

# Comp Arch - HW4 Group Project

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# Report

## Part 4.1 (dllist\_step1.S)

Input your RISC-V code here:

```
54 # This procedure clears 128 bits (16 bytes) of memory
55 # previously allocated to a dll_node
56 ##########
57 FREE_DLLIST:
58
59 # Must clear the memory pointed to by x12 (16 bytes)
60 sw x0, 0(x12)      # Clears bytes 0-3
61 sw x0, 4(x12)      # Clears bytes 4-7
62 sw x0, 8(x12)      # Clears bytes 8-11
63 sw x0, 12(x12)     # Clears bytes 12-15
64
65 jalr x0, x1, 0      # jump and link back to main
66
67 END:
68 nop
```

Reset Step Run CPU: 32 Hz ▾

```
[line 10]: lui x5, 0x7ffff  
[line 11]: addi x18, x5, 0x620  
[line 12]: add x5, x0, x0  
[line 16]: add x12, x18, x0  
[line 17]: jal x1, CALLOC_DLLIST
```

Input your RISC-V code here:

```

54 # This procedure clears 128 bits (16 bytes) of memory
55 # previously allocated to a dll_node
56 ######
57 FREE_DLLIST:
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59 # Must clear the memory pointed to by x12 (16 bytes)
60 sw x0, 0(x12)          # Clears bytes 0-3
61 sw x0, 4(x12)          # Clears bytes 4-7
62 sw x0, 8(x12)          # Clears bytes 8-11
63 sw x0, 12(x12)         # Clears bytes 12-15
64
65 jalr x0, x1, 0          # jump and link back to main
66
67 END:
68 nop

```

Reset Step Run CPU: 32 Hz ▾

```
[line 47]: jalr x0, x1, 0  
[line 18]: add x1, x0, x0  
[line 19]: add x12, x0, x0  
[line 22]: add x12, x18, x0  
[line 23]: jal x1, FREE_DLLIST
```

**Report Deliverable:** Describe how you call the function `calloc` for the `dllist` struct, allocate the memory for the `dllist` struct, free the `sllist` struct main call, and detail a `FREE_DLLIST` procedure. You should show how your code allocates the 128 bits using `calloc`, as well as freeing the struct.

Set x12 to the address of the head pointer(x18) and then jump and link to CALLOC\_DLLIST. In CALLOC\_DLLIST, we stored x0 to 16 sequential bytes of memory starting at wherever x12 was pointing. Freeing the dlist similarly stores x0 to the 16 sequential bytes of memory starting at x12. After freeing the dlist struct main call, we added x0 and x0 and stored it in x18 to clear/zero the pointer. Then we branched to END using beq x0, x0, END.

## Part 4.2 (dllist\_step2.S)

**Report Deliverable:** Describe how you set up the for loop, and how you would insert the first node into the doubly linked list. You should show the steps of how the node is created, and how both head\_ptr and tail\_ptr are pointing to the only node.

Hint: Remember that, in a Doubly Linked List, that if there is only node, then both the head\_ptr and tail\_ptr point to that node. Also, remember the previous and next pointers must initially be NULL (whereas in the Singly Linked List, the next pointer must initially be NULL)

Memory Address		Decimal	Hex	Binary
0x7ffff620	head_ptr	2147481136	0x7ffff630	0b01111111111111111111011000110000
0x7ffff624	tail_ptr	2147481136	0x7ffff630	0b01111111111111111111011000110000
0x7ffff628		0	0x00000000	0b00000000000000000000000000000000
0x7ffff62c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff630	node->value	3	0x00000003	0b00000000000000000000000000000011
0x7ffff634	node->prev_node	0	NULL	0b00000000000000000000000000000000
0x7ffff638	node->next_node	0	NULL	0b00000000000000000000000000000000
0x7ffff63c		0	0x00000000	0b00000000000000000000000000000000

The for loop starts by checking the conditional to see if the block should be executed or skipped. It uses bge to compare x0 and x19. Meaning that if the value in x19 is less than or equal to zero, it will stop running the for loop by branching to "END\_FOR".

To insert the first node, the program determines that the head\_ptr ( $x18 \rightarrow x7$ ) is NULL. This means that the first node has not been inputted and, as such, the code does not take the "ELSE" path. Inside the IF statement, x18 (head\_ptr) is stored in a temporary register and x28 is incremented by 16 bytes. This places the address in x28 at the location where the node->value will be stored for the first node we are about to insert. The following 4 addresses (or 16 bytes) are then cleared so as to account for the value, prev\_node, next\_node, and the segmentation between this current node and the eventual next node. Finally x7 and x28 are cleared, and y is decremented by one. The PC then jumps back to BEGIN\_FOR. Eventually END\_FOR is followed and x6 is cleared as well.

Note that prev\_node and next\_node are intentionally NULL and zero in the photo as there is no node that comes before it and no node that follows it yet.

## Part 4.3 (dllist\_step3.S)

**Report Deliverable:** Describe how you will insert the other nodes, including setting up the pointers to the previous and next nodes. Show examples of what the data memory looks like when as you step through the nodes.

Inserting the other nodes begins after the initial IF statement, which is reserved for inserting the first node. We start by loading the address of the head node from head\_ptr into x28. We then get the next node by loading the location of x28 offset by 8 (2 addresses later of the head node). At this point, we check if the next\_node is null and branch to END WHILE if it is. We now set x29 = curr\_ptr→next\_node by loading an offset of 8(x28), effectively an offset of 8 for x29 since it was just loaded into x28. We then update the tail\_ptr to x28, storing it in the obvious spot of x18(the head\_ptr) offset by 4.

y=2

Memory Address		Decimal	Hex	Binary
0x7fffff620	head_ptr	2147481136	0x7fffff630	0b0111111111111111111111011000110000
0x7fffff624	tail_ptr	2147481152	0x7fffff640	0b0111111111111111111111101100100000
0x7fffff628		0	0x00000000	0b00000000000000000000000000000000
0x7fffff62c		0	0x00000000	0b00000000000000000000000000000000
0x7fffff630	Node → Value	3	0x00000003	0b00000000000000000000000000000001
0x7fffff634	Node → preV_node	0	NULL	0b00000000000000000000000000000000
0x7fffff638	Node → next_node	2147481152	0x7fffff640	0b0111111111111111111111011000100000
0x7fffff63c		0	0x00000000	0b00000000000000000000000000000000
0x7fffff640	Node → Value	2	0x00000002	0b000000000000000000000000000000010
0x7fffff644	Node → preV_node	2147481136	0x7fffff630	0b01111111111111111111111011000110000
0x7fffff648	Node → next_node	0	NULL	0b00000000000000000000000000000000
0x7fffff64c		0	0x00000000	0b00000000000000000000000000000000
0x7fffff650		0	0x00000000	0b00000000000000000000000000000000
0x7fffff654		0	0x00000000	0b00000000000000000000000000000000

y=1

Memory Address		Decimal	Hex	Binary
0x7ffff620	head_ptr	2147481136	0x7ffff630	0b01111111111111111111011000110000
0x7ffff624	tail_ptr	2147481168	0x7ffff650	0b011111111111111111110110001010000
0x7ffff628		0	0x00000000	0b00000000000000000000000000000000
0x7ffff62c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff630	Node->value	3	0x00000003	0b00000000000000000000000000000011
0x7ffff634	Node->prev_node	0	NULL	0b00000000000000000000000000000000
0x7ffff638	Node->next_node	2147481152	0x7ffff640	0b01111111111111111111011001000000
0x7ffff63c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff640	Node->value	2	0x00000002	0b00000000000000000000000000000010
0x7ffff644	Node->prev_node	2147481136	0x7ffff630	0b01111111111111111111011000110000
0x7ffff648	Node->next_node	2147481168	0x7ffff650	0b01111111111111111111011001010000
0x7ffff64c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff650	Node->value	1	0x00000001	0b00000000000000000000000000000001
0x7ffff654	Node->prev_node	2147481152	0x7ffff640	0b01111111111111111111011001000000
0x7ffff658	Node->next_node	0	NULL	0b00000000000000000000000000000000
0x7ffff65c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff660		0	0x00000000	0b00000000000000000000000000000000
0x7ffff664		0	0x00000000	0b00000000000000000000000000000000
0x7ffff668		0	0x00000000	0b00000000000000000000000000000000

y=0

Memory Address		Decimal	Hex	Binary
0x7ffff620	head_ptr	2147481136	0x7ffff630	0b01111111111111111111011000110000
0x7ffff624	tail_ptr	2147481184	0x7ffff660	0b01111111111111111111011000110000
0x7ffff628		0	0x00000000	0b00000000000000000000000000000000
0x7ffff62c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff630	Node->value	3	0x00000003	0b00000000000000000000000000000011
0x7ffff634	Node->prev_node	0	NULL	0b00000000000000000000000000000000
0x7ffff638	Node->next_node	2147481152	0x7ffff640	0b01111111111111111111011001000000
0x7ffff63c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff640	Node->value	2	0x00000002	0b00000000000000000000000000000010
0x7ffff644	Node->prev_node	2147481136	0x7ffff630	0b01111111111111111111011000110000
0x7ffff648	Node->next_node	2147481168	0x7ffff650	0b01111111111111111111011001010000
0x7ffff64c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff650	Node->value	1	0x00000001	0b00000000000000000000000000000001
0x7ffff654	Node->prev_node	2147481152	0x7ffff640	0b01111111111111111111011001000000
0x7ffff658	Node->next_node	2147481184	0x7ffff660	0b01111111111111111111011001100000
0x7ffff65c		0	0x00000000	0b00000000000000000000000000000000
0x7ffff660		0	0x00000000	0b00000000000000000000000000000000
0x7ffff664	Node->value	2147481168	0x7ffff650	0b01111111111111111111011001010000
0x7ffff668	Node->prev_node	0	NULL	0b00000000000000000000000000000000
0x7ffff66c	Node->next_node	0	0x00000000	0b00000000000000000000000000000000

#### Part 4.4 (dllist\_step4.S)

**Report Deliverable:** Describe how you will find and free every `dll_node*`, including seeking and freeing the nodes.

We load the head node and the next subsequent node into `curr_ptr` and `next_ptr`, respectively accomplishing the following C code:

```
sll_node* curr_ptr = the_list→head_ptr;  
sll_node* next_ptr = curr_ptr→next_ptr;
```

We then create a “FREE\_NODE\_LOOP” which will iterate through all the nodes in memory, freeing them as necessary with `sw` instructions. We continue to loop until the `curr_ptr(x7)` is `NULL` (equal to `x0`). We then free the `curr_ptr` (its subsequent 16 bytes). We then update the pointer to `next_ptr` and set the `next_ptr` equal to the `curr_ptr→next_ptr` through `x7` offset by 8 accomplishing the following C code:

```
curr_ptr = next_ptr;  
next_ptr = curr_ptr→next_ptr;
```

We then continue the loop until all nodes have been cleared from memory.

## Part 4.5 (dllist\_final.S)

**Report Deliverable:** Show how you would find and delete one element in the Linked List based on the value. Note how this is called before the free() in main. This means that the find code should delete one node, and then free the other nodes after the call.

We will load word a current pointer(x7) with the head of the list and then run the equivalent of a while loop until we find the pointer that contains the value or the pointer points to null(using beq x7, x0, END\_DELETE\_NODE\_LOOP). If the pointer that contains the value is found then there are a series of if statements depending on if the pointer is the head, somewhere in the middle, the tail, or if the head pointer is equal to the tail pointer.

If the head pointer is equal to the tail pointer, then we set both of them to zero using store word and x0. Else if the pointer is the head, then we use lw and sw to make the head pointer point to the curr→next, and then we make the head pointer's prev pointer's value zero. Else if the pointer is the tail, then the tail pointer becomes the curr\_ptr→prev and then the tail pointer→next value is set to zero using lw's and sw's like in the previous else if. Finally, the else statement is for if the pointer is a middle node. Thus, we use a series of lw's, addi's, and sw's to 'bridge the gap' that the deletion will make. The curr→prev→next is set to curr→next and the curr→next→prev is set to curr→prev. At the end of each of the if and elif statements, there is a beq x0, x0, END\_DELETE\_IF. This clears all of the registers that were used and also deletes the node from memory.

If the value is not found on the current pointer, then curr = curr→next. We did this by getting the curr→next from memory (lw x28, 8(x7)) and then setting x7 equal to it with add x7, x28, x0.

See subsequent pages for images.



Memory Address		Decimal	Hex	Binary	delete_val = 0
0x7ffff620	head_ptr	2147481136	0x7ffff630	0b0111111111111111011000110000	
0x7ffff624	tail_ptr	2147481168	0x7ffff650	0b0111111111111111011001010000	
0x7ffff628		0	0x00000000	0b00000000000000000000000000000000	
0x7ffff62c		0	0x00000000	0b00000000000000000000000000000000	
0x7ffff630	Node->value	3	0x00000003	0b00000000000000000000000000000011	
0x7ffff634	Node->prev_node	0	0x00000000	0b00000000000000000000000000000000	
0x7ffff638	Node->next_node	2147481152	0x7ffff640	0b01111111111111110110010000000000	
0x7ffff63c		0	0x00000000	0b00000000000000000000000000000000	
0x7ffff640	Node->value	2	0x00000002	0b00000000000000000000000000000010	
0x7ffff644	Node->prev_node	2147481136	0x7ffff630	0b01111111111111110110001100000000	
0x7ffff648	Node->next_node	2147481168	0x7ffff650	0b01111111111111110110010100000000	
0x7ffff64c		0	0x00000000	0b00000000000000000000000000000000	
0x7ffff650	Node->value	1	0x00000001	0b00000000000000000000000000000001	
0x7ffff654	Node->prev_node	2147481152	0x7ffff640	0b01111111111111110110010000000000	
0x7ffff658	Node->next_node	0	0x00000000	0b00000000000000000000000000000000	
0x7ffff65c		0	0x00000000	0b00000000000000000000000000000000	
0x7ffff660		0	0x00000000	0b00000000000000000000000000000000	
0x7ffff664		0	0x00000000	0b00000000000000000000000000000000	
0x7ffff668		0	0x00000000	0b00000000000000000000000000000000	
0x7ffff66c		0	0x00000000	0b00000000000000000000000000000000	