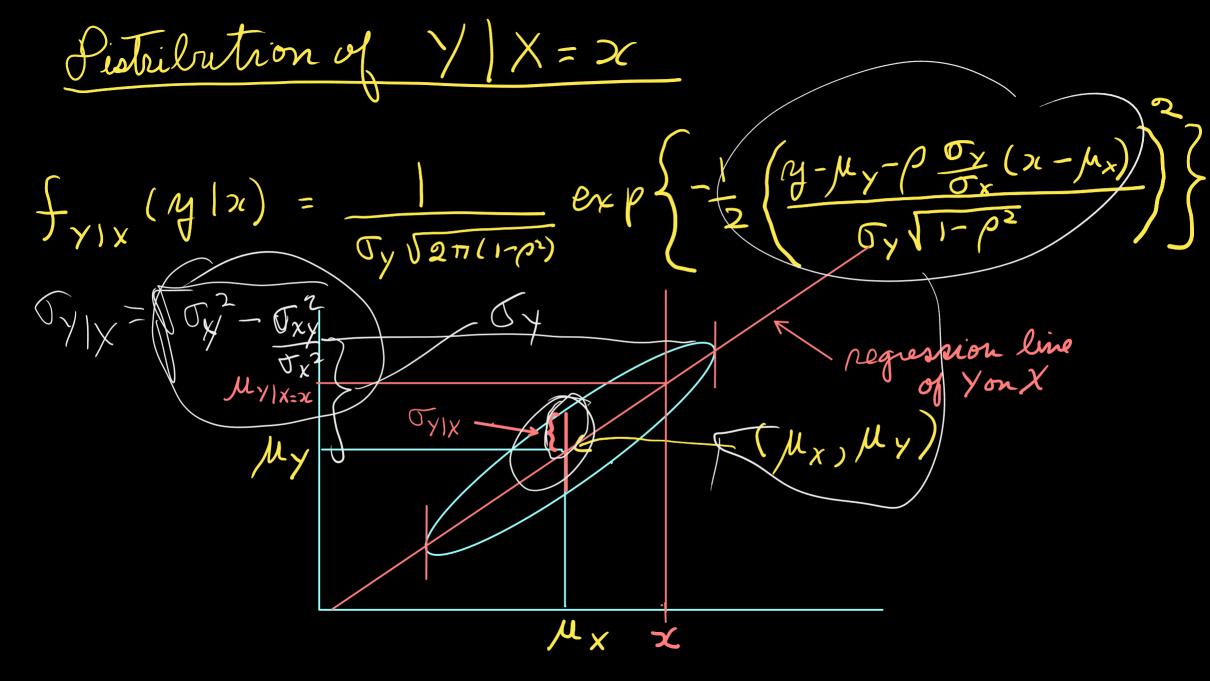
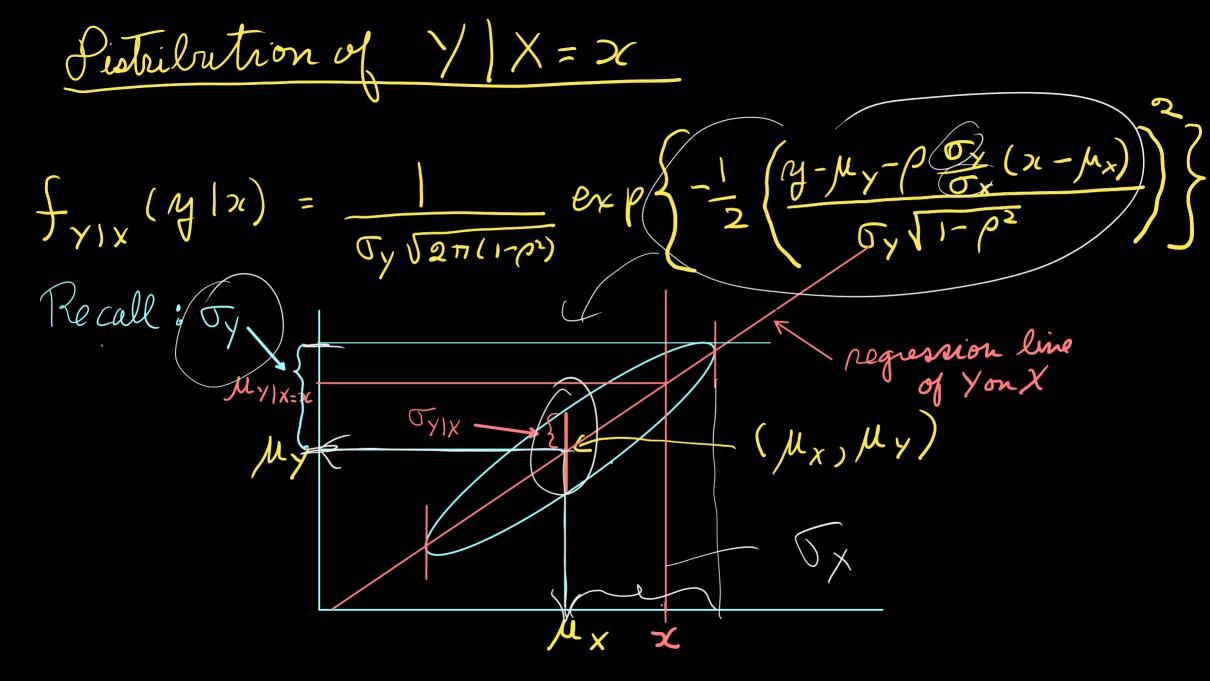
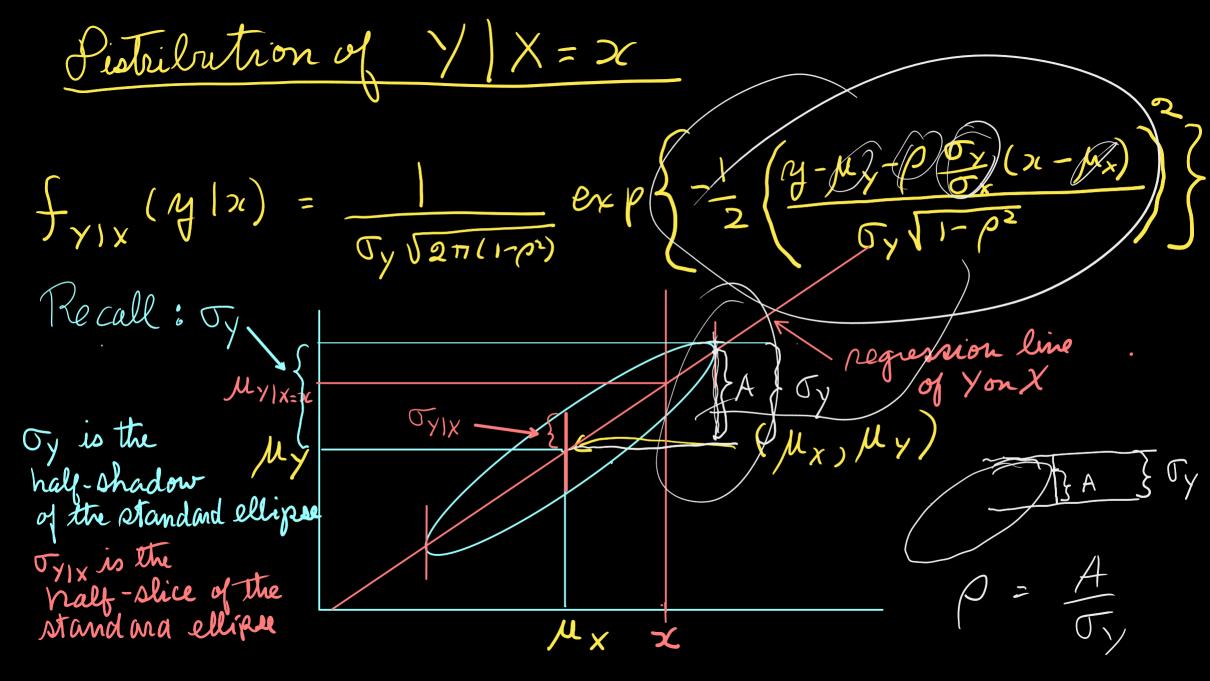
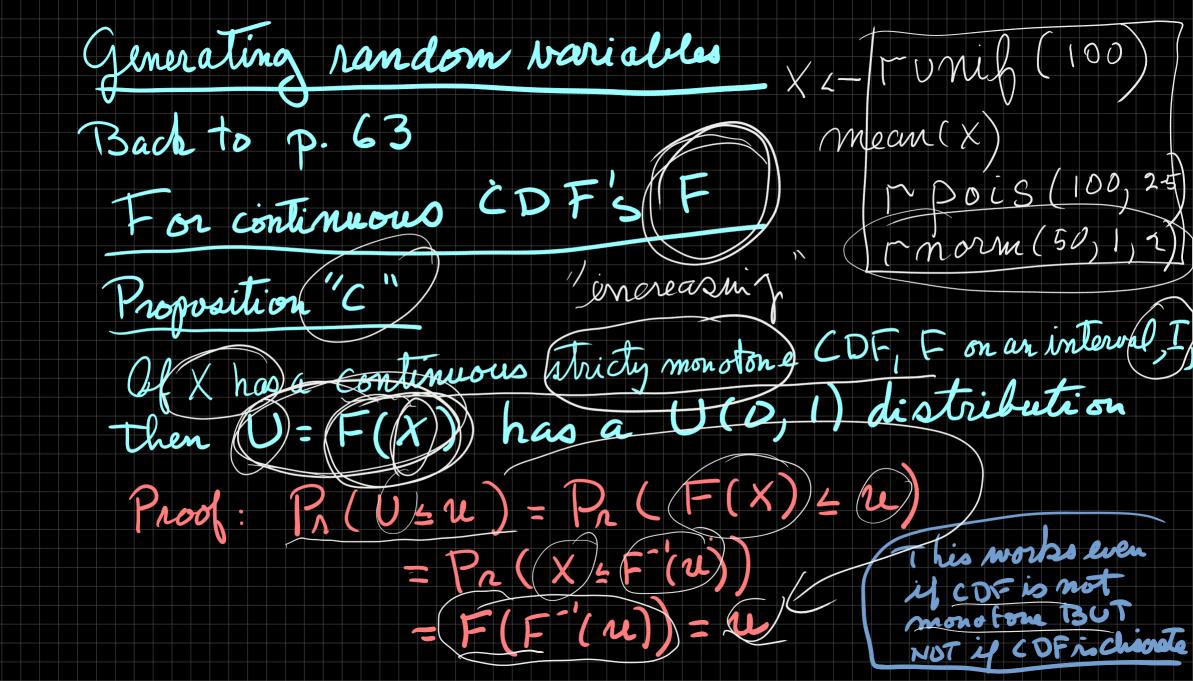


Pistribution of / X = x $= \frac{1}{\sqrt{5}\sqrt{2\pi(1-p^2)}} \exp\left(-\frac{1}{2}\right)$ + (y /2)









Proposition "D" Let F be continuous, stricty mono for Fet (V~V(0,1). Let X=F'(V). on an interval T Then COFO(X is F.) Proof: Pr(X=x) = Pr(F-(U) = x) $= P_{\mathcal{L}}(U \leq F(x)) = F(x)$ Why is this important and useful? - We know how to generate (U(0,1) sandom variables. So if we can program F, then we can generate random observations of X = F(U)

Rejection Method: p.92
Suppose X has density & with
f(x) = 0 outside an interval $[a, b]$
- Could use $X = F'(U)$ el F'is easy to
ONAGO AMO:
But sometimes, we can program to but For F' very hard.
- Of f has a maximum on [a, b]
ne can use the rejection method.

Simple algorithm 4 Un U(0,1) 0) Choose M2 max fix) then Z = (b-a)U+a ~ Orieform (a, b) 1) Jet U, ~ Uniform (a, 6) 2) Let U2 ~ Uniform (0, 11) repeat 3) Jet X = {U, if U2 \(\psi \), Laso buch to (1) otherwise neve enough

eve enough

Let's check whether this works: $P_{\lambda}(X \leq x) = P_{\lambda}(U_{1} \leq x) U_{2} \leq f(U_{1})$

$$P_{n}(x \in x \pm dx) = P_{n}(U_{1} \in x \pm dx) U_{2} \pm f(U_{1})$$

$$= P_{n}(U_{2} \pm f(U_{1})) U_{1} \in x \pm dx) P_{n}(U_{1} \in x \pm dx)$$

$$= P_{n}(U_{2} \pm f(U_{1}))$$

$$= \frac{f(x)}{h} \frac{dx}{h-a} / \left(\int_{a}^{b} f(x) dx / hx(h-a)\right)$$

$$= f(x) / 1 = f(x)$$

$$= f(x)$$

Example E Bayesian Onference 20 sided die Y = # on face 9 fair (Pr(Y=10))= (1/2) But maybe not. 04041 Let Pr(Y410) = 0 We don't know 0

Baysian approach: - Use a probability distribution to represent your uncertainty about O - Tough act. But let's pretend Prior (6~U(0,1))is reasonable distri-- get data: Toss die n Cimos! bution - Let Sche # of Times Y 4 (0) What do we know about 6 now? Posterior Distribution of a given X = x distribution

$$p(x|\theta) = \binom{n}{x} \theta^{x} (1-\theta)^{n-x}$$

$$p(x|\theta) = p(x|\theta) p(\theta)$$

$$= \binom{n}{x} \theta^{x} (1-\theta)^{n-x} \times 1, \quad 0.10 \times 1$$

$$= \binom{n}{x} \theta^{x} (1-\theta)^{n-x} \times 1, \quad 0.10 \times 1$$

$$= \binom{n}{x} \theta^{x} (1-\theta)^{n-x} \times 1, \quad 0.10 \times 1$$

$$= \binom{n}{x} \theta^{x} (1-\theta)^{n-x} d\theta$$

So $P(\theta|x) = (n+1)\binom{n}{x}\theta^{x}(1-\theta)^{n-x}$ 02041 This is a Beta (2011, n-2011) and has mean $\frac{x+1}{n+2}$ The "frequentist" estimate of θ is $\frac{2c}{n}$ Very close to x+1 if n large and x not close to 161h