



# DAA ASSIGNMENT

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# QUESTION

The task is to implement a recursive algorithm in C++ that generates the  $n U r$  permutations with unrestricted repetitions in reverse lexical order.

The algorithm should generate all possible permutations of  $r$  elements from a set of  $n$  distinct elements, allowing repetitions.



# Algo

1. Define a recursive function generate Permutations( $n, r, \text{perm}$ ) that takes three arguments:
  - a.  $n$ : the number of distinct elements in the set
  - b.  $r$ : the desired length of the permutation
  - c.  $p$ : a vector to store the permutation
2. Check if  $r == 0$ . If it is, print the permutation in reverse lexical order and return.
3. Otherwise, for each element  $i$  from 1 to  $n$ , do the following:
  - a. Add  $i$  to the end of the permutation vector.
  - b. Recursively call `permute( $p, n, r - 1$ )`.
  - c. Remove the last element from the permutation vector.
4. In the main function:
  - a. Initialize  $n$  and  $r$ .
  - b. Call `Permute( $p, n, r$ )` to generate the permutations.



Code :

```
#include <iostream>
#include <vector>

using namespace std;

void permute(vector<int> &p, int n, int r)
{
    if (r == 0)
    {
        for (int i = p.size() - 1; i >= 0; i--)
        {
            cout << p[i] << " ";
        }
        cout << endl;
    }
    else
    {
        for (int i = 1; i <= n; i++)
        {
            p.push_back(i);
            permute(p, n, r - 1);
            p.pop_back();
        }
    }
}

int main()
{
    int n, r;
    cout << "Enter n and r: ";
    cin >> n >> r;

    vector<int> p;
    permute(p, n, r);

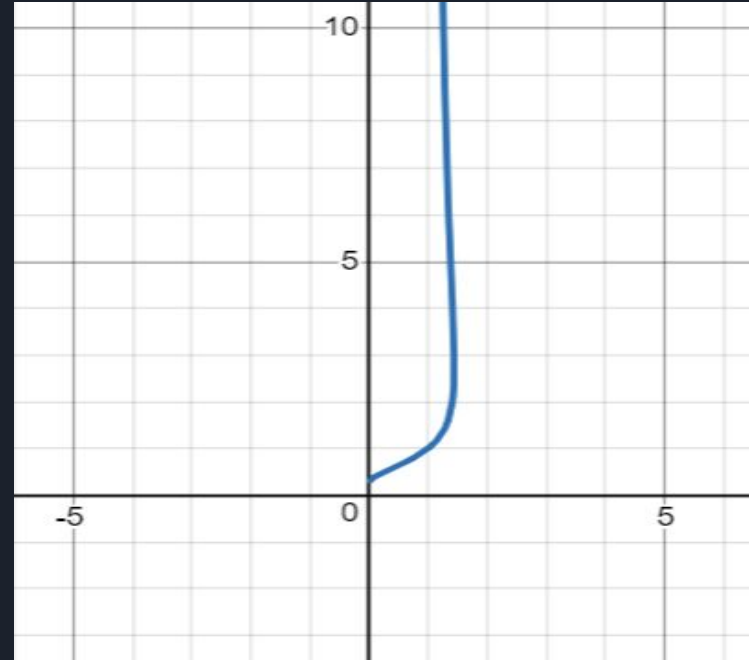
    return 0;
}
```

# Time Complexity

The time complexity of the given code is  $O(n^r)$ , where  $n$  is the number of elements to choose from and  $r$  is the number of elements in each permutation.

The reason for this complexity is that the code generates all possible permutations with unrestricted repetitions, and there are  $n^r$  such permutations.

The recursive function `permute` is called  $r$  times (the depth of the recursion), and at each level of the recursion, there are  $n$  branches (the for loop with  $i$  from 1 to  $n$ ). Therefore, the total number of function calls is  $n^r$ .

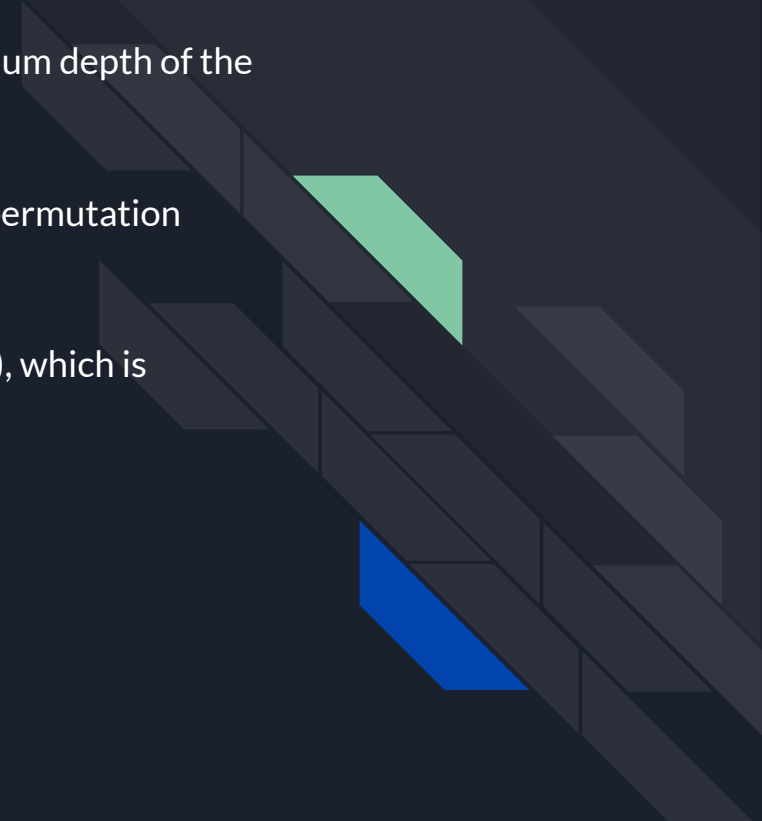


# Space Complexity

The space used by the algorithm depends on the maximum depth of the recursive tree, which is equal to  $r$ .

At each level of the tree, a new integer is added to the permutation vector, which has a maximum size of  $r$ .

Therefore, the space complexity of the algorithm is  $O(r)$ , which is relatively small compared to the time complexity.





# APOSTERIORI ANALYSIS

We are using a variable  $t$  to count the number of operations required for a particular value of  $n$ ;

For  $n=3$   $r=2$  number of Operations=144

For  $n=3$   $r=3$  number of Operations=558

For  $n=4$   $r=3$  number of Operations=1272

For  $n=4$   $r=4$  number of Operations=6136

For  $n=5$   $r=3$  number of Operations=2430

We observe that as  $n$  increases number of Operations also increase proportional to  $O(n^r)$ .

# Conclusion

The C++ implementation of the recursive algorithm for generating  $n$  U  $r$  permutations with unrestricted repetitions in reverse lexical order provides a useful tool for enumerating all possible permutations of  $r$  elements chosen from a set of  $n$  distinct objects.

While the time and space complexity of the algorithm can be limiting for large values of  $n$  and  $r$ , it is efficient and accurate for small to medium-sized inputs. The algorithm can be easily adapted to other problems that require generating all possible combinations of a set of objects, making it a useful reference for researchers, engineers, and programmers.

