# Anonymous Key Agreements for V2X Communication

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May 1, 2024

Introduction

Preliminaries

Our Proposition





# V2X Related Terminology

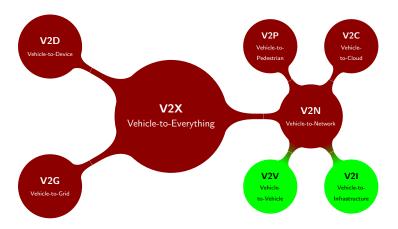


Figure 1: A breakdown of V2X.

# Message Types in V2X

- Cooperative Awareness Messages (CAMs)<sup>1</sup> and Basic Safety Messages (BSMs)<sup>2</sup>.
  - Exchanged between vehicles to create awareness and support cooperative performance of vehicles in the road network.
  - Includes status information such as time, position, speed, active systems, vehicle dimensions, etc.
  - Broadcasted unencrypted in 5.9 GHz channel (ETSI ITS-G5).
  - 4 Huge privacy concerns and threats!

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<sup>&</sup>lt;sup>1</sup>European Telecommunications Standards Institute. "Intelligent Transport Systems (ITS); Vehicular Communications; Basic Set of Applications; Part 2: Specification of Cooperative Awareness Basic Service". In: ETSI EN 302 637-2 V1.4.1 (2019). URL: https://www.etsi.org/deliver/etsi en/302600 302699/30263702/01.04.01 60/en 30263702v010401p.pdf.

<sup>&</sup>lt;sup>2</sup> J2735\_202309: V2X Communications Message Set Dictionary - SAE International. URL: https://www.sae.org/standards/content/j2735\_202309/ (visited on 04/15/2024). □ ▶ ◀ 를 ▶ ◀ 를 ▶ ◀ 를 ▶

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- Other types of messages
  - Signal Phase and Timing (SPaT)
  - Roadside Infrastructure Information (MAP)

<sup>&</sup>lt;sup>1</sup>European Telecommunications Standards Institute, "Intelligent Transport Systems (ITS); Vehicular Communications; Basic Set of Applications; Part 2: Specification of Cooperative Awareness Basic Service".

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#### Motivation and Goals

- Do we really need to encrypt CAMs?
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  - Focus on encrypting more sensitive messages and information sent less frequently.
- Unlimited privacy.
- Better security guarantees (privacy, authenticity, confidentiality).

# **Pairings**

#### Definition 1

Pairing<sup>a</sup> Let  $\mathbb{G}_0 = \langle g_0 \rangle$ ,  $\mathbb{G}_1 = \langle g_1 \rangle$ ,  $\mathbb{G}_{\mathcal{T}}$  be three cyclic groups of prime order q. A *pairing* is an efficiently computable function  $e : \mathbb{G}_0 \times \mathbb{G}_1 \to \mathbb{G}_{\mathcal{T}}$  satisfying the following properties:

**1** bilinear: for all  $u, u' \in \mathbb{G}_0$  and  $v, v' \in \mathbb{G}_1$ , we have

$$e(uu',v) = e(u,v)e(u',v)$$
 (1)

$$e(u, vv') = e(u, v) e(u, v')$$
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② non-degenerate:  $g_T := e(g_0, g_1)$  is a generator of  $\mathbb{G}_T$ .

<sup>&</sup>lt;sup>a</sup> A Graduate Course in Applied Cryptography. URL: https://toc.cryptobook.us/ (visited on 04/30/2024).

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- **1** Here,  $\mathbb{G}_0$  and  $\mathbb{G}_1$  are called source groups and  $\mathbb{G}_T$  is called the target group.
- ② When  $\mathbb{G}_0 = \mathbb{G}_1$ , the pairing is said to be *symmetric*.

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• A key agreement protocol where two parties agree on a shared secret key, without being able to determine the other party.

<sup>&</sup>lt;sup>3</sup>Aniket Kate, Greg Zaverucha, and Ian Goldberg. "Pairing-Based Onion Routing". In: *Privacy Enhancing Technologies*. Ed. by Nikita Borisov and Philippe Golle. Vol. 4776. Berlin, Heidelberg: Springer Berlin Heidelberg, 2007, pp. 95–112. ISBN: 978-3-540-75551-7\_7. URL: http://link.springer.com/10.1007/978-3-540-75551-7\_7. (visited on 04/04/2024): ▶ ◀ ● ▶ ◀ ■ ▶ ▲ ■ ▶ ▼

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- **③** Clients can now create **psuedonyms** or fake identities  $id \rightarrow (\mathcal{H}(id))^r$ ,  $\mathcal{H}: \mathcal{ID} \rightarrow \mathcal{G}, r \in \mathbb{Z}_q$ .



• Attributes: Labels associated with a user that describe them fully, such as role of a user.

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<sup>&</sup>lt;sup>4</sup>Jan Camenisch et al. Zone Encryption with Anonymous Authentication for V2V Communication. 2020. URL: https://eprint.iacr.org/2020/043 (visited on 02/04/2024). preprint.

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- For V2X
  - Anonymous credentials issued to vehicles regularly.
  - We use DGSA (Digital Group Signatures with Attributes)<sup>4</sup>, which gives us a **randomizable** group element as the credential  $\sigma \to \sigma^r$ ,  $r \in \mathbb{Z}_a$ .

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## Proposed Message Flow Diagram

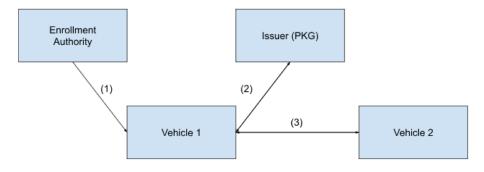


Figure 2: Message flow of the proposed scheme.

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  - DGSA credentials guarantee authenticity.
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- Vehicles exchange DGSA-signed randomized psuedonyms to generate shared key for futher communication.
  - Used in verifying legitimacy of the other party.



# **Analysis**

#### Advantages

- Fully anonymous communication, unlimited privacy between communicating parties.
- Third parties cannot identify who is communicating.
- Useful for sending extremely sensitive data.
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#### Oisadvantages

- Lots of pairing computations, for DGSA and for anonymous key agreement. Incurs computational overheads.
- Works for single-hop connections only.
- May not be scalable to communicating with many vehicles simultaneously in terms of storage overhead.

#### **Future Work**

Encrypt V2X messages like CAMs.

<sup>&</sup>lt;sup>5</sup>Camenisch et al., Zone Encryption with Anonymous Authentication for V2V Communication.

#### **Future Work**

- Encrypt V2X messages like CAMs.
- Improve efficiency of the present work.
  - Use one of DGSA or anonymous key agreement, but not both?

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<sup>&</sup>lt;sup>5</sup>Camenisch et al., Zone Encryption with Anonymous Authentication for V2V Communication.

<sup>&</sup>lt;sup>6</sup>Yue et al., "A Practical Privacy-Preserving Communication Scheme for CAMs in C=ITS": → → ↓ ≥ → ↓ ≥ → ↓ ≥ → ○ ↓ ○

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- 3 A new workflow for encryption using zones<sup>5</sup> and zone managers<sup>6</sup>

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