#### OUTLINES



# Parallel Programming(CDSC 604)

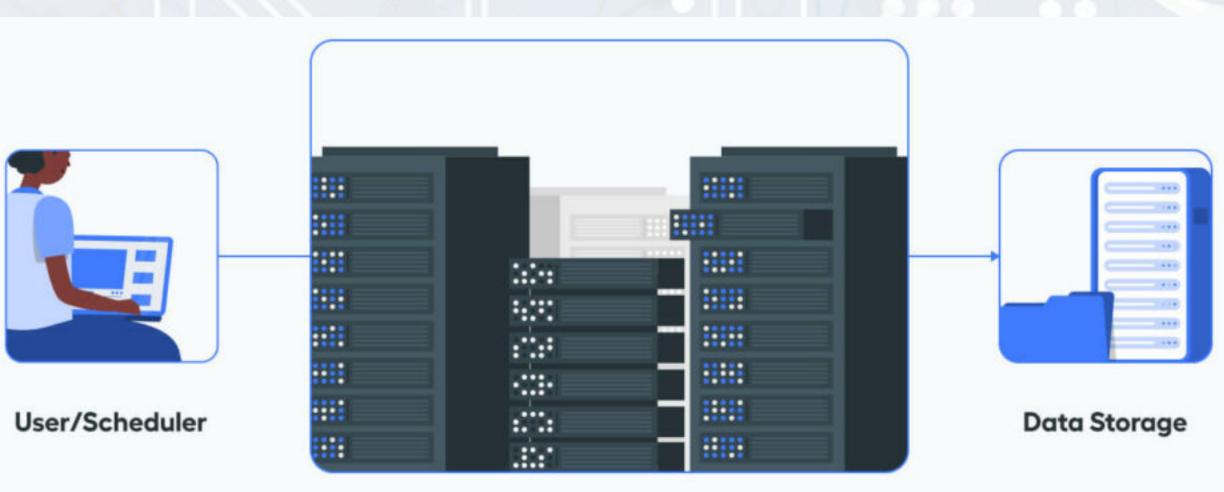
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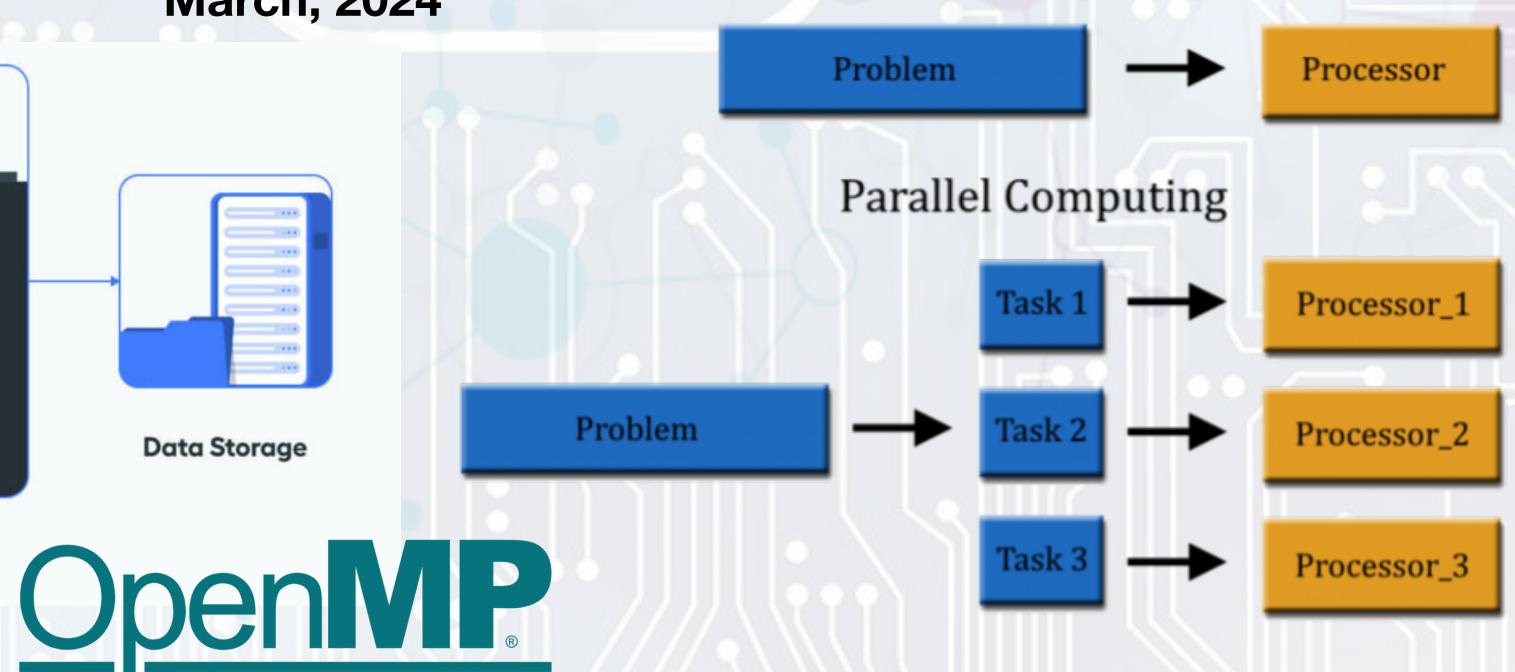






**Compute Cluster** 

#### Serial Computing



### Parallel computing

Part 1: Lecture 3

Cluster computing: Factors to measure performance of parallel programs

- Question:
  - 1. What do you say if Absolute Speedup ( $S_{an}$ ) is less than one?
  - 2. What do you say if Absolute Efficiency  $(E_{an})=1$
  - 3. What do you say if Absolute Efficiency( $E_{an}$ ) > 1
  - 4. What do you observe by looking at Relative Speedup ( $S_{rn}$ ) for different value of N
  - 5. What do you observe by looking at Relative Efficiency ( $E_{\it rn}$ ) for different value of N

Part 1: Lecture3

#### Introduction to MPI

- In this section we will cover:
  - Introduction to MPI
  - Steps in MPI implementation
  - Installing OpenMPI
  - Testing OpenMPI

Part 1: Lecture 3

### Message Passing Programming Paradigm

- Parallel program = program composed of tasks(processes) which communicate to accomplish an overall computational goal
- Each processor in a message-passing program runs a sub-program
  - Written in a conferential sequential language
  - All variables are private
  - Communication via special subroutine calls
- MPI refers primarily the message-passing parallel programming model, in which data is moved from the address space of one process to that of another process through cooperative operations on each process.
- It is a language-independent communications protocol used to program parallel computers.

Part 1: Lecture 3

#### Why Use MPI?

- Message passing now mature as programming paradigm:
  - Well understood
  - Efficient match to hardware
  - Many application
- Full range of desired features:
  - Modularity
  - Access to peak performance
  - Portability
  - Heterogeneity
  - Subgroups
  - Topologies
  - Performance measurement tools

Part 1: Lecture 3

#### Introduction to MPI

- MPI is the first standardized, vendor independent, message passing library.
- MPI closely match the design goals of portability, efficiency, and flexibility
- It become the "industry standard" for writing message passing programs on HPC platforms.
- The first standard for Message passing Interface (MPI-1) is released in 1994 and followed with subsequent revisions.
- MPI Venders that played a significant role in advancing the adoption and optimization of MPI across a wide range of computing platforms:
  - IBM, Intel, TMC, SGI, Seiko, Cray, Convex, Ncube, ....
- Library contributed to the MPI ecosystem:
  - PVM, P4, Zipcode, TCGMSG, Chameleon, Express, Linda, DP(HKU), PM(Japan),
  - AM(Berkeley), FM(HPVM at Illinois)

Part 1: Lecture 3

#### Introduction to MPI

- The message passing model demonstrates the following characteristics:
  - A set of tasks that uses their own local memory during computation.
  - Multiple tasks can reside on the same physical machine as well as across an arbitrary number of machines.
  - Tasks exchange data through communications by sending and receiving messages.
  - Data transfer usually requires cooperative operations to be performed by each process. For example, a send operation must have a matching receive operation.

Part 1: Lecture 3

#### Reason for using MPI

- Standardization MPI is the only message passing library that can be considered as a standard. It is supported on virtually by all HPC platforms
- Portability There is little or no need to modify your source code when you port your application to a different platform that supports (and is compliant with) the MPI standard
- Performance Opportunities Vendor implementations should be able to exploit native hardware features to optimize performance. Any implementation is free to develop optimized algorithms.
- Functionality There are over 430 routines defined in MPI-3, which includes the majority of those in MPI-2 and MPI-1.
- Availability A variety of implementations are available, both vendor and public domain.

#### Part 1: Lecture 3

#### **MPI Implementations**

- From a programming perspective, message passing implementations commonly comprise a library of subroutines that are embedded in the source code.
- The programmer is responsible for determining all parallelism.
- A large number of message passing libraries are available since the 1980s.
- Message Passing usually thought of in the context of distributed memory, parallel computers
- MPI program can be Single program or Multiple program
- However, the same code can be run well on
  - A shared memory parallel computer
  - A network of workstations
  - A collection of heterogeneous (different architecture) processors
  - Even on a single workstation

Part 1: Lecture 3

### Steps in MPI Implementation

- In manual approach of implementing parallel program, there are steps to follow which fairly guide the programmer to successfully accomplish the job.
- · However, there is no single best rule (method) to do manual parallelization.

Part 1: Lecture 3

### Step 1: Understand the Problem and the Program

- Undoubtedly, the first step in developing parallel software is to understand the problem that you wish to solve in parallel.
- If you are starting with a serial program, this necessitates to understand the implementation logic of the existing code (program)





- Monte Carlo Method for π Approximation:
  - 1. Initialize Variables
    - Set the total number of random sample N.
    - Initialize the count of points inside the unit circle inside circle to zero
  - 2. Generate random points:
    - Repeat the following step N times:
      - Generate a random point (x, y) within the unit square  $[0,1] \times [0,1]$
  - 3. Determine if the points are inside the Unit Circle:
    - For each generated point(x,y):
      - If  $x^2 + y^2 \le 1$  increment inside circle by one
  - **4.** Compute π Approximation:
    - The ratio of points inside the unit circle to the total number of points is proportional to

the area of the unit circle to the area of the unit square:

$$\frac{\text{inside circle}}{N} \approx \frac{\text{Area of Circle}}{\text{Area of Square}} = \frac{\pi \cdot (1^2)}{(1^2)} = \pi$$

• Hence, the approximation of  $\pi$  is given by:

. 
$$\pi pprox \frac{4 \cdot \text{inside circle}}{N}$$

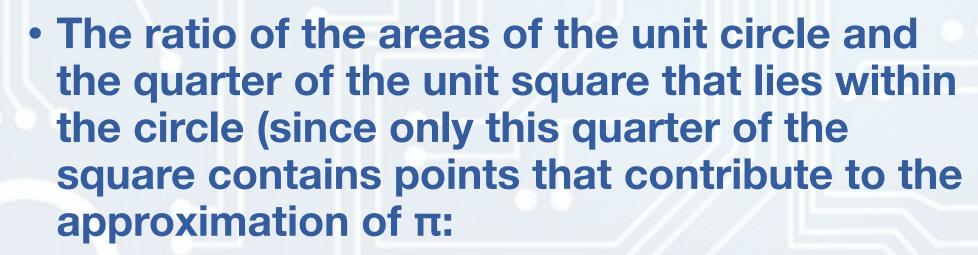
- 5. Output Result:
  - Print or return the approximation of  $\pi$  obtained from step 4.

This problem is solved in parallel!



### Example of Parallelizable Problem: (Approximating PI)

- Monte Carlo Method for  $\pi$  Approximation:
  - 1. Initialize Variables
    - Set the total number of random sample N.
    - Initialize the count of points inside the unit circle inside circle to zero
  - 2. Generate random points:
    - Repeat the following step N times:
      - Generate a random point (x, y) within the unit square where both x and y are between 0 and 1.
  - 3. Determine if the points are inside the Unit Circle:
    - For each generated point(x,y) calculate the distance from origin:
      - If  $x^2 + y^2 \le 1$  increment inside circle by one
  - 4. Compute π Approximation:



Area of Circle 
$$=\frac{\pi}{4}$$
Area of quarter Square

• Hence, the approximation of  $\pi$  is given by:

. 
$$\pi \approx \frac{4 \cdot \text{inside circle}}{N}$$

- 5. Output Result:
  - Print or return the approximation of  $\pi$  obtained from step 4.

This problem is solved in parallel!

#### Part 1: Lecture 3

#### Parallelizable Problem: (Approximating PI) Serial version

```
import random
def approximate pi(num samples):
    points inside circle = 0
        for in range(num samples):
        # Generate random point coordinates within the unit square
        x = random.uniform(0, 1)
        y = random.uniform(0, 1)
        # Check if the point lies within the unit circle
        if x^{**2} + y^{**2} <= 1:
            points inside circle += 1
     # Approximate pi using the ratio of points inside the circle to total points
    pi approx = 4 * (points inside circle / num samples)
    return pi approx
     name == " main
    num samples = 10000000 # Number of random samples
    # Compute the approximation of pi
    pi approx = approximate pi(num samples)
    # Output the result
    print("Approximation of pi:", pi approx)
```

Part 1: Lecture 3

#### Example of a Non-parallelizable Problem:

- Calculation of the Fibonacci series (1,1,2,3,5,8,13,21,...) by use of the formula: F(k + 2) = F(k + 1) + F(k).
  - This is a non-parallelizable problem because the calculation of the Fibonacci sequence as shown above, entail dependent calculations rather than independent ones.
  - The calculation of the k + 2 value uses those of both k + 1 and k.
  - These three terms cannot be calculated independently and therefore, not in parallelizable.

### Part 1: Lecture 3

#### Step 3: Design the parallel Solution

• In order to write programmer directed parallel program we need to identify the program's *hotspots* location, the *bottle necks*, *inhibitors*, and the *algorithm* used.

#### Hot Spots

- Knowing the hot spots includes identifying where most of the real work is being done
- The majority of scientific and technical programs usually accomplish most of their work in a few places)
- Profilers and performance analysis tools can help in identifying the hot spots if we have the equivalent serial codes.
- Once identified we need to focus on parallelizing the hotspots and ignore those sections of the program that account for little CPU usage.

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Part 1: Lecture 3

#### Step 3: Design the parallel Solution

- In order to write programmer directed parallel program we need to identify the program's *hotspots* location, the *bottle necks*, *inhibitors*, and the *algorithm* used.
- Bottlenecks
  - are the areas that are disproportionately slow, or cause parallelizable work to halt or be deferred. (for example, I/O is usually something that slows a program down)
  - it may be possible to restructure the program or use a different algorithm to reduce or eliminate unnecessary slow area
- Inhibitors refers to issues that affect parallelization.
  - Some inhibitors include data dependency as demonstrated by the Fibonacci sequence above.
  - Finally we should design the parallel algorithm (data parallel or instruction parallel or both)

Part 1: Lecture 3

Step 4: Implement the parallel Algorithm

 The next step is to implement the parallel solution using the selected parallel programming model

Part 1: Lecture 3

Step 5: Debug and test the Program

- This is the time that we need to sit and observe logical flaws that will lead into a wrong logical conclusion
- We should carefully see all the logic implemented are correct as per the requirement stated in the algorithm design

Part 1: Lecture 3

Step 6: Submit the job to the parallel program manager

- This enable you to run your program in the HPC environment.
- We should be able to know how to submit parallel jobs into the HPC resources.

Part 1: Lecture 3

Step 7: Analyze the result

This enable you to know how successful you are

Part 1: Lecture 3

#### **Installing OpenMPI**

- OpenMPI is Open source implementation of MPI standards for use in cluster environment
- Got the the openMPI download site(<a href="https://www.open-mpi.org/software/ompi/v5.0/">https://www.open-mpi.org/software/ompi/v5.0/</a>
- Download the stable version for your machine say openmpi-5.0.2.tar.bz2
- Tar the file
  - tar -xjvf openmpi-5.0.2.tar.bz2
- Go to the folder openmpi-5.0.2
  - cd openmpi-5.0.2

#### Part 1: Lecture 3

#### **Installing OpenMPI**

- if you have access to /usr/local (root previlage) you don't need to specify the option —prefix
  - ./configure
  - make
  - make install
- Otherwise
  - Make installation directory if you don't have access to /usr/local
  - ./configure —prefix ~/openmpi
  - make
  - make install

#### Part 1: Lecture 3

### Test Program with OpenMPI in Fortran

- Modify the environment variable PATH and LD\_LIBRARY\_PATH
- For CSHELL users open the file ~/.cshrc and add the following into the file
  - Open file `~/.cshrc` for editing
  - Add the following lines to the file:
    - setenv PATH \${PATH}:BinPath
    - setenv LD\_LIBRARY\_PATH \${LD\_LIBRARY\_PATH}:LibPath
  - Close the file and execute the command
    - Source ~/.cshrc
- For Bash Shell open the file ~/.bashrc and add the following into the file
  - export PATH=\${PATH}: BinPath
  - export LD\_LIBRARY\_PATH=\${LD\_LIBRARY\_PATH}:LibPath
  - Close the file and execute the command
    - source ~/.bashrc

Part 1: Lecture 3

Test Program with OpenMPI in Fortran

```
program sample
  include 'mpif.h'
  integer :: ierr
 call MPI Init (ierr)
  write(*,*) 'Hello world'
  call MPI Finalize (ierr)
end program sample
```

#### Part 1: Lecture 3

### Test Program with OpenMPI in Fortran

- Compile the program as
  - Mpif90 sample.f90 -o sample.x
- Running the program with N number of processors
  - mpiexec -n 4 sample.x
  - mpiexec —use-hwthread-cpus -n 8 ./sample.x
  - Mpirun oversubscribe -np 16 ./sample.x

#### Part 1: Lecture 3

### Parallelizable Problem: (Approximating PI) MPI version

```
from mpi4py import MPI
import random
def approximate pi(num samples):
    # Initialize MPI
   comm = MPI.COMM WORLD
   rank = comm.Get rank()
   size = comm.Get size()
    # Initialize variables to count points inside the unit circle
   local inside circle = 0
    # Perform Monte Carlo simulation
   for in range(num samples // size):
        # Generate random points within the unit square
       x = random.random()
       y = random.random()
       # Check if the point lies inside the unit circle
       if x**2 + y**2 <= 1:
           local inside circle += 1
    # Gather the count of points inside the unit circle from all processes
    total inside circle = comm.reduce(local inside circle, op=MPI.SUM, root=0)
    # Master process calculates the final approximation of pi
    if rank == 0:
       pi_approx = (4 * total_inside_circle) / num_samples
        return pi approx
    else:
        return None
if name == " main ":
    num samples = 1000000 # Number of random samples
    # Compute the approximation of pi
   pi_approx = approximate_pi(num_samples)
    # Output the result
    if pi approx is not None:
       print("Approximation of pi:", pi approx)
```

