

hw5

● Graded

23 Hours, 55 Minutes Late

Student

Giancarlos Marte

Total Points

26 / 30 pts

Question 1

CFG derivations and parse trees

9 / 9 pts

✓ - 0 pts Correct

Question 2

design a CFG

10 / 10 pts

✓ - 0 pts Correct

Question 3

design a PDA

6 / 10 pts

✓ - 1.5 pts PDA still has content left at end

✓ - 1.5 pts Used incorrect cfg->pda conversion

✓ - 1 pt Combined too many transitions (minor infraction)

💬 You never popped the start state, and your implementation of the algorithm is a bit off. Separating into multiple rules would've made this problem easier for you, because with one rule, you only needed a single push and pop. You seem to be repeating yourself with your pushes and pops

Question 4

readme

1 / 1 pt

✓ - 0 pts Correct

Question assigned to the following page: [1](#)

1) Context-free grammars and string derivations

1.1

String 1:

i = 3; if i == 0: print((1, i)) else: print(i + 7)

Derivation of string 1:

step	derivation	rule
1	<STMTS>	start state
2	<STMT>; <STMTS>	<STMTS> → <STMT>; <STMTS>
3	<ID> = <EXPR>; <STMT>	<STMT> → <ID> = <EXPR> <STMTS> → <STMT>
4	i = <NUM>; if <EXPR>: <STMT> else: <STMT>	<ID> → i <EXPR> → <NUM> <STMT> → if <EXPR>: <STMT> else: <STMT>
5	i = 3; if <EXPR> == <EXPR>: print(<EXPR>) else: print(<EXPR>)	<NUM> → 3 <EXPR> → <EXPR> == <EXPR> <STMT> → print(<EXPR>) <STMT> → print(<EXPR>)
6	i = 3; if <ID> == <NUM>: print(<TUP>) else: print(<EXPR> + <EXPR>)	<EXPR> → <NUM> <EXPR> → <ID> <EXPR> → <TUP> <EXPR> → <EXPR> + <EXPR>
7	i = 3; if i == 3: print((<EXPRS>)) else: print(<ID> + <NUM>)	<ID> → i <NUM> → 3 <TUP> → (<EXPRS>) <EXPR> → <NUM> <EXPR> → <ID>
8	i = 3; if i == 3: print((<EXPR>, <EXPRS>)) else: print(i + 7)	<ID> → i <NUM> → 7 <EXPRS> → <EXPR>, <EXPRS>
9	i = 3; if i == 3: print((<NUM>, <EXPR>)) else: print(i + 7)	<EXPRS> → <EXPR> <EXPR> → <NUM>

Question assigned to the following page: [1](#)

10	i = 3; if i == 3: print((1, <ID>)) else: print(i + 7)	<EXPR> → <ID> <NUM> → 1
11	i = 3; if i == 3: print((1, i)) else: print(i + 7)	<ID> → i <NUM> → 1

String 2:

x = 4; y = lambda(z): z + 10; print(y(x))

Derivation of string 2:

step	derivation	rule
1	<STMTS>	start state
2	<STMT>; <STMTS>	<STMTS> → <STMT>; <STMTS>
3	<ID> = <EXPR>; <STMT>; <STMTS>	<STMT> → <ID> = <EXPR> <STMTS> → <STMT>; <STMTS>
4	x = <NUM>; <ID> = <EXPR>; <STMT>	<ID> → x <EXPR> → <NUM> <STMT> → <ID> = <EXPR> <STMTS> → <STMT>
5	x = 4; y = lambda(<IDS>) : <EXPR>; print(<EXPR>)	<NUM> → 4 <ID> → y <EXPR> → lambda(<IDS>) : <EXPR> <STMT> → print(<EXPR>)
6	x = 4; y = lambda(<ID>) : <EXPR> + <EXPR>; print(<ID> (<EXPR>))	<IDS> → <ID> <EXPR> → <EXPR> + <EXPR> <EXPR> → <ID> (<EXPR>)
7	x = 4; y = lambda(z) : <ID> + <NUM>; print(y (<ID>))	<ID> → z <EXPR> → <ID> <EXPR> → <NUM> <ID> → y <EXPR> → <ID>
8	x = 4; y = lambda(z) : z + 10; print(y (x))	<ID> → z <NUM> → 10 <ID> → x

Question assigned to the following page: [1](#)

1.2

Formal description of the grammar PY

A context free grammar for PY is a 4-tuple (V, Σ , R, S), where:

$V = \{ \langle \text{STMTS} \rangle, \langle \text{STMT} \rangle, \langle \text{EXPRS} \rangle, \langle \text{EXPR} \rangle, \langle \text{TUP} \rangle, \langle \text{NUM} \rangle, \langle \text{IDS} \rangle, \langle \text{ID} \rangle \}$

$\Sigma = \{ ;, =, \text{if}, :, \text{else}, \text{print}, (,), \epsilon, ==, +, \text{lambda}, 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, \text{a}, \text{b}, \text{c}, \text{d}, \text{e}, \text{f}, \text{g}, \text{h}, \text{i}, \text{j}, \text{k}, \text{l}, \text{m}, \text{n}, \text{o}, \text{p}, \text{q}, \text{r}, \text{s}, \text{t}, \text{u}, \text{v}, \text{w}, \text{x}, \text{y}, \text{z} \}$

R = RULES, where RULES is a set with the given rules

S = $\langle \text{STMTS} \rangle$

1.3a

String:

if x == 9: print((a, b, c)) else: print(0)

Derivation:

step	derivation	rule
1	$\langle \text{STMTS} \rangle$	start state
2	$\langle \text{STMT} \rangle$	$\langle \text{STMTS} \rangle \rightarrow \langle \text{STMT} \rangle$
3	if $\langle \text{EXPR} \rangle$: $\langle \text{STMT} \rangle$ else: $\langle \text{STMT} \rangle$	$\langle \text{STMT} \rangle \rightarrow \text{if } \langle \text{EXPR} \rangle : \langle \text{STMT} \rangle$ else: $\langle \text{STMT} \rangle$
4	if $\langle \text{EXPR} \rangle == \langle \text{EXPR} \rangle$: print($\langle \text{EXPR} \rangle$) else: print($\langle \text{EXPR} \rangle$)	$\langle \text{EXPR} \rangle \rightarrow \langle \text{EXPR} \rangle == \langle \text{EXPR} \rangle$ $\langle \text{STMT} \rangle \rightarrow \text{print}(\langle \text{EXPR} \rangle)$ $\langle \text{STMT} \rangle \rightarrow \text{print}(\langle \text{EXPR} \rangle)$
5	if $\langle \text{ID} \rangle == \langle \text{NUM} \rangle$: print($\langle \text{TUP} \rangle$) else: print($\langle \text{NUM} \rangle$)	$\langle \text{EXPR} \rangle \rightarrow \langle \text{ID} \rangle$ $\langle \text{EXPR} \rangle \rightarrow \langle \text{NUM} \rangle$ $\langle \text{EXPR} \rangle \rightarrow \langle \text{TUP} \rangle$ $\langle \text{EXPR} \rangle \rightarrow \langle \text{NUM} \rangle$
6	if x == 9: print(($\langle \text{EXPRS} \rangle$)) else: print(0)	$\langle \text{ID} \rangle \rightarrow \text{x}$ $\langle \text{NUM} \rangle \rightarrow 9$ $\langle \text{TUP} \rangle \rightarrow (\langle \text{EXPRS} \rangle)$ $\langle \text{NUM} \rangle \rightarrow 0$
7	if x == 9: print(($\langle \text{EXPR} \rangle, \langle \text{EXPRS} \rangle$)) else: print(0)	$\langle \text{EXPRS} \rangle \rightarrow \langle \text{EXPR} \rangle, \langle \text{EXPRS} \rangle$
8	if x == 9: print(($\langle \text{ID} \rangle, \langle \text{EXPR} \rangle, \langle \text{EXPRS} \rangle$)) else: print(0)	$\langle \text{EXPR} \rangle \rightarrow \langle \text{ID} \rangle$ $\langle \text{EXPRS} \rangle \rightarrow \langle \text{EXPR} \rangle, \langle \text{EXPRS} \rangle$

Question assigned to the following page: [1](#)

Question assigned to the following page: [1](#)

1.3b.

String: a = 1; f = lambda(x): x + a; f(5)

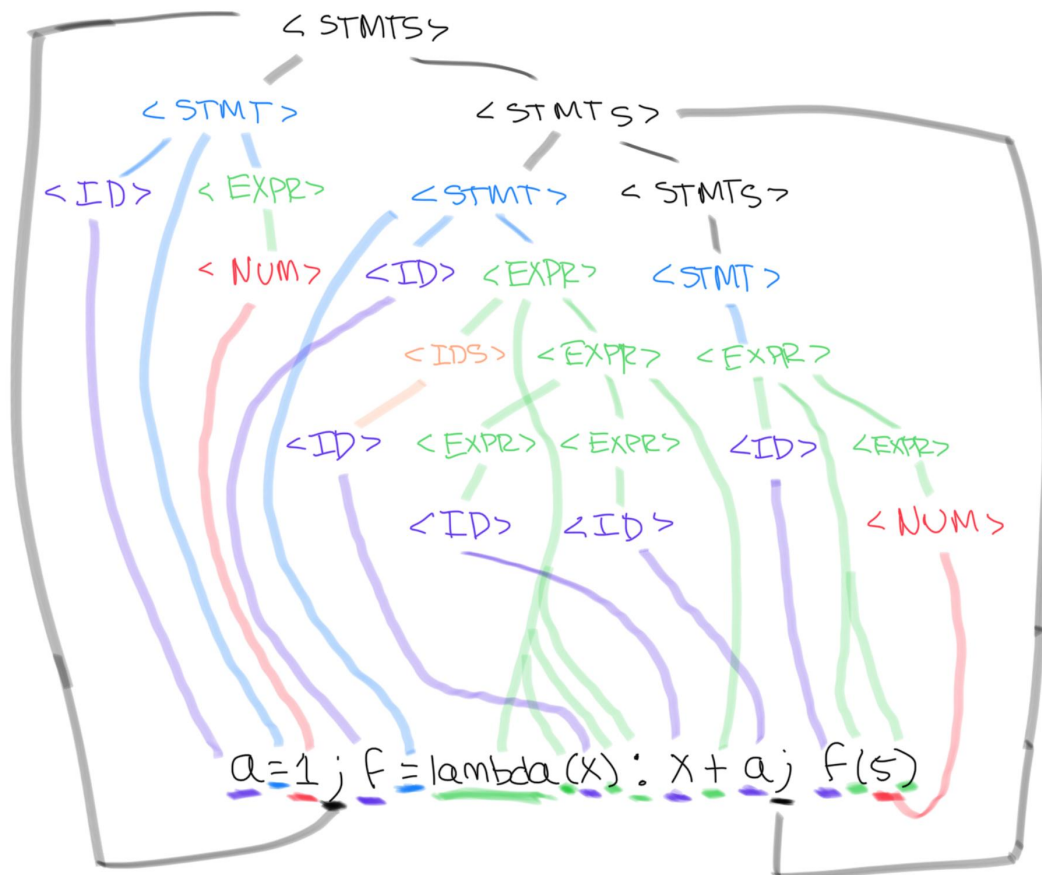
Derivation:

step	derivation	rule
1	<STMTS>	start state
2	<STMT>; <STMTS>	<STMTS> → <STMT>; <STMTS>
3	<ID> = <EXPR>; <STMT>; <STMTS>	<STMT> → <ID> = <EXPR> <STMTS> → <STMT>; <STMTS>
4	a = <NUM>; <ID> = <EXPR>; <STMT>	<ID> → a <EXPR> → <NUM> <STMT> → <ID> = <EXPR> <STMTS> → <STMT>
5	a = 1; f = lambda(<IDS>) : <EXPR>; <EXPR>	<NUM> → 1 <ID> → f <EXPR> → lambda(<IDS>) : <EXPR> <STMT> → <EXPR>
6	a = 1; f = lambda(<ID>) : <EXPR> + <EXPR>; <ID> (<EXPR>)	<IDS> → <ID> <EXPR> → <EXPR> + <EXPR> <EXPR> → <ID> (<EXPR>)
7	y = 1; f = lambda(x) : <ID> + <ID>; f (<NUM>)	<ID> → x <EXPR> → <ID> <EXPR> → <ID> <ID> → f <EXPR> → <NUM>
8	y = 1; f = lambda(x) : x + a; f (5)	<ID> → x <ID> → a <NUM> → 5

Question assigned to the following page: [1](#)

Parse tree:

3 B)



Question assigned to the following page: [2](#)

2) Design a CFG, including whitespace?

step	statement	justification
1	<p>prove that the following language B is a CFL:</p> $B = \{w \mid \text{if_var} == \text{num} : \text{indent}_1 \text{ num}_1 \text{ else} : \text{indent}_2 \text{ num}_2\}$ <p>B is a language over the alphabet $\Sigma = \{a, \dots, z, 0, 1, \dots, 9, _, =, :\}$</p>	given
2	A language is a CFL if a CFG describes it.	theorem and given
3	A CFG is a 4-tuple (V, Σ, R, S) , where V represents a set of variables, Σ is a set of terminals, R is a set of rules and S is the start variable.	definition of a CFG
4	$V = \{\text{var}, \text{num}, \text{num}_1, \text{num}_2, \text{indent}_1, \text{indent}_2\}$	(3) and (1)
5	$\Sigma = \{a, b, \dots, z, 0, 1, \dots, 9, _, =, :\}$	(1) and (3)
6	$R = \{$ $\text{start} \rightarrow \text{if_var} == \text{num} : \text{indent}_1 \text{ num}_1 \text{ else} : \text{indent}_2 \text{ num}_2,$ $\text{var} \rightarrow a \mid b \mid c \mid d \mid e \mid f \mid g \mid h \mid i \mid j \mid k \mid l \mid m \mid n \mid o \mid p \mid q \mid$ $r \mid s \mid t \mid u \mid v \mid w \mid x \mid y \mid z,$ $\text{num} \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9,$ $\text{num}_1 \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9,$ $\text{num}_2 \rightarrow 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9,$ $\text{indent}_1 \rightarrow _ _ *,$ $\text{indent}_2 \rightarrow \text{indent}_1$ $\}$	(4), (5) and (3)
7	$S = \text{start}$	(3) and (6)
8	A CFG (V, Σ, R, S) describes the language B, therefore it is a CFL.	(4), (5), (6) and (7)

Question assigned to the following page: [3](#)

3) Design of a PDA, including whitespace?

Let:

$X = \text{if_var} == \text{num} : \text{indent}_1 \text{ num}_1 \text{ else} : \text{indent}_2 \text{ num}_2,$

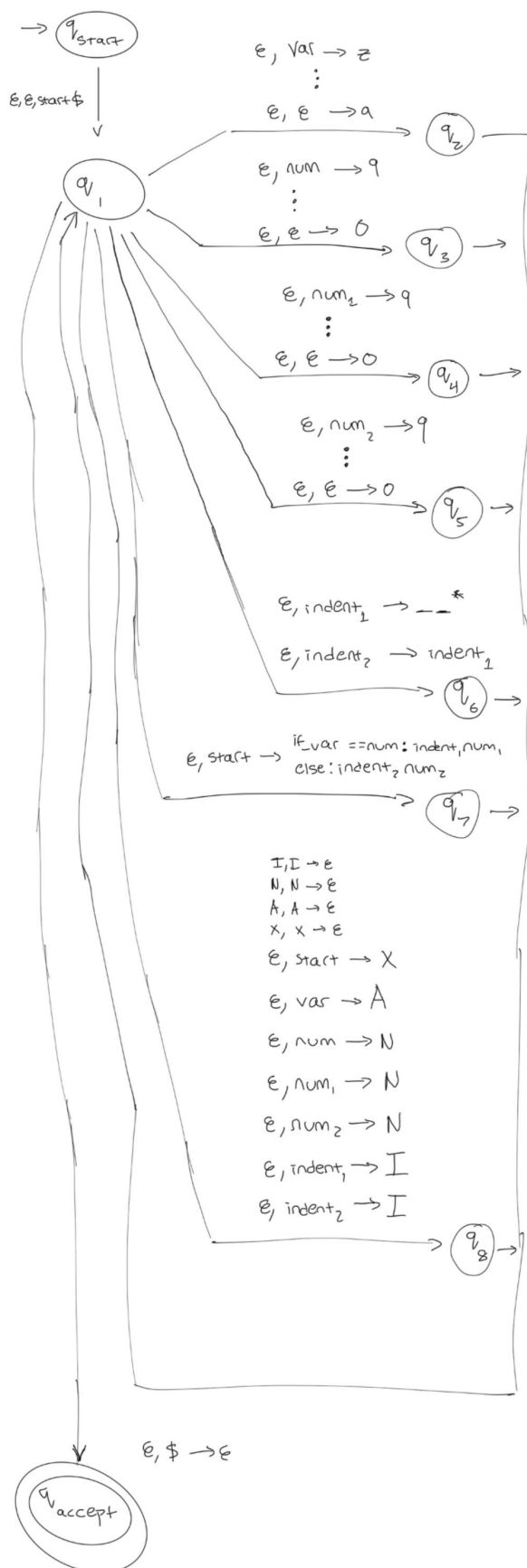
$A = a | b | c | d | e | f | g | h | i | j | k | l | m | n | o | p | q | r | s | t | u | v | w | x | y | z,$

$N = 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9,$

$I = _ _ *$

(graph is on the next page)

Question assigned to the following page: [3](#)



Question assigned to the following page: [4](#)

README

names of others: none

books/websites used: class slides

time spent: 5 hours