

Inversions

Problem

Given an array of size n . We need to count the number of indices j for each index i ($0 \leq i < n$) following the condition

$$\begin{aligned} j &< i \\ a[j] &> a[i] \end{aligned}$$

Note: Above condition is known as the property of inversion.

Constraints

$$1 \leq n \leq 10^5$$

Example input

```
5
4 1 3 5 2
```

Output

```
0 1 1 0 3
```

Brute force approach

Pseudo Code

```
for(int i=0; i<n; i++){
    for(int j=i-1; j>=0; j--){
        if(a[j] > a[i]){
            ans++;
        }
    }
}
```

Optimized Approach (Present Sir approach)

1. After input of each element x , mark that element on the number line as “Present sir” by marking 1 at the index x .
2. Create a segment tree of sum.
3. For each x , just count the number of 1's from $x+1$ till n .

Code

```
#include <bits/stdc++.h>
using namespace std;
#define int long long
const int N = 1e5+2, MOD = 1e9+7;

int tree[4*N], a[N];

int query(int node, int st, int en, int l, int r){
    if(st>r || en<l)
        return 0;

    if(l<=st && en<=r)
        return tree[node];

    int mid = (st + en)/2;
    int q1 = query(2*node, st, mid, l, r);
    int q2 = query(2*node+1, mid+1, en, l, r);

    return q1 + q2;;
}

void update(int node, int st, int en, int idx, int val){
    if(st == en){
        tree[node] = val;
        return;
    }

    int mid = (st+en)/2;
    if(idx <= mid){
        update(2*node, st, mid, idx, val);
    }
}
```

```
else
{
    update(2*node+1, mid+1, en, idx, val);
}

tree[node] = tree[2*node] + tree[2*node+1];
}

signed main()
{
    for(int i=0; i<4*N; i++){
        tree[i] = 0;
    }

    int n;
    cin >> n;
    int x;
    for(int i=0; i<n; i++){
        cin >> x;

        int ans = query(1,1,n,x,n);
        cout << ans <<" ";
        update(1,1,n,x,1);
    }
    return 0;
}
```