

*W4111 – Introduction to Databases
Section 002, Spring 2024
Lecture 8
Module II (1), NoSQL (1)*



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Lecture 8
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We will start in a couple of minutes.

Module II Kickoff

Course Modules – Reminder

Course Overview

Each section of W4111 is slightly different based on student interest and professor's focus. There is a common, core syllabus. Professors cover topics in different orders and grouping based on teaching style.

This section of W4111 has four modules:

- **Foundational concepts (50% of semester):** This module covers concepts like data models, relational model, relational databases and applications, schema, normalization, ... The module focuses on the relational model and relational databases. The concepts are critical and foundational for all types of databases and data centric applications.
- **Database management system architecture and implementation (10%):** This module covers the software architecture, algorithms and implementation techniques that allow [databases management systems](#) to deliver functions. Topics include memory hierarchy, storage systems, caching/buffer pools, indexes, query processing, query optimization, transaction processing, isolation and concurrency control.
- **NoSQL – “Not Only SQL” databases (20%):** This module provides motivation for [“NoSQL”](#) data models and databases, and covers examples and use cases. The module also includes cloud databases and databases-as-a-service.
- **Data Enabled Decision Support (20%):** This module covers data warehouses, data import and cleanse, OLAP, Pivot Tables, Star Schema, reporting and visualization, and provides an overview of analysis techniques, e.g. clustering, classification, analysis, mining.

Module II – DBMS Architecture and Implementation Overview and Reminder

Module II – DBMS Architecture and

What is a database? In essence a database is nothing more than a collection of information that exists over a long period of time, often many years. In common parlance, the term *database* refers to a collection of data that is managed by a DBMS. The DBMS is expected to:

- 
1. Allow users to create new databases and specify their *schemas* (logical structure of the data), using a specialized *data-definition language*.

Covered for the relational model.

Database Systems: The Complete Book (2nd Edition)

by [Hector Garcia-Molina](#) (Author), [Jeffrey D. Ullman](#) (Author), [Jennifer Widom](#) (Author)

Module II – DBMS Architecture and

- 
- 
2. Give users the ability to *query* the data (a “query” is database lingo for a question about the data) and modify the data, using an appropriate language, often called a *query language* or *data-manipulation language*.
 3. Support the storage of very large amounts of data — many terabytes or more — over a long period of time, allowing efficient access to the data for queries and database modifications.
 4. Enable *durability*, the recovery of the database in the face of failures, errors of many kinds, or intentional misuse.
 5. Control access to data from many users at once, without allowing unexpected interactions among users (called *isolation*) and without actions on the data to be performed partially but not completely (called *atomicity*).

Database Systems: The Complete Book (2nd Edition)

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Purpose of Database Systems

In the early days, database applications were built directly on top of file systems, which leads to:

- Data redundancy and inconsistency: data is stored in multiple file formats resulting in duplication of information in different files
- Difficulty in accessing data
 - Need to write a new program to carry out each new task
- Data isolation
 - Multiple files and formats
- Integrity problems
 - Integrity constraints (e.g., account balance > 0) become “buried” in program code rather than being stated explicitly
 - Hard to add new constraints or change existing ones

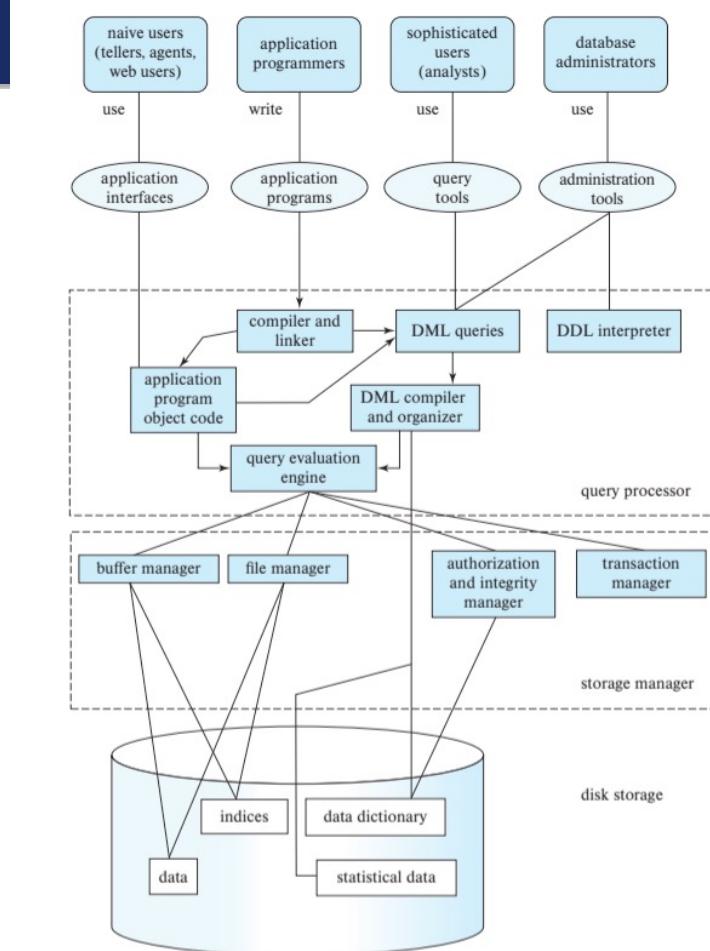
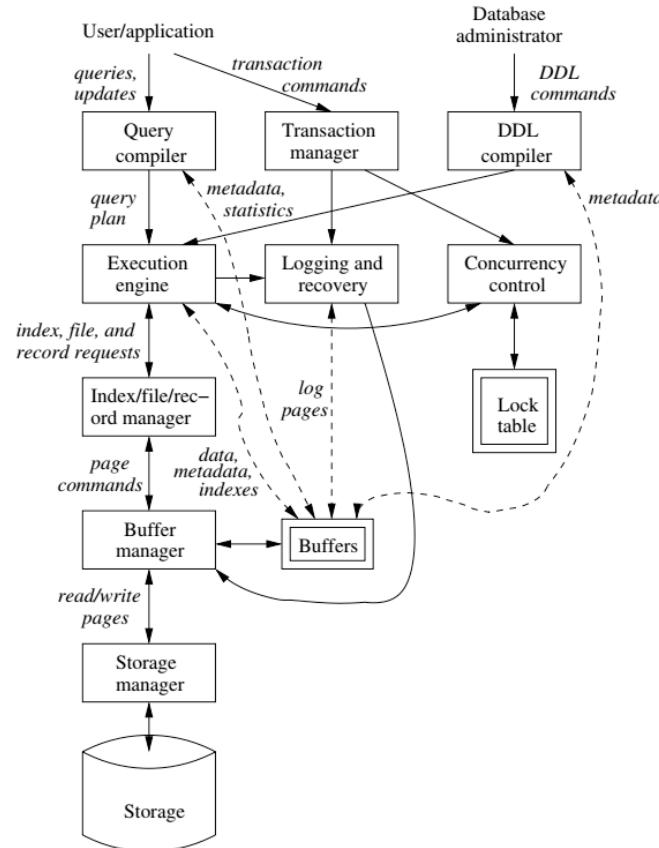


Purpose of Database Systems (Cont.)

- Atomicity of updates
 - Failures may leave database in an inconsistent state with partial updates carried out
 - Example: Transfer of funds from one account to another should either complete or not happen at all
- Concurrent access by multiple users
 - Concurrent access needed for performance
 - Uncontrolled concurrent accesses can lead to inconsistencies
 - Ex: Two people reading a balance (say 100) and updating it by withdrawing money (say 50 each) at the same time
- Security problems
 - Hard to provide user access to some, but not all, data

Database systems offer solutions to all the above problems

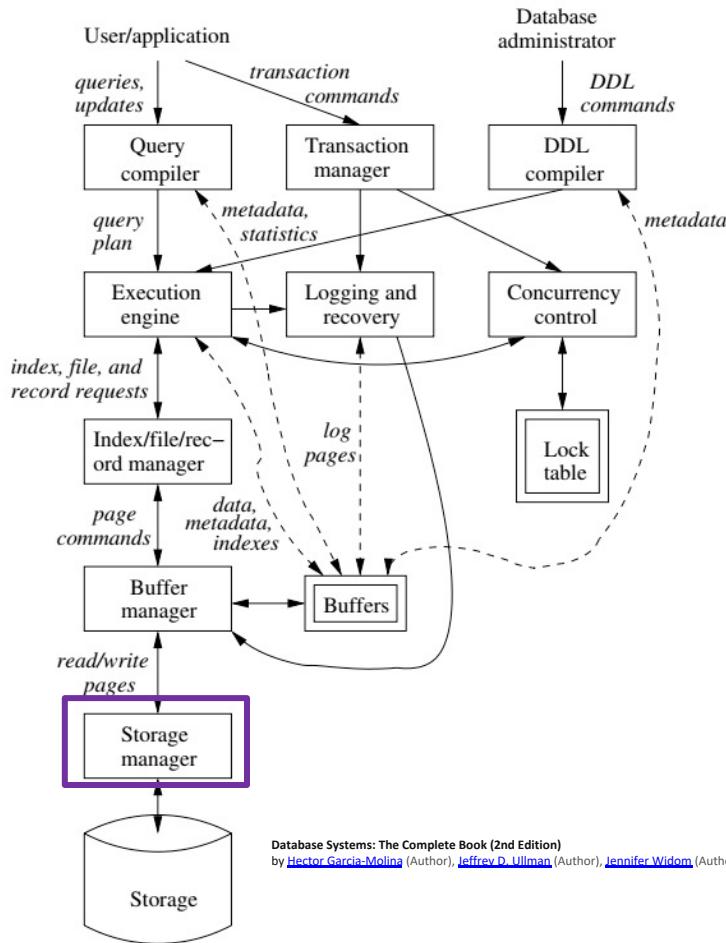
DBMS Arch.



Data Management

Today

- Load/save things quickly.



Disks

Input/Output (IO)

Disks as Far as the Eye can See

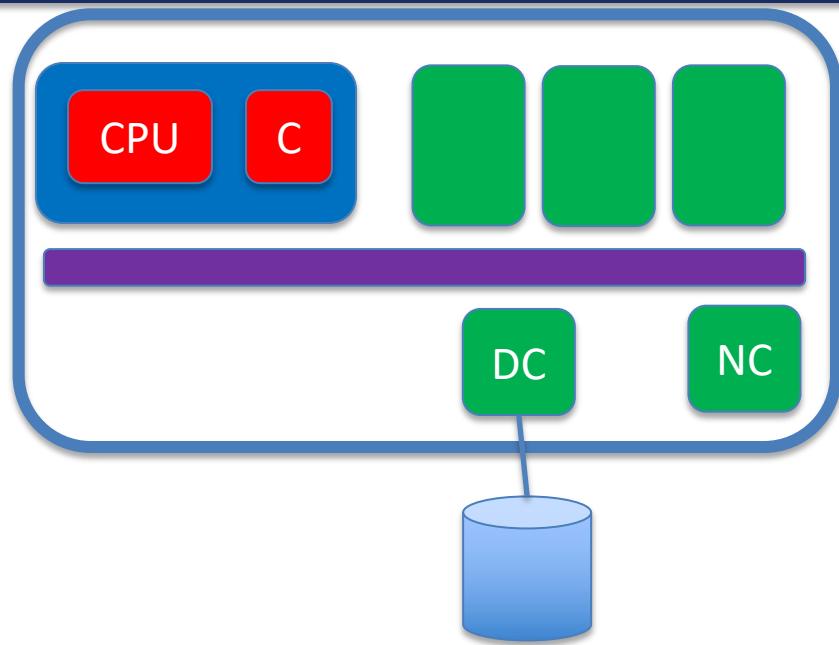
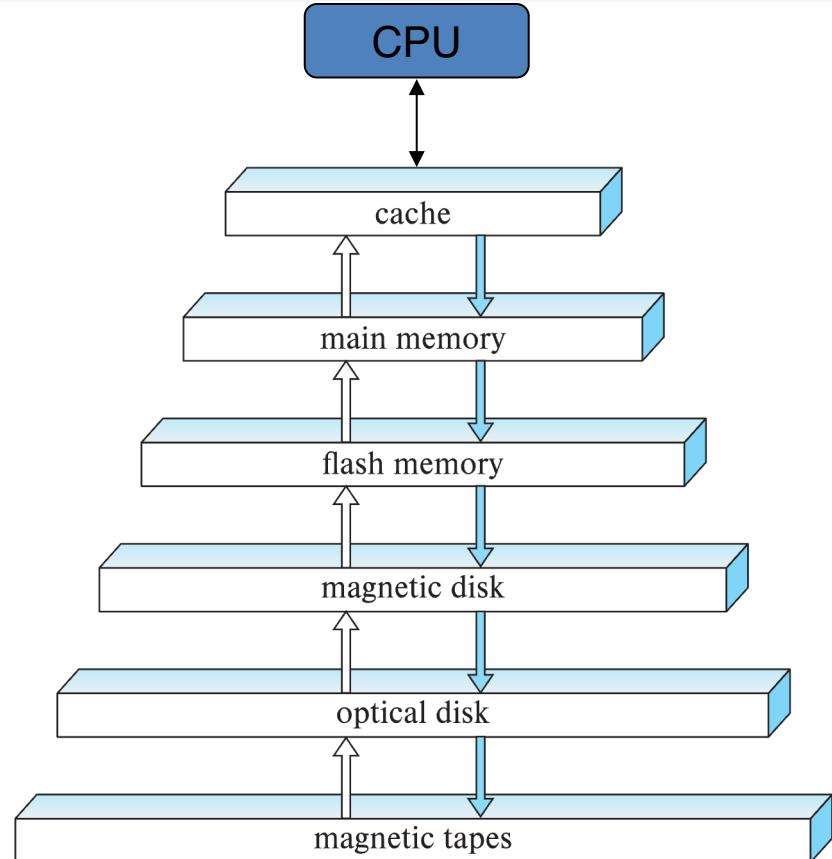




Classification of Physical Storage Media

- Can differentiate storage into:
 - **volatile storage:** loses contents when power is switched off
 - **non-volatile storage:**
 - Contents persist even when power is switched off.
 - Includes secondary and tertiary storage, as well as battery-backed up main-memory.
- Factors affecting choice of storage media include
 - Speed with which data can be accessed
 - Cost per unit of data
 - Reliability

Storage Hierarchy



From: Database System Concepts, 7th Ed.

Memory Hierarchy

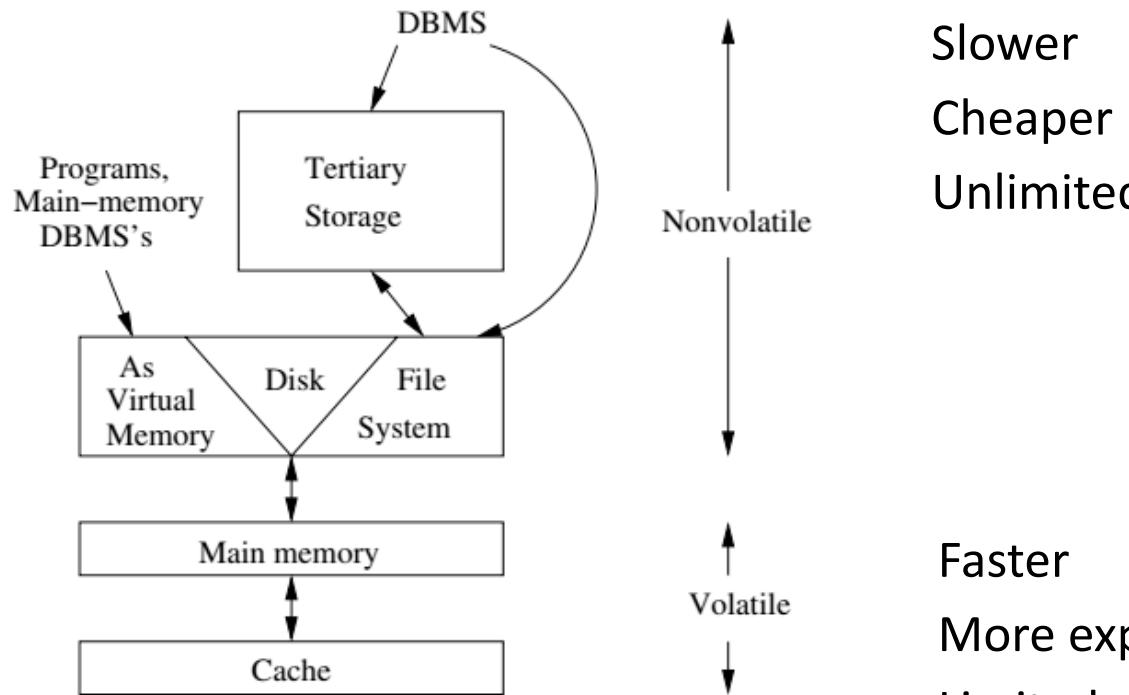


Figure 13.1: The memory hierarchy

From: Database System Concepts, 7th Ed.

Memory Hierarchy (Very Old Numbers – Still Directionally Valid)

Storage Technology

Price, Performance & Capacity

Technologies	Capacity (GB)	Latency (microS)	IOPs	Cost/IOPS (\$)	Cost/GB (\$)
Cloud Storage	Unlimited	60,000	20	17c/GB	0.15/month
Capacity HDDs	2,500	12,000	250	1.67	0.15
Performance HDDs	300	7,000	500	1.52	1.30
SSDs (write)	64	300	5000	0.20	13
SSDs (read only)	64	45	30,000	0.03	13
DRAM	8	0.005	500,000	0.001	52

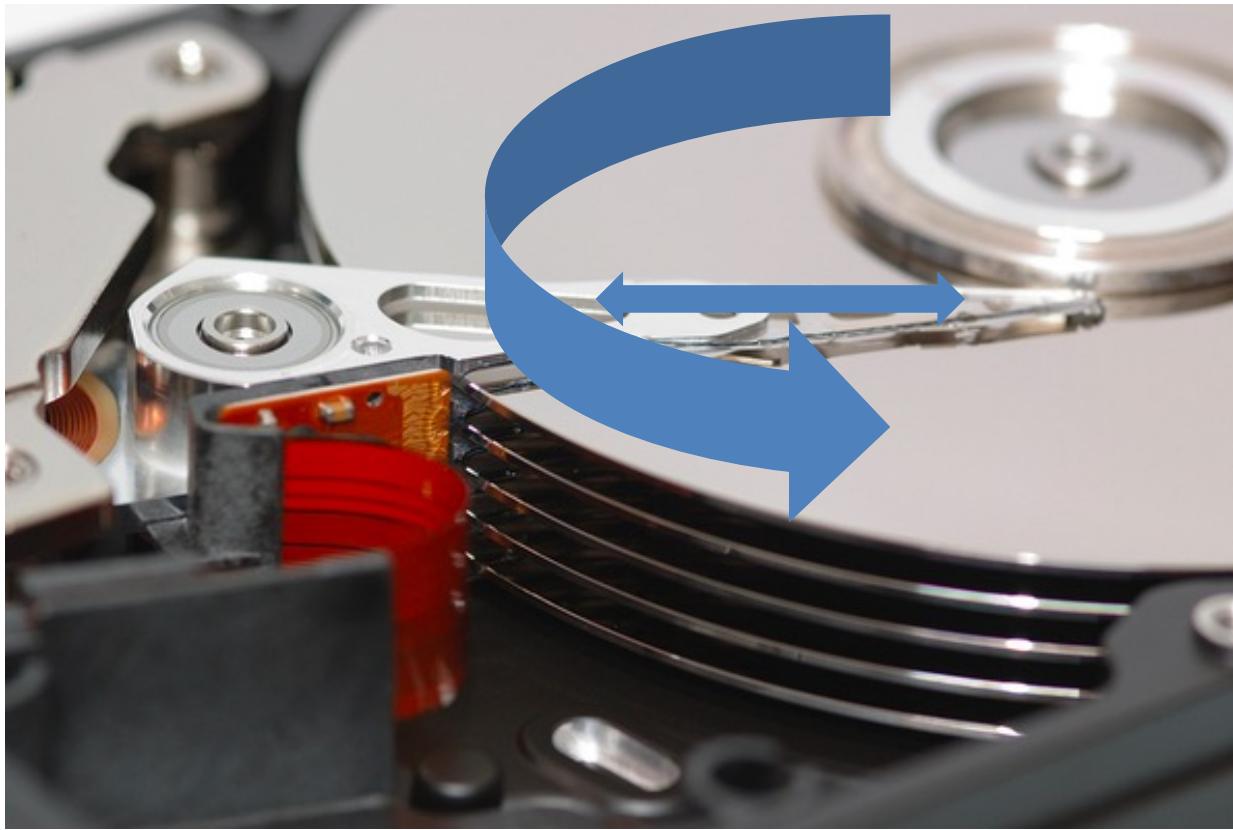
- These numbers are ancient.
- Looking for more modern numbers.
- But, does give an idea of
 - Price
 - Performance
- The general observation is that
 - Performance goes up 10X/level.
 - Price goes up 10x per level.
- Note: One major change is improved price performance of SSD relative to HDD for large data.



Storage Hierarchy (Cont.)

- **primary storage:** Fastest media but volatile (cache, main memory).
- **secondary storage:** next level in hierarchy, non-volatile, moderately fast access time
 - Also called **on-line storage**
 - E.g., flash memory, magnetic disks
- **tertiary storage:** lowest level in hierarchy, non-volatile, slow access time
 - also called **off-line storage** and used for **archival storage**
 - e.g., magnetic tape, optical storage
 - Magnetic tape
 - Sequential access, 1 to 12 TB capacity
 - A few drives with many tapes
 - Juke boxes with petabytes (1000's of TB) of storage

Hard Disk Drive



Disk Configuration

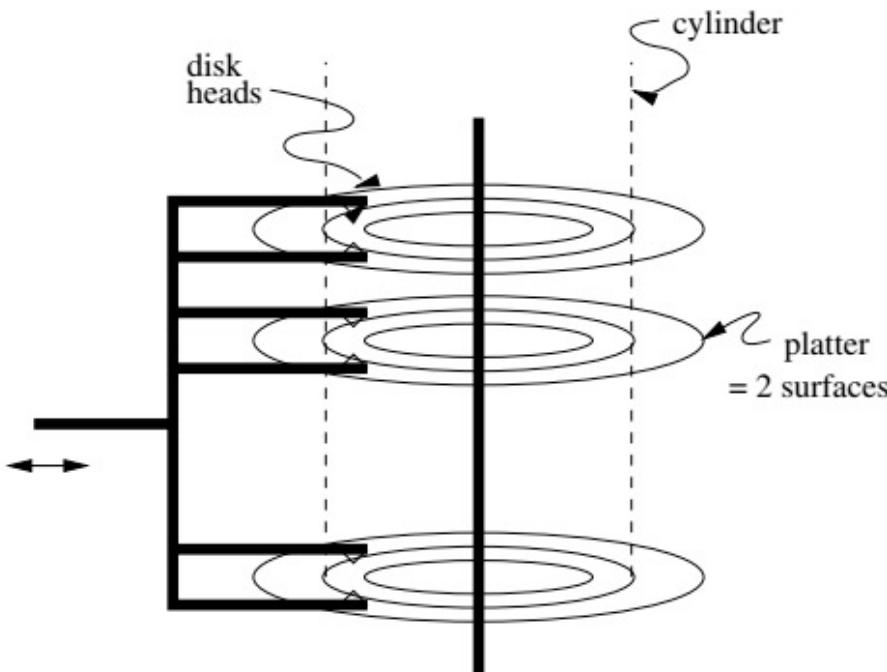


Figure 13.2: A typical disk

Components of disk I/O delay

Seek: Move head to cylinder/track.

Rotation: Wait for sector to get under head

Transfer: Move data from disk to memory.

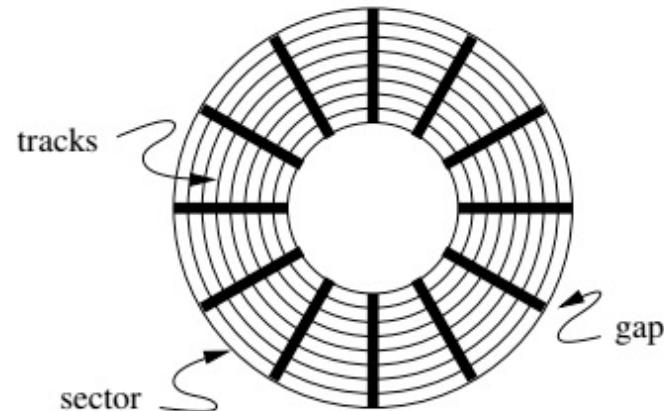
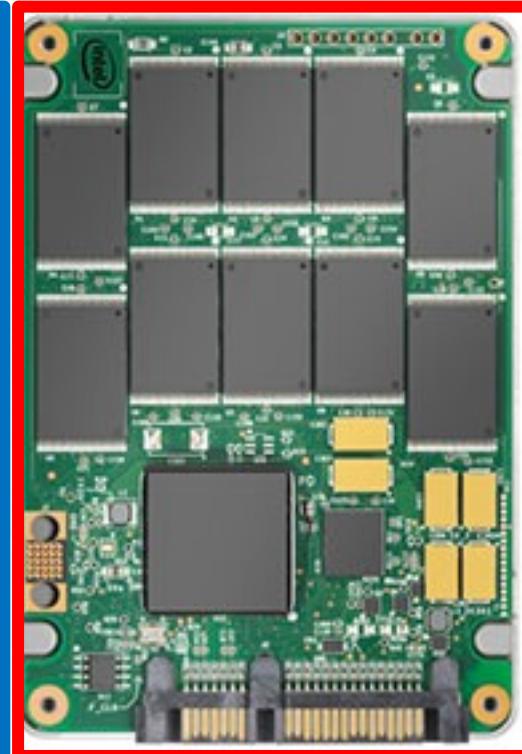


Figure 13.3: Top view of a disk surface

Database Systems: The Complete Book (2nd Edition)
by [Hector Garcia-Molina](#) (Author), [Jeffrey D. Ullman](#) (Author), [Jennifer Widom](#) (Author)

Hard Disk versus Solid State Disk

Hard
Disk
Drive



Solid
State
Drive



Flash Storage

- NOR flash vs NAND flash
- NAND flash
 - used widely for storage, cheaper than NOR flash
 - requires page-at-a-time read (page: 512 bytes to 4 KB)
 - 20 to 100 microseconds for a page read
 - Not much difference between sequential and random read
 - Page can only be written once
 - Must be erased to allow rewrite
- **Solid state disks**
 - Use standard block-oriented disk interfaces, but store data on multiple flash storage devices internally
 - Transfer rate of up to 500 MB/sec using SATA, and up to 3 GB/sec using NVMe PCIe

Logical Block Addressing

- Concept:
 - The *unit of transfer* from a “disk” to the computers memory is a “block.
Blocks are usually relatively large, e.g. 16 KB, 32 KB,
 - A program that reads or write a single byte, requires the database engine (or file system) to read/write the entire block.
- The address of a block in the entire space of blocks is:
 - (Device ID, Block ID)
 - Block ID is simple 0, 1, 2,
- The disk controller and disk implementation translate the *logical block address* into the *physical address of blocks*.
- The physical address changes over time for various reasons, e.g. performance optimization, internal HW failure, etc.

Logical/Physical Block Addressing

The mapping from LBA to physical block address can change over time.

- Internal HW failure.
- SSD writes in a funny way.
 - You have to erase before writing.
 - So, the SSD (for performance)
 - Writes to an empty block.
 - Erase the original block.
- Performance optimization on HDD
 - Based on block access patterns.
 - Place blocks on cylinder/sector/head in a way to minimize:
 - Seek
 - Rotate

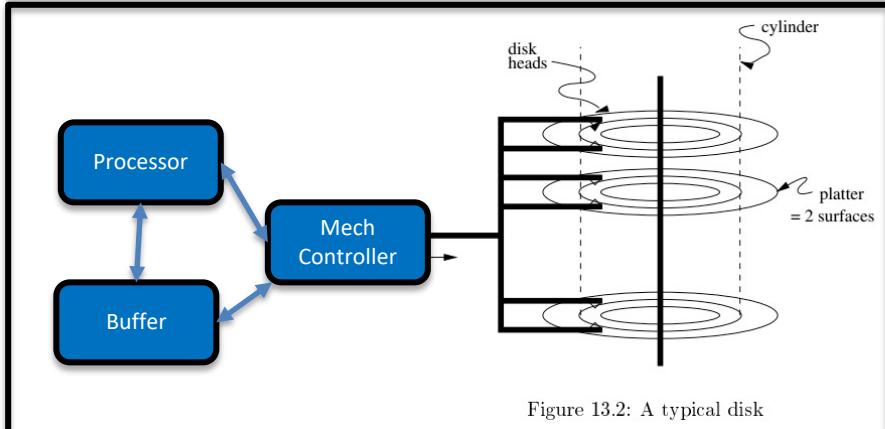
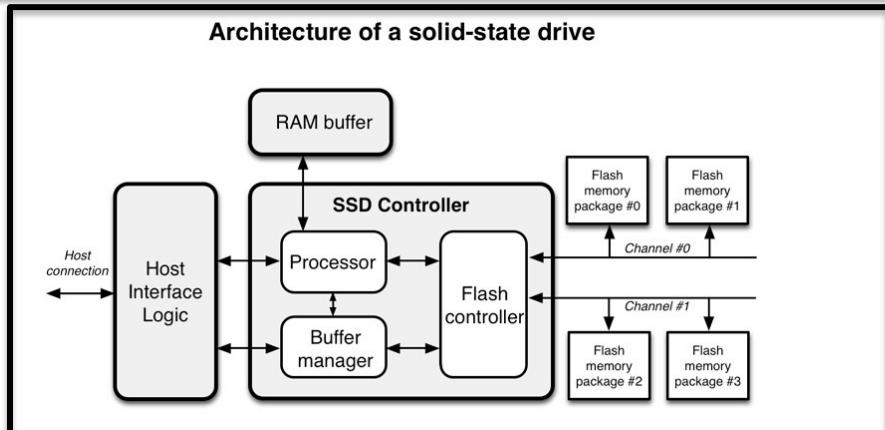
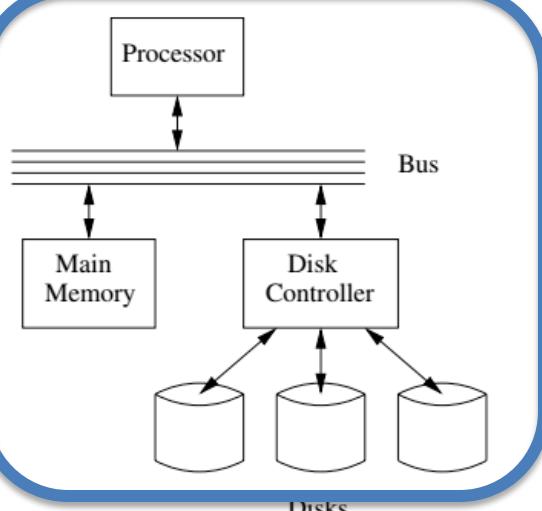


Figure 13.2: A typical disk

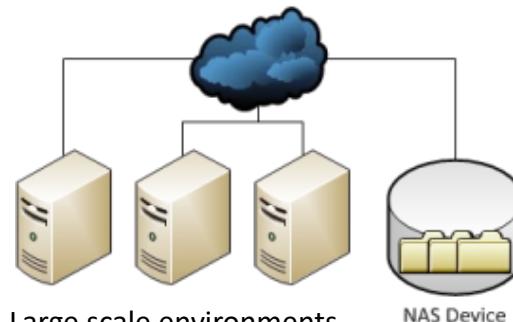
I/O Architecture



How we normally think of disks and I/O.

Network Attached Storage

- Shared storage over shared network
- File system
- Easier management

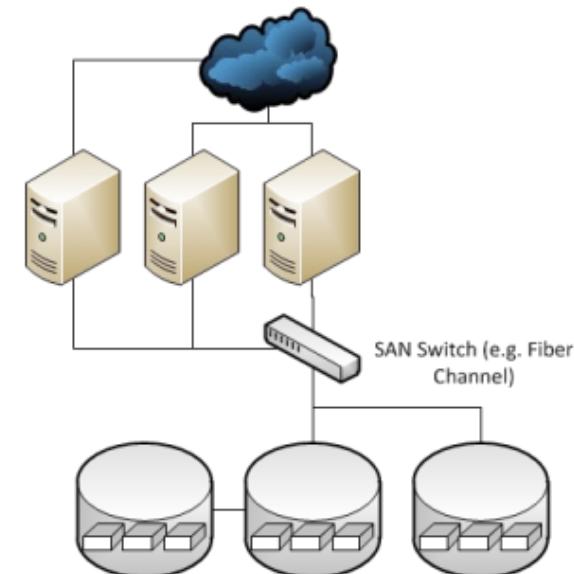


Large scale environments

- The bus-controller connection is over some kind of network.
- The disk controller is at the disks, and basically a “computer” with SW.
- Network is either
 - Standard communication network, or
 - Highly optimized I/O network.

Storage Area Network

- Shared storage over dedicated network
- Raw storage
- Fast, but costly





Magnetic Disks

- **Read-write head**
- Surface of platter divided into circular **tracks**
 - Over 50K-100K tracks per platter on typical hard disks
- Each track is divided into **sectors**.
 - A sector is the smallest unit of data that can be read or written.
 - Sector size typically 512 bytes
 - Typical sectors per track: 500 to 1000 (on inner tracks) to 1000 to 2000 (on outer tracks)
- To read/write a sector
 - disk arm swings to position head on right track
 - platter spins continually; data is read/written as sector passes under head
- Head-disk assemblies
 - multiple disk platters on a single spindle (1 to 5 usually)
 - one head per platter, mounted on a common arm.
- **Cylinder i** consists of i^{th} track of all the platters



Magnetic Disks (Cont.)

- **Disk controller** – interfaces between the computer system and the disk drive hardware.
 - accepts high-level commands to read or write a sector
 - initiates actions such as moving the disk arm to the right track and actually reading or writing the data
 - Computes and attaches **checksums** to each sector to verify that data is read back correctly
 - If data is corrupted, with very high probability stored checksum won't match recomputed checksum
 - Ensures successful writing by reading back sector after writing it
 - Performs **remapping of bad sectors**



Performance Measures of Disks

- **Access time** – the time it takes from when a read or write request is issued to when data transfer begins. Consists of:
 - **Seek time** – time it takes to reposition the arm over the correct track.
 - Average seek time is 1/2 the worst case seek time.
 - Would be 1/3 if all tracks had the same number of sectors, and we ignore the time to start and stop arm movement
 - 4 to 10 milliseconds on typical disks
 - **Rotational latency** – time it takes for the sector to be accessed to appear under the head.
 - 4 to 11 milliseconds on typical disks (5400 to 15000 r.p.m.)
 - Average latency is 1/2 of the above latency.
 - Overall latency is 5 to 20 msec depending on disk model
- **Data-transfer rate** – the rate at which data can be retrieved from or stored to the disk.
 - 25 to 200 MB per second max rate, lower for inner tracks



Performance Measures (Cont.)

- **Disk block** is a logical unit for storage allocation and retrieval
 - 4 to 16 kilobytes typically
 - Smaller blocks: more transfers from disk
 - Larger blocks: more space wasted due to partially filled blocks
- **Sequential access pattern**
 - Successive requests are for successive disk blocks
 - Disk seek required only for first block
- **Random access pattern**
 - Successive requests are for blocks that can be anywhere on disk
 - Each access requires a seek
 - Transfer rates are low since a lot of time is wasted in seeks
- **I/O operations per second (IOPS)**
 - Number of random block reads that a disk can support per second
 - 50 to 200 IOPS on current generation magnetic disks



Performance Measures (Cont.)

- **Mean time to failure (MTTF)** – the average time the disk is expected to run continuously without any failure.
 - Typically 3 to 5 years
 - Probability of failure of new disks is quite low, corresponding to a “theoretical MTTF” of 500,000 to 1,200,000 hours for a new disk
 - E.g., an MTTF of 1,200,000 hours for a new disk means that given 1000 relatively new disks, on an average one will fail every 1200 hours
 - MTTF decreases as disk ages

Performance Measures of Disks

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From: Database System Concepts, 7th Ed.

Logical Block Addressing (https://gerardnico.com/wiki/data_storage/lba)

3 - The LBA scheme

LBA	C	H	S
0	0	0	0
1	0	0	1
2	0	0	2
3	0	0	3
4	0	0	4
5	0	0	5
6	0	0	6
7	0	0	7
8	0	0	8
9	0	0	9
10	0	1	0
11	0	1	1
12	0	1	2
13	0	1	3
14	0	1	4
15	0	1	5
16	0	1	6
17	0	1	7
18	0	1	8
19	0	1	9
Cylinder 0			

LBA	C	H	S
20	1	0	0
21	1	0	1
22	1	0	2
23	1	0	3
24	1	0	4
25	1	0	5
26	1	0	6
27	1	0	7
28	1	0	8
29	1	0	9
30	1	1	0
31	1	1	1
32	1	1	2
33	1	1	3
34	1	1	4
35	1	1	5
36	1	1	6
37	1	1	7
38	1	1	8
39	1	1	9
Cylinder 1			

lowing formula:

tor number

Redundant Array of Independent Disks (RAID)



“RAID (redundant array of independent disks) is a data [storage virtualization](#) technology that combines multiple physical [disk drive](#) components into a single logical unit for the purposes of [data redundancy](#), performance improvement, or both. (...)

[RAID 0](#) consists of [striping](#), without [mirroring](#) or [parity](#). (...)

[RAID 1](#) consists of data mirroring, without parity or striping. (...)

[RAID 2](#) consists of bit-level striping with dedicated [Hamming-code](#) parity. (...)

[RAID 3](#) consists of byte-level striping with dedicated parity. (...)

[RAID 4](#) consists of block-level striping with dedicated parity. (...)

[RAID 5](#) consists of block-level striping with distributed parity. (...)

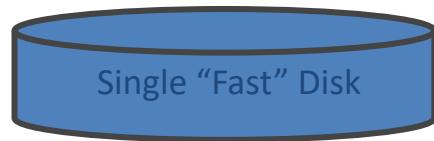
[RAID 6](#) consists of block-level striping with double distributed parity. (...)

Nested RAID

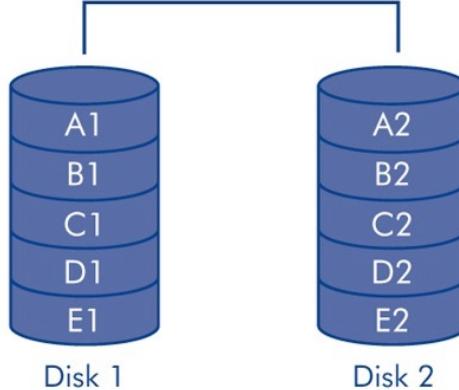
- RAID 0+1: creates two stripes and mirrors them. (...)
- RAID 1+0: creates a striped set from a series of mirrored drives. (...)
- **[JBOD RAID N+N](#)**: With JBOD (*just a bunch of disks*), (...)"

RAID-0 and RAID-1

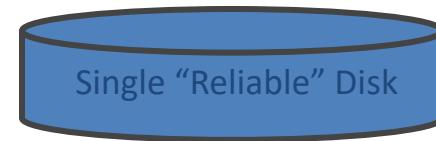
Two physical disks make
one single, logical **fast** disk



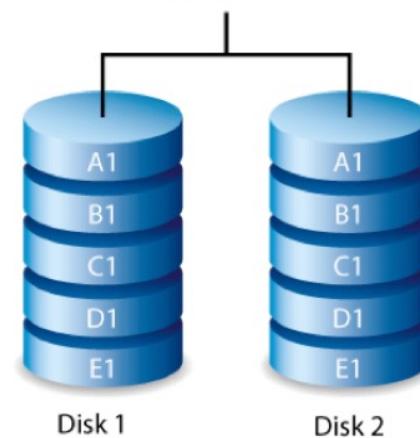
RAID 0



Two physical disks make
one single, logical **reliable** disk

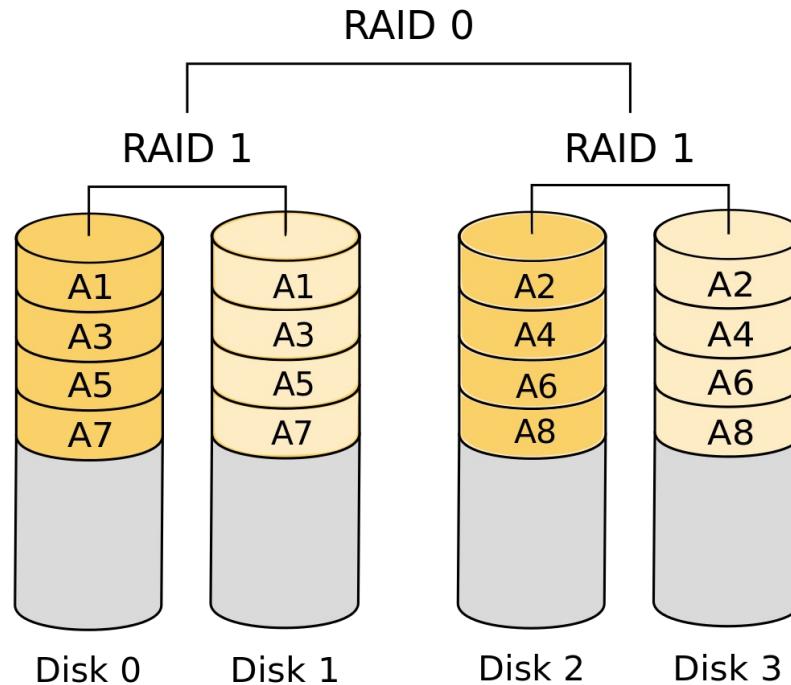


RAID 1



Mixed RAID Modes

RAID 1+0



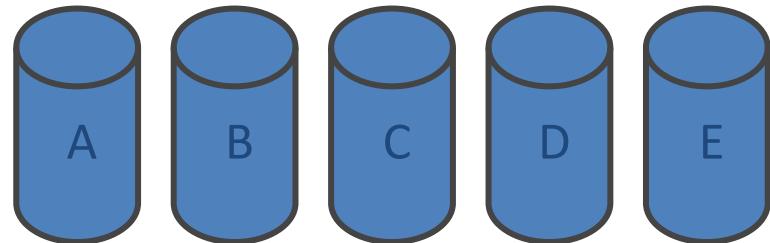
Stripe
And
Mirror

RAID-5

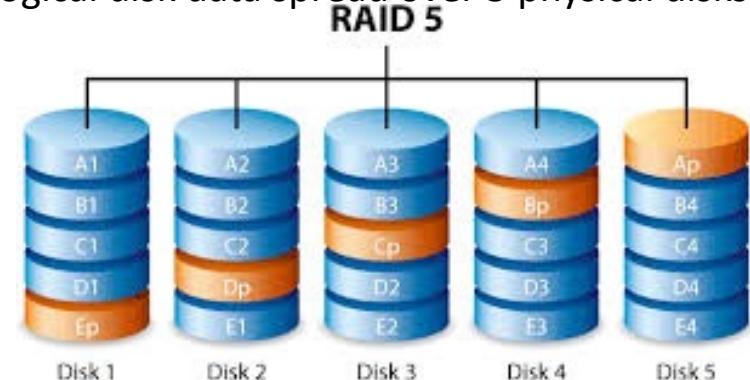
- Improved performance through parallelism
 - Rotation/seek
 - Transfer
- Availability uses *parity blocks*
 - Suppose I have 4 different data blocks on the logical drive A: A1, A2, A3, A4.
 - Parity function: $A_p = P(A_1, A_2, A_3, A_4)$
 - Recovery function: $A_2 = R(A_p, A_1, A_3, A_4)$
- During normal operations:
 - Read processing simply retrieves block.
 - Write processing of A2 updates A2 and Ap
- If an individual disk fails, the RAID
 - Read
 - Continues to function for reads on non-missing blocks.
 - Implements read on missing block by recalculating value.
 - Write
 - Updates block and parity block for non-missing blocks.
 - Computes missing block, and calculates parity based on old and new value.
 - Over time
 - “Hot Swap” the failed disk.
 - Rebuild the missing data from values and parity.



Is actually 5 smaller “logical” disks.

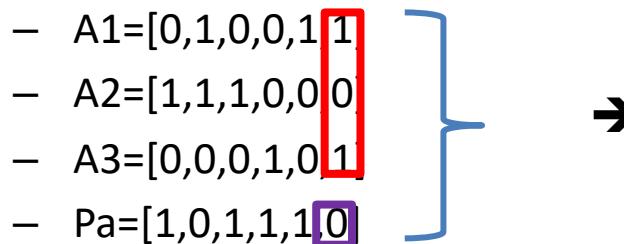


Logical disk data spread over 5 physical disks.



Very Simple Parity Example

- Even-Odd Parity
 - $b[i]$ is an array of bits (0 or 1)
 - $P(b[i]) =$
 - 0 if an even number of bits = 1. $\{P([0,1,1,0,1,1])=0\}$
 - 1 if an odd number of bits = 1. $\{P(0,0,1,0,1,1)=1\}$
 - Given an array with one missing bit and the parity bit, I can re-compute the missing bit.
 - Case 1: $[0,?,1,0,1,1]$ has $P=0$. There must be an EVEN number of ones and $?=1$.
 - Case 2: $[0,?,1,0,1,1]$ has $P=1$. There must be an ODD number of ones and $?=0$.
- Block Parity applies this to a set of blocks bitwise



If I am missing a block and have the parity block, I can re-compute the missing block bitwise from remaining blocks and parity block.

NoSQL

Concepts

Overview (I) (<https://en.wikipedia.org/wiki/NoSQL>)

A **NoSQL** (originally referring to "non SQL" or "non relational")^[1] database provides a mechanism for storage and retrieval of data that is modeled in **means other than the tabular relations used in relational databases**. Such databases have existed since the late 1960s, but did not obtain the "NoSQL" moniker until a surge of popularity in the early twenty-first century,^[2] triggered by the needs of Web 2.0 companies such as Facebook, Google, and Amazon.com.^{[3][4][5]} NoSQL databases are increasingly used in big data and real-time web applications.^[6] NoSQL systems are also sometimes called "**Not only SQL**" to emphasize that they may support SQL-like query languages.^{[7][8]}

Motivations for this approach include: simplicity of design, simpler "horizontal scaling" to clusters of machines (which is a problem for relational databases),^[2] and finer control over availability. The data structures used by NoSQL databases (e.g. key-value, wide column, graph, or document) are different from those used by default in relational databases, making **some operations faster in NoSQL**. The particular suitability of a given NoSQL database depends on the problem it must solve. Sometimes the **data structures used by NoSQL databases are also viewed as "more flexible"** than relational database tables.^[9]

Overview (I) (<https://en.wikipedia.org/wiki/NoSQL>)

Many NoSQL stores compromise [consistency](#) (in the sense of the [CAP theorem](#)) in favor of availability, partition tolerance, and speed. Barriers to the greater adoption of NoSQL stores include the use of low-level query languages (instead of SQL, for instance the lack of ability to perform ad-hoc joins across tables), lack of standardized interfaces, and huge previous investments in existing relational databases.^[10] Most NoSQL stores lack true [ACID](#) transactions,

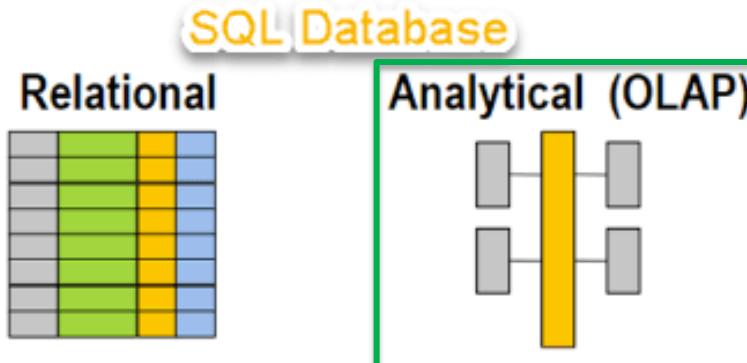
Instead, most NoSQL databases offer a concept of "eventual consistency" in which database changes are propagated to all nodes "eventually" (typically within milliseconds) so queries for data might not return updated data immediately or might result in reading data that is not accurate, a problem known as stale reads.^[11] Additionally, some NoSQL systems may exhibit lost writes and other forms of [data loss](#).^[12] Fortunately, some NoSQL systems provide concepts such as [write-ahead logging](#) to avoid data loss.^[13] For [distributed transaction processing](#) across multiple databases, data consistency is an even bigger challenge that is difficult for both NoSQL and relational databases. Even current relational databases "do not allow referential integrity constraints to span databases."^[14]

Simplistic Classification

(<https://medium.com/swlh/4-types-of-nosql-databases-d88ad21f7d3b>)

Relational is the foundational model.

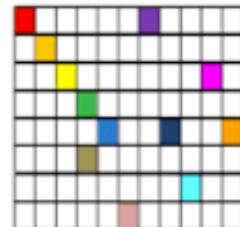
We covered graphs and examples.



We will see OLAP in a future lecture.

Subject of this lecture and part of HW3

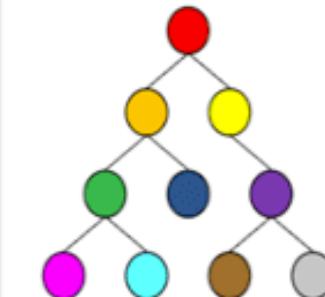
Column-Family



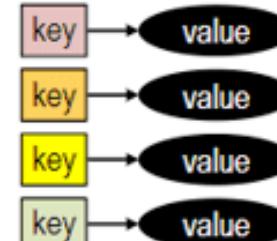
Graph



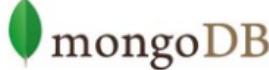
Document



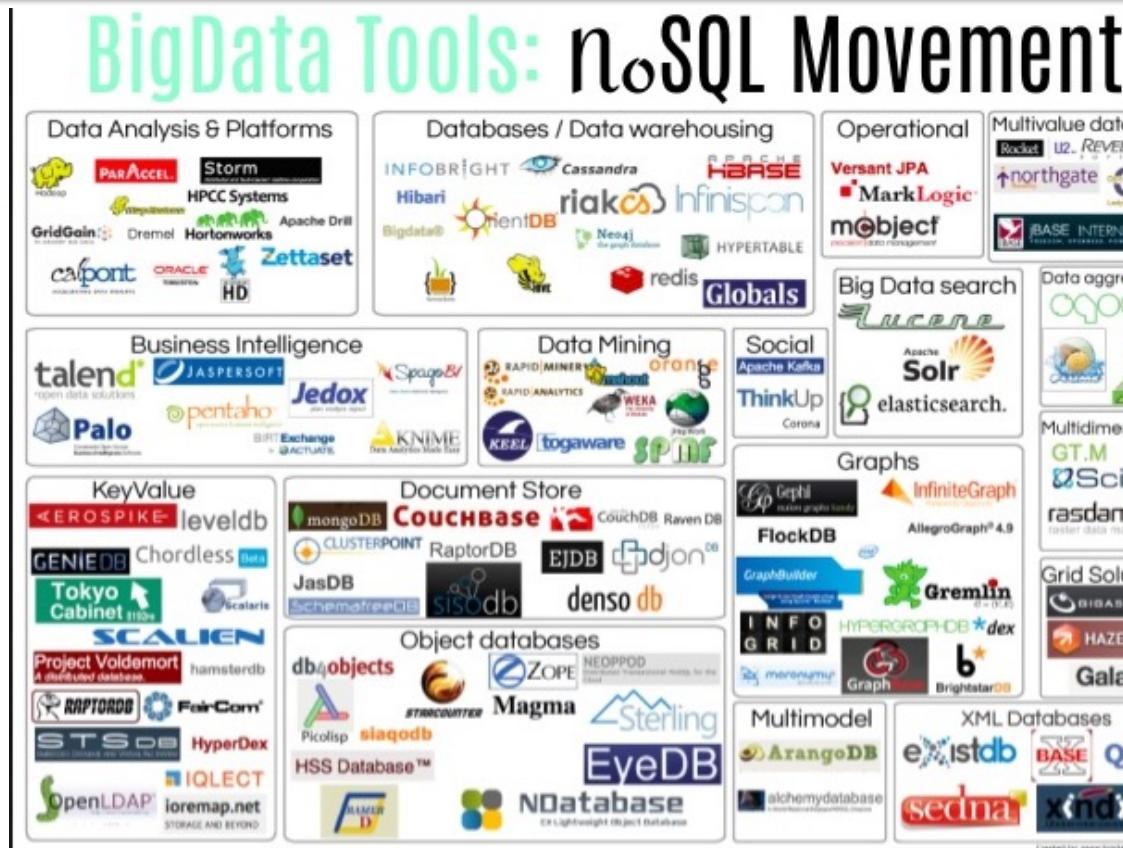
Key-Value



One Taxonomy

Document Database	Graph Databases
  	 
Wide Column Stores	Key-Value Databases
  	    

Another Taxonomy



Use Cases

Motivations

- Massive write performance.
- Fast key value look ups.
- Flexible schema and data types.
- No single point of failure.
- Fast prototyping and development.
- Out of the box scalability.
- Easy maintenance.

What is wrong with SQL/Relational?

- Nothing. One size fits all? Not really.
- Impedance mismatch. – Object Relational Mapping doesn't work quite well.
- Rigid schema design.
- Harder to scale.
- Replication.
- Joins across multiple nodes? Hard.
- How does RDMS handle data growth? Hard.
- Need for a DBA.
- Many programmers are already familiar with it.
- Transactions and ACID make development easy.
- Lots of tools to use.

Documents

Example Document – An Order

FOOD ORDER FORM TEMPLATE

Company Name
123 Main Street
Hamilton, OH 44416
(321) 456-7890
Email Address
Point of Contact
web address

YOUR LOGO

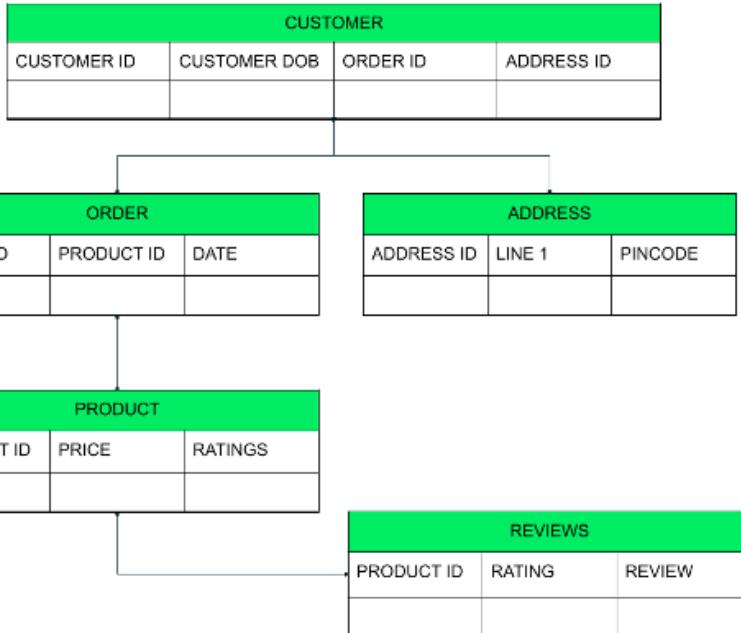
ORDER FORM

CUSTOMER		ORDER NO.	ORDER DATE
ATTN: Name / Dept		DATE NEEDED	TIME NEEDED
Company Name		ORDER RECEIVED BY	
123 Main Street			
Hamilton, OH 44416			
(321) 456-7890			
Email Address			
DESCRIPTION			
		UNIT PRICE	QUANTITY
		\$ -	
		\$ -	
		\$ -	
		\$ -	
		\$ -	
		\$ -	
		\$ -	
		\$ -	
		\$ -	
		\$ -	
		\$ -	
		\$ -	
TOTALS			
ENTER PERCENTAGE		TAX RATE	0.000%
ENTER INITIAL PAYMENT AMOUNT		TOTAL TAX	\$ -
		DELIVERY	\$ -
		GRAND TOTAL	\$ -
		LESS PAYMENT	\$ -
THANK YOU!			

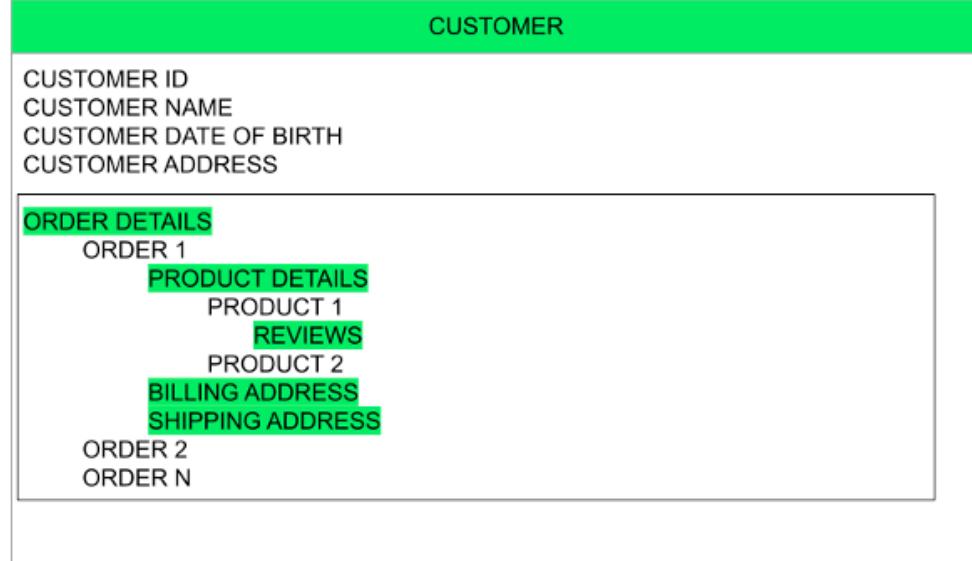
- There are 5 entity types on the form:
 - A “copy of” customer information.
 - Links (via productCode) to products
 - Order
 - OrderDetails
 - Comments
- But OrderDetails and Comments are somehow different from the others.
 - These are arrays of objects
 - That are sort of “inside” the order.
- These are weak entities, but relational does not always handle these well.

Relational vs Document

Relational



Relational



The key difference?

- Relational DB attributes are atomic.
- Document DB attributes may be *multi-valued* and *complex*.

UI Interface and Logical Data Model

Columbia University

CUSTOMER NUMBER	10001				
CONTACT NAME	CONTACT NUMBER				
Lord Voldemort	+1 212-555-6666				
Customer Address (THIS WAS INSANELY PAINFUL TO LAYOUT)					
STREET 1	STREET 2				
520 W 120th St.	4th Floor				
CITY	STATE	COUNTRY	POSTAL CODE		
New York	NY	US	10027		
I nearly went "postal" doing the previous layout and almost failed all of you!"					
NUMBER	STATUS	ORDER DATE	REQUIRED DATE	SHIPPED DATE	DETAILS
21	Shipped	2022-01-01	2022-02-01	2022-01-28	<input type="button" value="Expand"/>
LINE NUMBER	PRODUCT CODE	QUANTITY	PRICE EACH		
0	P0	0	0		
1	P1	2	3		
2	P2	4	6		
NUMBER	STATUS	ORDER DATE	REQUIRED DATE	SHIPPED DATE	DETAILS
22	Pending	2022-03-01	2022-04-01	0000000	<input type="button" value="Expand"/>

Note to DFF: Start

- [~/Dropbox/000/000-A-Current-Examples/W4111-FastAPI-IMDB-Solution](#)
- [~/Dropbox/000/000-A-Current-Examples/current-dashboard](#)

```
{  
    customerName: "Columbia University",  
    contact: {  
        name: "Lord Voldemort",  
        phone: "+1 212-555-6666"  
    },  
    address: {  
        line1: "520 W 120th St."  
        line2: "Floor 4",  
        ... ...  
    },  
    orders: [  
        {  
            number: ...,  
            ... ...  
            details: [  
                ... ...  
            ]  
        }  
    ]  
}
```

Switch to Notebook



MongoDB

MongoDB Concepts

RDBMS	MongoDB
Database	Database
Table	Collection
Tuple/Row	Document
column	Field
Table Join	Embedded Documents
Primary Key	Primary Key (Default key <code>_id</code> provided by MongoDB itself)
Database Server, Client, Tools, Packages	
mysqld/Oracle	<code>mongod</code>
mysql/sqlplus	<code>mongo</code>
DataGrip	Compass
pymysql	<code>pymongo</code>

Core Operations

Basic Operations:

- Create database
- Create collection
- Create-Retrieve-Update-Delete (CRUD):
 - Create: insert()
 - Retrieve:
 - find()
 - find_one()
 - Update: update()
 - Delete: remove()

More Advanced Concepts:

- Limit
- Sort
- Aggregation Pipelines
 - Merge
 - Union
 - Lookup
 - Match
 - Merge
 - Sample
 -

We will just cover the basics for now and may cover more things in HW or other lectures.

find()

- Note:
 - MongoDB uses a more `pymysql` approach, e.g. an API, than pure declarative languages like SQL.
 - The parameters for `find()` are where the declarative language appears.
- The basic forms of `find()` and `find_one()` have two parameters:
 - *filter expression*
 - *Project expression*
- You can use the Compass tool and screen captures for some HW and exam answers.
- What if I want the answer in a Jupyter Notebook?

The screenshot shows the MongoDB Compass interface with the following details:

- Collection:** GOT.seasons
- Documents:** 8 (TOTAL SIZE 1.0MB, AVG. SIZE 130.2KB)
- Indexes:** 1 (TOTAL SIZE 20.0KB, AVG. SIZE 20.0KB)
- Filter:** `{"episodes.scenes.location": "The Dothraki Sea"}`
- Project:** `{ season: 1, "episodes.episodeNum":1, "episodes.episodeLink": 1, "episodes.episodeTitle": 1 }`
- Sort:** `{ field: -1 }`
- Collation:** `{ locale: 'simple' }`
- Options:** MAX TIME MS 60000, SKIP 0, LIMIT 0
- Results:** Displays three documents (1-3 of 3). Each document has an _id, season (1), and episodes array. The first document's episodes array is expanded to show 10 individual episode objects, each with episodeNum, episodeTitle, and episodeLink.

Generate Code

The screenshot shows the MongoDB Compass interface. On the left, the sidebar lists databases (Local, 7 DBs) and collections (6 COLLECTIONS). The 'GOT.seasons' collection is selected. In the main area, a document is displayed with fields: _id, season, and episodes. Below the document, two results are shown, each with an _id, season, and episodes field. A modal window titled 'Export Query To Language' is open. It contains a 'My Query:' section with the following query:

```
1 {  
2   "episodes.scenes.location": "The Dothraki Sea"  
3 }
```

Below this, the 'Export Query To:' dropdown is set to 'PYTHON 3'. The generated Python code is:

```
1 # Requires the PyMongo package.  
2 # https://api.mongodb.com/python/current  
3  
4 client = MongoClient('mongodb://localhost:27017/?  
5 filter={  
6   'episodes.scenes.location': 'The Dothraki Sea'  
7 }  
8 project={  
9   'season': 1,  
10  'episodes.episodeNum': 1,  
11  'episodes.episodeLink': 1,  
12  'episodes.episodeTitle': 1  
13 }  
14  
15 result = client['GOT']['seasons'].find(  
16   filter=filter,
```

At the bottom of the modal, there are two checkboxes: 'Include Import Statements' (unchecked) and 'Include Driver Syntax' (checked). A 'CLOSE' button is at the bottom right.

- Choose Export to Language.
- Copy into the notebook.
- The export has an option to include all the connection setup, choosing DB,
- Switch to Notebook

Result is not Quite You Expect

GOT.seasons

DOCUMENTS 8 TOTAL SIZE 1.0MB AVG. SIZE 130.2KB INDEXES 1 TOTAL SIZE 20.0KB AVG. SIZE 20.0KB

Documents Aggregations Schema Explain Plan Indexes Validation

FILTER { "episodes.scenes.location": "The Dothraki Sea" }
PROJECT { \$season: 1, "episodes.episodeNum": 1, "episodes.episodeLink": 1, "episodes.episodeTitle": 1 }
SORT { field: -1 }
COLLATION { locale: 'simple' }

FIND RESET ...
MAX TIME MS 60000
SKIP 0 LIMIT 0

VIEW Refresh

Displaying documents 1 - 3 of 3 < > C REFRESH

```
_id:ObjectId("60577e50c68b67110968b6d1")
season:"1"
episodes:Array
  ▾ 0:Object
    episodeNum: 1
    episodeTitle: "Winter Is Coming"
    episodeLink: "/title/tt1480055/"
  ▾ 1:Object
  ▾ 2:Object
  ▾ 3:Object
    episodeNum: 4
    episodeTitle: "Cripples, Bastards, and Broken Things"
    episodeLink: "/title/tt1829963/"
  ▾ 4:Object
  ▾ 5:Object
  ▾ 6:Object
  ▾ 7:Object
  ▾ 8:Object
  ▾ 9:Object

_id:ObjectId("60577e50c68b67110968b6d5")
season:"5"
episodes:Array

_id:ObjectId("60577e50c68b67110968b6d6")
season:"6"
episodes:Array
```

- The query returns documents that match.
 - The document is “Large” and has seasons and episodes and seasons.
 - If you do a \$project requesting episodes/episode content,
 - You get all episodes in the documents that match.
 - Not just the episodes with the scene/location.
 - Projecting array elements from arrays whose elements are arrays is complex and baffling.
- You also get back something (a cursor) that is iterable.

Result is not Quite You Expect

The screenshot shows the MongoDB Compass interface. At the top, there's a navigation bar with tabs for 'Documents', 'Aggregations', 'Schema', 'Explain Plan', 'Indexes', and 'Validation'. Below the navigation bar, there are sections for 'DOCUMENTS' (8 documents, 1.0MB total size, 130.2KB avg size) and 'INDEXES' (1 index, 20.0KB total size, 20.0KB avg size). The main area displays a table of query results. The first row of the table has a 'FILTER' column containing the query: { "episodes.scenes.location": "The Dothraki Sea" }, and a 'PROJECT' column containing the projection: { \$project: { season: 1, episodes.episodeNum: 1, episodes.episodeLink: 1, episodes.episodeTitle: 1 } }. The table also includes columns for 'SORT' (with a field of -1), 'COLLATION' (with a locale of 'simple'), and 'OPTIONS' (with MAX TIME MS set to 60000). At the bottom of the table, there are buttons for 'FIND', 'RESET', and '...', along with 'SKIP' (0) and 'LIMIT' (0) controls.

Net for HWs and exams:

- We will keep the queries simple.
- The language is as complex as SQL, and we spent several weeks on the language.
- Let's take a look in the Jupyter Notebook on some create and insert functions.
- But, first, the datamodel impedance mismatch concept.

Raw JSON documents from the results table:

```
_id:ObjectId("60577e50c68b67110968b6d1")
season:""
episodes:Array
  ▾ 0:Object
    episodeNum: 1
    episodeTitle: "Winter Is Coming"
    episodeLink: "/title/tt1480055/"
  ▾ 1:Object
  ▾ 2:Object
  ▾ 3:Object
    episodeNum: 4
    episodeTitle: "Cripples, Bastards, and
    episodeLink: "/title/tt1829963/"
  ▾ 4:Object
  ▾ 5:Object
  ▾ 6:Object
  ▾ 7:Object
  ▾ 8:Object
  ▾ 9:Object

_id:ObjectId("60577e50c68b67110968b6d5")
season:""
episodes:Array

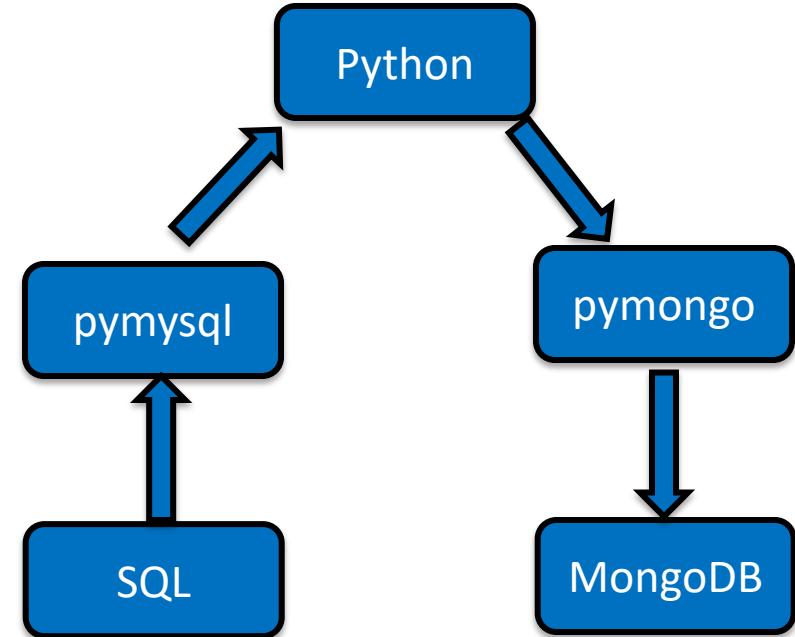
_id:ObjectId("60577e50c68b67110968b6d6")
season:""
episodes:Array
```

- The query returns documents that match.
 - The document is “Large” and has many episodes and seasons.
- You do a \$project requesting episodes/episode content,
 - You get all episodes in the documents that match.
- Not just the episodes with the scene/location.
- Projecting array elements from arrays whose elements are arrays is complex and baffling.
- You also get back something (a cursor) that is iterable.

More Fun – Data Types

MongoDB supports many datatypes. Some of them are –

- **String** – This is the most commonly used datatype to store the data. String in MongoDB must be UTF-8 valid.
- **Integer** – This type is used to store a numerical value. Integer can be 32 bit or 64 bit depending upon your server.
- **Boolean** – This type is used to store a boolean (true/ false) value.
- **Double** – This type is used to store floating point values.
- **Min/ Max keys** – This type is used to compare a value against the lowest and highest BSON elements.
- **Arrays** – This type is used to store arrays or list or multiple values into one key.
- **Timestamp** – timestamp. This can be handy for recording when a document has been modified or added.
- **Object** – This datatype is used for embedded documents.
- **Null** – This type is used to store a Null value.
- **Symbol** – This datatype is used identically to a string; however, it's generally reserved for languages that use a specific symbol type.
- **Date** – This datatype is used to store the current date or time in UNIX time format. You can specify your own date time by creating object of Date and passing day, month, year into it.
- **Object ID** – This datatype is used to store the document's ID.
- **Binary data** – This datatype is used to store binary data.
- **Code** – This datatype is used to store JavaScript code into the document.
- **Regular expression** – This datatype is used to store regular expression.



(Some) MongoDB CRUD Operations

- Create:
 - db.collection.insertOne()
 - db.collection.insertMany()
- Retrieve:
 - db.collection.find()
 - db.collection.findOne()
 - db.collection.findOneAndUpdate()
 -
- Update:
 - db.collection.updateOne()
 - db.collection.updateMany()
 - db.collection.replaceOne()
- Delete:
 - db.collection.deleteOne()
 - db.collection.deleteMany()

pymongo maps the camel case to _, e.g.

- findOne()
- find_one()

There are good online tutorials:

- https://www.tutorialspoint.com/python_data_access
- <https://www.tutorialspoint.com/mongodb/index.htm>

(Some) MongoDB Pipeline Operators

<https://www.slideshare.net/mongodb/s01-e04-analytics>

Aggregation operators

- Pipeline and Expression operators

Pipeline	Expression	Arithmetic	Conditional
\$match	\$addToSet	\$add	\$cond
\$sort	\$first	\$divide	\$ifNull
\$limit	\$last	\$mod	
\$skip	\$max	\$multiply	
\$project	\$min	\$subtract	
\$unwind	\$avg		Variables
\$group	\$push		
\$geoNear	\$sum		
\$text			\$let
\$search			\$map

Tip: Other operators for date, time, boolean and string manipulation

MongoDB Checkpoint

- You can see that MongoDB has powerful, sophisticated
 - Operators
 - Expressions
 - Pipelines
- We have only skimmed the surface. There is a lot more:
 - Indexes
 - Replication, Sharding
 - Embedded Map-Reduce support
 -
- We will explore a little more in subsequent lectures, homework,
- You will have to install MongoDB and Compass for HW4 and final exam.