

# Algorithmic Game Theory

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October 11, 2022

# Overview

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## 1. Week 1 - Normal Form Games

- 1.1 Definition
- 1.2 Strategies

## 2. Week 2 - Solution Concepts

- 2.1 Maximin
- 2.2 Nash
- 2.3 Non-Rational Player
- 2.4 Evaluating Strategies

## 3. Week 3 - Zero Sum, First Algorithms

- 3.1 Fictitious Play
- 3.2 LP prerequisites
- 3.3 Nash and LP
- 3.4 Correlated Equilibrium

## 4. Week 4 - Regret Minimization

- 4.1 External regret
- 4.2 Application to Games

## 5. Week 6 - Counterfactual Regret Minimization

- 5.1 Motivation

# About the Course

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## Class

- Simultaneous and sequential decision making
- Solution concepts and optimal policies
- Practical algorithm for finding the optimal policies

## Homeworks

- You will get to implement the games and algorithms!

# Understand These!



(a) AlphaZero



(b) AlphaStar



(c) DeepStack

# Game Theory - Reinforcement Learning

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## Reinforcement Learning

- Single agent settings
- Maximize reward
- Scalable practical algorithms

## Game Theory

- Multi agent settings
- Analyzes agent interaction, incentives
- Optimal solution concepts
- Algorithms (historically) tabular and not scalable

# Terminology

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## Reinforcement Learning

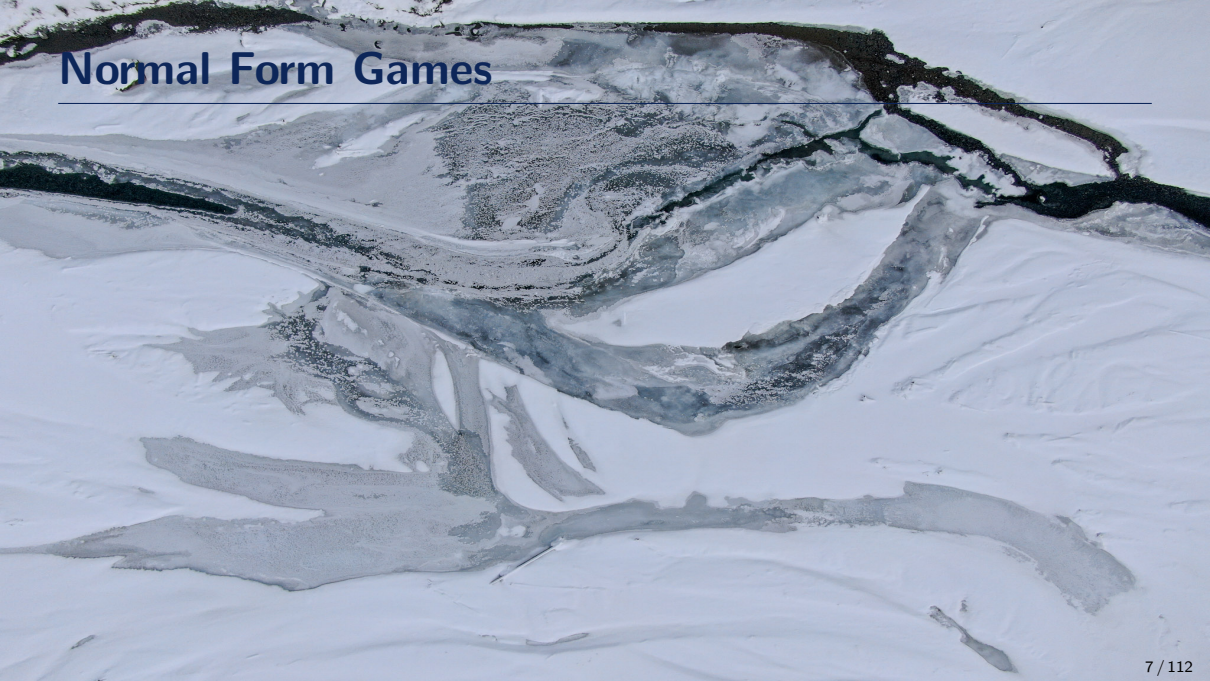
1. Environment
2. Agent
3. Policy
4. Reward

## Game Theory

1. Game
2. Player
3. Strategy
4. Utility

# Normal Form Games

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# Normal Form Games

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The normal form games is a model in which each player chooses his strategy, and then all players play simultaneously. The outcome depends on the actions chosen by the players.

## Definition: Normal Form Game

is a tuple  $\langle N, (A_i), (u_i) \rangle$ , where

- $N$  is the **finite** set of players
- $A_i$  is the nonempty set of actions available to the player  $i$
- $u_i$  is a **payoff/utility** function for the player  $i$ . Let  $A = \times_{i \in N} A_i$ .  
 $u_i : A \rightarrow \mathbb{R}$



# Normal Form Games

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- If there are only two players ( $|N| = 2$ ), we can conveniently describe the game using a table

	Rock	Paper	Scissors
Rock	(0, 0)	(-1, 1)	(1, -1)
Paper	(1, -1)	(0, 0)	(-1, 1)
Scissors	(-1, 1)	(1, -1)	(0, 0)

(a) Rock-Paper-Scissors

	Confess	Be quiet
Confess	(8, 8)	(0, 10)
Be quiet	(10, 0)	(2, 2)

(c) Prisoner's dilemma

- Rows/columns correspond to actions of player one/two
- The cell  $(i, j)$  contains the players' payoffs  $u_1(i, j)$  and  $u_2(i, j)$

# Constant Sum Games

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- Constant-sum game is a game for which  $u_1 + u_2 = c$
- Zero-sum game is a constant-sum game for  $c = 0$ , so  $u_1 = -u_2$
- Critical implications!

# Zero Sum Games

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	Rock	Paper	Scissors
Rock	(0, 0)	(-1, 1)	(1, -1)
Paper	(1, -1)	(0, 0)	(-1, 1)
Scissors	(-1, 1)	(1, -1)	(0, 0)

	Rock	Paper	Scissors
Rock	0	-1	1
Paper	1	0	-1
Scissors	-1	1	0

# Normal Form Game Strategies

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## Definition: Pure Strategy

$a_i \in A_i$  is player  $i$ 's pure strategy. This strategy is referred to as pure, because there's no probability involved. For example, the player can always play Scissors.

## Definition: Mixed Strategy

is a probability measure over the player's pure strategies. The set of player  $i$ 's mixed strategies is denoted as  $\Pi_i$ . Given  $\pi_i \in \Pi_i$ , we denote the probability that the player chooses the action  $a_j \in A_i$  as  $\pi^{\pi_i}(a_j)$ . Mixed strategies allow a player to probabilistically choose actions.

## Definition: Strategy profile

Is the set of all players' strategies (one for every player), denoted as  $\pi = (\pi_0, \pi_1 \dots \pi_n)$ . Finally,  $\pi_{-i}$  refers to all the strategies in  $\pi$  except  $\pi_i$ .

# Normal Form Game Strategies II

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- **Pure Strategy**  $a_i \in A_i$  is player  $i$ 's pure strategy.
- **Mixed Strategy** is a probability measure over the player's pure strategies. The set of player  $i$ 's mixed strategies is denoted as  $\Pi_i$ . Given  $\pi_i \in \Pi_i$ , we denote the probability that the player chooses the action  $a_j \in A_i$  as  $\pi_i(a_j)$
- **Strategy profile** Is a tuple of all players' strategies, denoted as  $\pi = (\pi_0, \pi_1 \dots \pi_n)$ . Finally,  $\pi_{-i}$  refers to all the strategies in  $\pi$  except for  $\pi_i$ .

# Outcome

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- Given a pure strategies of all players, we can easily compute the utilities. Player  $i$ 's utility =  $u_i(a)$
- How to compute the outcome if the players use mixed strategy (they randomize among the pure strategies)? We simply compute the expected value given the probability measure.
- Since the players choose the actions simultaneously, the events are independent and consequently  $\pi^\pi((a_0, a_1, \dots, a_n)) = \pi^{\pi_0}(a_0)\pi^{\pi_1}(a_1) \dots \pi^{\pi_n}(a_n)$
- Using this fact, computing the expected value is easy

$$u_i(\pi) = \sum_{a \in A} \pi^\pi(a) u_i(a)$$

# Best Response

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- One of the key concepts, that you will see throughout the class
- Given the strategies  $\pi_{-i}$  of the opponents, the **best response** is the strategy that maximizes the utility for the player.

## Definition: Best Response

Best response against a policy  $\pi_i$  is:

$$\arg \max_{\pi_{-i} \in \Pi_{-i}} R_{-i}(\pi_i, \pi_{-i})$$

We use  $\mathbb{BR}(\pi_i)$  to denote the set of best response policies against the policy  $\pi_i$ .

# Best Response

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Note that for zero-sum games, opponent maximizing their reward is equivalent to opponent minimizing our reward.

$$\arg \max_{\pi_{-i}} R_{-i}(\pi_i, \pi_{-i}) = \arg \min_{\pi_{-i}} R_i(\pi_i, \pi_{-i})$$

As this means the player's value against any best-response strategy is unique, we denote this unique value as  $BRV_i(\pi_i)$ .

$$BRV_i(\pi_i) = \min_{\pi_{-i}} R_i(\pi_i, \pi_{-i}) = - \max_{\pi_{-i}} R_{-i}(\pi_i, \pi_{-i})$$



# Best Response

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## Lemma

For any best response strategy  $\pi_i \in \mathbb{BR}_i(\pi_{-i})$ , all the actions in the support have the same expected value.

## Lemma

The best response set  $\mathbb{BR}(\pi_{-i})$  is convex.

# Dominated Strategies

- Some actions can be clearly poor choices, and it makes no sense for a rational player to take.
- Strategy  $\pi_i^a$  **strictly dominates**  $\pi_i^b$  iff for any  $\pi_{-i}$   
$$u_i(\pi_i^a, \pi_{-i}) > u_i(\pi_i^b, \pi_{-i})$$
- Strategy  $\pi_i^a$  **weakly dominates**  $\pi_i^b$  iff for any  $\pi_{-i}$   
$$u_i(\pi_i^a, \pi_{-i}) \geq u_i(\pi_i^b, \pi_{-i})$$
- Strategy is **strictly/weakly** dominated if there's a strategy that strictly/weakly dominates it.
- Strategies  $\pi_i^a, \pi_i^b$  are **intransitive** iff one neither dominates nor is dominated by the other.

## Examples

Can a weakly/strictly dominated strategy be a best response?

# Iterated elimination of dominated strategies

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- A rational player does not play dominated strategy
- Iterated elimination of dominated strategies

	Left	Center	Right
Top	(13, 3)	(1, 4)	(7, 3)
Middle	(4, 1)	(3, 3)	(6, 2)
Up	(-1, 9)	(2, 8)	(8, -1)

## Examples

Can a weakly/strictly dominated strategy that we found during the iterated elimination be a best response in the original game?

# Week 1 Homework

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1. Python and notebooks
2. Strategy pair evaluation for a matrix game
3. Best response calculation
4. Strategy evaluation against a best response
5. Iterated removal of dominated strategies
6. OpenSpiel