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MILANO 1863

DIPARTIMENTO DI ELETTRONICA
INFORMAZIONE E BIOINGEGNERIA

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Automotive Security

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Program of these two lessons:

- Introduction to the automotive context
- Automotive (CAN) Attacks
- Automotive (CAN) Countermeasures
- Secure Positioning Systems
- Exercises

Automotive Ecosystem Players:

Automotive Manufacturers

Original Equipment Manufacturers

Processor/Semiconductor/Board Companies

Regulations and State requirements

Fleet management companies (Uber, Taxis, Truck hauling companies, Postal services, Rental etc...)

All these players have a role in the development of new vehicles

Services



UBER



zipcar.
wheels when you want them



TESLOOP



RideCell

Cars



Public Transportation



Commercial



Air Taxis



Security

ARGUS
CYBER SECURITY

Arilou

Karamba
Security

CYMICTIVE
TECHNOLOGIES
Cyber security on the move

Symantec™

Data Analytics

IBM

Microsoft

Airbiquity®

QUALCOMM

Sas

Hortonworks

CISCO™

Software

MOBILEYE®

TESLA

NVIDIA®

Baidu 百度

WAYMO

Carnegie
Mellon
University

Hardware and Peripherals



BOSCH

Continental



DELPHI

QUALCOMM® QUANERGY

NVIDIA®

Valeo

Velodyne

LeddarTech®

intel

Google

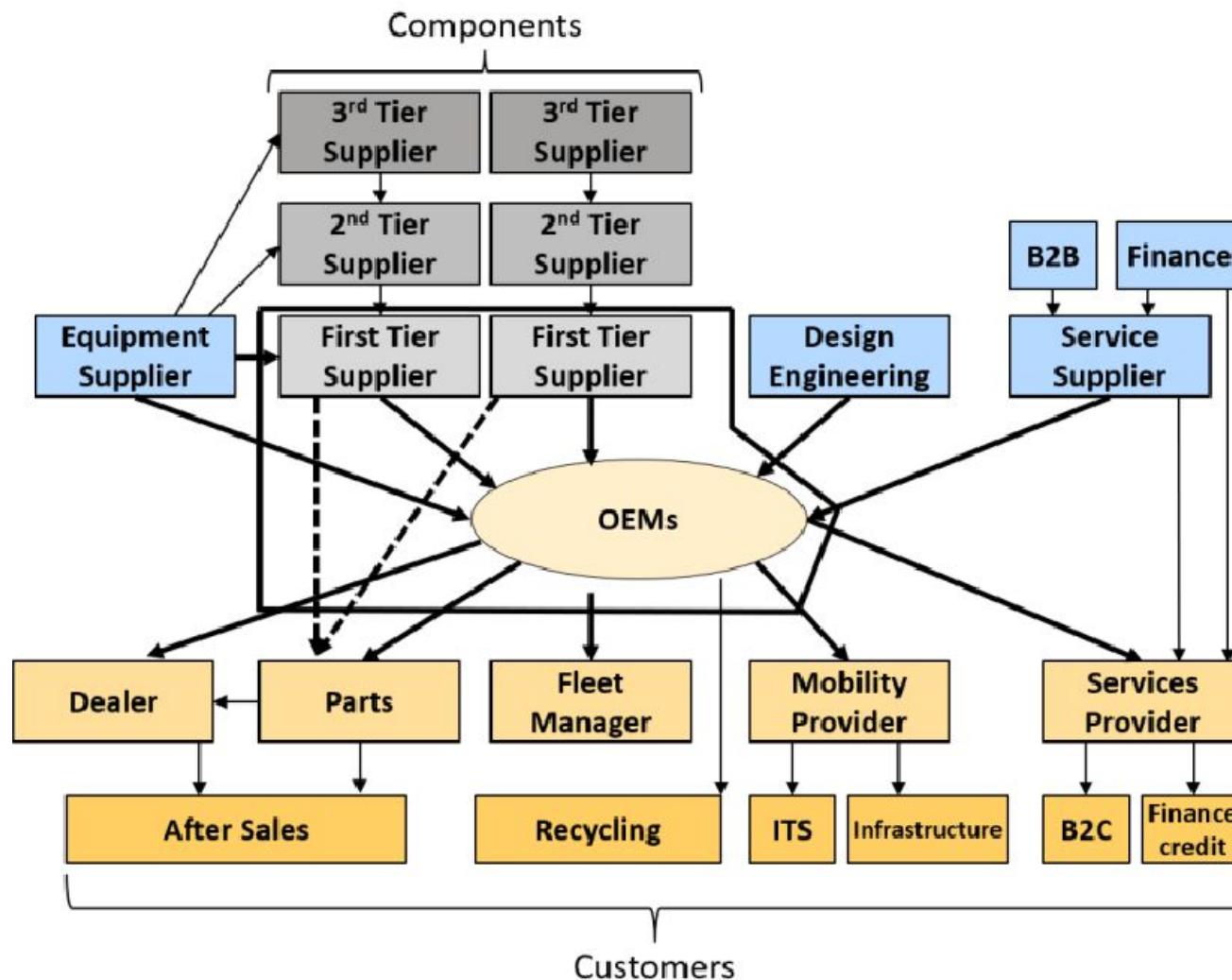


QNX™

AUTOMOTIVE
GRADE LINUX

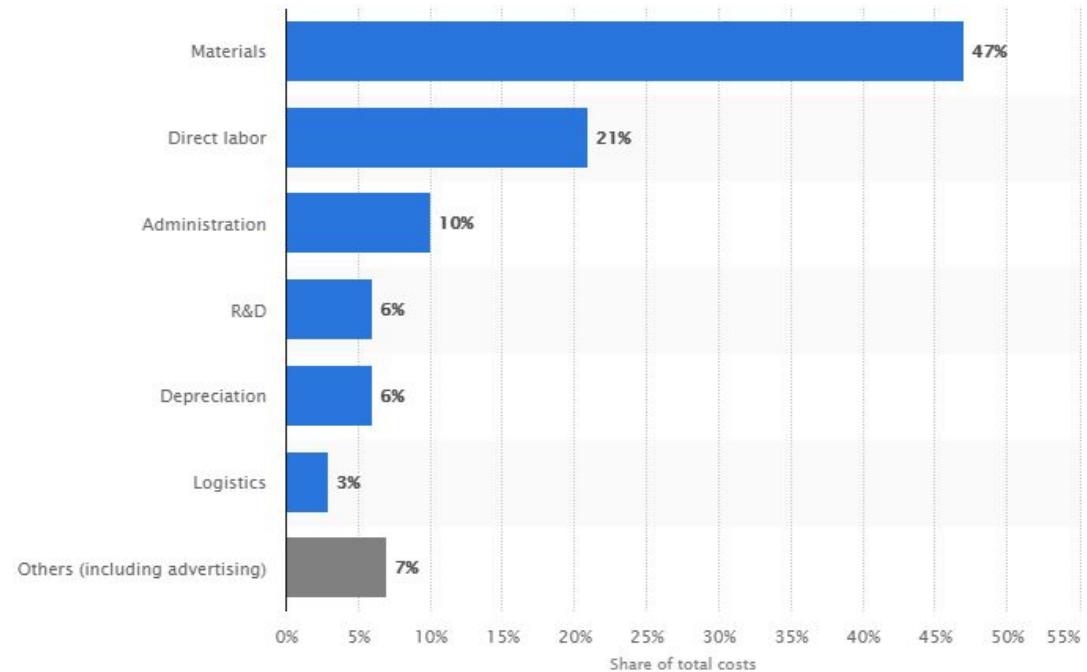
HARMAN™

Connected Car Features



Costs and profits

Profit per vehicle: around 3k\$



Details: Worldwide; 2015

<https://www.statista.com/statistics/744910/cost-breakdown-of-car-production-by-segment/>
<https://www.zaginvestor.com/tesla-profit-per-vehicle/>

Latest Vehicle Advancements: Safety

Pedestrian Detection and braking systems



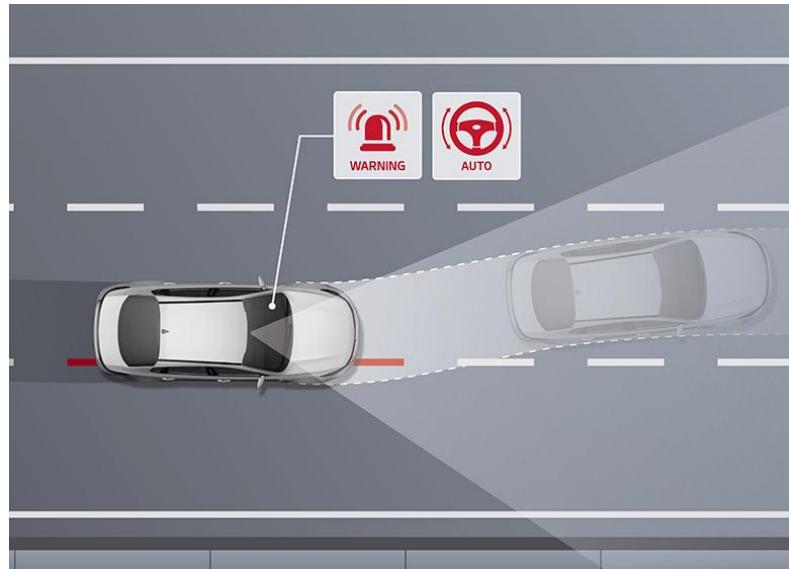
Latest Vehicle Advancements: Safety

Drowsiness detection



Latest Vehicle Advancements: Safety

Lane Assist



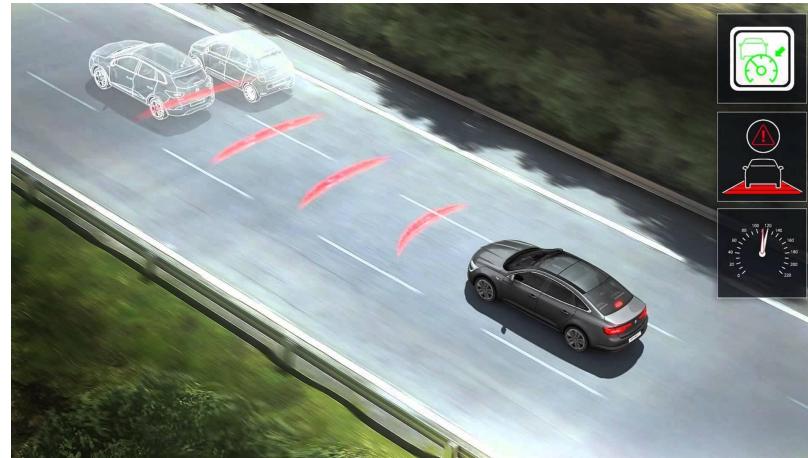
Latest Vehicle Advancements: Safety

Emergency Call



Latest Vehicle Advancements: Comfort

Adaptive Cruise Control



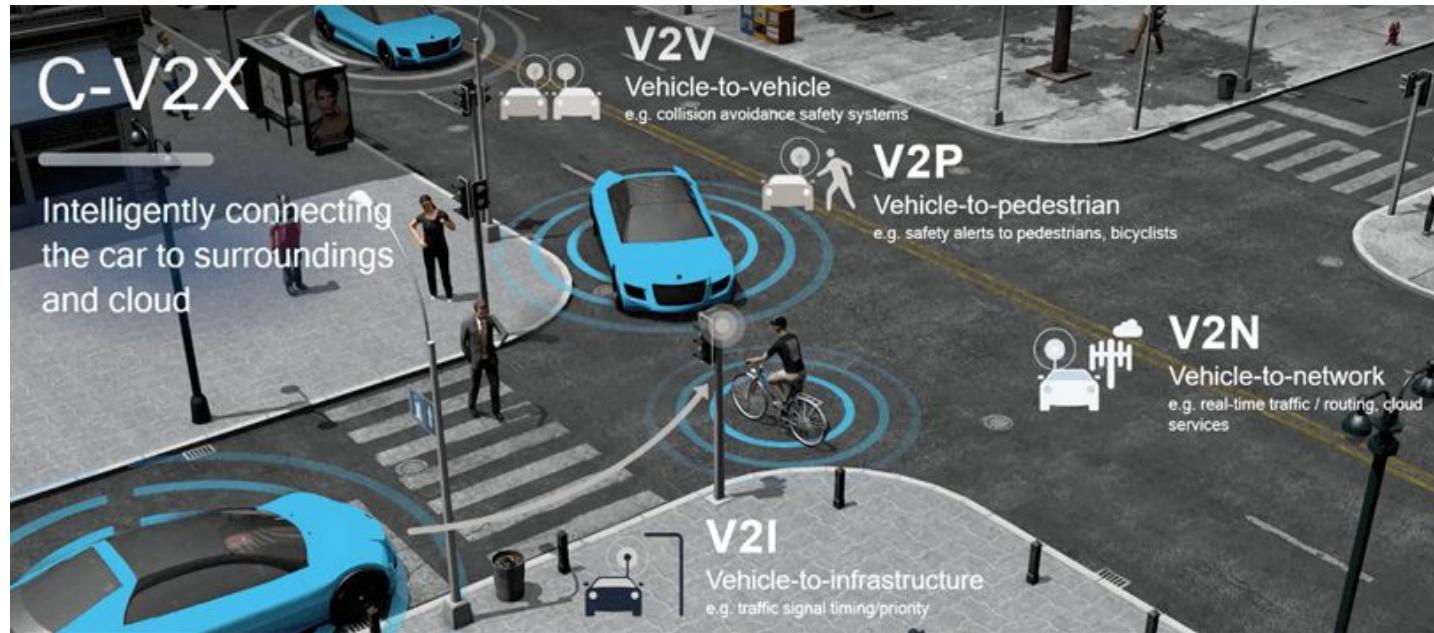
Latest Vehicle Advancements: Comfort

Personalized adjustable settings



Latest Vehicle Advancements: Communication

V2X



Latest Vehicle Advancements: Communication

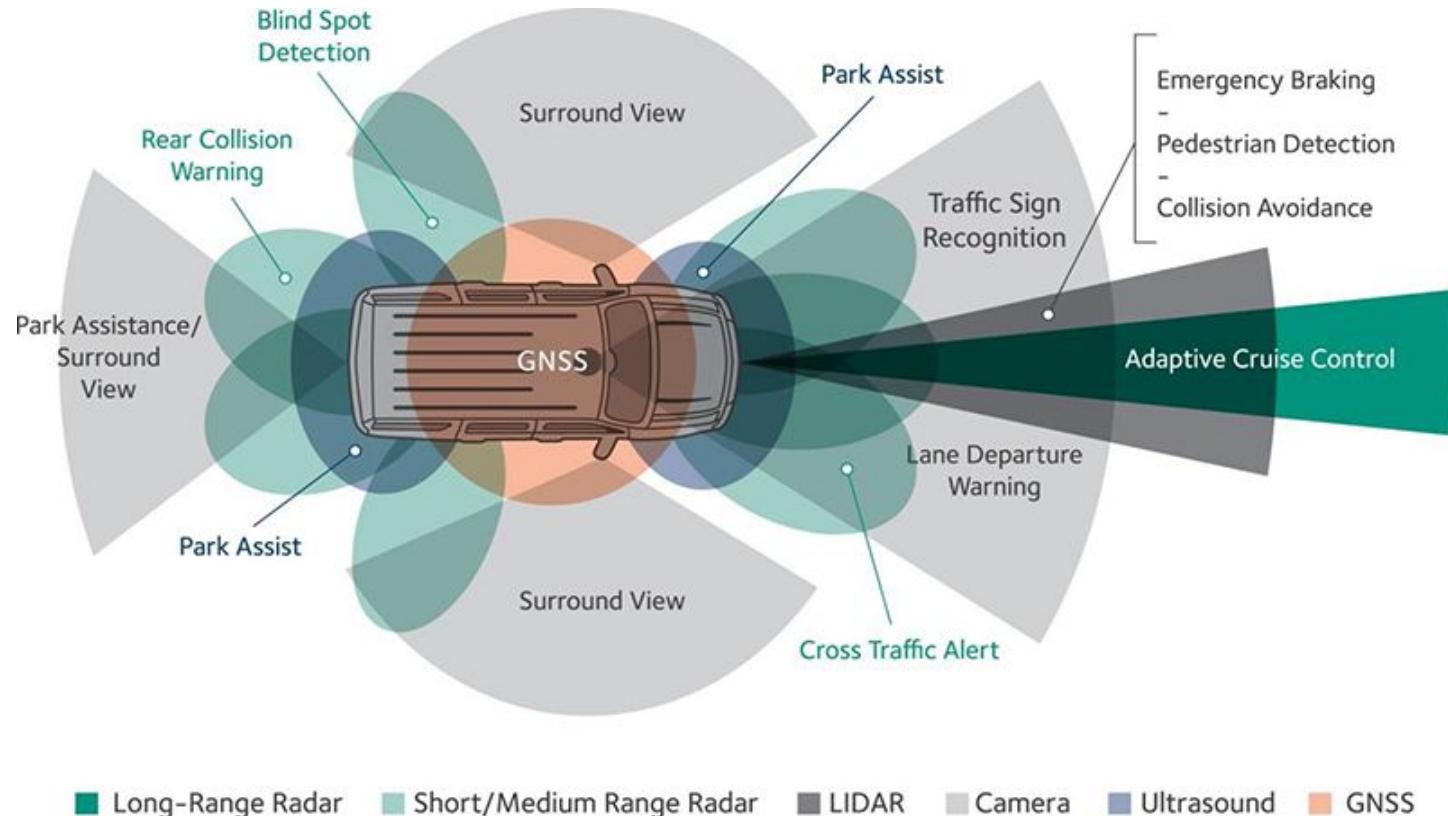
Aftermarket Devices



Future Vehicle Advancements

	SAE LEVEL 0	SAE LEVEL 1	SAE LEVEL 2	SAE LEVEL 3	SAE LEVEL 4	SAE LEVEL 5
What does the human in the driver's seat have to do?	You are driving whenever these driver support features are engaged – even if your feet are off the pedals and you are not steering	You must constantly supervise these support features; you must steer, brake or accelerate as needed to maintain safety		You are not driving when these automated driving features are engaged – even if you are seated in “the driver’s seat”	When the feature requests, you must drive	These automated driving features will not require you to take over driving
What do these features do?	These features are limited to providing warnings and momentary assistance	These features provide steering OR brake/acceleration support to the driver	These features provide steering AND brake/acceleration support to the driver	These features can drive the vehicle under limited conditions and will not operate unless all required conditions are met	These are driver support features	These are automated driving features
Example Features	<ul style="list-style-type: none">automatic emergency brakingblind spot warninglane departure warning	<ul style="list-style-type: none">lane centering ORadaptive cruise control	<ul style="list-style-type: none">lane centering ANDadaptive cruise control at the same time	<ul style="list-style-type: none">traffic jam chauffeur	<ul style="list-style-type: none">local driverless taxipedals/steering wheel may or may not be installed	<ul style="list-style-type: none">same as level 4, but feature can drive everywhere in all conditions

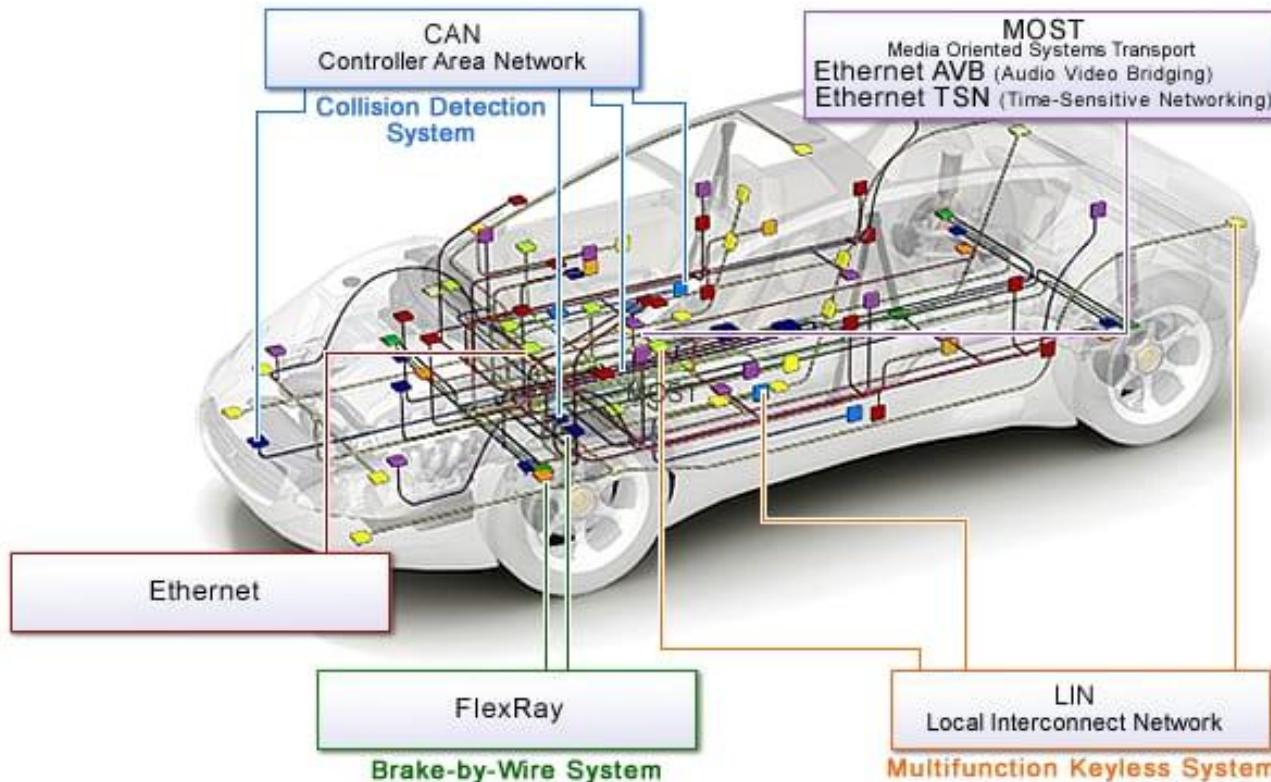
Technologies to get there: External Sensors



Technologies to get there: External Sensors

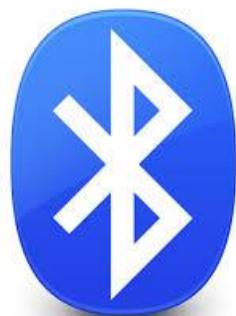


Technologies to get there: On-board networks



Technologies to get there: Communication Networks

Local Communication



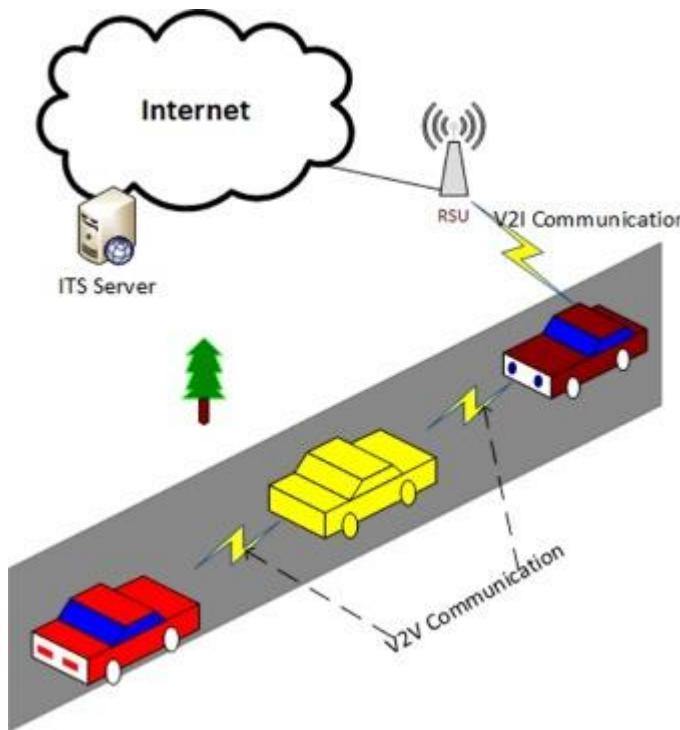
Remote Communication



DSRC



Technologies to get there: Communication Networks



Sending and receiving traffic data

Sending status updates

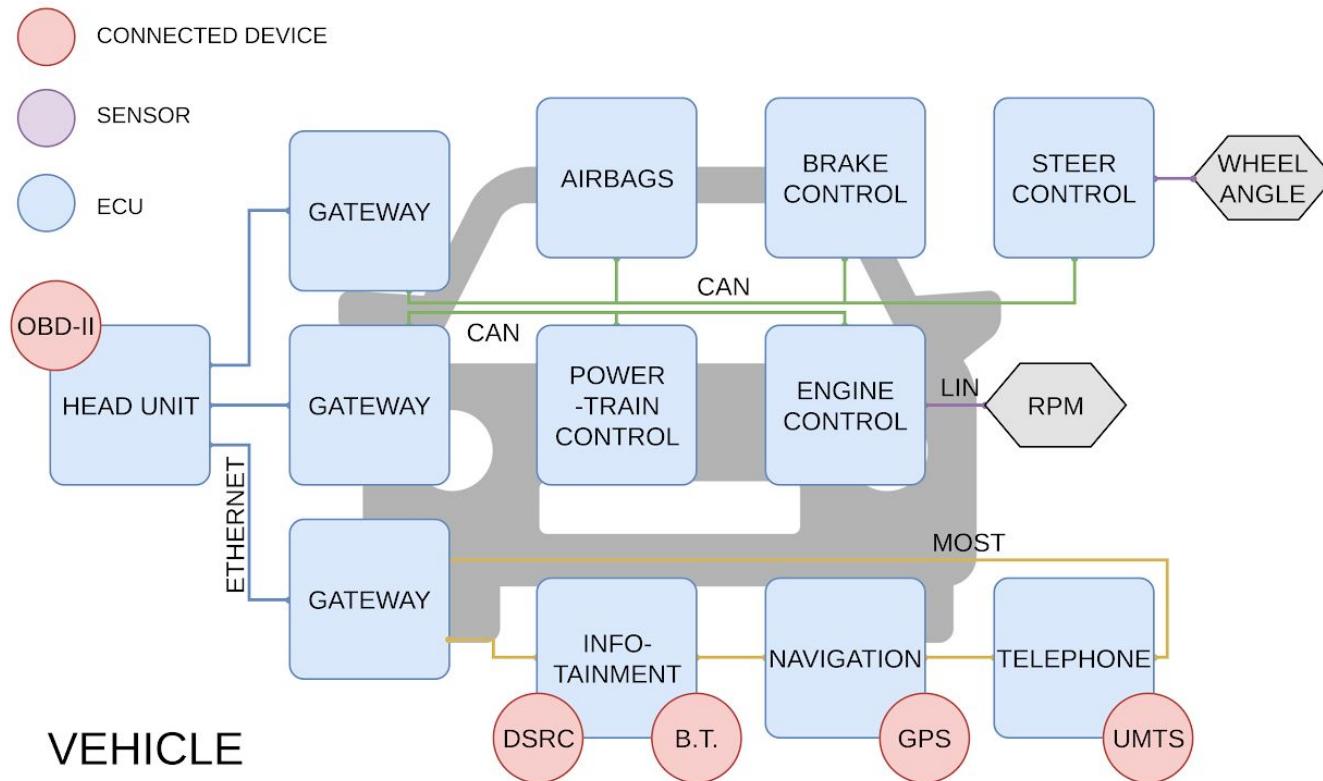
Receiving software updates

Phone calls

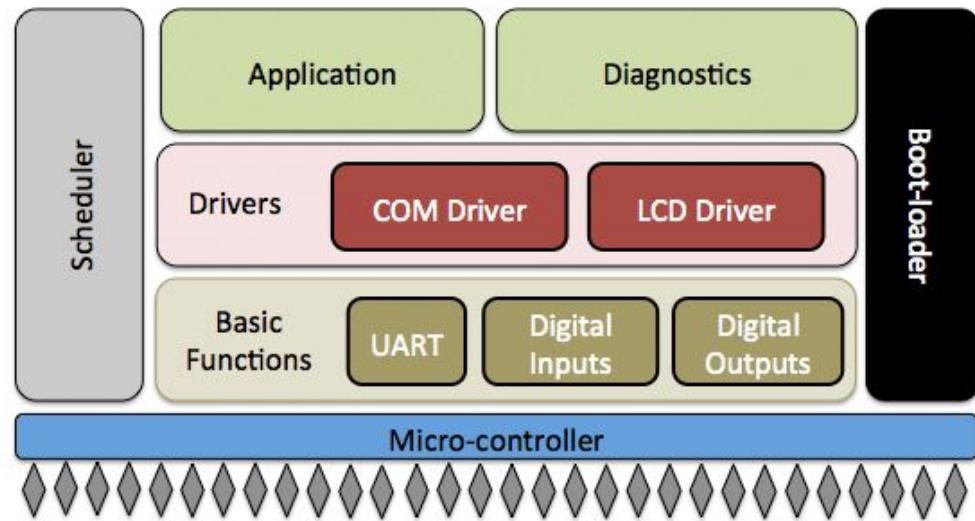
Internet navigation

...

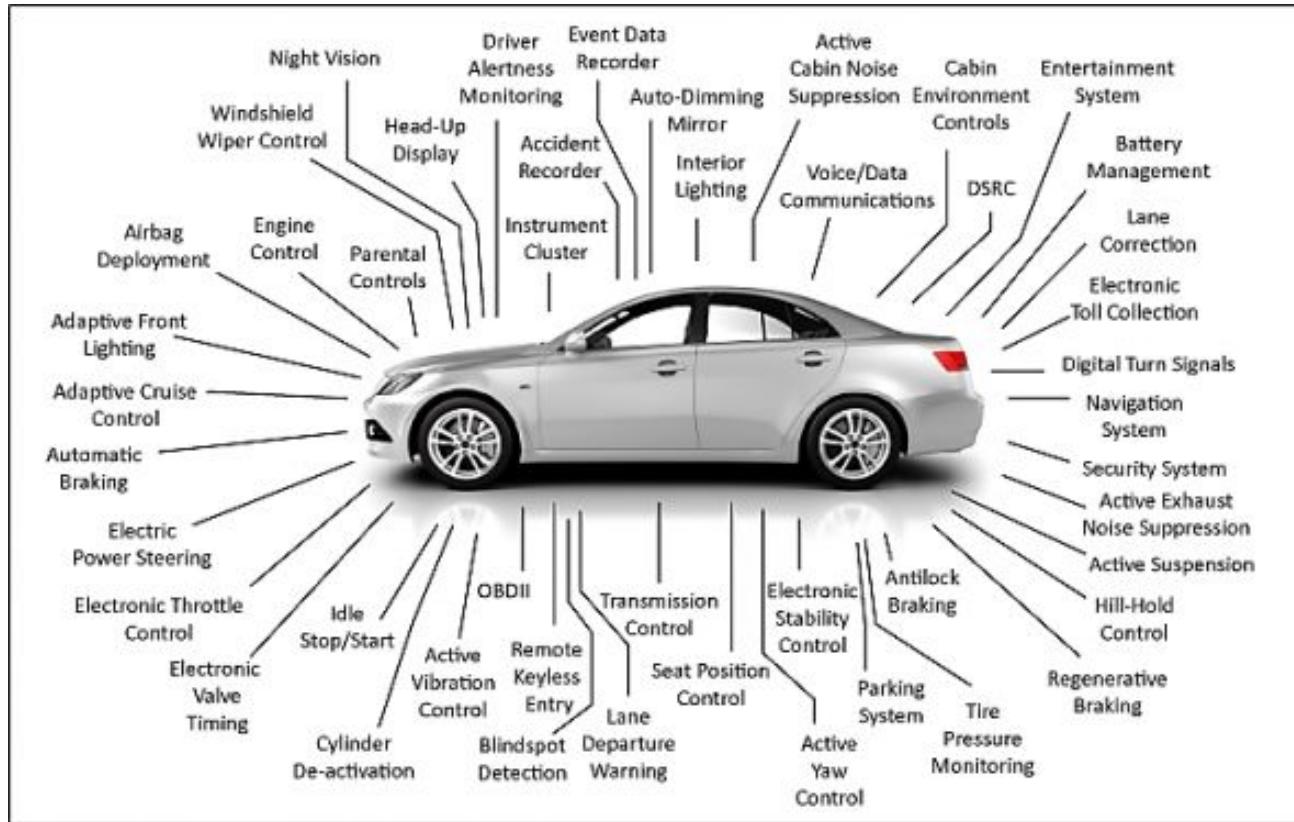
Technologies to get there: On-board networks

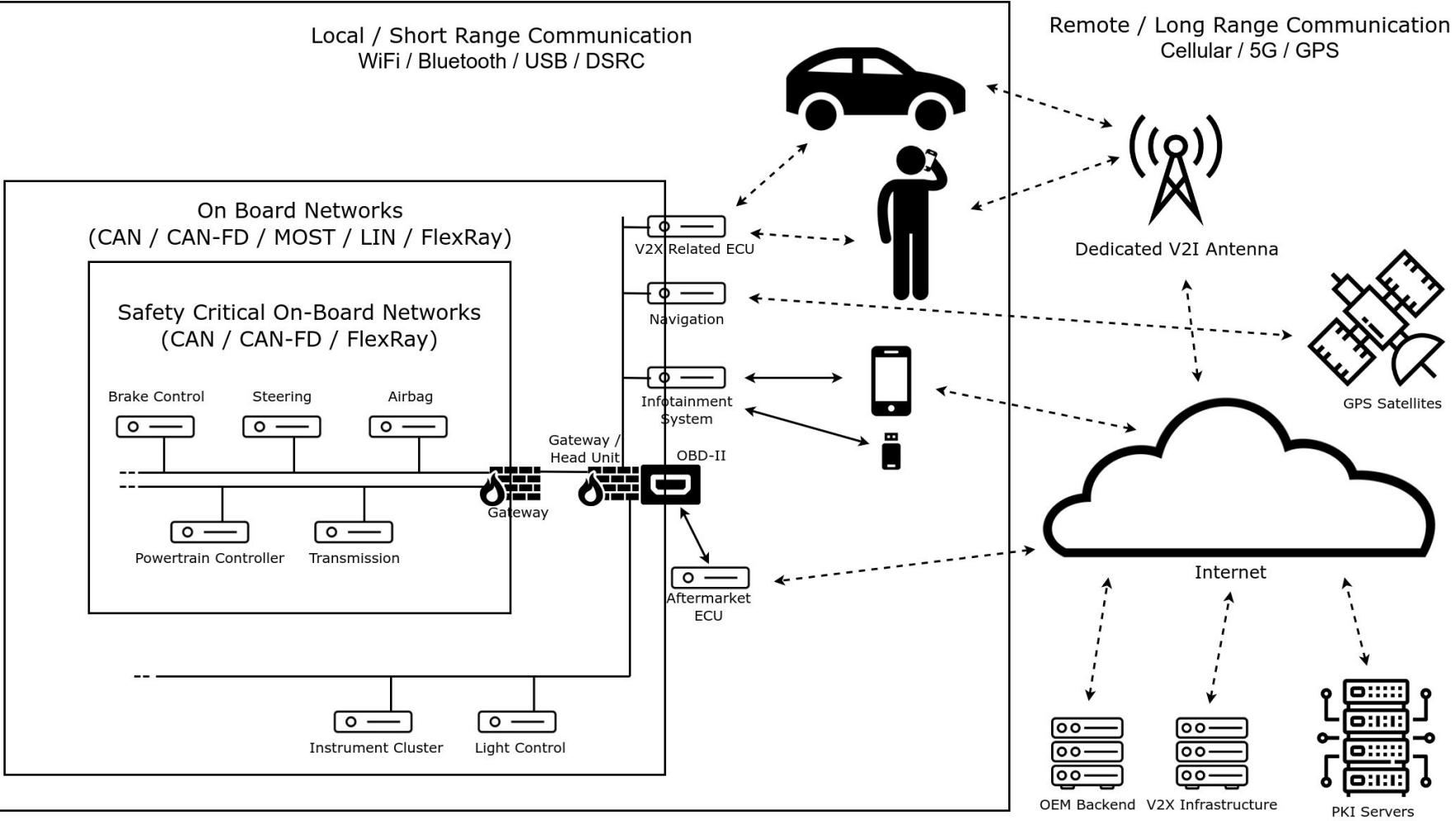


Technologies to get there: ECUs



Overview of intra vehicle devices







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CAN NETWORKS



Most common on-board subnetworks

CAN - Controller Area Network: Standard, Real-Time, Cheap

CAN-FD - “faster” CAN, basically identical, carries more data

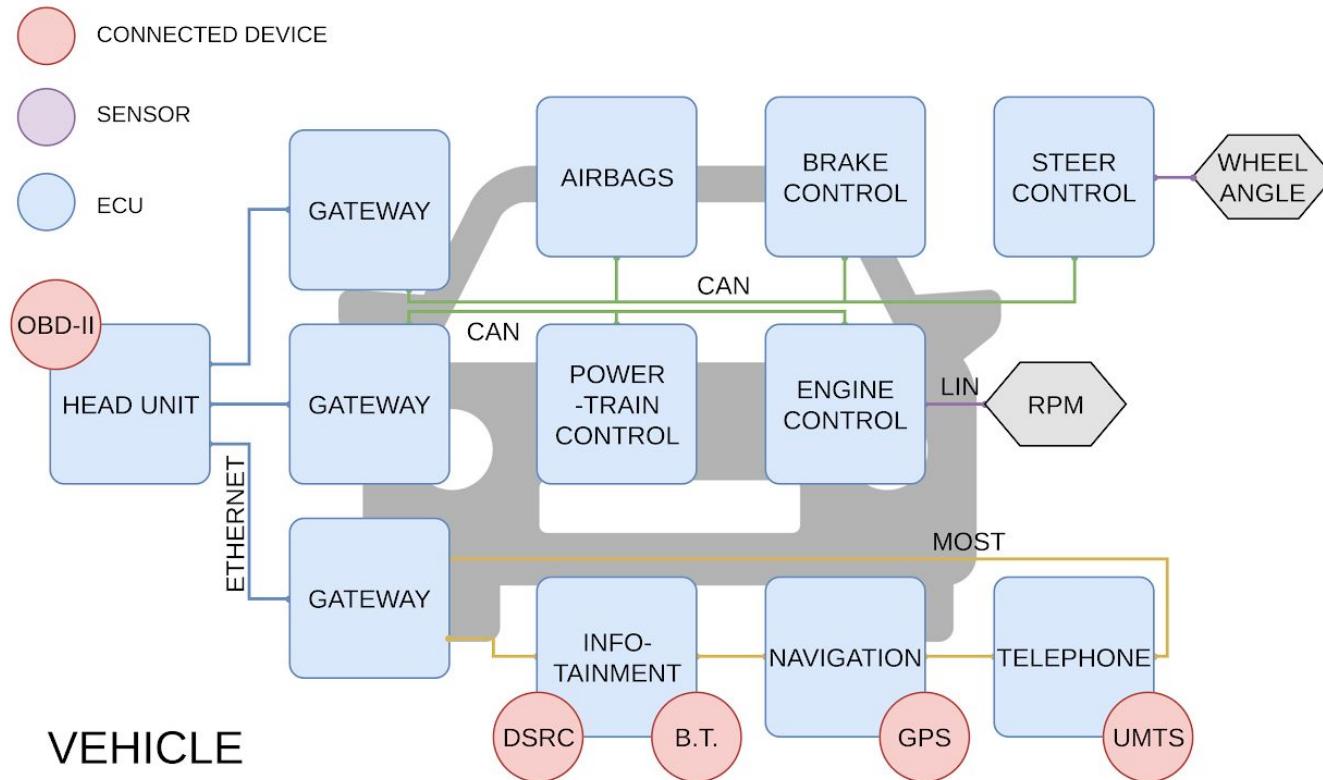
FlexRay: Expensive, Real-Time

MOST - Media Oriented System Network: High Bandwidth

LIN - Local Interconnect Network: Cheap, for sensors

Automotive Ethernet: Higher speed than CAN, already used in other contexts

Most common on-board subnetworks

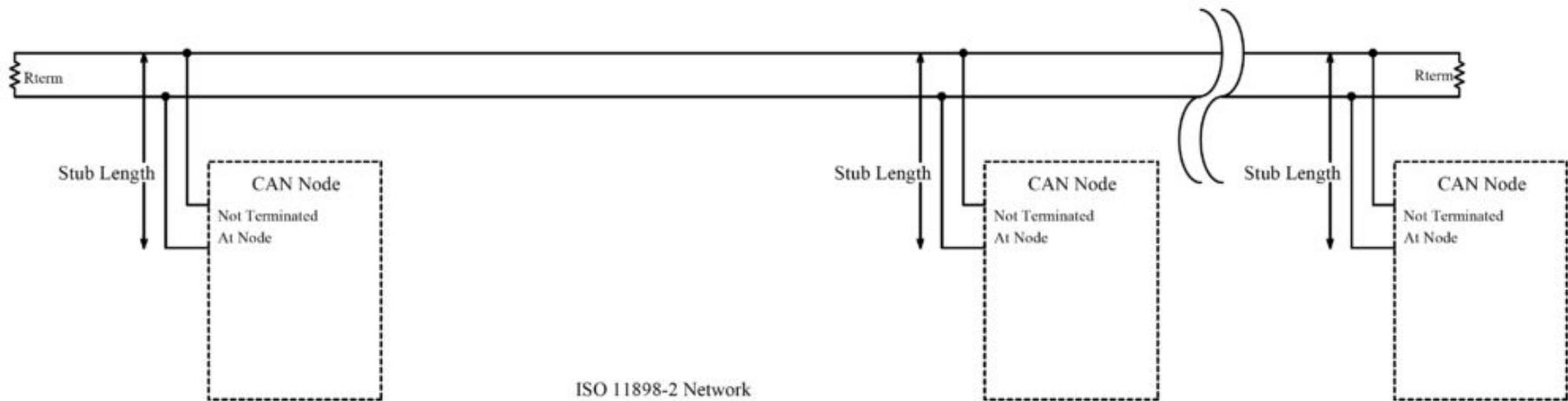


Controller Area Network

- Developed by Bosch in 1980's
- Current de-facto standard
- Data Link and Physical layers
- Developed with focus on safety:
 - No issues with electromagnetic interferences
 - Broadcast nature
 - Arbitration focused on favouring most important messages
- Max network length of 40 meters
- Max baud rate 1Mbps (usually 500/250kbps)

Controller Area Network

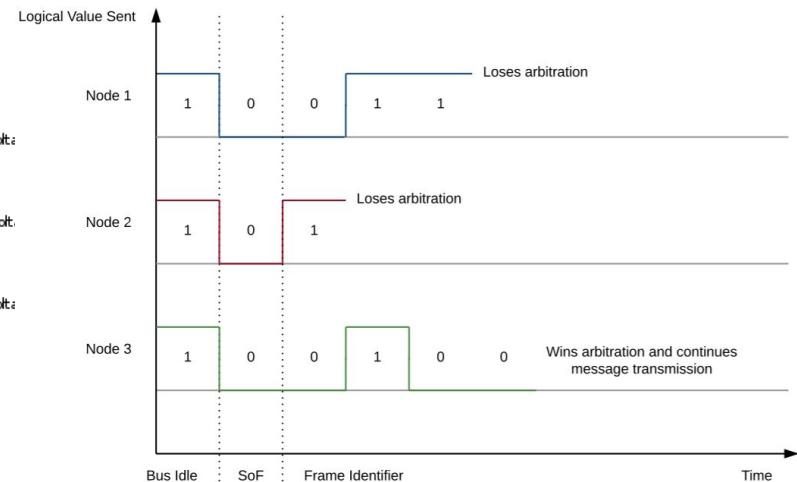
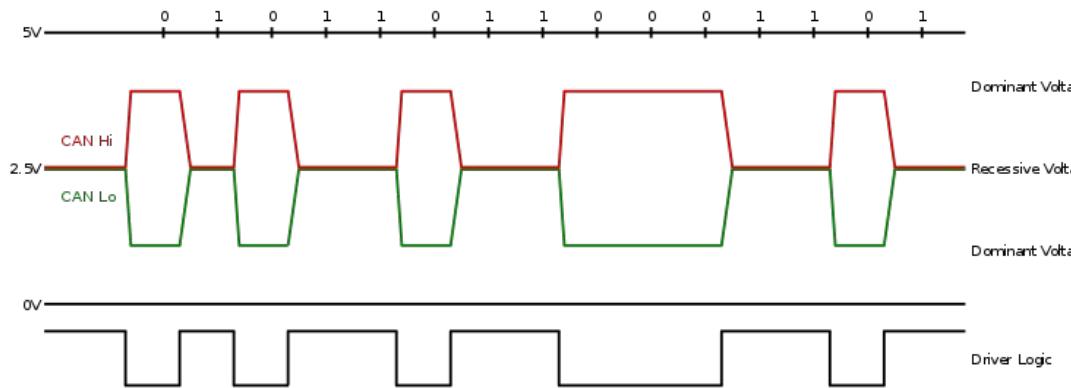
- Multi Master
- Broadcast
- Bus topology



Controller Area Network

Physical layer: Differential signaling over two wires

Data Link Layer: CSMA/BA -> This enables Real-Time



Controller Area Network: Frames

- Data Frames
- Remote Frames
- Error Frames
- Overload Frames

Controller Area Network: Data Frames

Base Frame format: 11 bit ID (2^{11} 2048 possible IDs)

Extended Frame format: 29 bit ID ($2^{29} = 536870912$ possible IDs)

8 data bytes (64 for CAN-FD)

S O F	11-bit Identifier	S R R	I D E	18-bit Identifier	R T R	r1	r0	DLC	0...8 Bytes Data	CRC	ACK	E O F	I F S
-------------	----------------------	-------------	-------------	----------------------	-------------	----	----	-----	------------------	-----	-----	-------------	-------------

S O F	11-bit Identifier	R T R	I D E	r0	DLC	0...8 Bytes Data	CRC	ACK	E O F	I F S
-------------	----------------------	-------------	-------------	----	-----	------------------	-----	-----	-------------	-------------

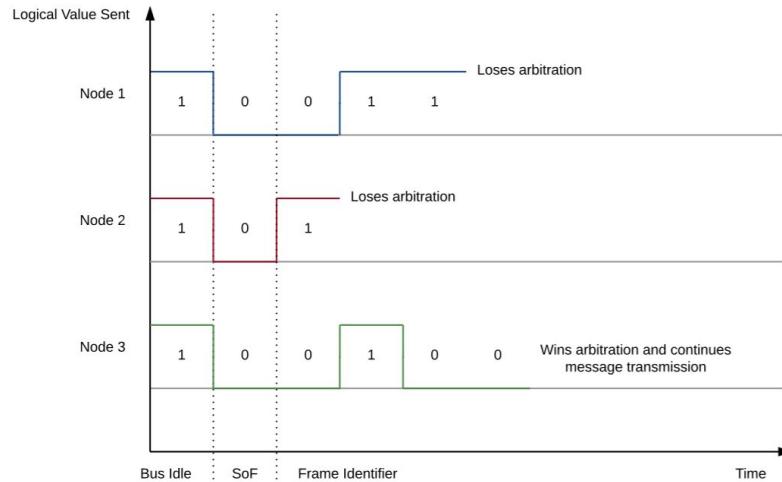
Controller Area Network: Data Frames

(Wikipedia)

Field name	Length (bits)	Purpose
Start-of-frame	1	Denotes the start of frame transmission
Identifier (green)	11	A (unique) identifier which also represents the message priority
Remote transmission request (RTR) (blue)	1	Must be dominant (0) for data frames and recessive (1) for remote request frames (see Remote Frame , below)
Identifier extension bit (IDE)	1	Must be dominant (0) for base frame format with 11-bit identifiers
Reserved bit (r0)	1	Reserved bit. Must be dominant (0), but accepted as either dominant or recessive.
Data length code (DLC) (yellow)	4	Number of bytes of data (0–8 bytes) ^[a]
Data field (red)	0–64 (0–8 bytes)	Data to be transmitted (length in bytes dictated by DLC field)
CRC	15	Cyclic redundancy check
CRC delimiter	1	Must be recessive (1)
ACK slot	1	Transmitter sends recessive (1) and any receiver can assert a dominant (0)
ACK delimiter	1	Must be recessive (1)
End-of-frame (EOF)	7	Must be recessive (1)

Controller Area Network: Data Frames

IDs enable priority as seen before



Lower ID = higher priority, hence messages with high real time requirements are associated to lower IDs

Controller Area Network: Data Frames

Each ECU “owns” a set of IDs

No other ECU can use those IDs

Data field of each ID has a fixed meaning

Each ECU “subscribes” to a list of IDs and knows the meaning of their data

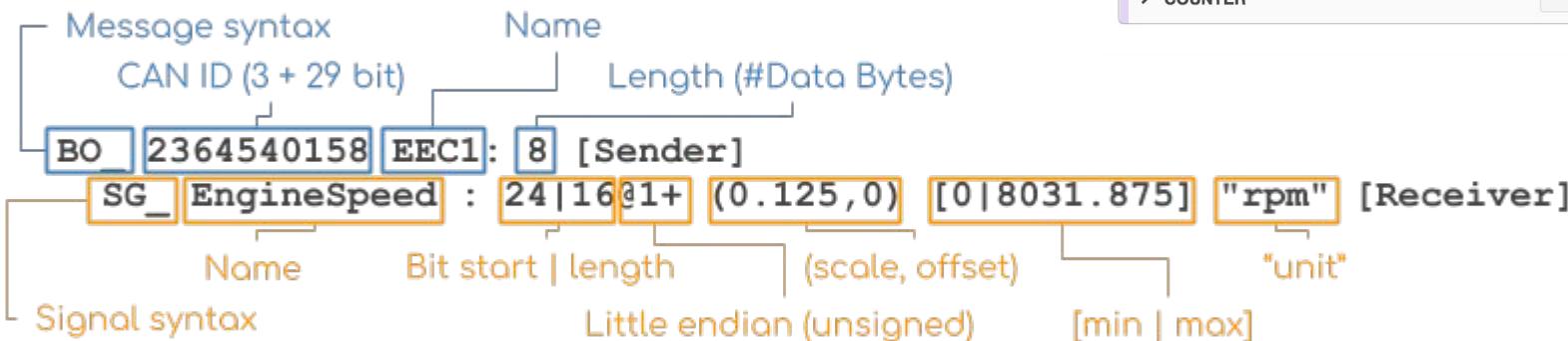
Majority of packets are periodic

▼ Edit Signals

time: 13.123

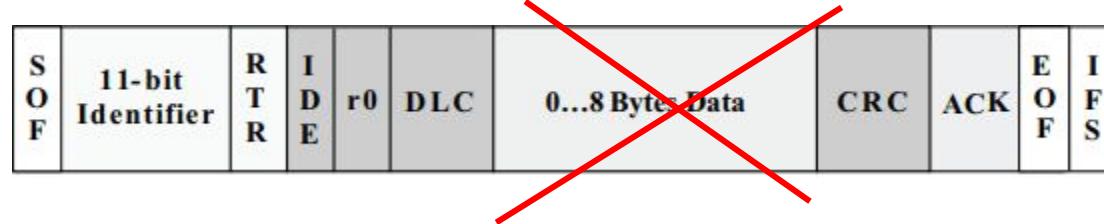
0	msb	0	0	0	0	0	0	0	0	00		
0	0	0	0	0	0	1	1	0	lsb	06		
0	msb	0	0	0	0	0	0	0	0	00		
0	0	0	0	0	0	0	0	0	lsb	00		
0	0	0	0	0	0	1	1	1	07			
0	0	1	msb	0	lsb	1	msb	1	0	1	lsb	2d

- > STEER_ANGLE [Show Plot](#)
- > STEER_ANGLE_RATE [Show Plot](#)
- > CHECKSUM [Show Plot](#)
- > COUNTER [Show Plot](#)



Controller Area Network: Remote Frames

Data frames used for requests. Not often seen.



Controller Area Network: Error Frames

6 bit time sequences of zeroes or ones. Used to say that the current data/remote frame has to be considered not valid

Reasons:

Sender ECU sent a wrong bit

Receiver ECU reads something non CAN-compliant

Conflict on the bus

Arbitration				Control				Data				Error Flag					
...	0	1	0	0	0	1	1	0	1	0	0	1	0	0	0	0	...
...	0	1	0	0	0	1	1	0	1	0	0	0	0	0	0	0	

Controller Area Network: Overload Frames

Error frames sent while bus is idle to request a pause from transmission -
not often used in newer ECUs

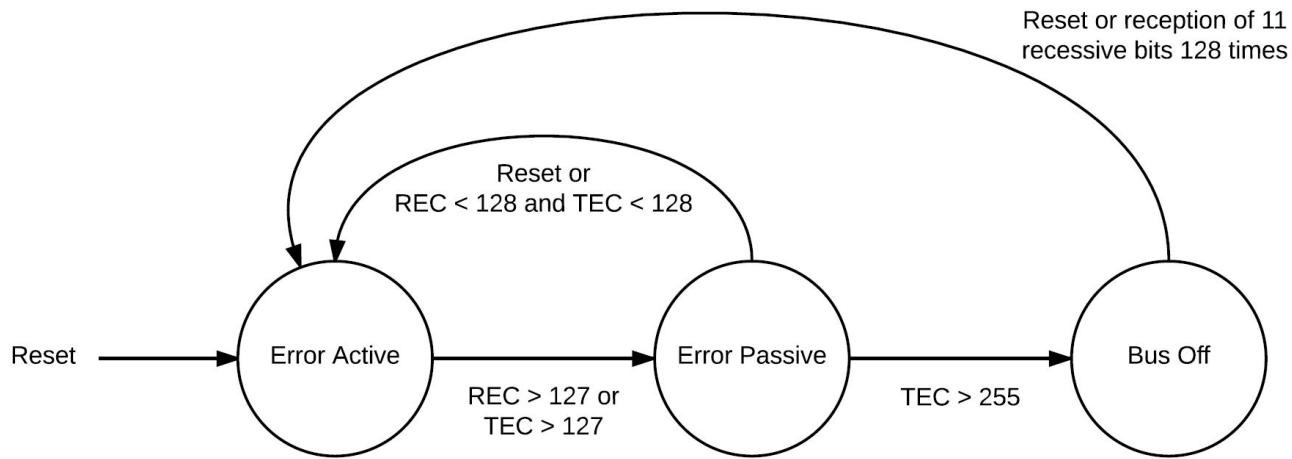
Controller Area Network: Error Handling

Bit, Stuff, CRC, Form or Acknowledgment Errors are possible

Fault Confinement:

- Units can be in Error Active, Error Passive or Bus Off state
- Two Error Counters, Transmit and Receive
- Error Counter increases by 8 every time an error occurs
- When any counter reaches 128 the unit goes into Error Passive state
- When the transmit error counter reaches 256 the units shuts down the CAN controller

Controller Area Network: Fault Confinement

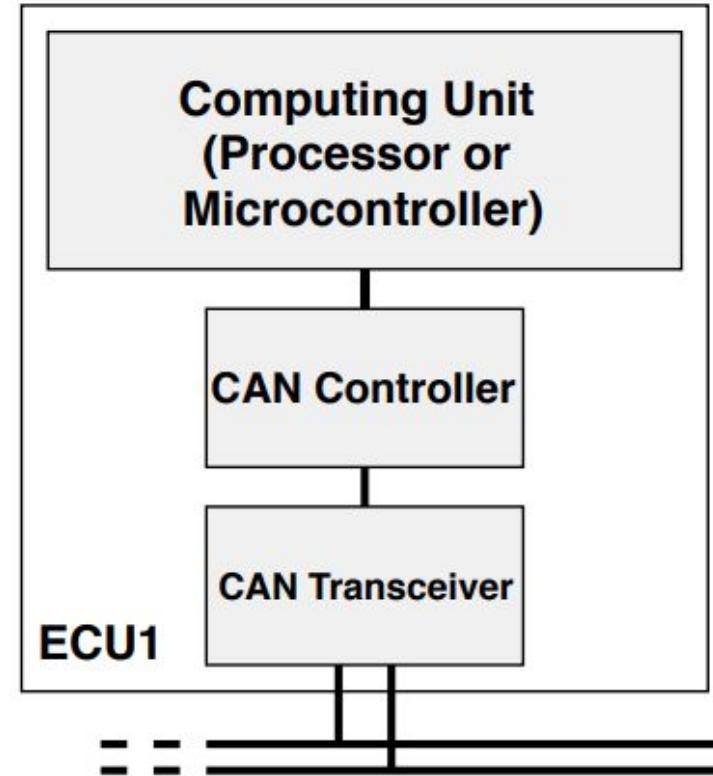


Controller Area Network

Processor is the application layer

CAN controller reads the bits and translates them in actual packets with meaning

CAN transceiver translates only electric signals in bits and vice-versa



Controller Area Network: Example Log

Only data frames (and remote, rarely). All others are “removed” by the CAN transceiver in advance.

<https://www.sciencedirect.com/science/article/pii/S2352340920300433>

CAN higher layers: OBD-II diagnostics

Mandatory in many countries,
almost always present

Enables reporting and
self-diagnostics of the whole
network



PID (hex)	PID (Dec)	Data bytes returned	Description	Min value	Max value	Units	Formula ^[a]
00	0	4	PIDs supported [01 - 20]				Bit encoded [A7..D0] == [PID \$01..PID \$20] See below
01	1	4	Monitor status since DTCs cleared. (Includes malfunction indicator lamp (MIL) status and number of DTCs.)				Bit encoded. See below
02	2	2	Freeze DTC				
03	3	2	Fuel system status				Bit encoded. See below
04	4	1	Calculated engine load	0	100	%	$\frac{100}{255} A \text{ (or } \frac{A}{2.55})$
05	5	1	Engine coolant temperature	-40	215	°C	$A - 40$
06	6	1	Short term fuel trim—Bank 1	-100	99.2 (Add Fuel: Too Lean)	%	$\frac{100}{128} A - 100$ (or $\frac{A}{1.28} - 100$)
07	7	1	Long term fuel trim—Bank 1	(Reduce Fuel: Too Rich)			
08	8	1	Short term fuel trim—Bank 2				
09	9	1	Long term fuel trim—Bank 2				
0A	10	1	Fuel pressure (gauge pressure)	0	765	kPa	$3A$
0B	11	1	Intake manifold absolute pressure	0	255	kPa	A
0C	12	2	Engine speed	0	16,383.75	rpm	$\frac{256A + B}{4}$
0D	13	1	Vehicle speed	0	255	km/h	A
0E	14	1	Timing advance	-64	63.5	° before TDC	$\frac{A}{2} - 64$
0F	15	1	Intake air temperature	-40	215	°C	$A - 40$
10	16	2	Mass air flow sensor (MAF) air flow rate	0	655.35	grams/sec	$\frac{256A + B}{100}$
11	17	1	Throttle position	0	100	%	$\frac{100}{255} A$

CAN higher layers: ISO-TP

Enables Multi-packet Messaging

Used for OBD-II diagnostics too

Type	Code	Description
Single frame	0	The single frame transferred contains the complete payload of up to 7 bytes (normal addressing) or 6 bytes (extended addressing)
First frame	1	The first frame of a longer multi-frame message packet, used when more than 6/7 bytes of data segmented must be communicated. The first frame contains the length of the full packet, along with the initial data.
Consecutive frame	2	A frame containing subsequent data for a multi-frame packet
Flow control frame	3	the response from the receiver, acknowledging a First-frame segment. It lays down the parameters for the transmission of further consecutive frames.
	4..15	Reserved

CAN higher layers: J1939

Higher Layer Protocol that uses CAN as a transfer protocol.

Used mainly by “bigger” vehicles (Agricultural, Buses, Ships, Military...)

Its main scope is to enable messaging longer than 8 bytes (i.e. sent in multiple CAN packets)

CAN higher layers: J1939

J1939 message (PGN & SPNs)



CAN higher layers: J1939

PGN61444 - Electronic Engine Controller 1 - EEC1

Transmission Repetition Rate: Engine speed dependent

Data Length: 8 bytes

Data Page: 0

PDU Format: 240

PDU Specific: 4

Default Priority: 3

Parameter Group Number: 61444 (0x00F004)

Bit Start/Byte	Length	SPN ID	SPN Description
1.1	4 bits	899	Engine Torque Mode
2	1 byte	512	Driver's Demand Engine - % Torque
3	1 byte	513	Actual Engine - Percent Torque
4-5	2 bytes	190	Engine Speed
6	1 byte	1483	SA of Controlling Device for Engine Control
7.1	4 bits	1675	Engine Starter Mode
8	1 byte	2432	Engine Demand - Percent Torque

CAN higher layers: J1939

CAN ID	Data bytes	PGN	61444
0CF00401	FF FF FF 68 13 FF FF FF	SPN	190
		Bit start	24
		Bit length	16
		Resolution	0.125 rpm/bit
		Offset	0
		Min	0
		Max	8031.875

CAN higher layers: J1939

Messages can be also multi-packet:

In the first packet we receive the layout (PGN + number of successive packets)

Each successive packet has the first data byte used to count and the other 7 carry data.

CAN higher layers: Unified Diagnostic Services

Diagnostic communication protocol widely used to ask specific ECUs to answer diagnostic questions, make diagnostic tests, and similar.

The majority of the “requests” are accessible only once authenticated with the ECU



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AUTOMOTIVE ATTACKS

Automotive Attacks: History

- Hoppe et al 2008 - Tire pressure monitoring
- Koscher et al. 2010 - On-board attacks
- Checkoway et al 2011 - Remote attack surfaces
- Miller & Valasek 2013/15/16 - Actual attacks



Automotive Attacks: History



Automotive Attacks: History



Automotive Attacks

Risks

Safety

Privacy

Financial

Goals (CIA triad)

Confidentiality

Integrity

Availability

Confidentiality

Information is not made available or disclosed to unauthorized individuals, entities, or processes

Integrity

Assuring the accuracy and completeness of data over its entire lifecycle. This means that data cannot be modified in an unauthorized or undetected manner.

Availability

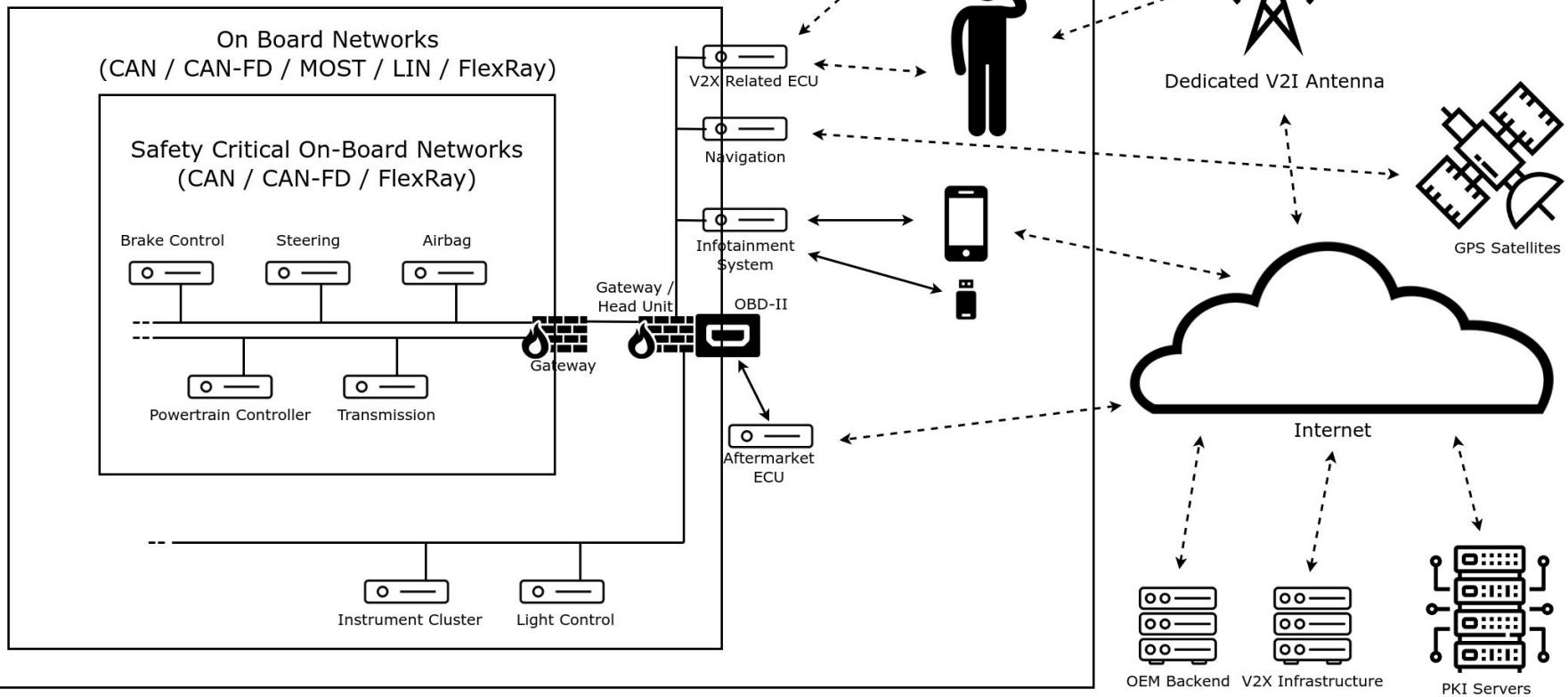
For any information system to serve its purpose, the information must be available when it is needed.

Attack Goals

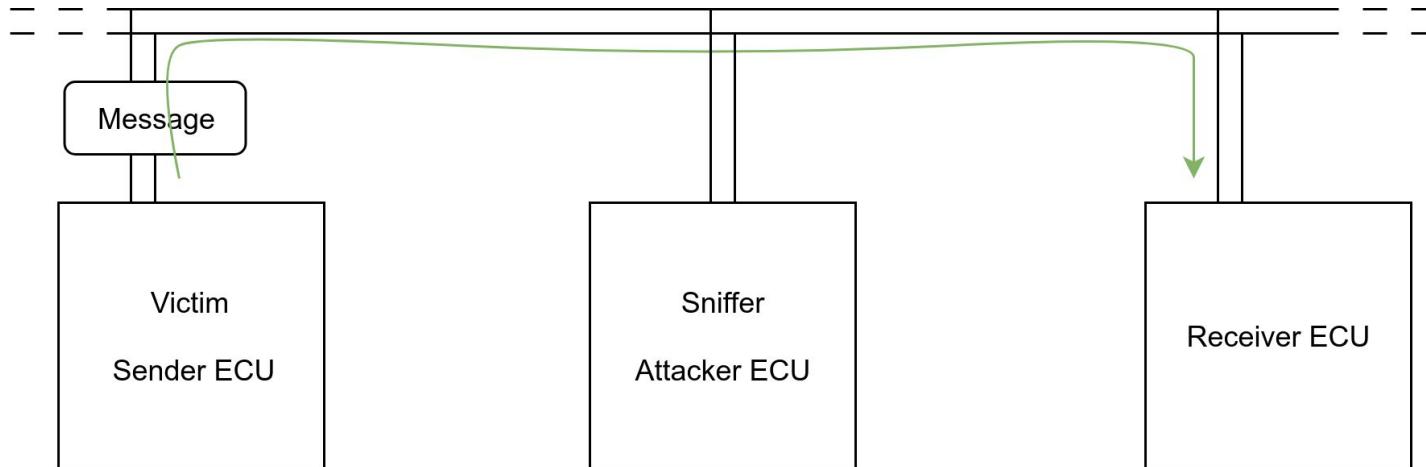
- Confidentiality → Sniffing
- Integrity → Spoofing
- Availability → Denial of Service

Local / Short Range Communication
WiFi / Bluetooth / USB / DSRC

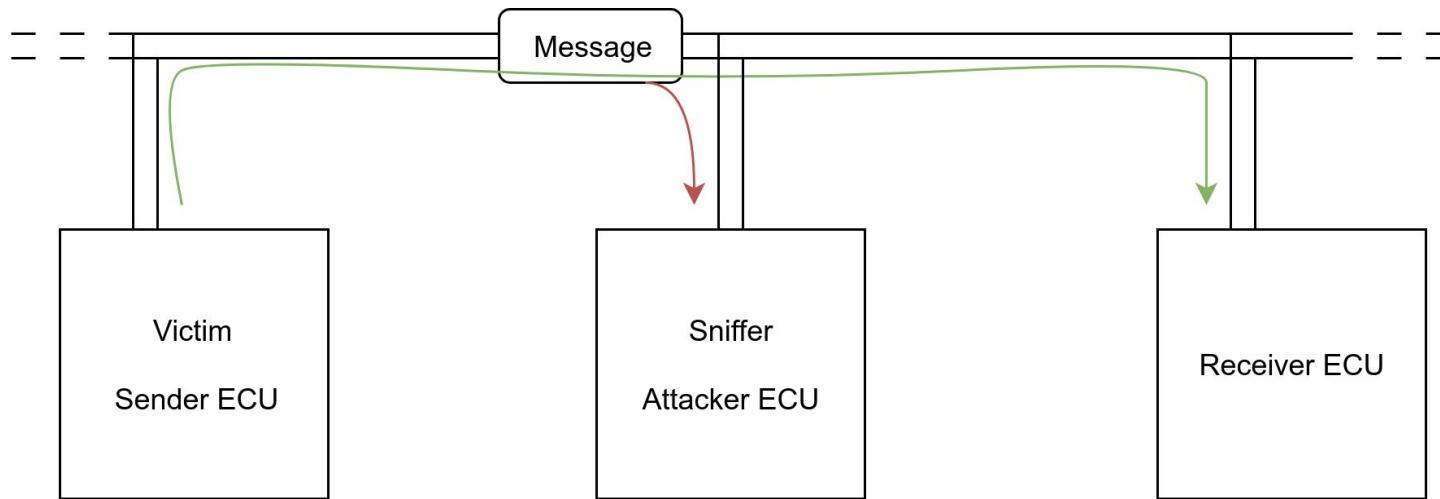
All three goals are available at different levels



CAN Sniffing



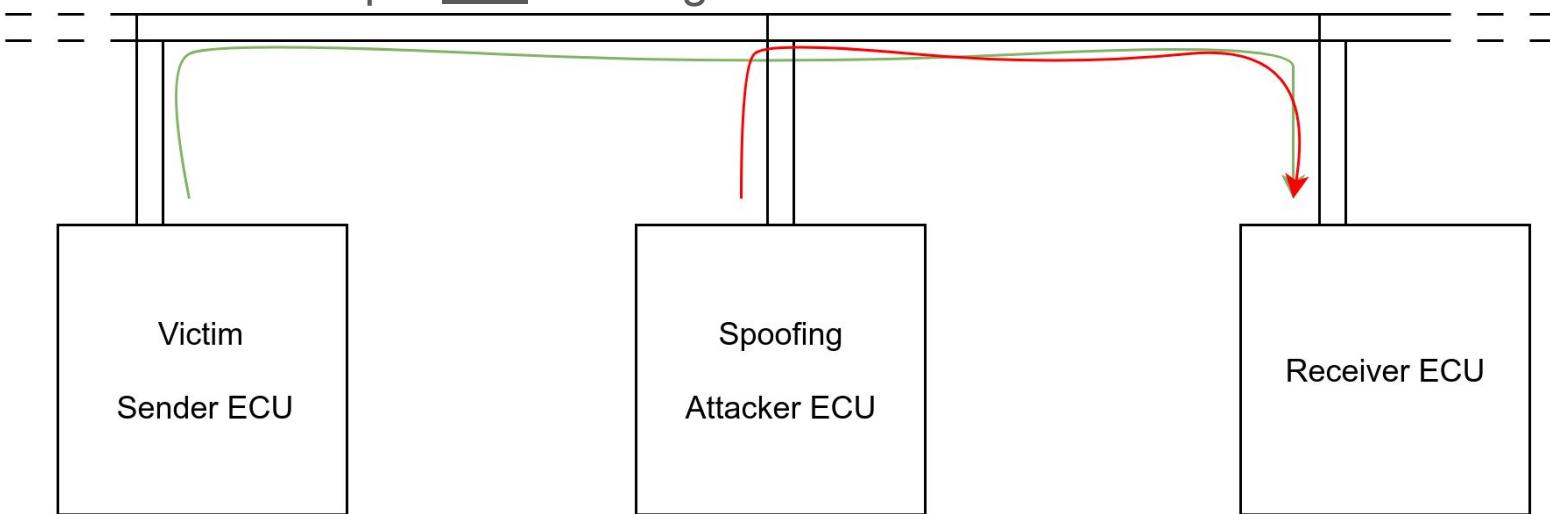
CAN Sniffing



CAN Spoofing

Pretty straightforward, sending messages with any ID is accepted in CAN. Even ones owned by other ECUs.

A basic receiver accepts ALL messages.



CAN Spoofing

Common spoofing attacks:

1) Masquerade

2) Replay

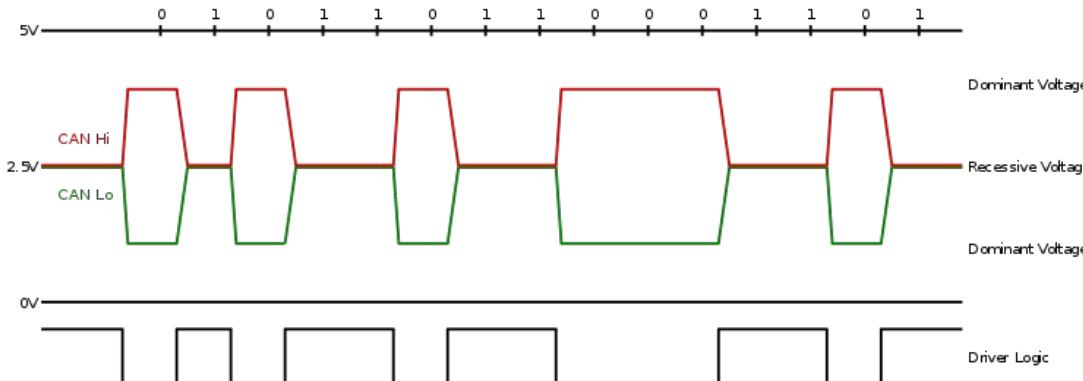
3) Fuzzing

CAN Denial of Service

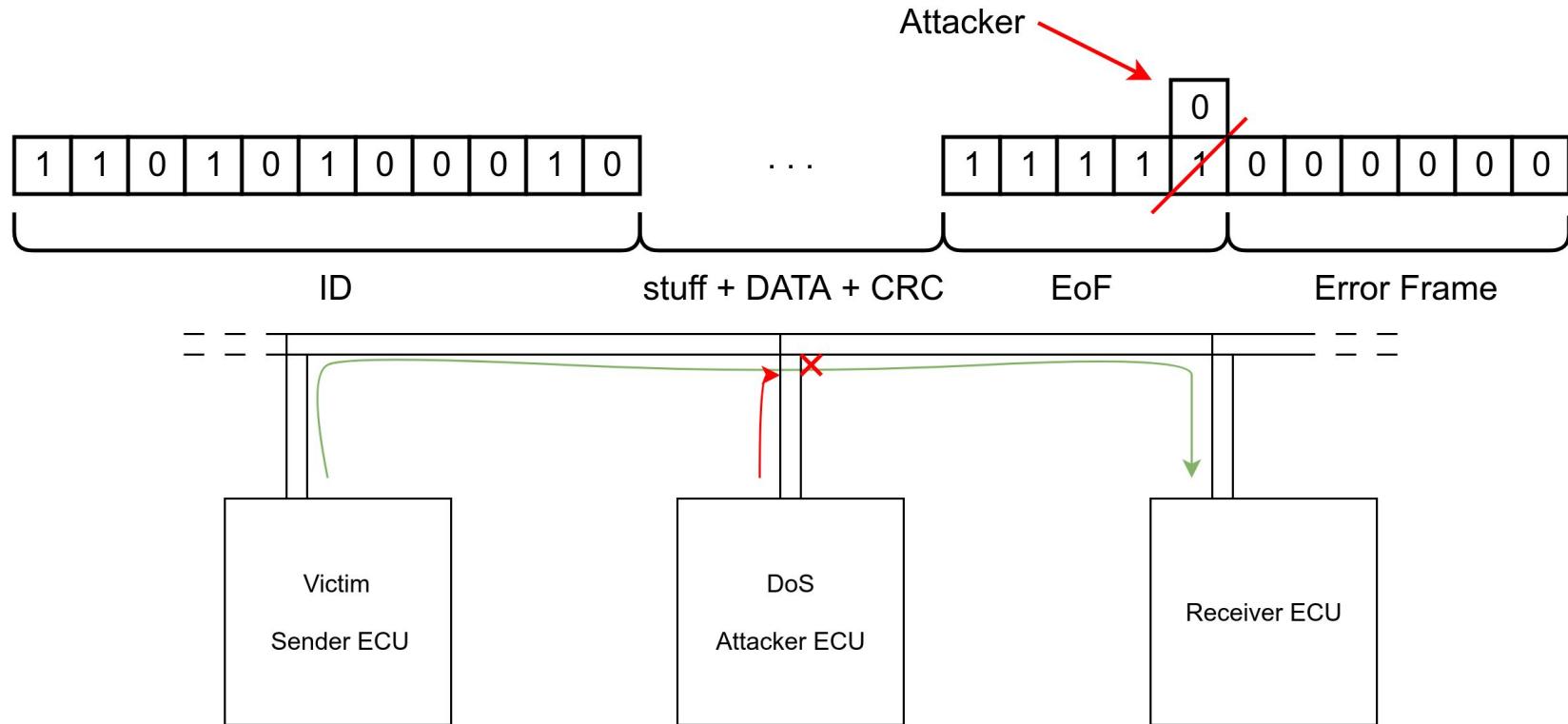
Recessive is deleted by Dominant Bit

CSMA/BA property:

An attacker can ALWAYS win arbitration



CAN Targeted Denial of Service



CAN Practical attack examples

http://illmatics.com/car_hacking.pdf

<http://illmatics.com/Remote%20Car%20Hacking.pdf>

<http://illmatics.com/can%20message%20injection.pdf>

<https://docs.google.com/presentation/d/1Js82wcBiDPJHvLL2d6lOWGITZh-tlF8E-OhghtL80PE/edit?usp=sharing>



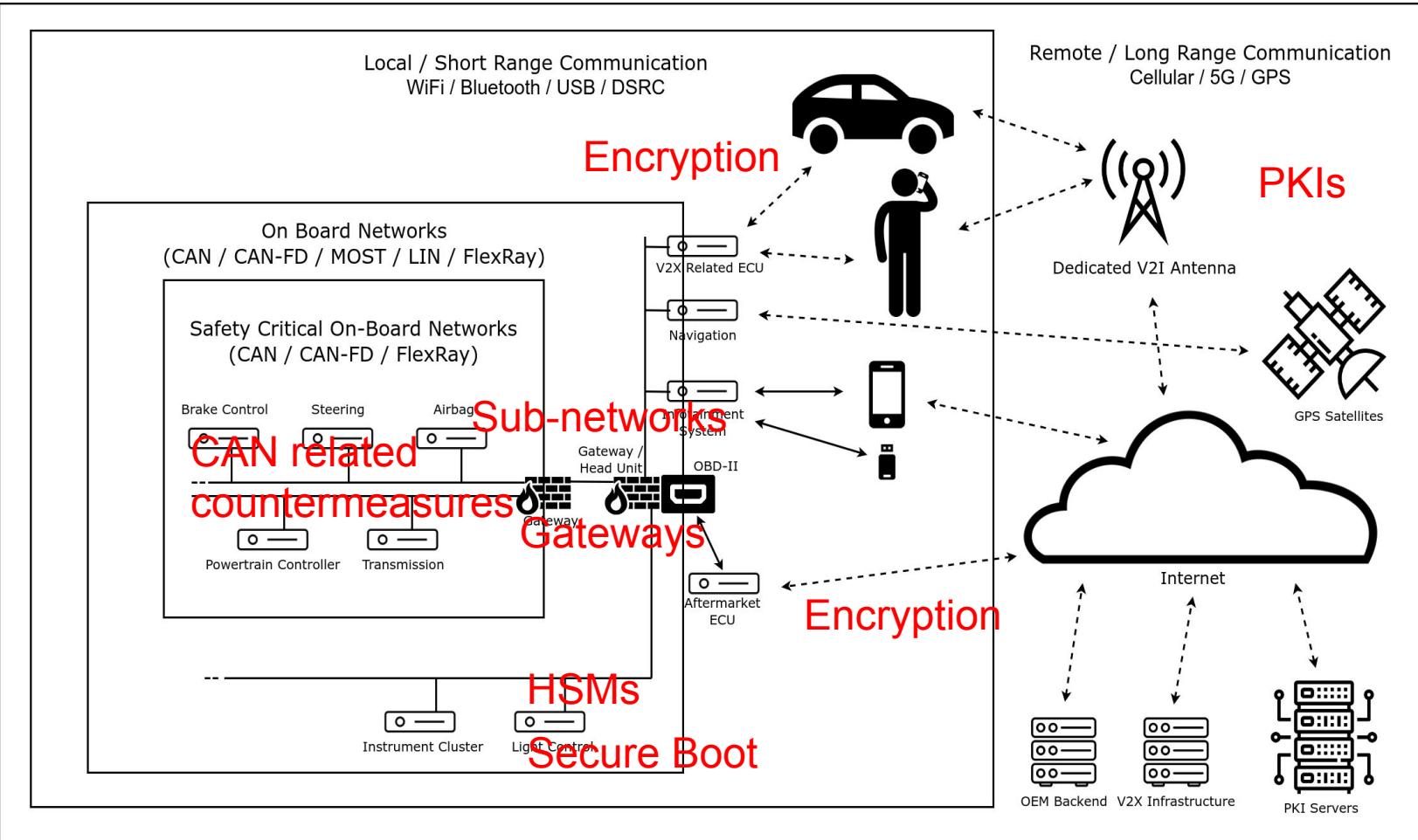
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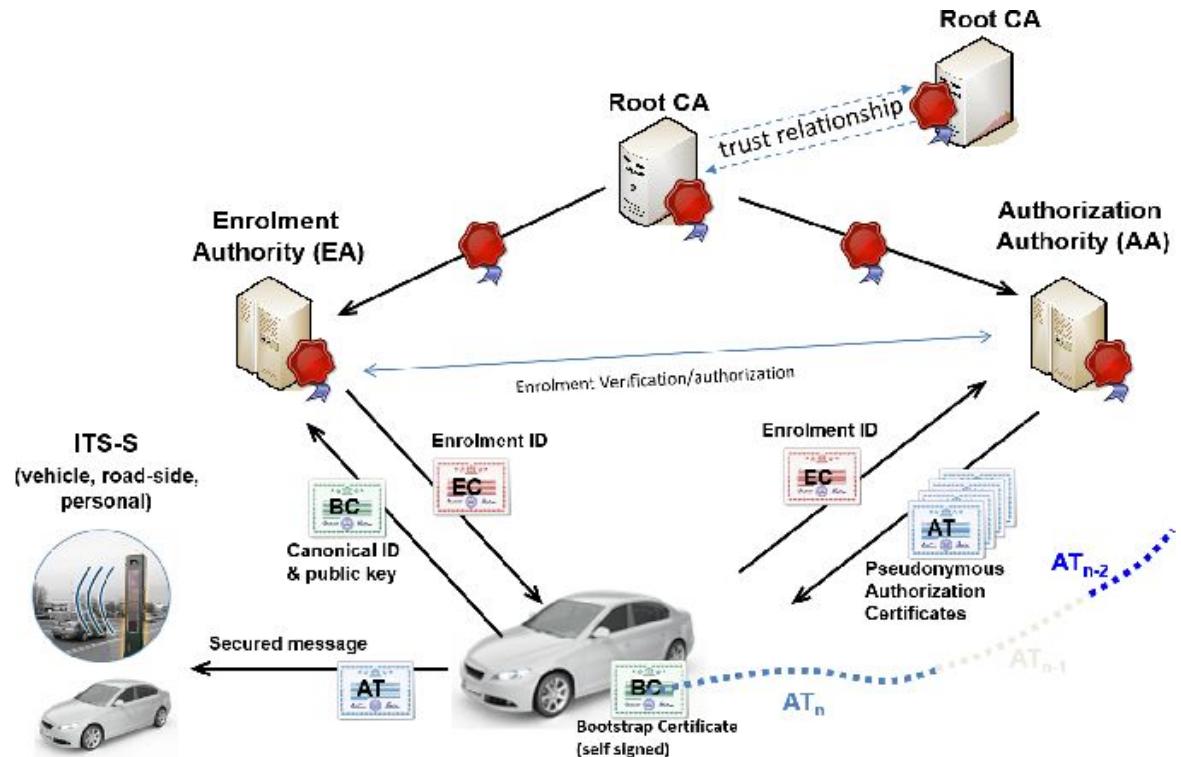
COUNTERMEASURES



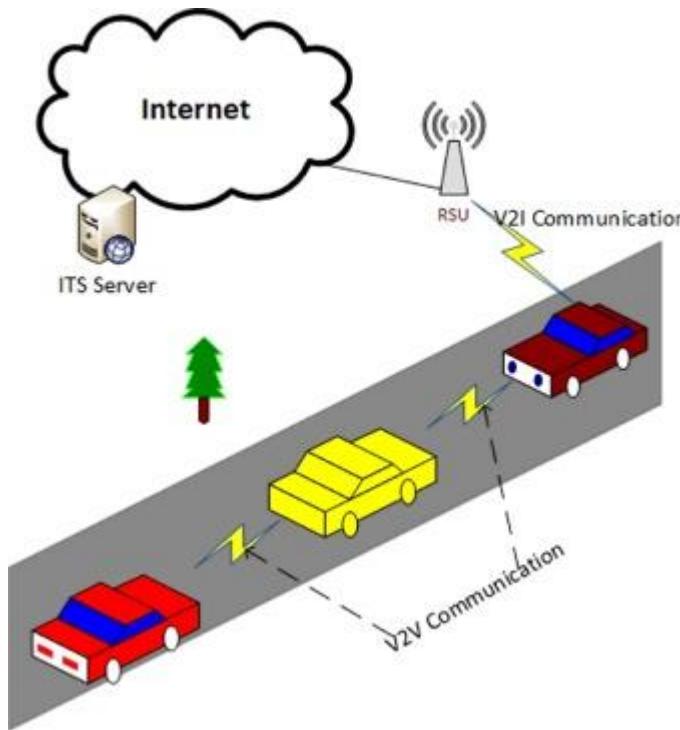
PKI

Requires Pseudonyms
(privacy)

Ensures sender's entity



PKI



Sending and receiving traffic data

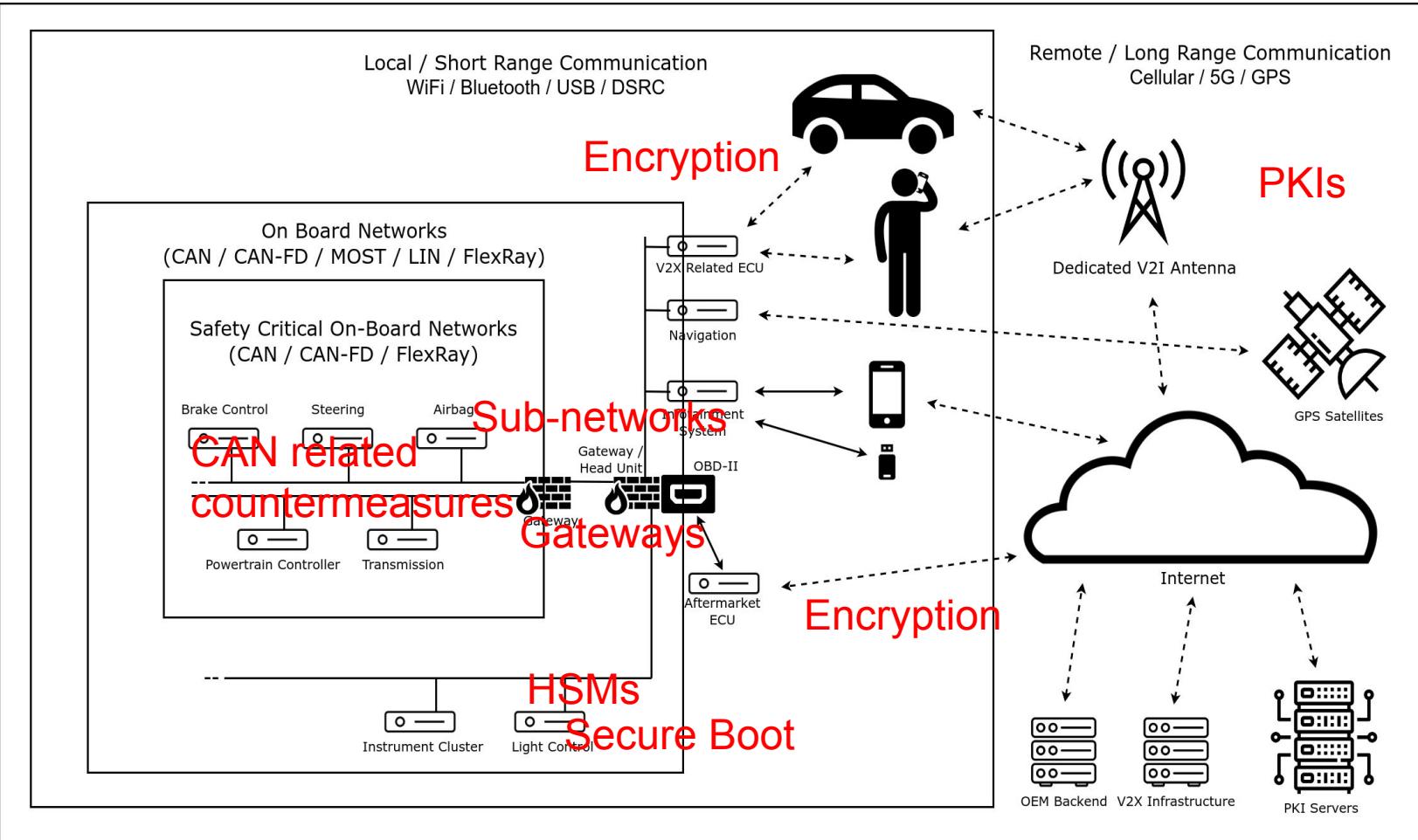
Sending status updates

Receiving software updates

Phone calls

Internet navigation

...



ECU security (Secure Boot)

Secure Boot:

- 1) Trusted code written in production in read only memory
- 2) Start of execution always from trusted code
- 3) Trusted code verifies other code through PKI or hash (potentially updated)
- 4) Other code gets executed

Pro: avoids easy flashing of ECUs

Cons: valid only for the specific ECU on which it is implemented

ECU security (Hardware Security Module)

- 1) Makes the cryptographic computation
- 2) Stores keys safely
- 3) Doesn't let the ECU communicate directly with the (untrusted) network

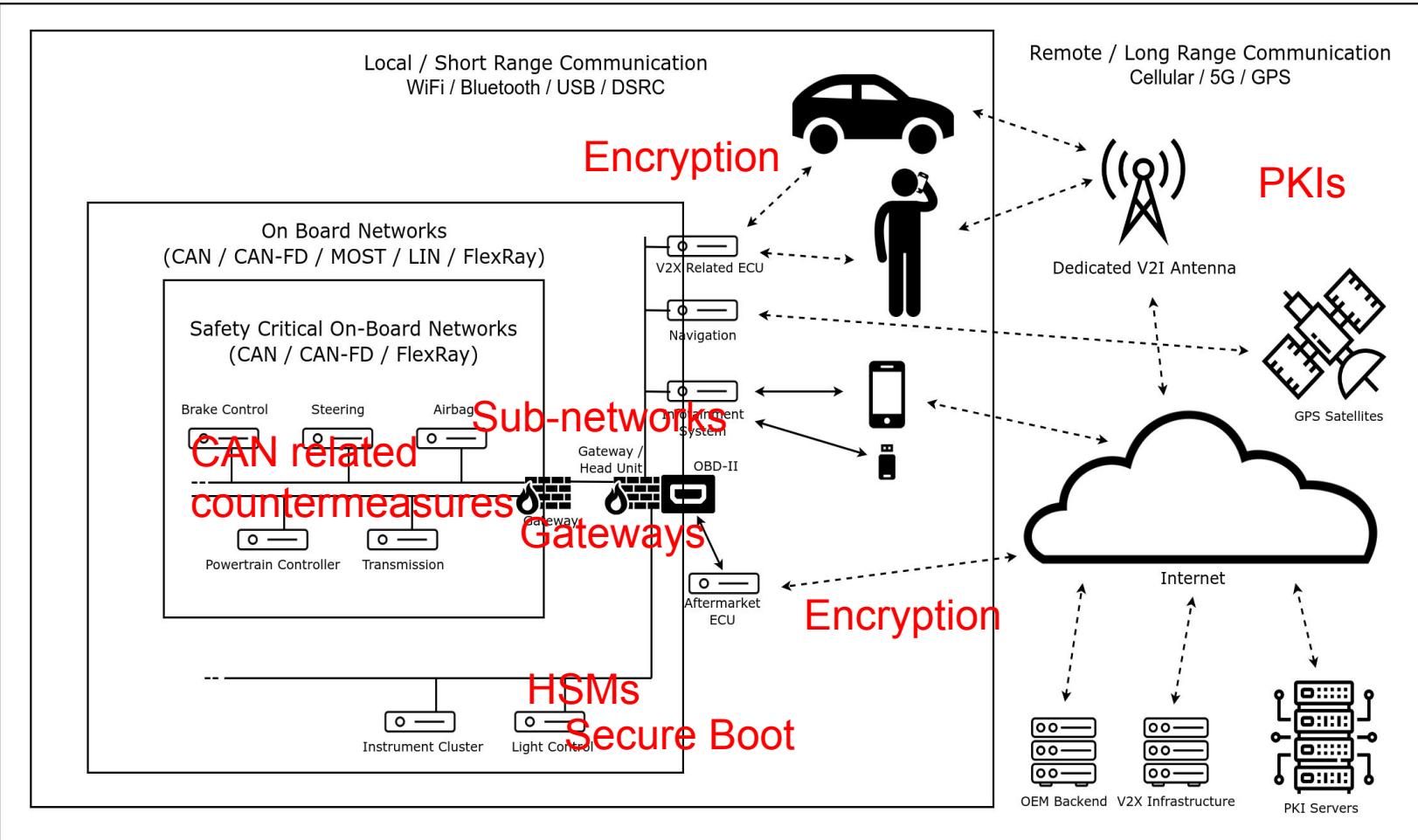
Pros: Increases security of the specific device by shielding it from external threats

Cons: Costs

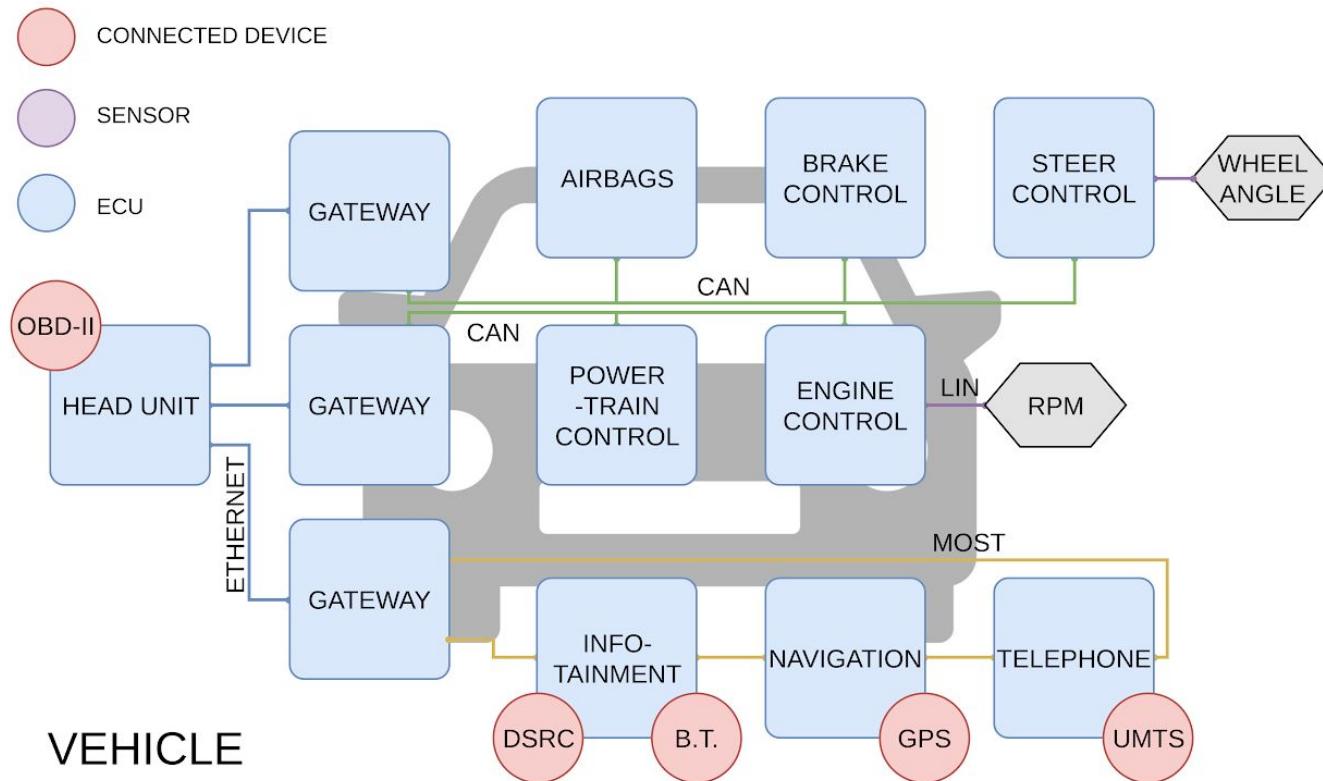
ECU security (OTA secure software updates)

Important to mention, although implemented by the previous technologies.

Enables patching of detected vulnerabilities without recalling the vehicles
(Jeep Hack costs were estimated in the tens to hundred millions \$ for FCA, since
1.4 million vehicles were recalled)



Gateways, subnetworks, firewalls



Gateways, subnetworks, firewalls

Gateways divide one subnetwork from the other, and firewalls (in the GW) avoid unwanted packets to reach critical locations

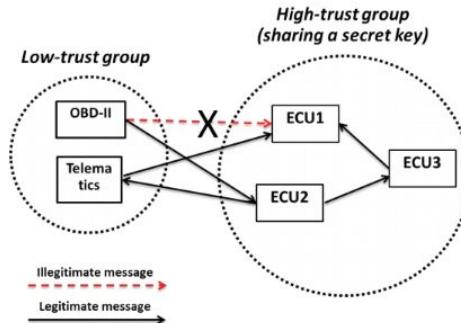
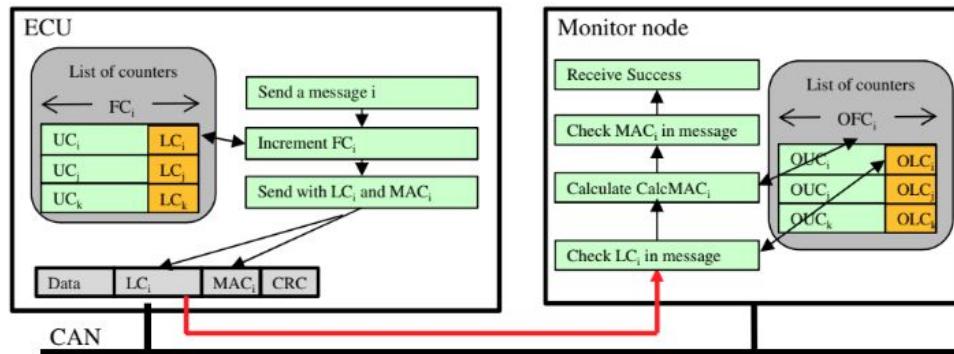
Subnetworks are divided so that critical ECUs are not in contact with externally connected ones

Pros: Gateways are required anyway to translate one kind of network packets to another. Only a small amount are required for the whole network

Cons: Do not defend from packets meant to pass in that GW (e.g.,)

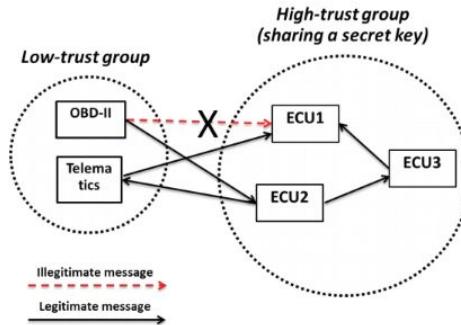
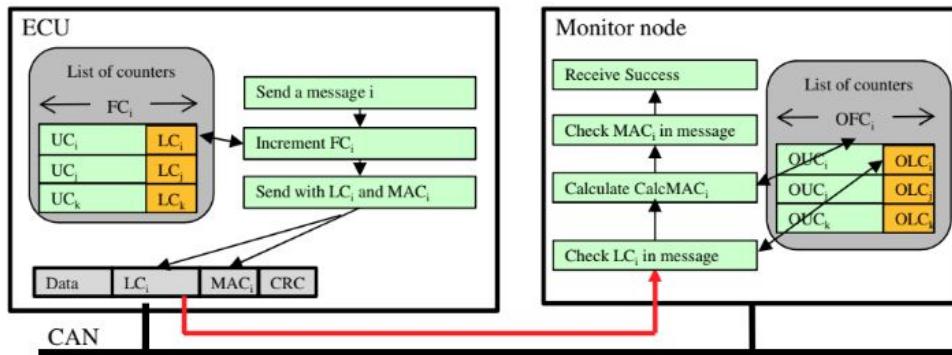
On-board Network Security: Authentication Protocols

Encrypt / Authenticate the message passing
on the CAN (or other) bus



On-board Network Security: Authentication Protocols

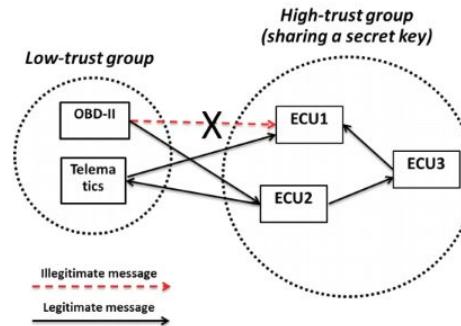
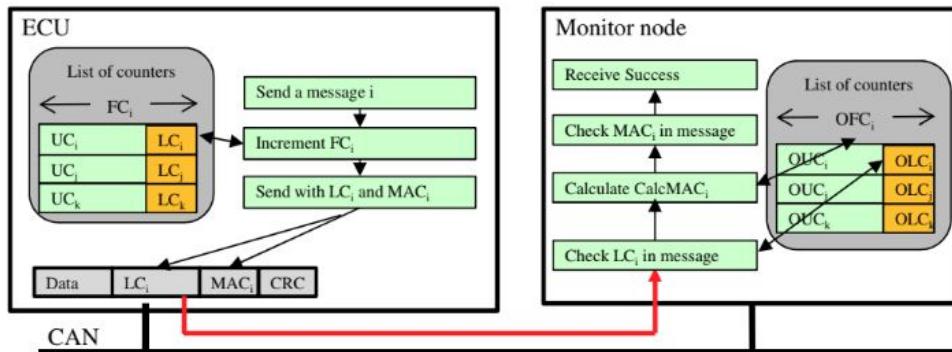
- 1) MAC/HMAC
- 2) Backward compatible
- 3) Centralized/Distributed
- 4) Key Distribution
- 5) Security Level



On-board Network Security: Authentication Protocols

Pros: Relatively high security measure,
especially if uses an HSM to store keys

Cons: Expensive, slows computation,
Increases message overhead/numbers



Generally used only for
firmware updates and
not even often.

On-board Network Security: Intrusion Detection

A singular (or a small amount) of devices that can read network traffic and analyze it to detect inconsistencies.

On-board Network Security: Intrusion Detection

Frequency Based IDSs:

- + CAN data is usually periodic, and when they are not, they are still sent at a specific frequency
- + Attackers often need to “overwrite” the messages sent by an ECU to convince another one of something (see Miller & Valasek speed)
- If the attacker can send the “writer” off the bus frequency can be kept identical

On-board Network Security: Intrusion Detection

Specification-based IDSs:

- a) Physical specifications (e.g., Voltages/Power, ...)
 - b) Logical specifications (ID ownership)
 - c) Protocol properties (Detect bus-off victim)
-
- + Make it impossible for the attacker to exploit some specific vulnerabilities (e.g. buffer overflows due to too high values, Sending an ID from an ECU that was not supposed to send it)
 - Extremely case-specific, hard to generalize. May require testing per-vehicle

On-board Network Security: Intrusion Detection

Payload-based IDSs:

- a) Data related rules (e.g., some bits may have fixed values)
 - b) Multi-value analysis (e.g., speed + rpm + gear)
 - c) Time-series analysis (e.g., speed increase not realistic)
 - d) Machine-Learning based rules
-
- + Impossible for the attacker to send non-protocol compliant requests
 - + Extremely complex for the attacker to comply to all the rules while being able to implement a meaningful attack
 - Computationally expensive
 - Hard to find rules strict enough and at the same time without false positives

On-board Network Security: Intrusion Reaction

- 1) Shut Down the attacker? -> Dangerous, also how do I tell?
- 2) Send Alert? -> to who?
- 3) Switch to a less “technology-reliant” driving mode?
- 4) Change IDs
- 5) Send data to data analysts to patch the problem ←

Defenses VS Attacks

- 1) Sniffing -> no real solution
- 2) Spoofing -> Authentication Protocols (but limited)
- 3) Spoofing -> Intrusion Detection Systems combinations
- 4) DoS -> No real solution :((yet kinda less useful than others)
- 5) DoS + Spoofing -> CopyCAN and similar (see slides)

<https://docs.google.com/presentation/d/1Js82wcBiDPJHvLL2d6lOWGITZh-tlF8E-OhghtL80PE/edit?usp=sharing>

Overall security solution (best case) :

- 1) Threat analysis is done in advance, while designing the vehicle
- 2) Penetration testing is done on all ECUs (unrealistic, the most safety critical ones at least)
- 3) Firewalls are always implemented in gateways, and threat analysis defines the best networks layouts
- 4) ECUs that expect to have software updates (infotainment system, safety critical ones) are either implemented with an HSM or with a secure implementation of key storage and PKIs
- 5) An IDS is implemented on board, or part of the data are sent remotely to be analyzed by a more powerful one
- 6) A team studies the alerts being sent by different vehicles, understand the current threats and implements patches.

Overall security solution (best case) :

Pros:

Costs are contained

Overall known attacks are (more or less) unfeasible

The reactions can be implemented also on vehicles already on the road

Cons:

The attack being implemented on a specific vehicle is not always stoppable

The patch for an attack may arrive after months



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Secure Positioning: RKEs and GPS

Remote Passive Keyless Entry/Start Systems

How do they work?

- 1) The vehicle periodically sends a message asking for the key
- 2) When the key receives it, it responds through RFID
- 3) The vehicle opens the doors/ignites the engine

Remote Passive Keyless Entry/Start Systems

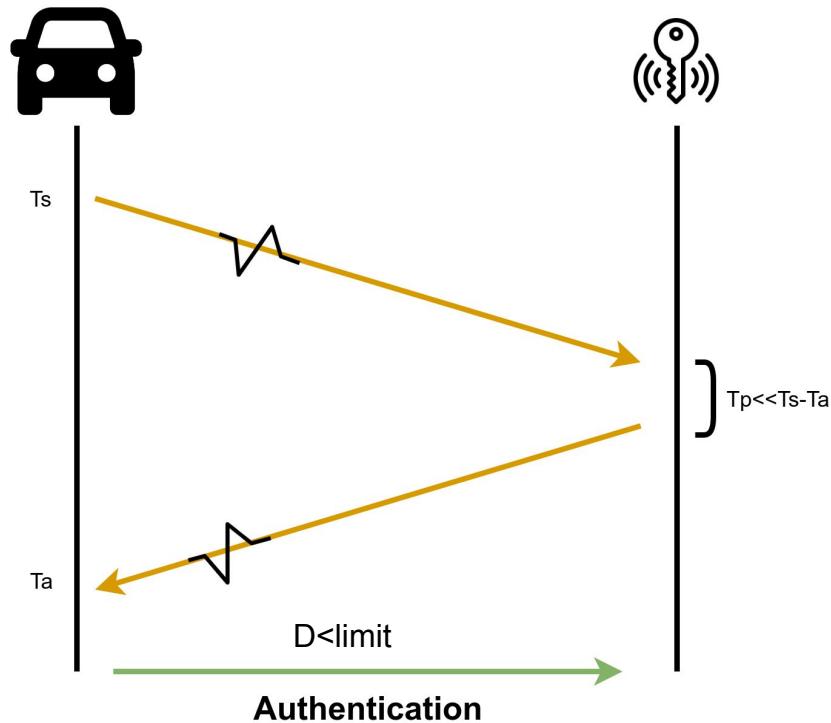
Does it need security? -> Well it is a substitute of the key, so yeah

Does it implement security measures? -> yes

Is it secure? ...



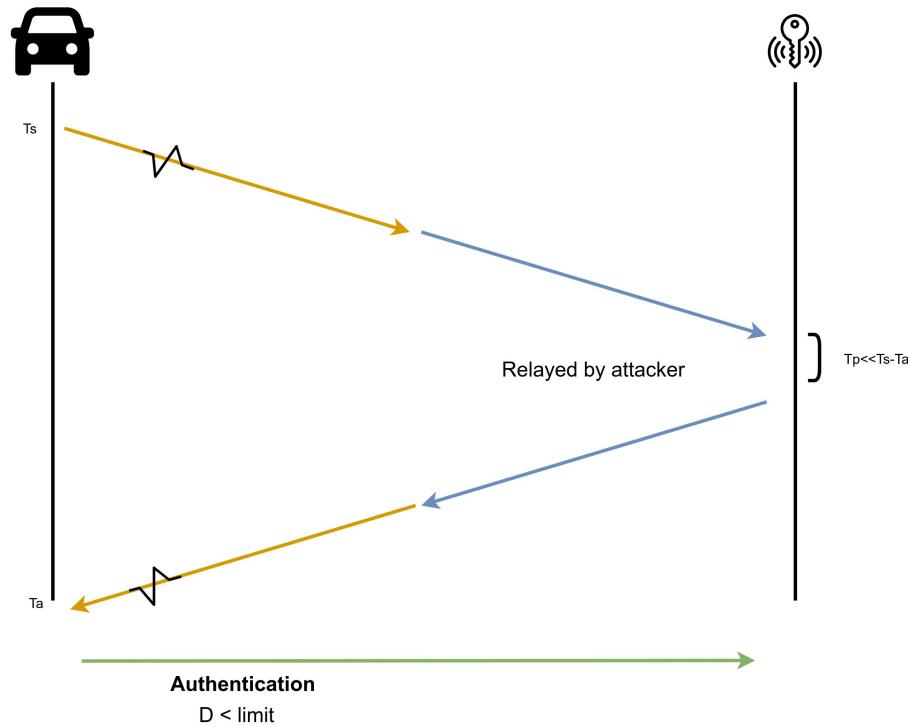
How does RKES work?



- 1) Single code
- 2) Rolling code
- 3) Manufacturer specific global codes
- 4) Challenge Response

Periodic
 $D = (T_a - T_s - T_p) * c / 2$

What does an attacker want?



Generally to fake the key to be close to the car even if it isn't. In this way the car will start.

How is it secured?

IDM - Indirect Distance Measurement

- 1) NFC/RFID
- 2) RSSI (Received Signal Strength Indication) Measurement
- 3) Phase Measurement
- 4) AoA (Angle of Arrival) Measurement

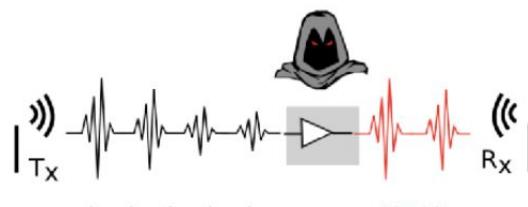
...

DDM - Direct Distance Measurement (Time-of-Flight)

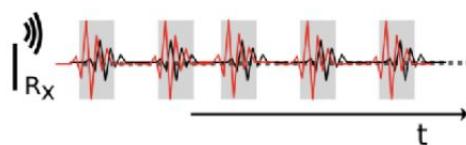
Most secure because basically uncheatable

Known Attacks

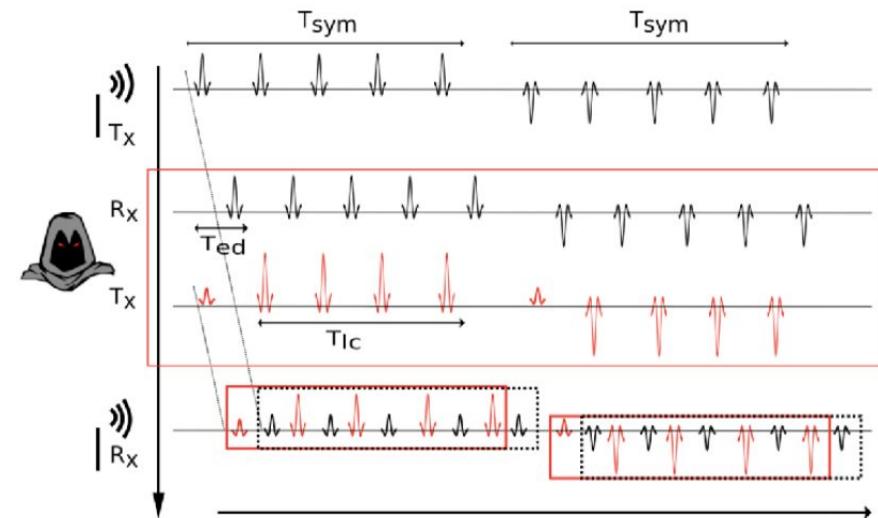
Simple Relay, Phase Relay, Signal Amplification, Early Detect / Late Commit, Cicada, Preamble Advance, ...



a) Relay Attack

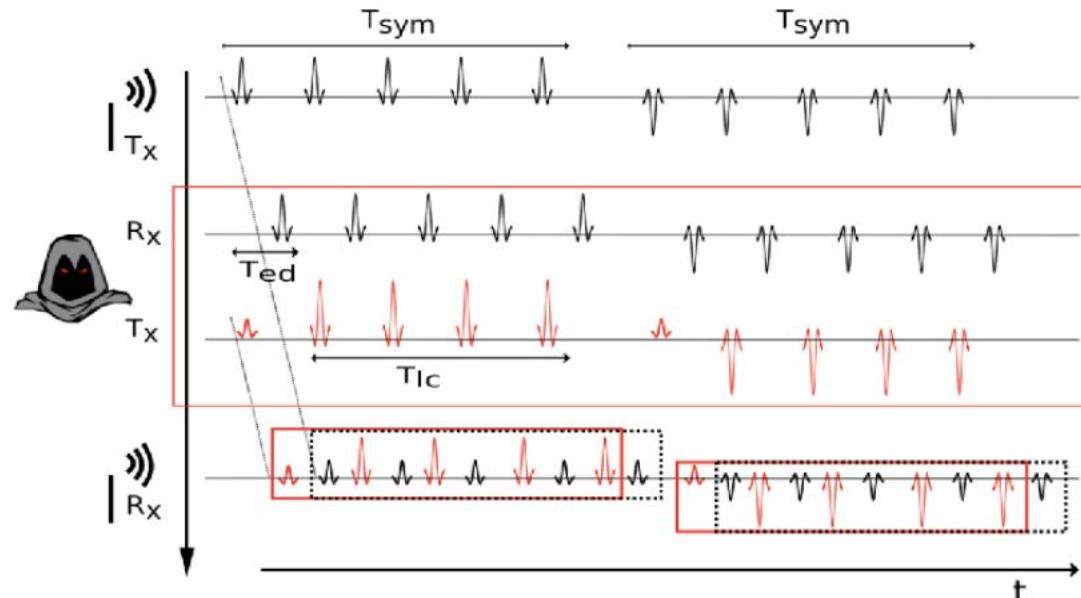


b) Cicada Attack



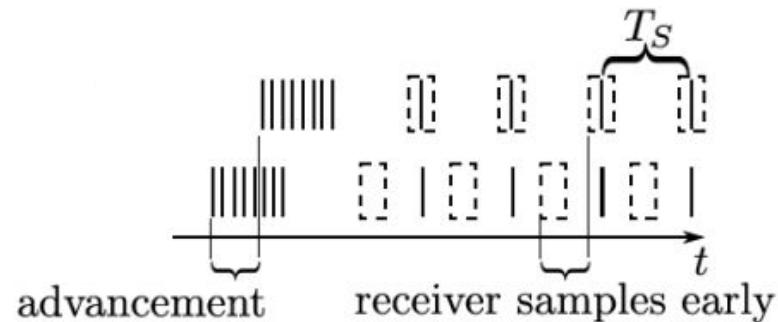
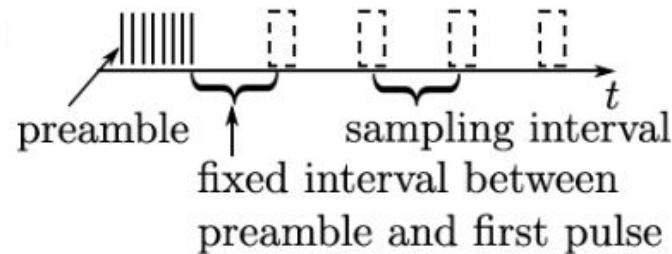
c) ED/LC Attack

Known Attacks: Early Detect/Late Commit

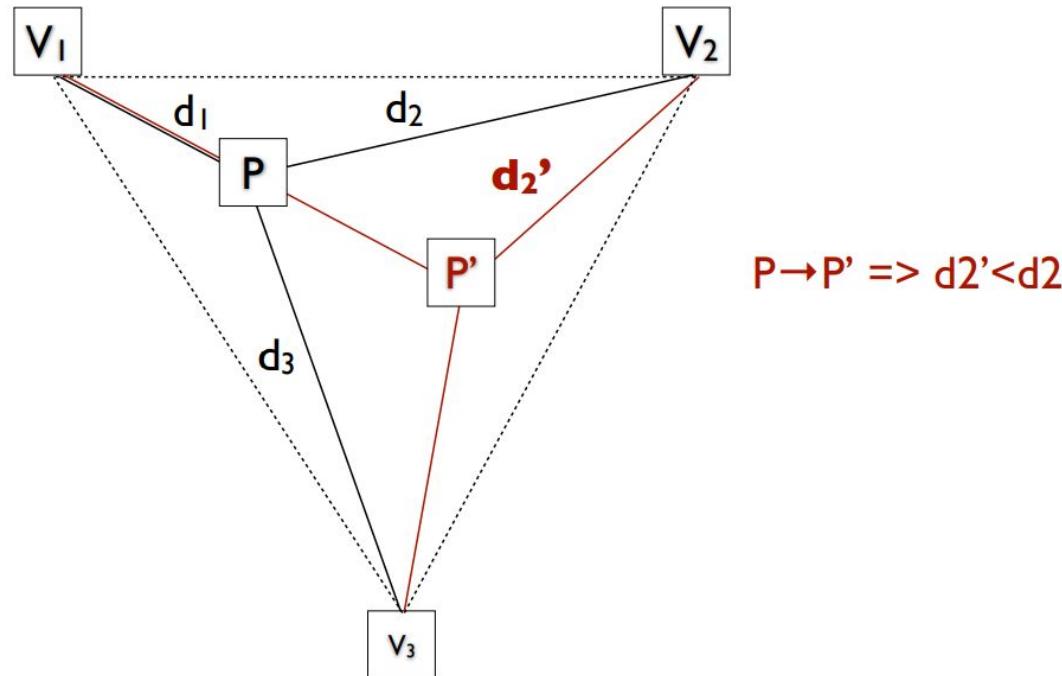


c) ED/LC Attack

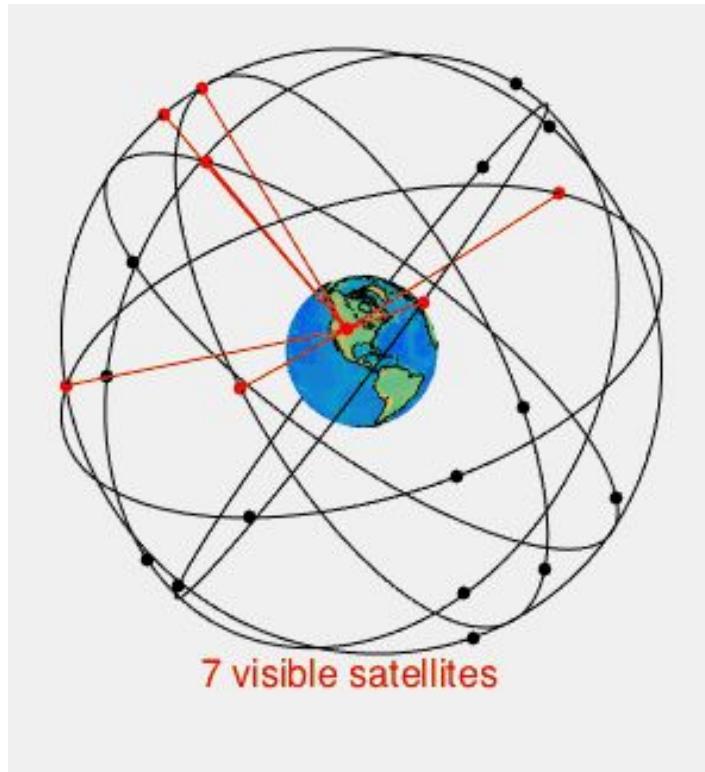
Solution: Preamble!



Triangulation OK! (Short Range)



Long Range Positioning: GPS / GNSS



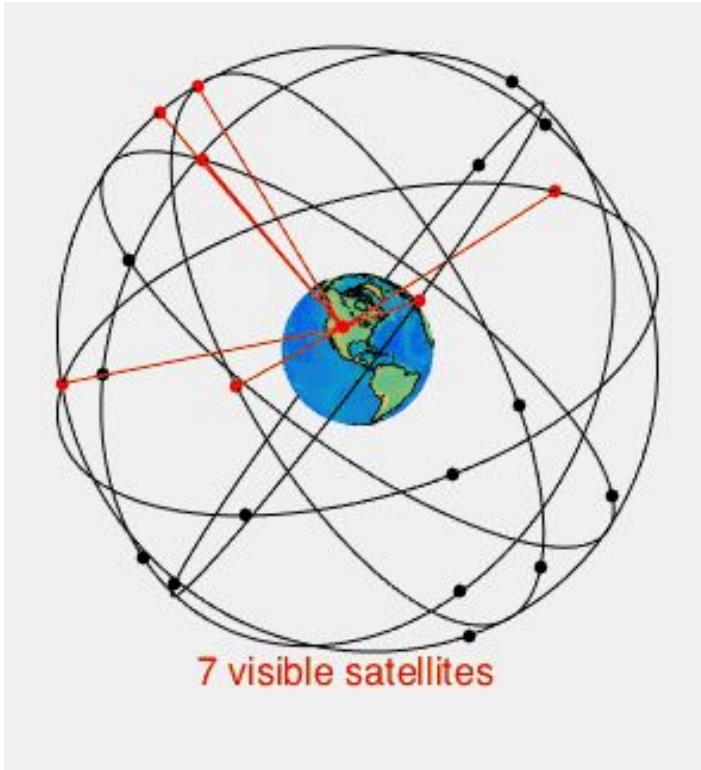
24 satellites at 20.200 Km

Each satellite sends its location and
PRECISE time of transmission

GPS receiver measures the distance
from each satellite

Receiver uses trilateration to calculate its
position

GPS Encryption

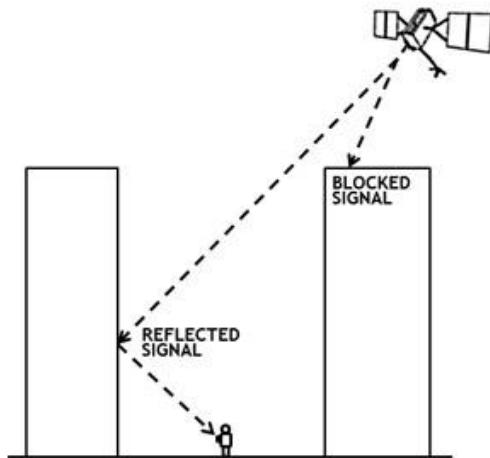


GPS is encrypted ONLY for military purposes.

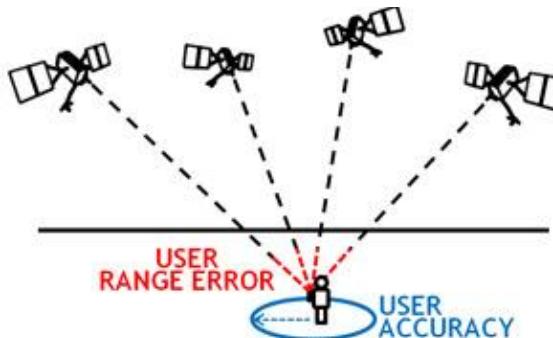
Why?

- hard to keep the secret a secret
- all satellites has to know all keys
- can't create 10000000000 keys

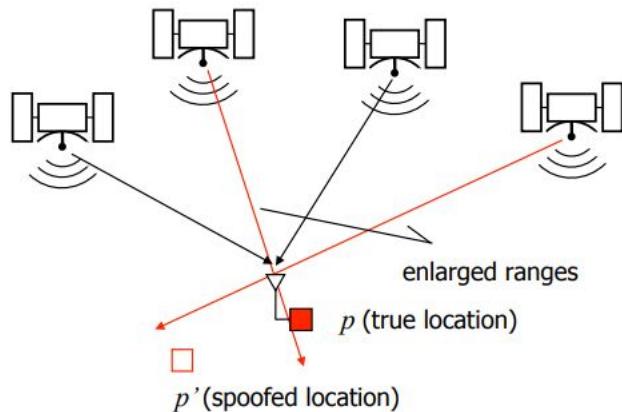
GPS Precision



The government commits to broadcasting the GPS signal in space with a global average user range error (URE) of ≤ 7.8 m (25.6 ft.), with 95% probability. Actual performance exceeds the specification. On May 11, 2016, the global average URE was ≤ 0.715 m (2.3 ft.), 95% of the time.



GPS Spoofing



The attacker:

Modifies the navigation message contents

or

Manipulates the time of arrival

(Military GPS can only be delayed)

GNSS Countermeasures

Changing the protocol:

- Authentication of messages -> can still be delayed
- Direct Sequence Spread Spectrum -> requires secret shared keys

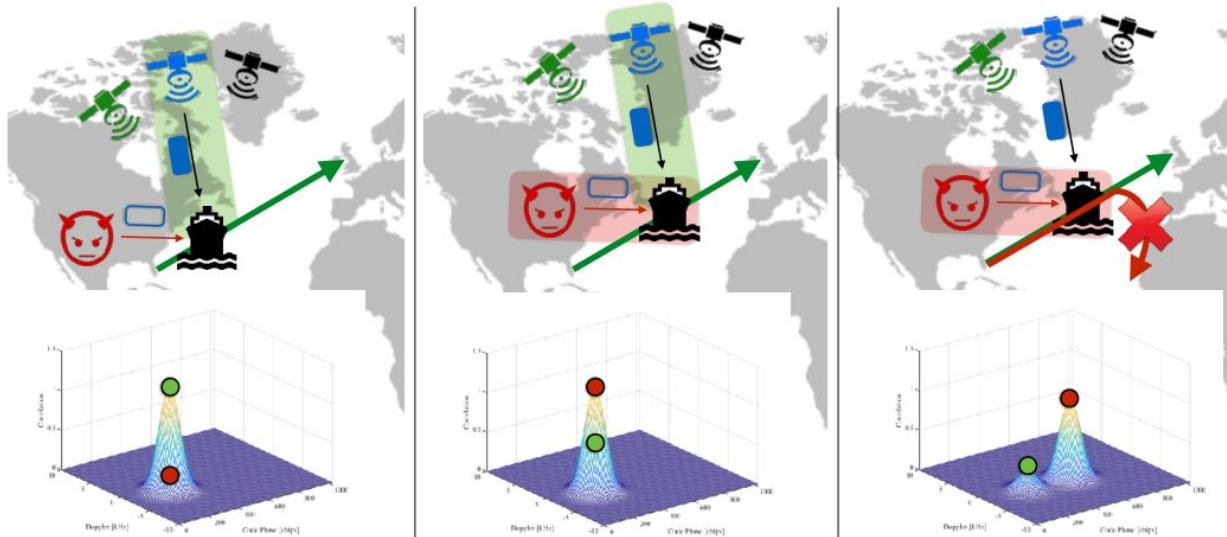
Not changing the protocol:

- Noise level, # of satellites...
- Spatial Diversity (AoA...)

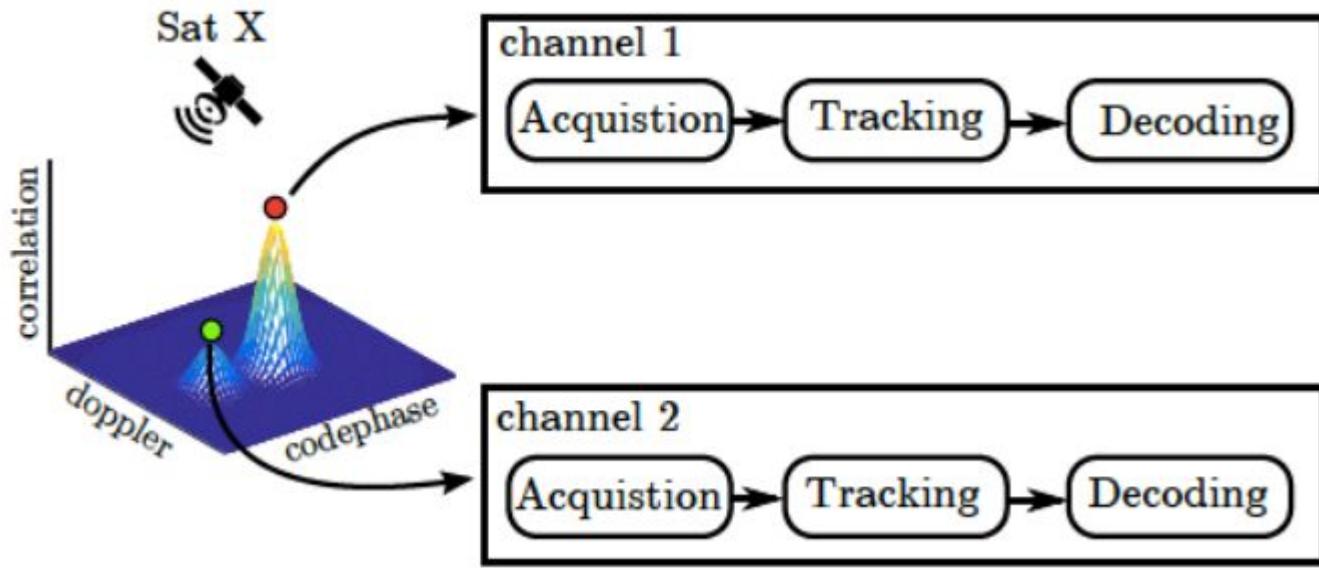
More or less feasible...

GNSS Seamless Takeover Attack

...But counterable

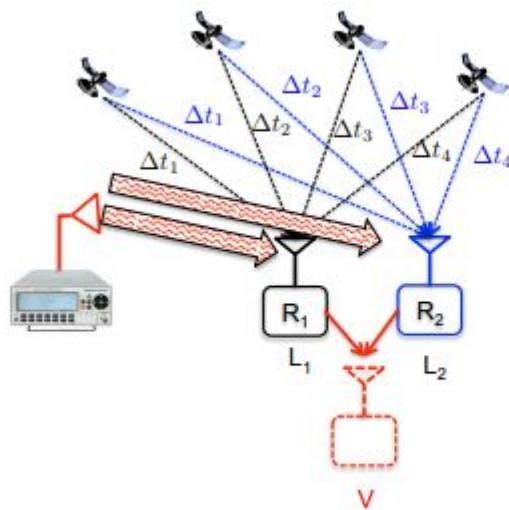


Detection with one receiver: SPREE



Detection with multiple receivers:

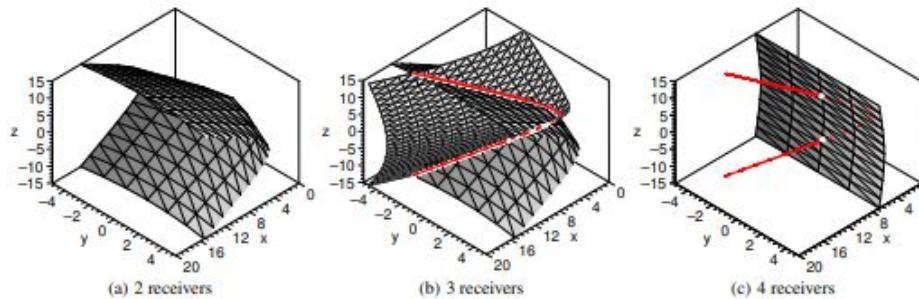
Leverage spatial diversity



If $d(R_1, R_2)$ is known
=> spoofing detection

Detection with multiple receivers:

Spatial diversity and number of nodes constraint the attacker



n	Spoofing to one location		Spoofing to multiple locations (preserved formation)	
	Civ. & Mil. GPS	Civilian GPS	Military GPS	
1	$P_i^A \in \mathbb{R}^3$	-	-	
2	$P_i^A \in \mathbb{R}^3$	set of hyperboloids	one hyperboloid	
3	$P_i^A \in \mathbb{R}^3$	set of intersections of two hyperboloids	intersection of two hyperboloids	
4	$P_i^A \in \mathbb{R}^3$	set of 2 points	2 points	
≥ 5	$P_i^A \in \mathbb{R}^3$	set of points	1 point	



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EXERCISE

Go to this link and download the dataset: (password to decrypt: cefriel2020)

<https://drive.google.com/file/d/19rYU4nNGR9YloQ3u7qmdbmQBYq6W9IR8/view?usp=sharing>

Answer the following questions:

- 1) which IDs are present in the datasets (same IDs for all datasets)
- 2) what is the frequency of each ID in the noattack.csv dataset?
- 3) which attack is being implemented on each log?
- 4) what IDs are being attacked?
- 5) what is the goal of the attacker? (if feasible to understand)
- 6) implement a detector for the specific kind of attack on the specific log