

Identify
Locality in a
Program
Hierar





The Memory Hierarchy

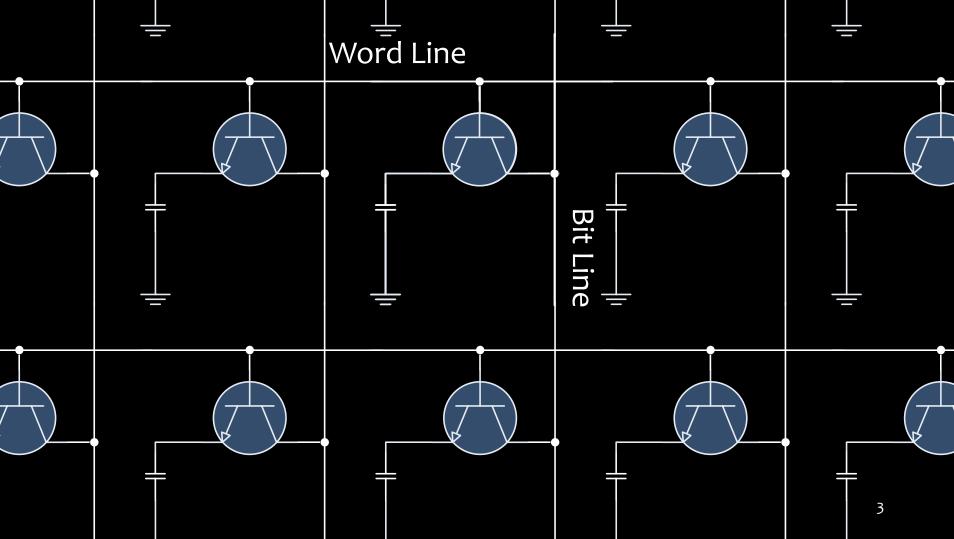
- ① Storage Technologies and Trends
 - ② Locality of Reference
- ③ Caching in the Memory Hierarchy

Random-access Memory (RAM)

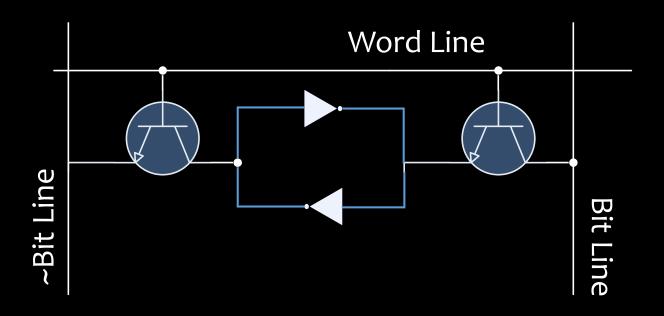


			Needs refresh		Applications
SRAM	4 or 6	1X	No	100x	Cache memories
DRAM	1	10X	Yes	1X	Main memories

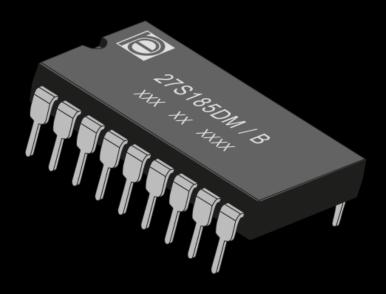
Dynamic RAM



Static RAM



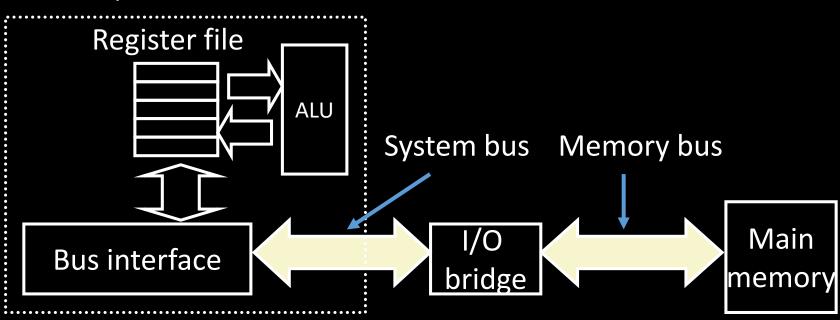
Nonvolatile Memories



Read-only memory (ROM)
Programmable ROM (PROM)
Erasable PROM (EPROM)
Electrically erasable PROM (EEPROM)
Flash memory

Bus

CPU chip

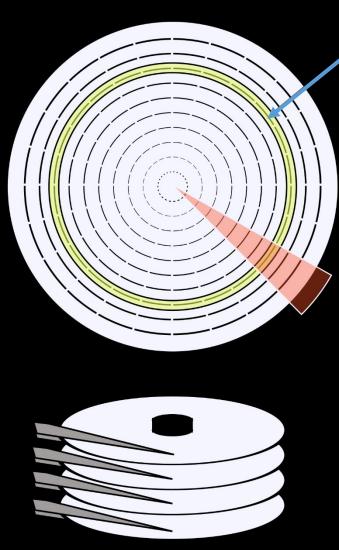


Rotating Disk

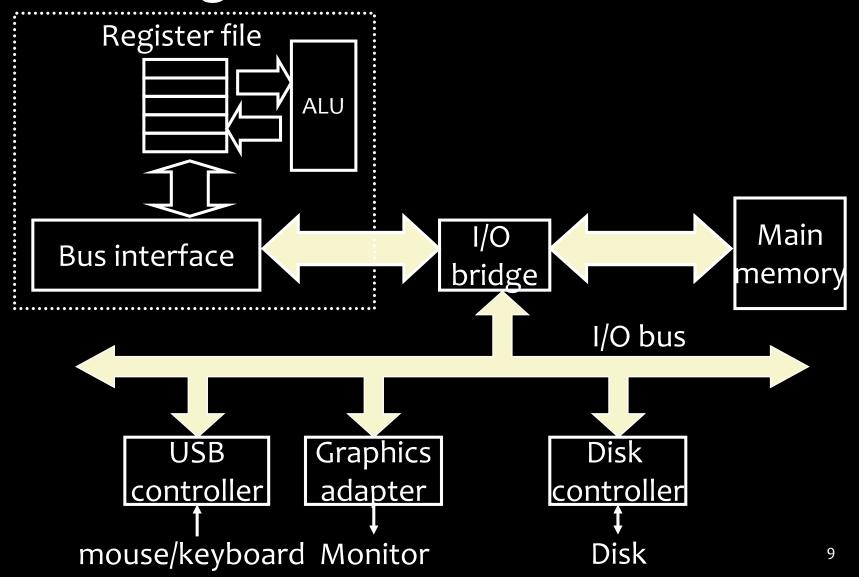


Disk Access

Track/Cylinder

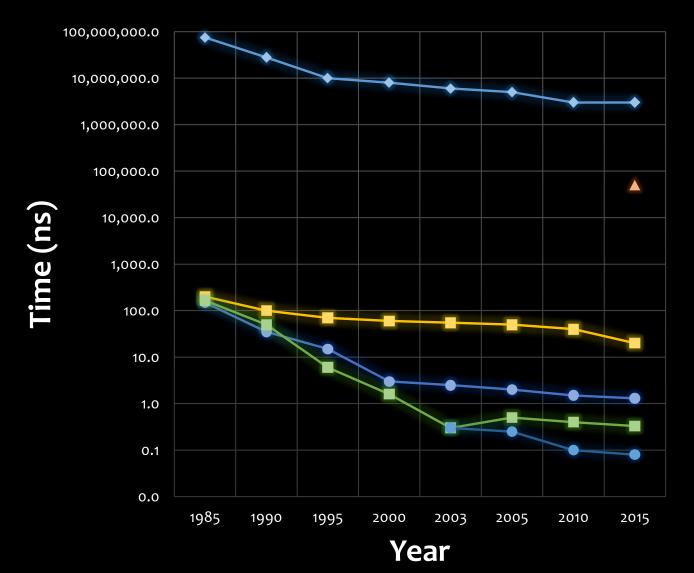


Reading a Disk Sector



Solid State Disks I/O bus Solid State Disk (SSD) Flash translation layer Flash memory Block o Block B-1 Page o Page 1 Page P-Page o Page 1 Page P-1

The CPU-Memory Gap

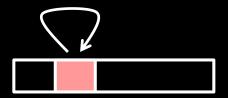


- → Disk seek time
- **→**SSD access time
- DRAM access time
- SRAM access time
- CPU cycle time
- Effective CPU cycle time

Locality

 Programs tend to use data and instructions with addresses near or equal to those they have used recently

Temporal locality



Spatial locality



Locality Example

```
sum = 0;
       for (i = 0; i < n; i++)
       → sum += a[i]; ←
      return sum;
Temporal locality
                           Spatial locality
```

Q: Does this function have good locality with respect to array a?

```
int sum array rows(int a[M][N])
    int i, j, sum = 0;
    for (i = 0; i < M; i++)
        for (j = 0; j < N; j++)
            sum += a[i][j];
    return sum;
```

Q: Does this function have good locality with respect to array a?

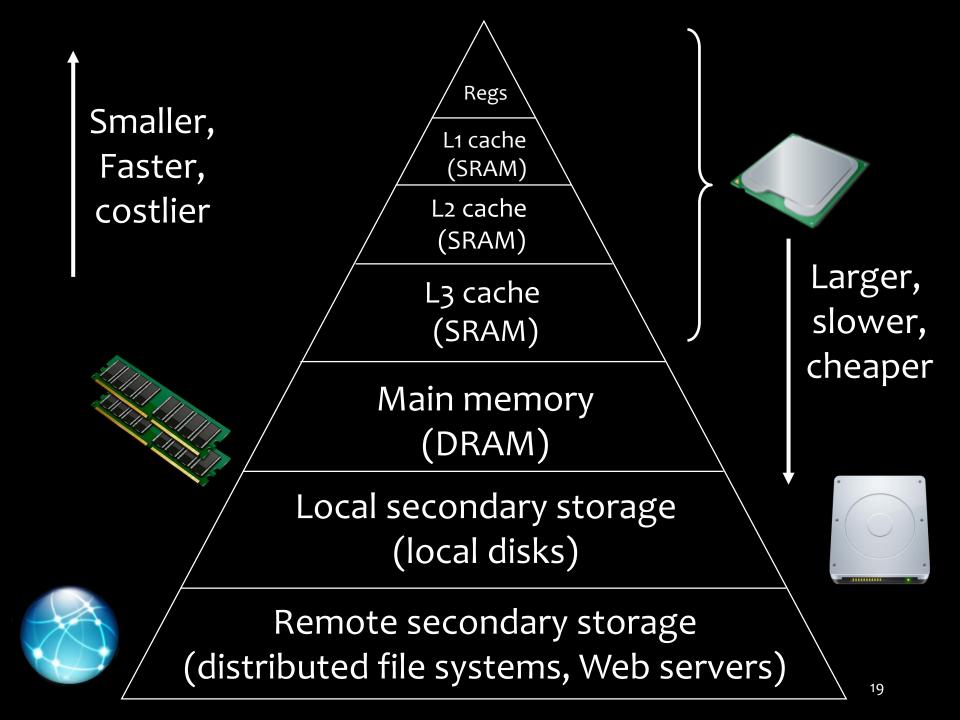
```
int sum array cols(int a[M][N])
    int i, j, sum = 0;
    for (j = 0; j < N; j++)
        for (i = 0; i < M; i++)
            sum += a[i][j];
    return sum;
```

Q: Can you permute the loops so that the function scans the 3-d array a with a stride-1 reference pattern (and thus has good spatial locality)?

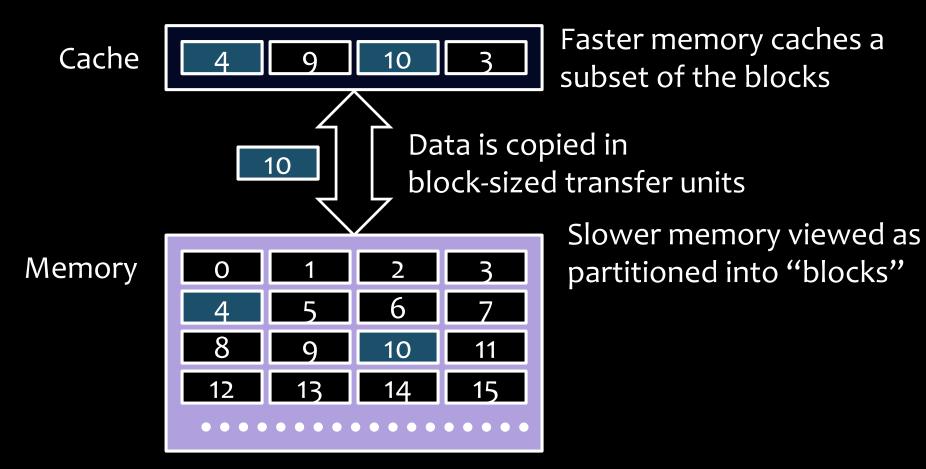
```
int sum array 3d(int a[M][N][N])
    int i, j, k, sum = 0;
    for (i = 0; i < M; i++)
        for (j = 0; j < N; j++)
            for (k = 0; k < N; k++)
                sum += a[k][i][j];
    return sum;
```

Fundamental properties

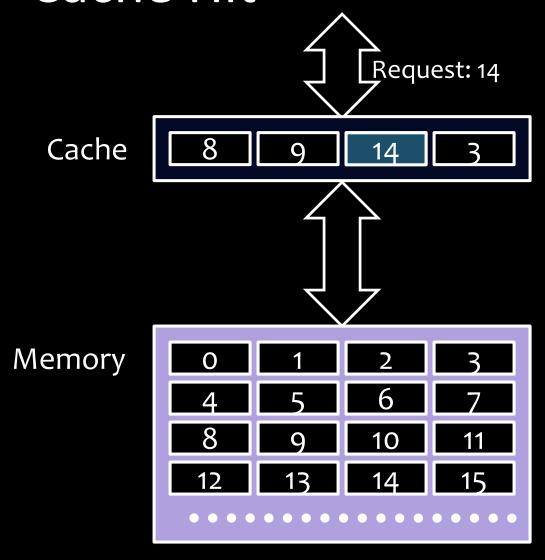
- Fast storage technologies cost more per byte, have less capacity, and require more power (heat!).
- The gap between CPU and main memory speed is widening.
- Well-written programs tend to exhibit good locality.



General Cache Concepts



Cache Hit

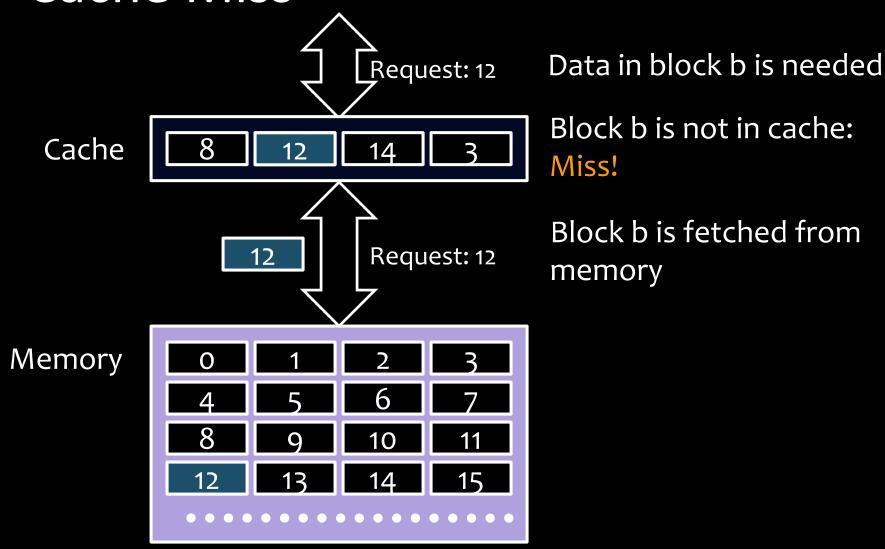


Data in block b is needed

Block b is in cache:

Hit!

Cache Miss



Cache Type	What is Cached?	Where is it Cached?	Latency (cycles)	Managed by
Registers	4-8 bytes words	CPU core	0	Compiler
TLB	Address translations	On-chip TLB	0	Hardware MMU
L1 cache	64-byte blocks	On-Chip L1	4	Hardware
L2 cache	64-byte blocks	On-Chip L2	10	Hardware
Virtual Memory	4-KB pages	Main memory	100	Hardware + OS
Buffer cache	Parts of files	Main memory	100	OS
Disk cache	Disk sectors	Disk controller	100K	Disk firmware
Network buffer cache	Parts of files	Local disk	10M	NFS client
Browser cache	Web pages	Local disk	10M	Web browser
Web cache	Web pages	Remote server disks	1B	Web proxy server

Summary

- Storage Technologies
 - RAM, ROM
 - Rotating disk, Solid state disk
- Locality
 - Temporal
 - Spatial
- Memory Hierarchy
- Cache
 - Hit
 - Miss





Bill Gates

Co-founder of Microsoft



No one will need more than 637Kb of memory for a personal computer.

