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Quiz 5

1) gcd(33, 24, X) ← not zero(24), gcd(24, 33 mod 24, X), gcd(33, 24, X).

gcd(33, 24, X).

gcd(33, 24, X).

gcd(24, 9, X) ← not zero(9), gcd(9, 24 mod 9, X), gcd(24, 9, X).

gcd(9, 6, X) ← not zero(6), gcd(6, 9 mod 6, X), gcd(9, 6, X).

gcd(6, 3, X) ← not zero(3), gcd(3, 6 mod 3, X), gcd(6, 3, X).

gcd(3, 0, X)

∴ Setting X to 3 gives the empty statement and matches the gcd(u, 0, u) statement.

1.

2) 1. P -> q	Premise
2. P-> r	Premise
3. P tq	$1(a)$ [Subsungtion. $(p \rightarrow q) \equiv (\bar{p} + q)$]
4. p + r	$2(a)$ [Subsumption.($p \rightarrow q$) = $(p+q)$]
5. (p+q)(p+r)	3,4(c)
6.Ptgr	5(a)[Subsumption. Ptqr=(ptq)(ptr)]
7.p→qr	6(a) Coloranolion. P->q = P+q]

2.

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3. 1.E_1 \rightarrow E_2 + T (The sum of the values of E2 and T will be assigned to E1)
   2.E_1 \rightarrow E_2 - T (The difference of the values of E2 and T will be assigned to E1)
   3.E \rightarrow T (The value of T will be assigned to E)
   4.T_1 \rightarrow T_2 * F (The product of T2 and F will be assigned to T1)
   5.T_1 \rightarrow T_2/F The division of T2 by F will be assigned to T1)
   6.T \rightarrow F (The value of F will be assigned to T)
   7.F_1 \rightarrow -F_2 (The value of F1 will be equal to the additive inverse of the value of F2)
   8.F \rightarrow (E) (The value of E will be assigned to F)
   9.F \rightarrow const (A constant will be assigned to F)
   1.E \rightarrow T TT (TT is the value of the T pointer and E is the value of the TT pointer)
   2.TT_1 \rightarrow +T \ TT_2 (The sum of T and TT_2 will be assigned to TT_1)
   3.TT_1 \rightarrow -T \ TT_2 ( The difference between T and TT_2 will be assigned to TT_1)
   4.TT \rightarrow \epsilon (The value of TT is stored)
   5.T \rightarrow F FT (FT is the value of the F pointer and T is the value of the FT pointer)
   6.FT_1 \rightarrow *F FT_2 (The product of F and FT_2 will be assigned to FT_1)
   7.FT_1 \rightarrow /F \ FT_2 (The division of F by FT_2 will be assigned to FT_1)
   8.FT \rightarrow \epsilon (The value of FT will be stored)
   9.F_1 \rightarrow -F_2 (The value of F1 will be equal to the additive inverse of the value of F2)
   10.F \rightarrow (E) (The value of E will be assigned to F)
   11.F \rightarrow const (A constant will be assigned to F)
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An action routine is a semantic function that we tell the compiler to execute at a particular point in the parse. When it predicts a production, the parser pushes all of the right-hand side onto the stack, including terminals to be matched, nonterminals (to drive future predictions), and pointers to action routines. When a pointer to an action routine at the top of the parse stack is found, the parser calls it and passes the appropriate attributes as arguments. In Fig 6 and Fig 7, the action routines would be embedded among the symbols of the right hand sides before performing the same task as in the syntax tree.

Questi	en-4.
	The following assumes the corponents in the input are all positive integers.
	P T more - Ts D more - Ts. St:= T. d DP-d: more - Ts. d
	T → rum T-tail D T-tal · e := rum · V DT · d := T-tail · d
•	T- tal -> x emp Demp, c: T-tail.c DT. tail.d: zemp.d
	T-bail → E D T. bail. d :=""
	enp -> * * num Dexp. d = float - to -string (cnp. c x num. V) + (6"11 x * * ") + int- to - string (num. V - 1)
	onp = E Desp. d: float - to - String (esp. c)
	more - Ts1 -> t T more - Ts2 s more - Ts1, St + "+" + T.d Ts1 d
	more - Ts, d: more - Ts2 d more - Ts > E D more - Ts, d: = more - Ts. st
-	D more - 15, d;

4.

Practice Problems:

- 4.7 3
- 4.8 3
- 4.9 2
- 4.10 2
- 4.13 2
- 4.16 3
- 4.17 2