

ECE 341 Midterm Exam Solution

Problem No. 1

(a) True, (b) False, (c) False, (d) True, (e) False

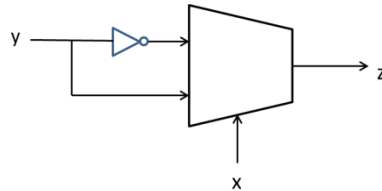
Problem No. 2

(a)

- At $t = 18$ ns, **$Q = 0$** (Q retains its initial value until there is a positive clock edge)
- At $t = 35$ ns, **$Q = 1$** (Positive clock edge arrives at $t = 20$ ns and $J = 1$, $K = 0$)
- At $t = 50$ ns, **$Q = 0$** (Q toggles its value because there is a positive clock edge at 40 ns and $J = 1$, $K = 1$)

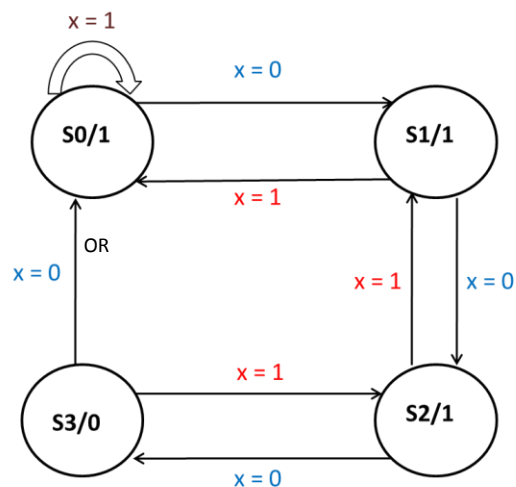
(b) $z = \overline{x \oplus y} = \overline{xy} + xy$

To implement the function with a 2-input multiplexer, we need one select input and two data inputs. We can choose either of the two variables as a select input. Then the second variable will appear in the data inputs as follows:



Problem No. 3

(a) The state machine for the given counter is as follows:



(b) The state-assigned state table is as follows:

Present State	Next State		Output z
	$x = 0$	$x = 1$	
$y_2 y_1$	$Y_2 Y_1$	$Y_2 Y_1$	
00	01	00	1
01	10	00	1
10	11	01	1
11	00	10	0

$$\begin{aligned}
 (c) \quad z &= \overline{y_2} \overline{y_1} + \overline{y_2} y_1 + y_2 \overline{y_1} \\
 &= \overline{y_2} (\overline{y_1} + y_1) + y_2 \overline{y_1} \\
 &= \overline{y_2} (1) + y_2 \overline{y_1} \\
 &= \overline{y_2} + y_2 \overline{y_1} \\
 &= (\overline{y_2} + y_2)(\overline{y_2} + \overline{y_1}) \\
 &= (1)(\overline{y_2} + \overline{y_1}) \\
 &= \overline{y_2} + \overline{y_1}
 \end{aligned}$$

$$\begin{aligned}
 Y_2 &= \bar{x} \overline{y_2} y_1 + \bar{x} y_2 \overline{y_1} + x y_2 y_1 \\
 &= \bar{x} (y_2 \oplus y_1) + x y_2 y_1
 \end{aligned}$$

Problem No. 4

$$\begin{aligned}
 (a) \quad X &= 1101, Y = 1010, c_0 = 1 \\
 G_0 &= x_0 \text{ AND } y_0 = 1 \text{ AND } 0 = 0 \\
 G_1 &= x_1 \text{ AND } y_1 = 0 \text{ AND } 1 = 0 \\
 G_2 &= x_2 \text{ AND } y_2 = 1 \text{ AND } 0 = 0 \\
 G_3 &= x_3 \text{ AND } y_3 = 1 \text{ AND } 1 = 1
 \end{aligned}$$

$$\begin{aligned}
 P_0 &= x_0 \text{ OR } y_0 = 1 \text{ OR } 0 = 1 \\
 P_1 &= x_1 \text{ OR } y_1 = 0 \text{ OR } 1 = 1 \\
 P_2 &= x_2 \text{ OR } y_2 = 1 \text{ OR } 0 = 1 \\
 P_3 &= x_3 \text{ OR } y_3 = 1 \text{ OR } 1 = 1
 \end{aligned}$$

$$\begin{aligned}
 c_3 &= G_2 + P_2 G_1 + P_2 P_1 G_0 + P_2 P_1 P_0 c_0 \\
 &= 0 \text{ OR } (1 \text{ AND } 0) \text{ OR } (1 \text{ AND } 1 \text{ AND } 0) \text{ OR } (1 \text{ AND } 1 \text{ AND } 1 \text{ AND } 1) \\
 &= 0 \text{ OR } 0 \text{ OR } 0 \text{ OR } 1 \\
 &= 1
 \end{aligned}$$

(b) Both these numbers can be represented as 5-bit binary numbers in 2's complement representation:

$$\begin{array}{r}
 \begin{array}{cccccc}
 & 0 & 1 & 0 & 0 & 0 & (+8) \\
 \times & 1 & 1 & 0 & 1 & 1 & (-5) \\
 \hline
 \end{array}
 & \Rightarrow &
 \begin{array}{cccccc}
 & 0 & 1 & 0 & 0 & 0 \\
 & 0 & -1 & +1 & 0 & -1 \\
 \hline
 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\
 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
 \hline
 1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & (-40)
 \end{array}
 \end{array}$$

(c) $\text{CPI} = (10\% * 3) + (30\% * 5) + (20\% * 4) + (40\% * 4) = 4.2$

Instruction Execution Rate = 500 million instructions per second = $5 * 10^8$ instructions per second

Since:

$$\text{Execution time} = (\text{Number of Instructions} * \text{CPI}) / \text{Clock Rate}$$

Therefore:

$$\text{Clock Rate} = (\text{Number of Instructions} * \text{CPI}) / \text{Execution Time} = \text{CPI} * \text{Instruction Execution Rate}$$
$$= 4.2 * 5 * 10^8 = 2.1 * 10^9 = 2.1 \text{ GHz}$$

Therefore the processor must operate at a clock speed of at least 2.1 GHz to satisfy the instruction execution rate target.

Problem No. 5

(a) (i) *MuxB* selects inputs **1**, **0**, and **1** for instructions 1, 2, and 3, respectively.

(ii)

For instruction 1:

Contents of $RY = \text{Contents of memory address } (28000 + 200) = \mathbf{500}$

The contents of register R4 become 500 after instruction 1.

For instruction 2:

Contents of $RY = [R2] + [R4] = 200 + 500 = \mathbf{700}$

The contents of register R5 become 700 after instruction 2.

(iii) *RF_write* has values of **1**, **1**, and **0** for instructions 1, 2, and 3, respectively.

(b) The “store” instruction writes the contents of R5 to the memory location at address $28000 + 400 = 28400$.

Therefore the final contents of memory location 28400 are **700**.

The final contents of memory location 28200 remain unchanged at **500**.