

DISCRETE STRUCTURES

Lecture 1. Propositional logic and equivalence

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- ① Propositional logic
 - Proposition
 - Compound proposition
 - Truth table
 - Conjunction, Disjunction,...
 - Converse, Inverse, Contrapositive,...
- ② Propositional equivalence

Review

Construct the truth table of the compound propositions:

$$(p \oplus q) \wedge \neg r \rightarrow p$$

Precedence of Logical Operators

Operators	Precedence
\neg	1
\wedge	2
\vee	3
\rightarrow	4
\leftrightarrow	5

Logical Equivalence

Definition 2.1

The propositions p and q are **logically equivalent** if they have the same truth tables. We also write $p \equiv q$.

Example 2.2

Using truth table, prove that the below propositions are logically equivalence.

$$\neg p \vee q \text{ and } p \rightarrow q$$

Tautology and Contradiction

Definition 2.3

- ① A **tautology** (denoted by T) is a compound proposition that is always true.
- ② A **contradiction** (denoted by F) is a compound proposition that is always false.

Thus p and q are logical equivalent if and only if $p \leftrightarrow q$ is a tautology.

Table: Truth tables for $p \wedge \neg p$ and $p \vee \neg p$

p	$\neg p$	$p \wedge \neg p$
T	F	F
F	T	F

p	$\neg p$	$p \vee \neg p$
T	F	T
F	T	T

Then, $p \wedge \neg p$ is a contradiction and $p \vee \neg p$ is tautology.

Tautology and Contradiction

Example 2.4

Check if the following is tautology or contradiction:

① $(p \wedge (p \rightarrow q)) \rightarrow q$

② $p \wedge (\neg p \rightarrow q) \rightarrow r$

De Morgan's Laws

De Morgan's Law 1

The compound propositions $\neg(p \vee q)$ and $\neg p \wedge \neg q$ are logically equivalent.

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

De Morgan's Law 2

The compound propositions $\neg(p \wedge q)$ and $\neg p \vee \neg q$ are logically equivalent.

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

Example 2.5

Using De Morgan's laws to find the negation of the following sentences:

- a) Jan is rich and happy.
- b) Mei walks or takes the bus to class.

Logical Equivalences

Equivalence	Name
$p \wedge T \equiv p$ $p \vee F \equiv p$	Identity laws
$p \vee T \equiv T$ $p \wedge F \equiv F$	Domination laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double Negation law
$p \vee \neg p \equiv T$ $p \wedge \neg p \equiv F$	Negation laws

Logical Equivalences (continued)

Equivalence	Name
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
$(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	Associative laws
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption laws

Logical Equivalences Involving Implications

$$p \rightarrow q \equiv \neg p \vee q$$

$$p \rightarrow q \equiv \neg q \rightarrow \neg p \text{ (contrapositive)}$$

$$p \vee q \equiv \neg p \rightarrow q$$

$$p \wedge q \equiv \neg(p \rightarrow \neg q)$$

$$\neg(p \rightarrow q) \equiv p \wedge \neg q$$

$$(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$$

$$(p \rightarrow r) \wedge (q \rightarrow r) \equiv (p \vee q) \rightarrow r$$

$$(p \rightarrow q) \vee (p \rightarrow r) \equiv p \rightarrow (q \vee r)$$

$$(p \rightarrow r) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$$

Logical Equivalences Involving Bi-Implications

$$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$

$$p \leftrightarrow q \equiv \neg p \leftrightarrow \neg q$$

$$p \leftrightarrow q \equiv (p \wedge q) \vee (\neg p \wedge \neg q)$$

$$\neg(p \leftrightarrow q) \equiv p \leftrightarrow \neg q$$

Disjunctive Normal Form

Definition 2.6

A compound proposition is said to be in **disjunctive normal form** if it is a disjunction of conjunctions of the variables or their negations.

Example 2.7

The compound proposition

$$(p \wedge q \wedge r) \vee (\neg p \wedge q \wedge \neg r) \vee (\neg p \wedge \neg q \wedge r)$$

Logical Equivalences

Question: How to prove two propositions are equivalent?

Two methods:

- ① Using truth tables:
 - Not good for long formula
- ② Using the logical equivalences
 - The preferred method.

Logically Equivalences

Example 2.8

Show that $(p \rightarrow r) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$

- ① Using the truth table
- ② Using the logical equivalences

$$\begin{aligned}
 (p \rightarrow r) \vee (q \rightarrow r) &\equiv (\neg p \vee r) \vee (\neg q \vee r) && \text{Implication} \\
 &\equiv \neg p \vee r \vee \neg q \vee r && \text{Associative law} \\
 &\equiv \neg p \vee \neg q \vee r \vee r && \text{Commutative law} \\
 &\equiv (\neg p \vee \neg q) \vee (r \vee r) && \text{Associative law} \\
 &\equiv \neg(p \wedge q) \vee r && \text{De Morgan and Idempotent} \\
 &\equiv (p \wedge q) \rightarrow r
 \end{aligned}$$

Exercises

- ① Show that $(\neg p \rightarrow r) \wedge (q \rightarrow r)$ and $(p \rightarrow q) \rightarrow r$ are logically equivalent.
- ② Show that $(p \wedge q) \rightarrow r$ and $(p \rightarrow r) \wedge (q \rightarrow r)$ are not logically equivalent.