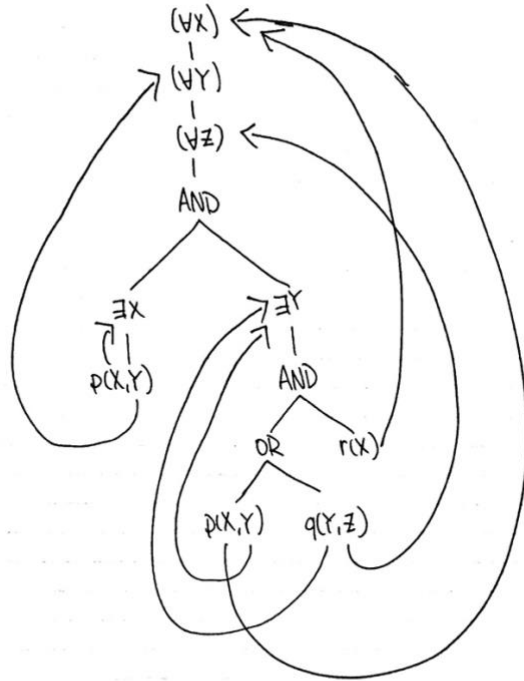


### Quiz 4

1)  
 (i)



(ii)  
 $(\forall x)(\forall y)(\forall z)((\exists x)p(x,y) \wedge (\exists y)((p(x,y) \vee q(y,z)) \wedge r(x)))$   
 $= (\forall x)(\forall y)(\forall z)((\exists x)p(x,y) \wedge (\exists A)((p(x,A) \vee q(A,z)) \wedge r(x)))$   
 $= (\forall x)(\forall y)(\forall z)((\exists B)p(B,y) \wedge (\exists A)((p(x,A) \vee q(A,z)) \wedge r(x)))$   
 $= (\forall x)(\forall y)(\forall z)(\exists B)(\exists A)(p(B,y) \wedge (p(x,A) \vee q(A,z)) \wedge r(x))$

1.

- 2) 1.  $E_1 \rightarrow E_2 + T$  (The sum of values of  $E_2$  and  $T$  will be assign as the value of  $E_1$ )
2.  $E_1 \rightarrow E_2 - T$  (Difference of values of  $E_2$  and  $T$  will be assign as the value of  $E_1$ )
3.  $E \rightarrow T$  (The value of  $T$  will be assign as the value of  $E$ )
4.  $T_1 \rightarrow T_2 * F$  (The multiplication of values of  $T_2$  and  $F$  will be assign as the value of  $T_1$ )
5.  $T_1 \rightarrow T_2 / F$  (The division of values of  $T_2$  and  $F$  will be assign as the value of  $T_1$ )
6.  $T \rightarrow F$  (The value of  $F$  will be assign as the value of  $T$ )
7.  $F_1 \rightarrow F_2$  (The value of  $F_1$  will be equal to the additive inverse of the value of  $F_2$ )
8.  $F \rightarrow (E)$  (The value of  $E$  will be assign as the value of  $F$ )
9.  $F \rightarrow \text{const}$  (A constant value will be assigned as the value of  $F$ )

def: Semantic rules attached to the productions of a tree grammar can be used to define the attribute flow of a syntax tree in exactly the same way that semantic rules attached to the productions of a context-free grammar are used to define the attribute flow of a parse tree

2.

3.  $C \rightarrow \text{dgt}_x \cdot \text{dgt}_y$   
 $C.\text{val} := \text{dgt}_x.\text{val} + \text{dgt}_y.\text{val} \times 10^{-\text{dgt}_x.\text{len}}$   
 $\text{dgt} \rightarrow \text{dgt more\_dgt}$   
 $\text{dgt.val} := \text{dgt.val} \times 10^{\text{more\_dgt.len}} + \text{more\_dgt.val}$   
 $\text{dgt.len} := \text{more\_dgt.len} + 1$   
 $\text{more\_dgt} \rightarrow \text{dgt}$   
 $\text{more\_dgt.val} := \text{dgt.val}$   
 $\text{more\_dgt.len} := \text{dgt.len}$   
 $\text{more\_dgt} \rightarrow \epsilon$   
 $\text{more\_dgt.val} := 0$   
 $\text{more\_dgt.len} := 0$   
 $\text{dgt} \rightarrow 0 \quad \text{dgt.val} := 0$   
 $\text{dgt} \rightarrow 1 \quad \text{dgt.val} := 1$   
 $\text{dgt} \rightarrow 2 \quad \text{dgt.val} := 2$   
 $\text{dgt} \rightarrow 3 \quad \text{dgt.val} := 3$   
 $\text{dgt} \rightarrow 4 \quad \text{dgt.val} := 4$   
 $\text{dgt} \rightarrow 5 \quad \text{dgt.val} := 5$   
 $\text{dgt} \rightarrow 6 \quad \text{dgt.val} := 6$   
 $\text{dgt} \rightarrow 7 \quad \text{dgt.val} := 7$   
 $\text{dgt} \rightarrow 8 \quad \text{dgt.val} := 8$   
 $\text{dgt} \rightarrow 9 \quad \text{dgt.val} := 9$

3.