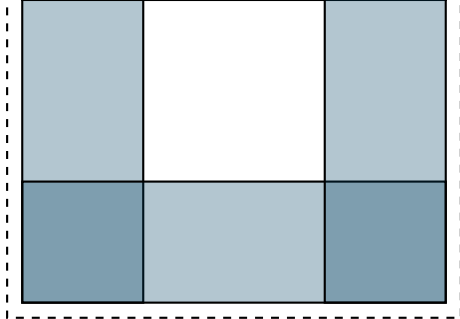
Composing Monitors [Dan16a]

- Spatial transformation
- Conjunctive composition
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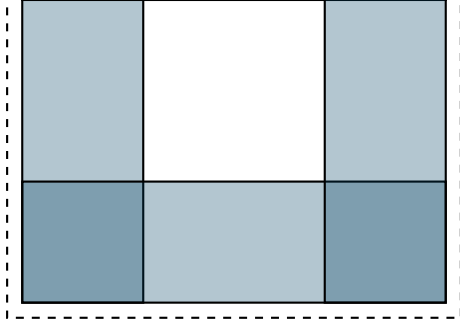
Composing Monitors [Dan16a]

- Spatial transformation
- Conjunctive composition
- Disjunctive composition



Composing Monitors [Dan16a]

- Spatial transformation
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- Disjunctive composition



Composing Monitors [Dan16a]

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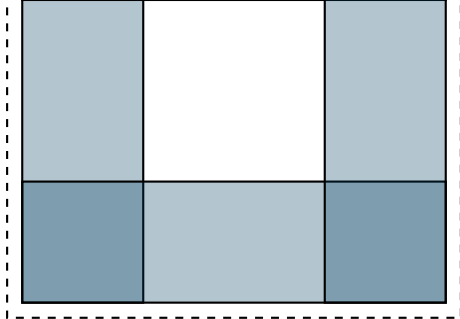
Transition region

Composing Monitors [Dan16a]

Thm

Thm

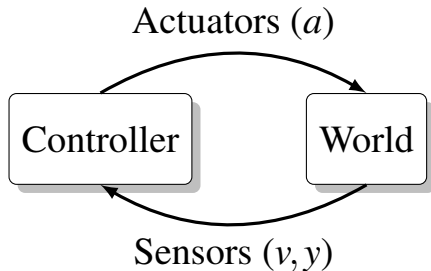
- Spatial transformation
- Conjunctive composition
- Disjunctive composition



Extending the Model: Robustness [Dan16b]

Tolerance to disturbances

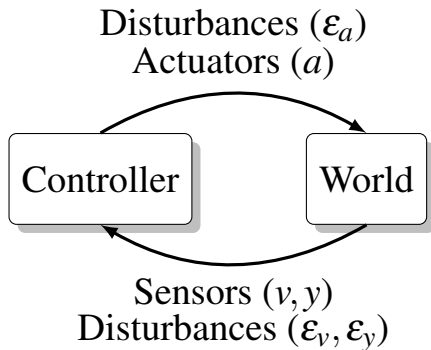
- Add disturbances
- Bound violations



Extending the Model: Robustness [Dan16b]

Tolerance to disturbances

- Add disturbances
- Bound violations

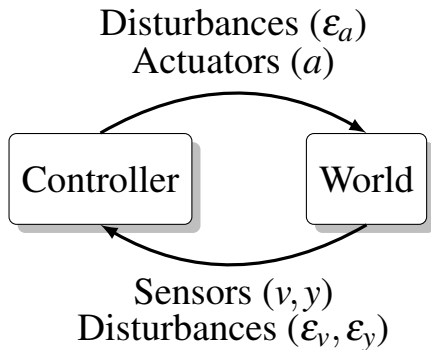


Extending the Model: Robustness [Dan16b]

Def

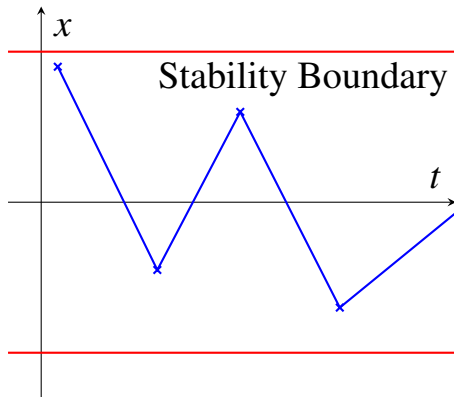
Tolerance to disturbances

- Add disturbances
- Bound violations



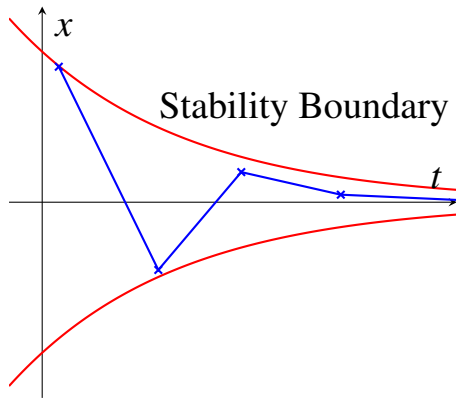
More Properties: Stability [Mat16]

- Boundedness over time



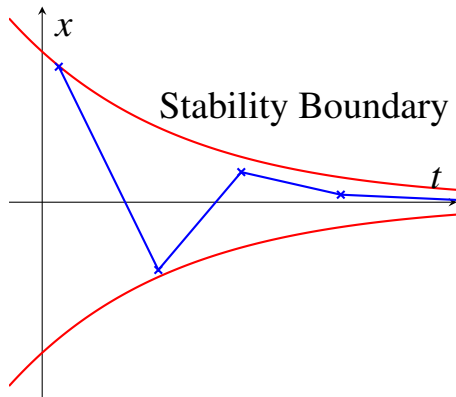
More Properties: Stability [Mat16]

- Boundedness over time
- Convergence to a goal



More Properties: Stability [Mat16]

- Boundedness over time
- Convergence to a goal
- 100 years of control theory!

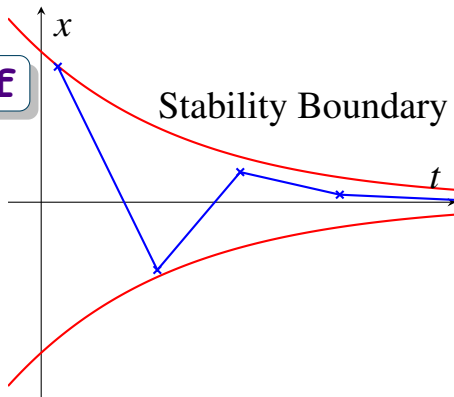


More Properties: Stability [Mat16]

Def

Def

- Boundedness over time
 - Convergence to a goal
- Thm Thm
- 100 years of control theory!



Outline

CPS Logic

Coq

Outline

∂ -Ind

Proof Rules

Checkers

CPS Logic

Coq

External Automation

Connect to other tools

External Automation

Connect to other tools

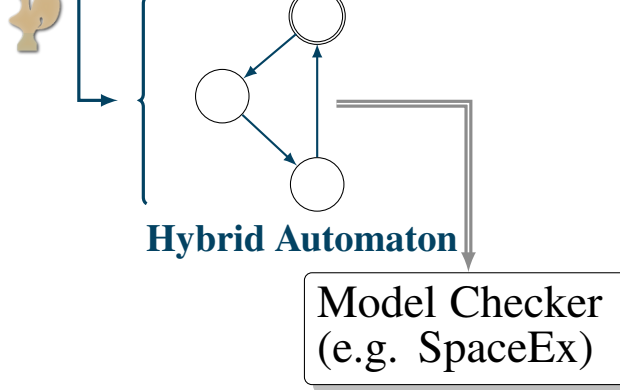
- Formalize in Coq

Model Checker
(e.g. SpaceEx)

External Automation

Connect to other tools

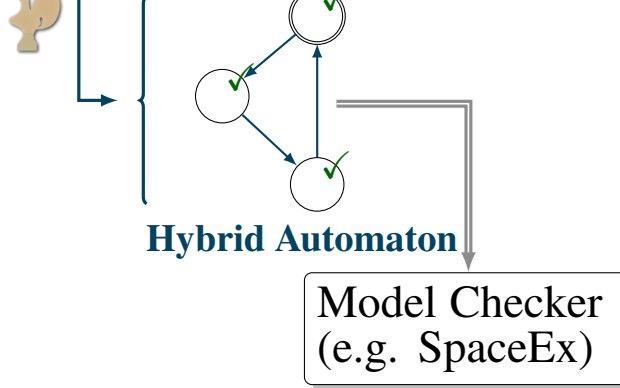
- Formalize in Coq
- Leverage automation



External Automation

Connect to other tools

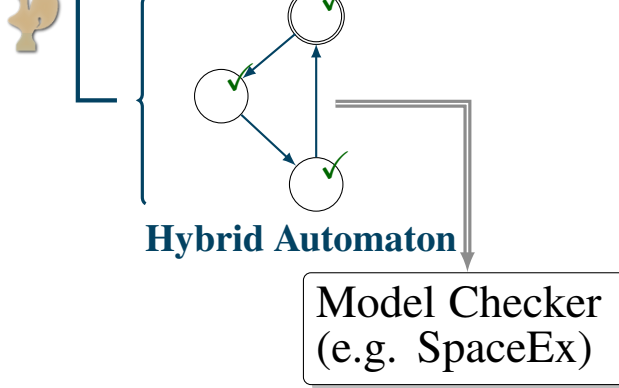
- Formalize in Coq
- Leverage automation



External Automation

Connect to other tools

- Formalize in Coq
- Leverage automation



External Automation

Connect to other tools

- Formalize in Coq
- Leverage automation

Trust or Verify?

Hybrid Automaton

**Model Checker
(e.g. SpaceEx)**

External Automation

Connect to other tools

- Formalize in Coq
- Leverage automation
- Combine with other reasoning

Thm?

Trust or Verify?

Hybrid Automaton

**Model Checker
(e.g. SpaceEx)**

Outline

∂ -Ind

Proof Rules

Checkers

CPS Logic

Coq

Outline

CPS Logic

Component

Hardware

Coq

End-to-End Guarantees

$$\text{Sys}_\Delta (\quad D, \quad W) \vdash P$$

Connect models & code

- Existing developments

End. Floating point & execution

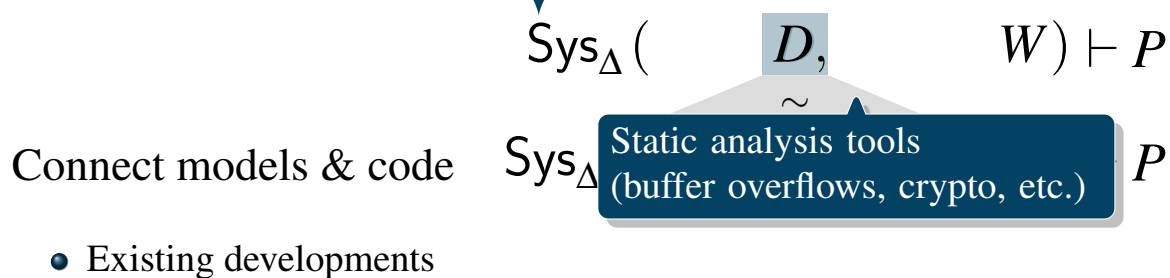
$$\text{Sys}_\Delta (\boxed{D}, W) \vdash P$$

Connect models & code

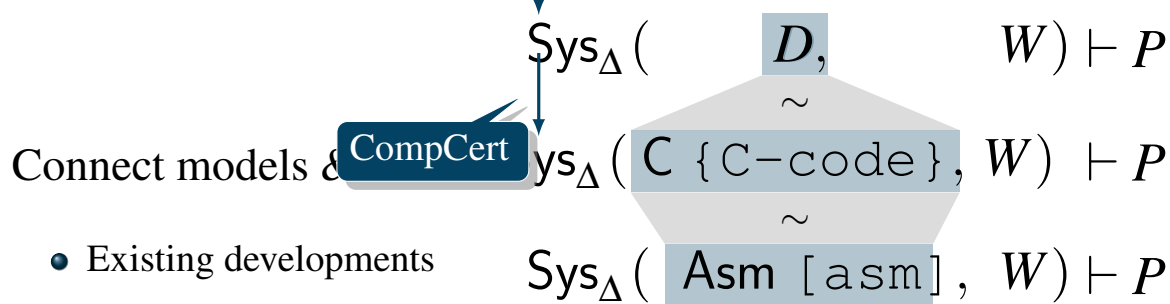
$$\text{Sys}_\Delta (\boxed{C \{ \text{C-code} \}}, W) \vdash P$$

- Existing developments

End-to-End Guarantees



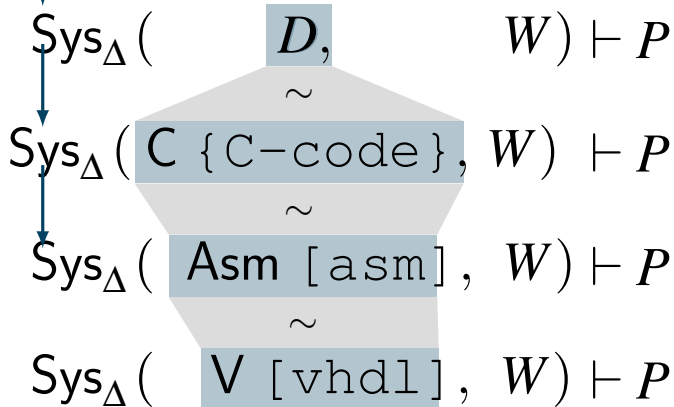
End-to-End Guarantees



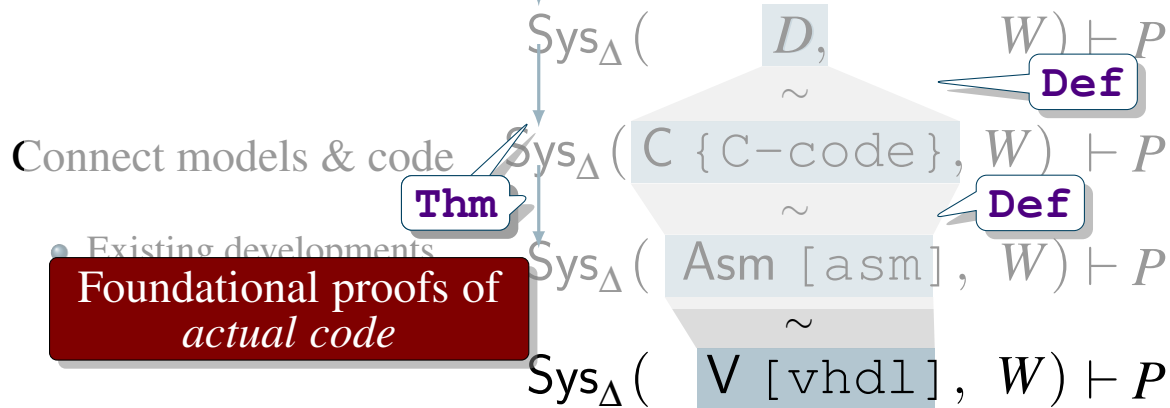
End-to-End Guarantees

Connect models & code

- Existing developments



End-to-End Guarantees



Outline

CPS Logic

Component

Hardware

Coq

Outline

CPS Logic

Discrete + Cont

Probabilistic

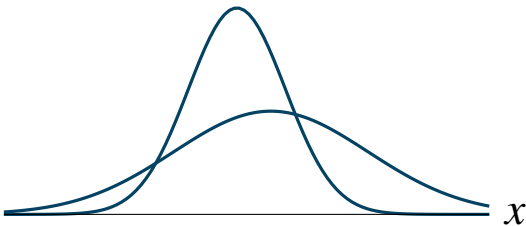
Coq

Uncertainty

gps

barometer

- Probabilistic processes
- Minimize uncertainty



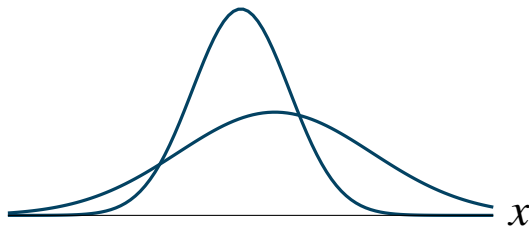
Uncertainty

gps

fused

barometer

- Probabilistic processes
- Minimize uncertainty

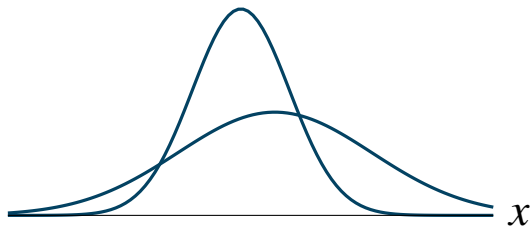


Uncertainty

Safe (80%)

Unsafe (20%)

- Probabilistic processes
- Minimize uncertainty
- Decide with uncertainty

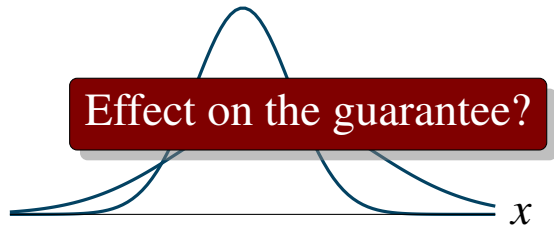


Uncertainty

Safe (80%)

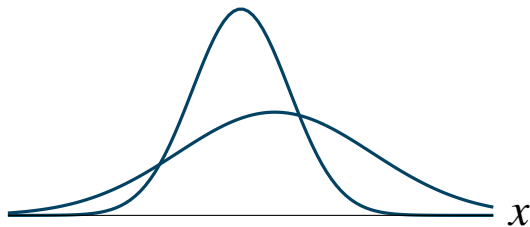
Unsafe (20%)

- Probabilistic processes
- Minimize uncertainty
- Decide with uncertainty



Uncertainty

- Probabilistic processes
- Minimize uncertainty
- Decide with uncertainty



Outline

CPS Logic

Coq

Out



<http://verification.ucsd.edu/>

CPS Logic

Coq

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