# K<sup>th</sup> One

## **Problem**

You are given a binary array that is  $a[i] \in \{0, 1\}$  of size n. You will be given m queries and point updates on the array.

Queries and updates will be of the type given below:

1 i : change the ith element with its opposite.

2 k: find the kth one (ones are numbered from 0, it is guaranteed that there are enough ones in the array).

#### Constraints

 $1 <= n, m <= 10^5$ 0 <= a[i] <= 1

## **Example Input**

5 7 1 1 0 1 0

2 0

2 2

1 2

2 3

1 0

2 0

## Output

0

3 3 1

#### **Approach**

- 1. In these kind of problems, we descend the segment tree
- 2. When we are standing in any segment, we have the decision ability to either go to the left child of the segment or right child of the segment.

#### Code

```
#include "bits/stdc++.h"
using namespace std;
#define int long long
const int N = 1e5+2, MOD = 1e9+7;
int tree[4*N], a[N];
void build(int node, int st, int en)
    int mid = (st + en)/2;
   build(2*node, st, mid);
    build(2*node+1, mid+1, en);
    tree[node] = tree[2*node] + tree[2*node+1];
    if(st == en) {
        return st;
    if(k<tree[2*node]){</pre>
       return kthOne(2*node, st, mid, k);
        return kthOne(2*node+1, mid+1, en, k-tree[2*node]);
```

```
void update(int node, int st, int en, int idx){
       if(a[st] == 1){
           a[st] = 0;
          tree[node] = a[st];
          a[st] = 1;
          tree[node] = a[st];
   int mid = (st+en)/2;
   if(idx <= mid) {</pre>
       update(2*node, st, mid, idx);
       update(2*node+1, mid+1, en, idx);
    tree[node] = tree[2*node] + tree[2*node+1];
   int n,m;
      cin >> a[i];
   build(1,0,n-1);
   while (m--) {
       int type;
       cin >> type;
```

```
if(type == 1) {
    int idx;
    cin >> idx;
    update(1,0,n-1,idx);
}
else if(type == 2) {
    int k;
    cin >> k;

    int ans = kthOne(1,0,n-1,k);
    cout << ans << endl;
}
return 0;
}</pre>
```