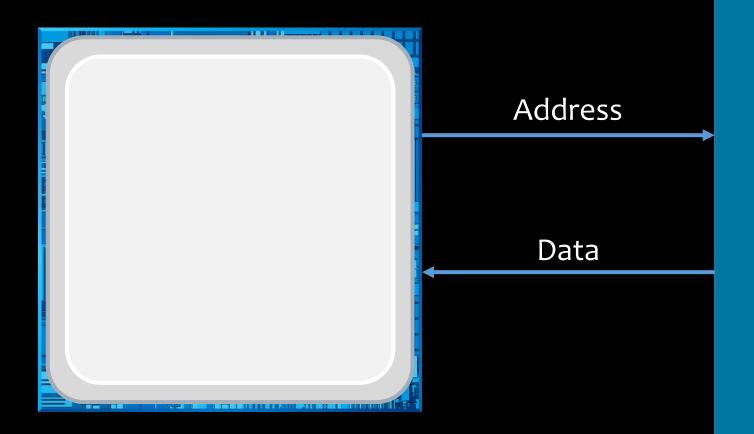




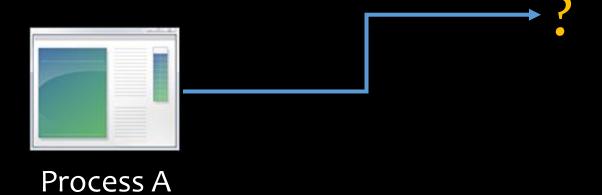
# Virtual Memory

- ① Address Spaces
- ② VM for Memory Management
  - 3 Address Translation

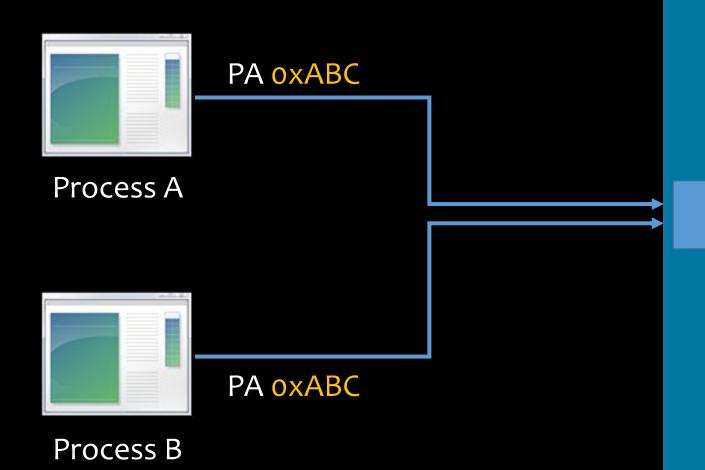
## Physical Address



### Problem



### Problem





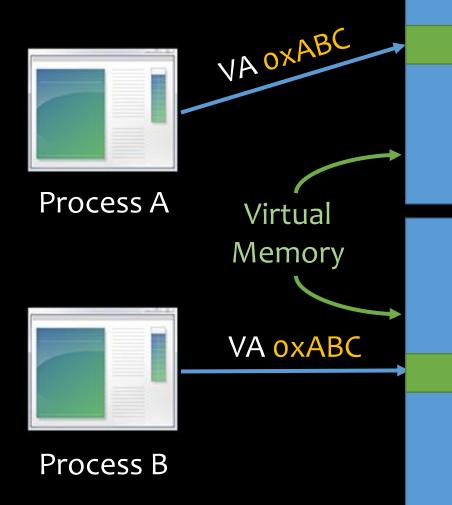
#### David John Wheeler

1985 Computer Pioneer Award for assembly language programming



All problems in computer science can be solved by another level of indirection ... except of course for the problem of too many indirections.

### Virtual Address

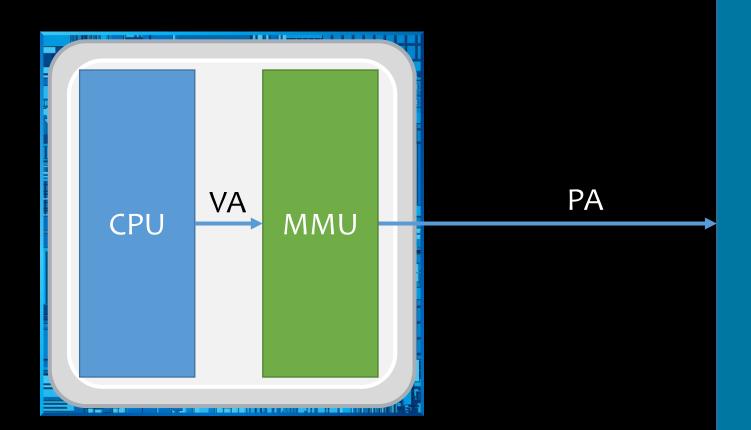


PA 0x200

PA

0X100

## Memory Management Unit



### Address Space

Virtual Address Space

$$\{0,1,...,N-1\}$$



Physical Address Space

$$\{0,1,\ldots,M-1\}$$

Virtual Address

n

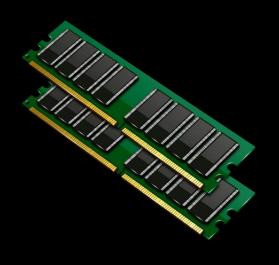
Physical Address

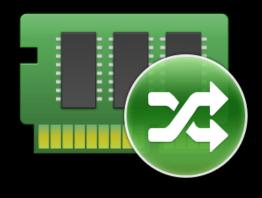
m

### Virtual Address

Number of virtual address bits (n)	Number of virtual addresses (N)	Largest possible virtual address
8		
	2 <sup>?</sup> =64K	
		2 <sup>32</sup> -1=?G-1
	2 <sup>?</sup> =256T	
64		

## Why Virtual Memory?







Efficiency

Simplification

Protection

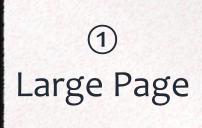
## Page Table Entries

n	P=2 <sup>p</sup>	Number of PTEs
16	4K	
16	8K	
32	4K	
32	8K	

# DRAM Cache

Regs L1 cache (SRAM) L<sub>2</sub> cache (SRAM) L<sub>3</sub> cache (SRAM) 10X slower Main memory 100,000X (DRAM) slower Local secondary storage (local disks) Remote secondary storage

(distributed file systems, Web servers)



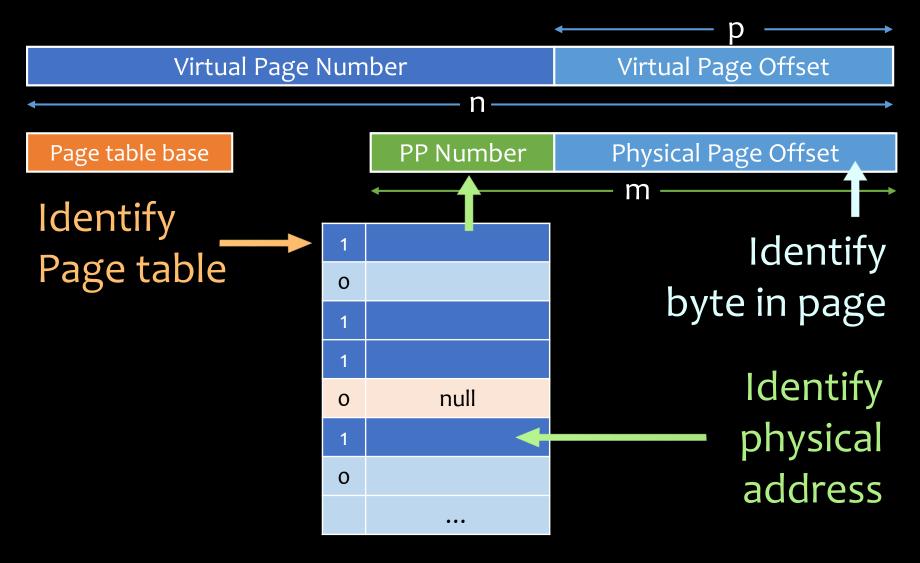


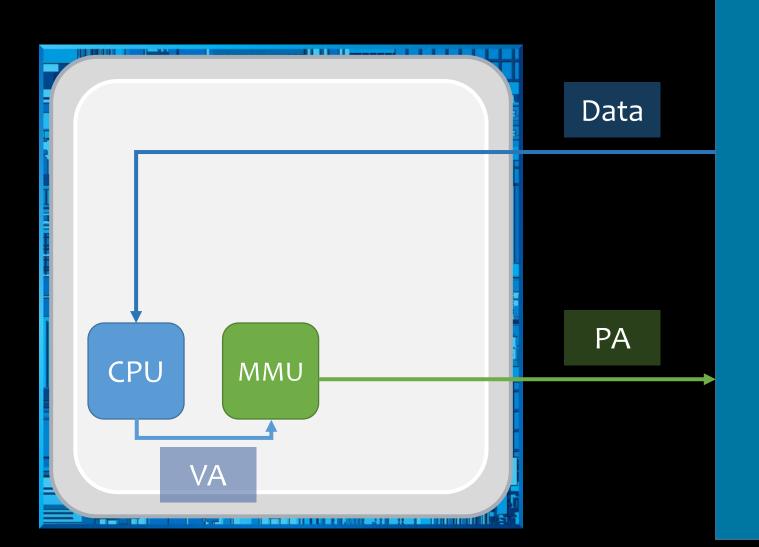
② Fully Associative

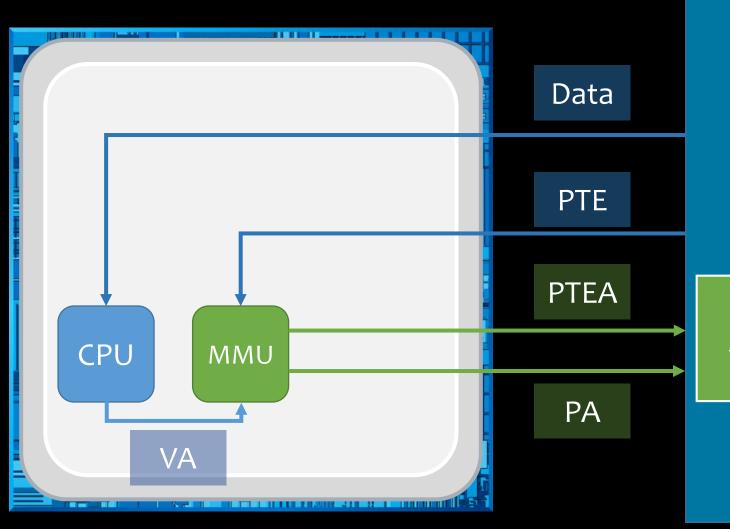


③ Write-back

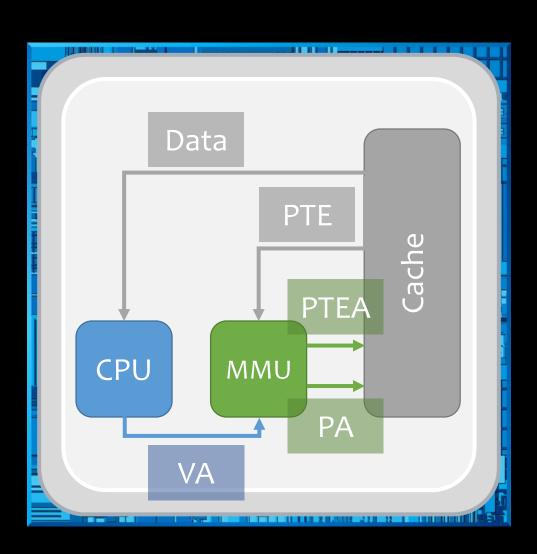
#### Address Translation



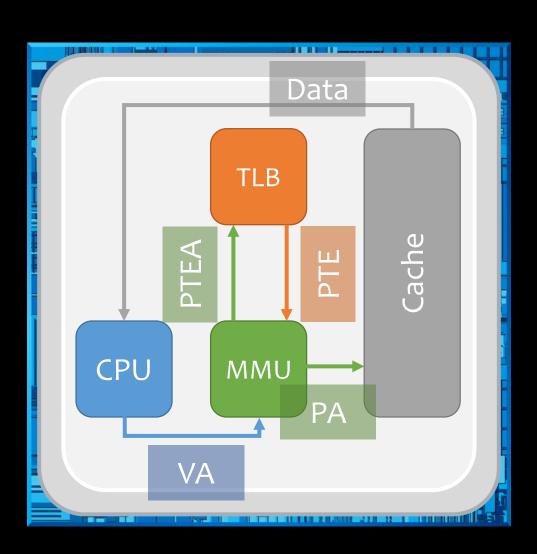




Page table

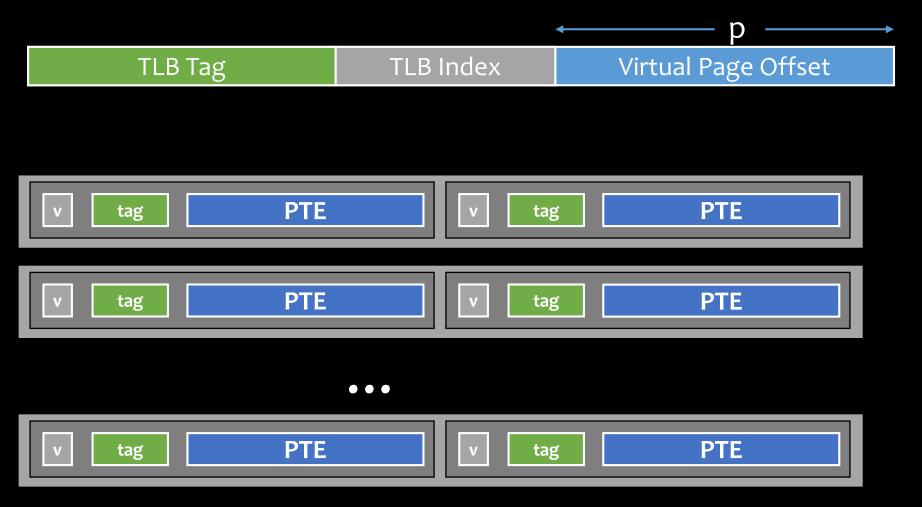


Page table



Page table

#### Translation Lookaside Buffer



#### Summary

- Address
  - Physical
  - Virtual
- Page Hit / Page Fault
- Translation Lookaside Buffer