

Notational conventions

- \mathcal{H} a complex hilbertspace
- Tr The trace of an endomorphism (or square matrix)
- $V \otimes W$ the tensor-product of V and W . This will be treated as the Kronecker-Product on all finite dimensional vector spaces.
- X, Y, Z Pauli matrices (see def. 1.4)
- $\mathcal{E}_n := \{\sigma_1 \otimes \dots \otimes \sigma_n, \text{ where } \sigma_i \in \{I, X, Y, Z\} \text{ for all } i \in \underline{n}\}.$
- $D_n = \langle a, b \mid a^2 b^n (ab)^2 \rangle$ the dihedral group on a regular polygon with n sides.
- V^* : the dual space to V .
- $\text{Hom}_G(A, B)$ the set of G-Homomorphisms from A to B (i.e., G-equivariant homomorphisms)
- $G \wr H$ the wreath product of G and H
- $\omega \in \mathbb{F}_4$ a root of $x^2 + x + 1 \in \mathbb{F}_2[x]$
- Im the image a mapping, Ker the kernel of a homomorphisms
- $\sigma|_S$ denotes $\sigma_{s_1} \otimes \dots \otimes \sigma_{s_k}$, where $S = \{s_1, \dots, s_k\} \subseteq \underline{n}$.

Contents

1	Preliminaries	3
1.1	Remarks on linear algebra	3
1.2	The Pauli Group	3
1.3	Random Unitary Matrices	4
1.4	Physical background: The postulates of quantum mechanics	6
2	Qubits and Quantum Codes	7
2.1	Error correction	7
2.2	Stabilizer codes	8
2.3	Logical Operators	9
2.4	Code distance	9
2.5	CSS Codes	10
3	Quantum Weight Enumerators	11
3.1	The Shor-Laflamme and Rains enumerators	11
3.2	The relationship between the enumerators	13
3.3	The MacWilliams transform	15
4	Applications of the weight enumerators: Bounds	15
4.1	The relationship between the code distance and the enumerators	15
4.2	Shadow enumerators	18

1 Preliminaries

1.1 Remarks on linear algebra

Throughout this thesis we will focus mainly on groups acting on tensor product spaces of a particular finite-dimensional complex Hilbert space. Recall that a complex Hilbert space \mathcal{H} is a complete complex vector space with a complex scalar product $\langle \cdot, \cdot \rangle$, i.e. a positive-definite, hermitian, bilinear form. Every finite dimensional complex vector space with a scalar product is complete, and as we will only work with such in these thesis, we need not worry about the completeness.

Remark 1.1. The trace function $\text{Tr} : \mathbb{C}^{n \times n} \rightarrow \mathbb{C}$ induces an associative, symmetric, bilinear form on $\mathbb{C}^{n \times n}$, the trace bilinear form: $\text{Tr} : \mathbb{C}^{n \times n} \times \mathbb{C}^{n \times n} \rightarrow \mathbb{C}, (A, B) \mapsto \text{Tr}(AB)$. This means the following:

- $\text{Tr}(kA + B, C) = k\text{Tr}(A, C) + \text{Tr}(B, C)$ for all $A, B, C \in \mathbb{C}^{n \times n}, k \in \mathbb{C}$.
- $\text{Tr}(A, B) = \text{Tr}(B, A)$ for all $\mathbb{C}^{n \times n}$.
- $\text{Tr}(AB, C) = \text{Tr}(B, CA)$ for all $A, B, C \in \mathbb{C}^{n \times n}$.

All of which immediately follow from the definition.

Definition 1.2. Let V, W be finite-dimensional Hilbert spaces, and $A \in \text{End}(V), B \in \text{End}(W)$ Hermitian. We define the partial trace of $A \otimes B$ over W as $\text{Tr}_W(A \otimes B) := \text{Tr}(B)A$, and extend the definition linearly. Since the trace is linear, this is well defined.

Remark 1.3. With the notation from def. 1.2, the partial trace of $A \in \text{End}(V \otimes W)$ Hermitian, is the unique Hermitian operator $\text{Tr}_W(A) \in \text{End}(V)$ which satisfies:

$$\text{Tr}(A(B \otimes \text{Id}_W)) = \text{Tr}(\text{Tr}_W(A)B) \text{ for all } B \in \text{End}(W) \text{ Hermitian} \quad (1)$$

Further for a Hilbert space with a particular decomposition $V_1 \otimes \dots \otimes V_n$, e.g. $(\mathbb{C}^2)^{\otimes n}$ with $V_i \cong \mathbb{C}^2$ for all $i \in \underline{n}$, and $S \cup S^c = \underline{n}, S = \{s_1 < \dots < s_j\}, S^c = \{s_1^c < \dots < s_k^c\}$ we define $V_S := V_{s_1} \otimes \dots \otimes V_{s_j}, V_{S^c} := V_{s_1^c} \otimes \dots \otimes V_{s_k^c}$. And use Tr_S as a short notation for $\text{Tr}_{V_S} : \text{End}(V_1 \otimes \dots \otimes V_n) \rightarrow \text{End}(V_{S^c})$. Note that for $s_j \geq s_1^c$ this is not entirely correct, since we cannot write V as $V_S \otimes V_{S^c}$, but we get a canonical isomorphism from reordering the tensor product components, and by abuse of notation we will always identify V with its image under it, such that we can always write $V = V_S \otimes V_{S^c}$. With this notation, and the definition of the partial trace it is easy to see that $\text{Tr}_{S \cup S'}(A) = \text{Tr}_S(\text{Tr}_{S'}(A))$ for all $S, S' \subseteq \underline{n}$ disjoint.

Proof. $\text{Tr}_W(A)$ is hermitian since A already is. Let $A = \sum_{i=1}^k V_i \otimes W_i$ with $V_i \in \text{End}(V), W_i \in \text{End}(W)$ Hermitian for all $i \in \underline{k}$. Then $\text{Tr}_W(A) = \sum_{i=1}^k \text{Tr}(W_i)V_i$ satisfies for all $B \in \text{End}(W)$ Herm.:

$$\text{Tr}(A(B \otimes \text{Id}_W)) = \text{Tr}\left(\sum_{i=1}^k V_i B \otimes W_i\right) = \sum_{i=1}^k \text{Tr}(V_i B) \text{Tr}(W_i) = \text{Tr}\left(\sum_{i=1}^k V_i B \text{Tr}(W_i)\right) = \text{Tr}(\text{Tr}_W(A)B)$$

Further is $\text{Tr}_W(A)$ unique with this property, since the restriction of trace bilinear form to hermitian operators is non-degenerate. Assuming there is a $T \in \text{End}(V)$ hermitian with $\text{Tr}(TB) = \text{Tr}(\text{Tr}_W(A)B)$ for all $B \in \text{End}(W)$ Herm., then it follows that $\text{Tr}((T - \text{Tr}_W(A))B) = 0$ for all such B which implies that $T - \text{Tr}_W(A) = 0$. \square

1.2 The Pauli Group

Definition 1.4. The Pauli matrices, or Pauli operators on \mathbb{C}^2 are defined by:

$$\text{Id}, X := \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, Z := \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}, Y := iXZ = \begin{bmatrix} 0 & i \\ -i & 0 \end{bmatrix} \quad (2)$$

We denote with

$$\mathcal{E}_n := \{\sigma_1 \otimes \dots \otimes \sigma_n, \text{ where } \sigma_i \in \{I, X, Y, Z\} \text{ for all } i \in \underline{n}\} \subseteq \text{End}((\mathbb{C}^2)^{\otimes n}) := \text{End}(\underbrace{\mathbb{C}^2 \otimes \mathbb{C}^2}_{n \text{ times}}) \quad (3)$$

the set of all so-called Pauli operators, and by $\mathcal{P}_n := \langle \mathcal{E} \rangle \leq \text{GL}_n(\mathbb{C})$, we denote the generated subgroup, which we will call the Pauli group on n qubits.

Remark 1.5. The set $\{\text{Id}, X, Y, Z\}$ is linearly independent and thus forms a basis of $\mathbb{C}^{2 \times 2}$. A simple calculation shows that this basis is orthogonal with respect to the trace bilinear form on $\mathbb{C}^{2 \times 2}$. Analogously, since $\text{Tr}(A \otimes B) = \text{Tr}(A)\text{Tr}(B)$ for all $A \in \mathbb{C}^{m \times m}, B \in \mathbb{C}^{m' \times m'}$, \mathcal{E} is an orthogonal basis of $\mathbb{C}^{2^n \times 2^n}$.

Proposition 1.6. Let $G = \mathcal{P}_n$ be the Pauli group on n qubits. Then the following is true:

- a) $Z(G) = \{\pm \text{Id}, \pm i \text{Id}\}$, where $Z(G)$ denotes the centre of G .
- b) $G' = \pm \{\text{Id}\} \leq Z(G)$, where G' denotes the commutator subgroup of G .
- c) $G/Z(G)$ is a $2n$ -dimensional \mathbb{F}_2 vectorspace
- d) The commutator $[\cdot, \cdot] : G \times G \rightarrow G$ induces a non-degenerate symplectic form on $G/Z(G)$: $[gZ(G), hZ(G)] \mapsto [g, h]$.

Proof. a) Let $A \in Z(G)$. Since $\langle G \rangle_{\mathbb{C}} = \mathbb{C}^{2^n \times 2^n}$ it means that $A \in \underbrace{Z(\mathbb{C}^{2^n \times 2^n})}_{\text{CId}} \cap G = \{\pm \text{Id}, \pm i \text{Id}\}$.

- b) It follows directly from the calculation for X, Y, Z and by induction on n .
- c) Since $G' \leq Z(G)$ is $G/Z(G)$ abelian. By induction on n it can be easily shown that $G^2 \subseteq Z(G)$, hence $G/Z(G)$ is elementary abelian. Since $|G| = 4^n \cdot 4^n = 2^{2n+2}$ it follows that $\dim_{\mathbb{F}_2}(G/Z(G)) = 2n$.
- d) Since $G' \leq Z(G)$ the induced mapping is well-defined. Since $[g, g] = 1$ for all $g \in G$ it follows that $[\cdot, \cdot]$ is alternating on $G/Z(G)$. \mathbb{F}_2 bilinearity follows directly from the definition. \square

Remark 1.7. Let $G = \mathcal{P}_n$, $G \ni g = i^k X_1^{a_1} \dots X_n^{a_n} Y_1^{b_1} \dots Y_n^{b_n}$, $k \in \underline{4}, a_i, b_i \in \{0, 1\}$ for all $i \in \underline{n}$. Note that by prop. 1.6 every element in G can be written than way. Then an \mathbb{F}_2 isomorphism $\varphi : G/G' \rightarrow \mathbb{F}_2^{2n}$ is given by:

$$\varphi(g = i^k X_1^{a_1} \dots X_n^{a_n} Y_1^{b_1} \dots Y_n^{b_n}) = (a_1, \dots, a_n, b_1, \dots, b_n) \in \mathbb{F}_2^{2n}$$

The image under φ is sometimes referred to as symplectic notation. Further let $\mathbb{F}_4 = \mathbb{F}_2[\omega]$, with $\omega^2 + \omega + 1 = 0$, equipped with the \mathbb{F}_2 -symplectic form $(x, y) \mapsto \text{Tr}(x\bar{y})$, where $(x \mapsto \bar{x})$ denotes the Galois automorphism from \mathbb{F}_4 : $(x \mapsto x^2)$. Note that since $\text{Tr}(x\bar{y}) = (\bar{x}\bar{y}) + (x\bar{y}) = x^2y + xy^2 = xy^2 + x^2y = \text{Tr}(y\bar{x}) (= -\text{Tr}(y\bar{x}))$ this is indeed a symplectic form. Then we can identify $\mathbb{F}_2^{2n} \cong_{\mathbb{F}_2} \mathbb{F}_4^n$ in a canonical way with respect to the symplectic notation: $(a_1, \dots, a_n, b_1, \dots, b_n) \mapsto (\omega a_1 + \bar{\omega} b_1, \dots, \omega a_n + \bar{\omega} b_n)$, and call the resulting isomorphism $\tilde{\varphi}$. $\tilde{\varphi}$ is an isometry with respect to the commutator and the \mathbb{F}_2 -symplectic form on \mathbb{F}_4 .

1.3 Random Unitary Matrices

Theorem 1.8 (Haar Measure). Let $\mathcal{U} := \mathcal{U}_n(\mathbb{C})$ denote the group of unitary $n \times n$ matrices over the complex field \mathbb{C} . Then there exists a unique measure μ on \mathcal{U} , called the Haar measure, which satisfies the following:

- μ is a countably additive, non-trivial measure on the Borel subsets of \mathcal{U} (i.e. on the σ -Algebra generated by the open subsets of \mathcal{U}).
- $\mu(\mathcal{U}) = 1$
- $\mu(gEh) = \mu(E)$ for all Borel sets $E \subseteq \mathcal{U}$, and for all $g, h \in \mathcal{U}$.

This last property is usually referred to as (left and right) translation invariance. A proof can be found in [?].

Remark 1.9. The Haar measure makes the group of unitary $n \times n$ matrices $\mathcal{U} := \mathcal{U}_n(\mathbb{C})$ over \mathbb{C} with its Borel sets into a measurable space, and through it, all familiar concepts of stochastics carry over. Since the Haar measure is translation invariant it is in particular uniform on \mathcal{U} , and thus we can interpret the identity mapping on \mathcal{U} as a random, uniformly distributed unitary matrix. The expected value of a measurable function f on \mathcal{U} can thus be explicitly be calculated through a Lebesgue integral:

$$E(f) = \int_{\mathcal{U}} f(x) d\mu(x) \tag{4}$$

Lemma 1.10. Let V be a finite dimensional \mathbb{C} vector space and V^* its dual space. Further let G be the unitary group acting on V . Then

$$\text{Hom}_G(V \otimes V^*, \mathbb{C}) \cong \text{Hom}_G(V, V) \quad (5)$$

as vector spaces, where G acts diagonally on $V \otimes V^*$.

Proof. We fix a \mathbb{C} basis e_1, \dots, e_n of V . With respect to this basis we see that both are basically just $n \times n$ matrices, and we get an isomorphism easily:

$$\text{End}_G(V, V) \ni (e_i \mapsto \sum_{j=0}^n a_{i,j} e_j) \mapsto \sum_{i,j=0}^n a_{i,j} (e_i \otimes e_j^*)^* \in \text{Hom}_G((V \otimes V^*), \mathbb{C}) \quad (6)$$

This linear mapping is well defined, since the condition of being G -equivariant translates to the set of linear equations on the coefficients $a_{i,j}$ given by the fact that the mapping described by $a_{i,j}$ has to commute with all representation matrices given by the action of G on V with respect to the basis e_1, \dots, e_n . From equation 6 we can immediately see the inverse mapping, which implies bijectivity. \square

Theorem 1.11. Let U be a random unitary $n \times n$ matrix over \mathbb{C} and $A, B \in \mathbb{C}^{n \times n}$, and let $\mathcal{U} := \mathcal{U}_n(\mathbb{C})$ be the set of unitary $n \times n$ matrices. Then:

$$E(UAU^\dagger) = \frac{1}{n} \text{Tr}(A) I_n \quad (7)$$

$$E(\text{Tr}(AU) \text{Tr}(BU^\dagger)) = \frac{1}{n} \text{Tr}(AB) \quad (8)$$

Where E is the expected value and I_n denotes the unit matrix $\in \mathbb{C}^{n \times n}$.

Proof. For the first one (eqn. 7) we note that for $W \in \mathcal{U}_n(\mathbb{C})$:

$$WE(UAU^\dagger)W^\dagger = W \int_{\mathcal{U}} (UAU^\dagger) d\mu(U) W^\dagger = \int_{\mathcal{U}} (WU) A (WU)^\dagger d\mu(WU) = E(UAU^\dagger) \quad (9)$$

, where in the second equality we used the fact that μ is translation invariant. In particular, $[E(UAU^\dagger), \mathcal{U}_n] = \{I_n\}$, and since \mathcal{U}_n is an irreducible representation of itself, it contains a \mathbb{C} basis for $\mathbb{C}^{n \times n}$ (for example, the set \mathcal{E} of pauli operators), which means that $E(UAU^\dagger) \in \{\lambda I_n \mid \lambda \in \mathbb{C}\}$, since it has to commute with every matrix in $\mathbb{C}^{n \times n}$.

Further we note that the mapping

$$\varphi : \mathbb{C}^{n \times n} \rightarrow \mathbb{C} I_n, A \mapsto E(UAU^\dagger)$$

is a \mathbb{C} -linear, \mathcal{U}_n -equivariant mapping, i.e. $\varphi \in \text{Hom}_{\mathcal{U}_n}(\mathbb{C}^n \otimes (\mathbb{C}^n)^*, \mathbb{C}) \cong \text{Hom}_{\mathcal{U}_n}(V, V)$ by lemma 1.10. Since \mathcal{U}_n is transitive on $V \setminus \{0\}$, $\dim_{\mathbb{C}}(\text{Hom}_{\mathcal{U}_n}(V, V)) = 1$, and thus $\text{Hom}_{\mathcal{U}_n}(\mathbb{C}^n \otimes (\mathbb{C}^n)^*, \mathbb{C})$ is also 1-dimensional. Since Tr is a \mathcal{U}_n -equivariant, non-trivial homomorphism, it follows that $\text{Hom}_{\mathcal{U}_n}(\mathbb{C}^n \otimes (\mathbb{C}^n)^*, \mathbb{C}) = \langle \text{Tr} \rangle_{\mathbb{C}}$, and in particular, $\varphi = \lambda \text{Tr}$, for some $\lambda \in \mathbb{C}$. That $\lambda = \frac{1}{n}$ follows immediately from $E(U I_n U^\dagger) = I_n = \lambda \text{Tr}(I_n) I_n = \lambda n I_n$.

For equation 8 we consider the mapping

$$\psi : \mathbb{C}^{n \times n} \times \mathbb{C}^{n \times n} \rightarrow \mathbb{C}, (A, B) \mapsto E(\text{Tr}(AU) \text{Tr}(BU^\dagger))$$

Since both the trace and the Lebesgue integral are linear mappings, it follows that ψ is a bilinear mapping. There are two crucial observations to complete the proof. First, $\psi(\cdot, I_n) \in \text{Hom}_{\mathcal{U}_n}(\mathbb{C}^n \otimes (\mathbb{C}^n)^*, \mathbb{C})$ and, by the same argument as above, there exists a $\lambda \in \mathbb{C}$ such, that $\psi(A, I_n) = \lambda \text{Tr}(A)$ for all $A \in \mathbb{C}^{n \times n}$. On the other hand, for $A \in \mathbb{C}^{n \times n}, B \in \mathcal{E}$

$$\psi(A, B) = \int_{\mathcal{U}_n} \text{Tr}(AU) \text{Tr}(BU^\dagger) d\mu(U) = \int_{\mathcal{U}_n} \underbrace{\text{Tr}(AU)}_{=\text{Tr}(ABBU)} \underbrace{\text{Tr}(BU^\dagger)}_{=\text{Tr}((BU)^\dagger)} d\mu(U) = \int_{\mathcal{U}_n} \text{Tr}(ABU) \text{Tr}(I_n U^\dagger) d\mu(U) = \psi(AB, I_n) \quad (10)$$

Since $B = B^\dagger = B^{-1}$ for all $B \in \mathcal{E} \subseteq \mathcal{U}_n$. But $\langle \mathcal{E} \rangle_{\mathbb{C}} = \mathbb{C}^{n \times n}$, so it holds for all $B \in \mathbb{C}^{n \times n}$ from the bilinearity of ψ . That $\lambda = \frac{1}{n}$ follows again from $\psi(I_n, I_n) = \int_{\mathcal{U}_n} \text{Tr}(U) \text{Tr}(U^\dagger) d\mu(U) = 1$, by a generalization of the Schur orthogonality relations for characters, see for example [?] \square

1.4 Physical background: The postulates of quantum mechanics

The motivation for studying quantum codes comes obviously from physics. To better understand quantum codes, it might therefore be a sensible thing to study the basic mathematical foundations of quantum mechanics. Quantum mechanics is a mathematical framework for describing physical phenomena, and can be given an axiomatic basis: this basis is usually referred to as 'the postulates of quantum mechanics'. The following formulation is mainly based on [?], section 2.2.

The postulates of Quantum Mechanics

- a) Every isolated physical system S is described by a tuple (\mathcal{H}, φ) , where \mathcal{H} is a complex Hilbert space, called the state space of S , and $\varphi : \mathbb{R} \rightarrow \mathcal{U}(\mathcal{H}), t \mapsto U_t$ a group homomorphism of $(\mathbb{R}, +)$ in the unitary group of \mathcal{H} is the mapping that gives the time evolution of the system. The concrete state of the space is described by a family of unit vectors $\{|\psi_t\rangle \mid t \in \mathbb{R}\}$, related by $|\psi_t\rangle = U_t U_{t'}^{-1} |\psi_{t'}\rangle$ for all $t, t' \in \mathbb{R}$. Alternatively the system can be seen as described by the projective space of \mathcal{H} . In this thesis \mathcal{H} will always be finite-dimensional, but the formalism does not require it to be.
- b) A measurement made on the system is described by a collection $\{M_m \mid m \in I\} \subseteq \text{End}(\mathcal{H})$ of operators, where I is some countable index set which represents the possible outcomes of the measurement. For a system in a state $|\psi\rangle$, the probability of measuring $m \in I$, $p(m)$ is given by $p(m) = \langle \psi | M_m^\dagger M_m | \psi \rangle$, and after a measurement with outcome $m \in I$ the system is in the state $\frac{1}{\|M_m |\psi\rangle\|} M_m |\psi\rangle$. Further, the set of measurement operators has to satisfy: $\sum_{m \in I} M_m^\dagger M_m = \text{Id}_{\mathcal{H}}$, which basically means that $\sum_{m \in I} p(m) = 1$.
- c) A system that is composed of various quantum systems with state spaces $\mathcal{H}_1, \dots, \mathcal{H}_n$ has the state space $\mathcal{H} = \mathcal{H}_1 \otimes \dots \otimes \mathcal{H}_n$

Remark 1.12. One does not have to think long to perceive an (apparent?) inconsistency in the postulates. The universe itself should be a closed quantum system, which means that the time evolution of the whole universe should be uniquely determined by its starting state by the 1. postulate. But the 2. postulate states that a measurement has an inherent probabilistic behaviour, which is inconsistent with this. This problem is called the measurement problem, and is in the end rather a philosophical problem, which is why we will not further pursue it in this thesis.

2 Qubits and Quantum Codes

Definition 2.1. A qubit is a two dimensional complex Hilbert space, $(\mathbb{C}^2, \langle \cdot, \cdot \rangle)$. It is an abstraction of a two state quantum system, whose physical realization we disregard. For convention we will fix an orthonormal basis of \mathbb{C}^2 and call the two vectors $|0\rangle$ and $|1\rangle$ in analogy to the two states of a bit, 0 and 1.

Remark 2.2. A many qubit system is, as by the postulates of quantum mechanics, described by the tensor product Hilbert space of its single qubits. This is the reason why an n -qubit system is described by $\mathcal{H} = \underbrace{\mathbb{C}^2 \otimes \dots \otimes \mathbb{C}^2}_{n \text{ times}} \cong \mathbb{C}^{2^n}$. The basis of this space given by the $|0\rangle, |1\rangle$ basis for each qubit can be abbreviated by

binary strings: $|b_1 b_2 \dots b_n\rangle := |b_1\rangle \otimes |b_2\rangle \dots \otimes |b_n\rangle$. For instance $|010\rangle = |0\rangle \otimes |1\rangle \otimes |0\rangle$. This notation also makes the analogy to bits further clear, hence the name, qubits.

Definition 2.3. A quantum code \mathcal{C} is a (linear) subspace of an n -qubit space \mathbb{C}^{2^n} . If the code is a 2^k dimensional subspace of \mathbb{C}^{2^n} for some integer $k \leq n$, then we call the code \mathcal{C} an $[[n, k]]$ -Code. In this case there is a bijective linear mapping from a k -qubit system to the code \mathcal{C} . We say the code encodes k qubits in n qubits.

2.1 Error correction

So, how does quantum error correction work? We will describe the basic theory of quantum error correction in analogy to the classical theory of linear codes. In particular, we will develop it in this section based on the very simple example of the $[3,1,3]$ binary linear code $0 \mapsto 000, 1 \mapsto 111$ (repetition code).

By the postulates of quantum mechanics, errors (possible time-evolutions of the physical space) are given by elements of the Unitary group on \mathcal{H} , which for this case is the group $\mathcal{U}_{2^n}(\mathbb{C})$. Because of its many useful properties, we will first consider the pauli operators (\mathcal{E}) as errors. These are, in particular, a unitary, hermitian basis of $\text{End}(\mathcal{H})$. There is also a very good physical motivation for choosing the pauli operators, since these are closely related to the spin of the system.

Example 2.4. A first, naive approach for building a quantum code is just the repetition code in qubits:

$$\begin{aligned} |0\rangle &\mapsto |000\rangle \\ |1\rangle &\mapsto |111\rangle \end{aligned}$$

It is not hard to see that this code will protect from one Pauli X error, for example, X_2 :

$$X_2|000\rangle = |010\rangle$$

And just as in the classical analog, we can see this (assuming only one X_i error) can only have been $|000\rangle$. It is not difficult to see that this code will protect from a single X_i error on an arbitrary qubit. Simple, right? The problem is, we have to measure in some way the state of the code in order to find out what the state is and what error happened. Fortunately, there is a way of covering all 'bit flip' (X_i) errors with a measurement: Define the following measurement operators:

$$\begin{aligned} P_0 &:= |000\rangle\langle 000| + |111\rangle\langle 111|, P_1 := |100\rangle\langle 100| + |011\rangle\langle 011| \\ P_2 &:= |010\rangle\langle 010| + |101\rangle\langle 101|, P_3 := |001\rangle\langle 001| + |110\rangle\langle 110| \end{aligned}$$

From the operators themselves it is easy to see that a measurement in this basis will immediately identify if exactly one 'bit flip' error (but nothing else) has happened, and leave everything as it is.

This example also shows another point where the analogies to classical codes have its limits: any unitary operator could be an error on the qubits. Consider the state $|+\rangle := 1/\sqrt{2}(|0\rangle + |1\rangle)$. Coding this would yield $1/\sqrt{2}(|000\rangle + |111\rangle)$. Since all error operators are linear, a bit flip would still be corrected with the measurement above. But how about a phase flip? Consider the error Z_1 , for example. $Z_1(1/\sqrt{2}(|000\rangle + |111\rangle)) = 1/\sqrt{2}(|000\rangle - |111\rangle)$, which would go undetected, but would be decoded to be $|-\rangle := 1/\sqrt{2}(|0\rangle - |1\rangle)$.

We can modify the code to detect phase flip (Z_i) errors: We have seen that the operator Z takes $|+\rangle$ to $|-\rangle$ and vice-versa. That is, it flips both states. If we were to rename our states $|0\rangle$ and $|1\rangle$, Z would act as an X . With this in mind, it is not difficult to see why coding $|0\rangle \mapsto |+++ \rangle$, $|1\rangle \mapsto |-- \rangle$ would be a good idea. If we make the same measurement as above, but in the $|+\rangle, |-\rangle$ basis instead, a phase flip would be the same as a bit flip in that basis, and would be detected. It can be immediately checked that the operator

$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

also called a Haddamard gate, takes $(|0\rangle, |1\rangle) \mapsto (|+\rangle, |-\rangle)$. So the measurement operators would be $HP_iH^\dagger = HP_iH, i = 0, 1, 2, 3$.

But this code, again, would not be able to detect a bit flip. We can try combining them, on a 9 qubit code:

Example 2.5. This 9 qubit code, also known as Shor code [?] can correct an X or Z error on a **single qubit**. This is achieved by combining the two codes previously mentioned. First the phase flip code, $|0\rangle \mapsto |+++\rangle, |1\rangle \mapsto |--\rangle$, then the bit flip code to each of the three qubits of the first encoding: $|+++\rangle = |+\rangle|+\rangle|+\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \mapsto \frac{1}{2\sqrt{2}}(|000\rangle + |111\rangle)(|000\rangle + |111\rangle)(|000\rangle + |111\rangle)$ and in an analogous fashion $|--\rangle \mapsto \frac{1}{2\sqrt{2}}(|000\rangle - |111\rangle)(|000\rangle - |111\rangle)(|000\rangle - |111\rangle)$

This, we have seen, can correct a single one of two different errors on a single qubit. It might seem like a lot of work for very little gain. Fortunately, as we will see, the correcting of these two errors is enough to correct an arbitrary single-qubit error! In fact, the following theorem (2.6) is one of the results that make quantum error correction possible in the first place, since it helps to bypass the problem of a continuous, infinite set of possible errors.

Theorem 2.6. Let $\{E_i \mid i \in I\}$ for some index set I be a set of unitary operators, and C be a quantum code that can correct errors on all these $E_i, i \in I$. Further let $\{F_j \mid j \in J\}$ for another index set J be another set of unitary operators, each of which is a linear combination of the E_i , i.e. $F_j = \sum_{i \in I} c_{i,j} E_i$ for all $j \in J$ where $c_{i,j} \in \mathbb{C}$ for all $i \in I, j \in J$, and almost all $c_{i,j} = 0$. Then, C can also correct all errors $F_j, j \in J$. This theorem actually holds for a wider class of errors, where E_i, F_j are so-called quantum operations; but this goes beyond the scope of this thesis. The more general statement, along with a proof, can be found in [?].

2.2 Stabilizer codes

We will now study a very important class of quantum codes, stabilizer or additive quantum codes.

Definition 2.7. Let V be a finite dimensional \mathbb{F}_4 vector space and let $(\cdot, \cdot) : V \times V \rightarrow \mathbb{F}_4$ be a symplectic, non-degenerate bilinear form on V . Further let $C \subset V$. We define the orthogonal complement of C to be $C^\perp := \{v \in V \mid (v, c) = 0 \text{ for all } c \in C\}$. If $C \subseteq C^\perp$ we say C is self-orthogonal.

Definition 2.8. Let $V, (\cdot, \cdot)$ be as above and let $C \leq_{\mathbb{F}_2} V$ be an \mathbb{F}_2 subspace of V , i.e. C is an additive subset (closed under addition) of V . If $C \subseteq C^\perp$ we call C an additive code.

Proposition 2.9. Let $C \leq_{\mathbb{F}_2} V = \mathbb{F}_4^n$ be an additive code, and the symplectic form $(x, y) = \text{Tr}(xy)$ for all $x, y \in V$ as in rem. 1.7. Then C defines a quantum code on n qubits via a lifting ψ of the inverse mapping to the induced by the symplectic notation defined in rem. 1.7:

$$\psi : V \rightarrow G, (\omega a_1 + \bar{\omega} b_1, \dots, \omega a_n + \bar{\omega} b_n) \mapsto X_1^{a_1} Y_1^{b_1} \dots X_n^{a_n} Y_n^{b_n}$$

$S := \psi(C)$ is an abelian subgroup of \mathcal{P}_n , called the stabilizer group of C , which defines a quantum code \tilde{C} on \mathbb{C}^{2^n} through $\tilde{C} := \{v \in \mathbb{C}^{2^n} \mid sv = v \text{ for all } s \in S\}$. The dimension of this code is $2^{n - \dim_{\mathbb{F}_2}(C)}$.

Proof. From rem. 1.7 we know that symplectic notation is an isometry to $\mathcal{P}_n/\mathcal{P}'_n$. Further since C is self-orthogonal it follows that S is abelian (since two elements of \mathcal{P}_n commute iff their images are orthogonal over \mathbb{F}_4) and since $C \leq_{\mathbb{F}_2} V$ it follows that S is an \mathbb{F}_2 vector space. Define $R_S := \frac{1}{|S|} \sum_{g \in S} g \in \text{End}(\mathbb{C}^{2^n})$. $R_S = R_S^2$ is a projection onto \tilde{C} , then for $s \in S, v \in V, c \in \tilde{C}$ it holds that:

$$sR_S(v) = \frac{1}{|S|} \sum_{g \in S} (sg)v = R_S(v), R_S(c) = \frac{1}{|S|} \sum_{g \in S} gc = \frac{1}{|S|} \sum_{g \in S} c = c$$

and in particular $c \in R_S^{-1}(c)$. It follows that

$$\dim_{\mathbb{C}}(\tilde{C}) = \text{Tr}(R_S) = \frac{1}{|S|} \sum_{g \in S} \text{Tr}(g) = \frac{1}{|S|} \text{Tr}(\text{Id}) = 2^n / |S| = 2^{n - \dim_{\mathbb{F}_2}(S)}$$

since for $\text{Id} \neq g \in G : \text{Tr}(g) = 0$ which follows directly from $\text{Tr}(X) = \text{Tr}(Y) = \text{Tr}(Z) = 0$ and $\text{Tr}(A \otimes B) = \text{Tr}(A)\text{Tr}(B)$ for all $A \in \mathbb{C}^{k_1 \times k_1}, B \in \mathbb{C}^{k_2 \times k_2}, k_1, k_2 \in \mathbb{N}$. \square

Definition 2.10. A subgroup $S \leq \mathcal{P}_n$ is called a stabilizer group iff $-I \notin S$. Note that this implies that S is abelian, since $\mathcal{P}'_n = \{\pm \text{Id}\}$. This definition is also consistent with the notation introduced in prop.[?]: Through symplectic notation, S defines an \mathbb{F}_2 vectorspace (equivalent to S being closed under multiplication) that is self orthogonal with respect to the commutator, since S is abelian.

Example 2.11. With stabilizer codes now properly defined, we can return to the previous example (2.5), the 9-qubit Shor code. The Shor code is generated by the 8 \mathbb{F}_2 -linearly independent generators in table 1 (see [?]). This is, in general, a much more compact way to write down the code. We see that, of course, the code is stabilized by all operators and there is 8 of them (what is in agreement with theorem ??). From their symplectic notation it can be immediately checked they are indeed \mathbb{F}_2 -linearly independent generators.

Name	Operator
g_1	$ZZIIIIII$
g_2	$I ZZIIIIII$
g_3	$III ZZIIII$
g_4	$IIII ZZII$
g_5	$IIII I ZZI$
g_6	$IIII I I ZZ$
g_7	$XXXXXX III$
g_8	$III XXXXX$

Table 1: Stabilizer Generators for the 9-Qubit Shor Code

2.3 Logical Operators

There is another class of operators which is very important for stabilizer codes: the so-called logical operators.

Definition 2.12. Let $S \leq \mathcal{P}_n$ be a stabilizer group. Let $C_{\mathcal{P}_n}(S)$ be the centralizer of S in \mathcal{P}_n , i.e.:

$$C_{\mathcal{P}_n}(S) = \{A \in \mathcal{P}_n \mid [A, B] = 0 \text{ for all } B \in S\}$$

Note that since S is abelian, $S \leq C_{\mathcal{P}_n}(S)$. An element $L \in C_{\mathcal{P}_n}(S) \setminus S$ is called a logical operator.

Logical operators have the following interpretation: An $[[n, k]]$ qubit code encodes k qubits in $n > k$. Once encoded, we are confronted with n physical qubits, but know that logically, only k are represented. If we were to do an operator A on the unencoded k qubits, we would obtain a different state for these k qubits. But, how would this new state look like when encoded in the n qubits? There has to be an operator $A^{(n)}$ on n qubits that applies A on the encoded qubits: $A^{(n)}$ is a logical operator, as it acts as A on the logical qubits.

But why are these exactly the elements in $C_{\mathcal{P}_n}(S) \setminus S$? Well, for a logical operator to act on the logical, encoded qubits, it has to map codewords to codewords; after all, the new encoded state is part of the code as well. It is precisely the elements of $C_{\mathcal{P}_n}(S)$ that do that. On the other hand, the elements of S , per definition, act as the identity on the code, leaving it unchanged. We want to define only non-trivial logical operators as such, and leave therefore S out.

Remark 2.13. Let $f : \mathbb{C}^{2^k} \rightarrow \mathbb{C}^{2^n}$ be a \mathbb{C} -linear, injective mapping that encodes k qubits in n qubits, i.e. $f(\mathbb{C}^{2^k}) =: C$ is an additive code. Then, $\text{Aut}_{\mathbb{C}}(C)$ acts on \mathbb{C}^{2^k} via $\sigma v := f^{-1}(\sigma f(v))$, and the stabilizer group S is the stabilizer of this action. Hence, for $A^{(n)}, B^{(n)}$ which both act as A on k qubits (as described above), there exists a $T \in S$ such, that $A^{(n)} = TB^{(n)}$.

Example 2.14. Again, the 9-qubit Shor code (ex. 2.5): It encodes only one qubit, so we will only search for the two logical operators \bar{X} and \bar{Z} . These are given by $\bar{X} = ZZZZZZZZZ$ and $\bar{Z} = XXXXXXXXX$, as can be checked with a simple calculation. Note that it is not as one might naively expect at first, but exactly the opposite instead.

2.4 Code distance

There is an easy way of quantifying how many errors a code can correct: the code distance. It is analogous to what is used in classical coding theory. It is also a good reason for studying logical operators. Before introducing it however, we need a further definition:

Definition 2.15. Let $V = (\mathbb{C}^2)^{\otimes n}$, and $\sigma = i^k \sigma_1 \otimes \dots \otimes \sigma_n \in \mathcal{P}_n$. We define the support of σ to be $\text{Supp}(\sigma) = \{i \in \underline{n} \mid \sigma_i \neq \text{Id}\}$ and its weight as $\text{wt}(\sigma) = |\text{Supp}(\sigma)|$.

Similarly, we define a weight on $\mathbb{F}_2^{2n} = (e_1, \dots, e_n, f_1, \dots, f_n)$ through

$$\left| \sum_{i=1}^n a_i e_i + b_i f_i \right| := |\{i \in \underline{n} \mid a_i \neq 0 \text{ or } b_i \neq 0\}|, a_i, b_i \in \mathbb{F}_2 \text{ for all } i \in \underline{n} \quad (11)$$

This coincides with the Hamming-weight on \mathbb{F}_4 with the identification from remark 1.7.

Definition 2.16. Let $S \leq \mathcal{P}_n$ be a stabilizer group with corresponding code C . Then, the minimal distance (or code distance) ρ of C is defined as:

$$\rho := \min\{|A| \in \mathcal{P}_n \mid A \text{ is a logical operator}\} = \min\{|A| \mid A \in C_{\mathcal{P}_n}(S) \setminus S\} \quad (12)$$

An $[[n, k]]$ code with distance ρ is also called an $[[n, k, \rho]]$ code.

The basic idea for the definition of ρ is that it marks the number of qubits an operator has to minimally affect for it to change one codeword to another. An argument similar to that of the classical coding theory (see [?]) can be made to prove that a code with distance ρ can correct all errors E which satisfy $|E| < \frac{\rho}{2}$. A direct proof for quantum coding theory can be found in [?].

Example 2.17. We can calculate the code distance of ex. 2.5. We know two logical operators, $\bar{X} = ZZZZZZZZ$ and $\bar{Z} = XXXXXXXX$, and know that they represent the only two equivalency classes of logical operators, since the Shor code encodes a single qubit. This means we can get all logical operators this way, by going through the equivalency classes:

$$C_{\mathcal{P}_n}(S) = \{A\bar{X} \mid A \in S\} \cup \{A\bar{Z} \mid A \in S\} \quad (13)$$

From table 1 we get the generators of S , so we know in principle all logical operators. We can look at, for example $\bar{Z}g_8 = ZZZIIIII$, or $g_2g_3g_6\bar{X} = ZIIIZZII$. It is not hard to see that these operators are of minimal weight in the two classes from eqn. 13. This means that $\rho = 3$, and hence, the Shor code corrects $1 < \frac{3}{2}$ errors.

Definition 2.16 can be generalized for any quantum code, and it will have the same properties as it has for additive codes, as well as in classical coding theory. See [?]

Definition 2.18. For a quantum code C on n qubits, define C to have minimal distance of at least d iff $\langle v|U_d|v\rangle = \langle v'|U_d|v'\rangle$ for all unit vectors $|v\rangle, |v'\rangle \in C$, where U_d ranges over all unitary operators with support on exactly d qubits.

2.5 CSS Codes

A class of codes worth mentioning, a subclass of stabilizer codes, are CSS codes. CSS codes are named after their inventors: A. R. Calderbank, Peter Shor and Andrew Steane. The original construction actually uses classical codes, and is one of the advantages of CSS codes, since you can derive much information about the code from the classical codes used on its construction.

Definition 2.19. Let $D_1, D_2 \leq \mathbb{F}_2^n$ be two (classical) linear codes, which encode k_1 and k_2 bits respectively, i.e. one is an $[n, k_1]$ code and the other one $[n, k_2]$, such that $D_2 \subseteq D_1^\perp$. Further let D_1 and D_2^\perp have both minimal distance $d \geq 2t + 1$ for some positive integer t . For $x \in \mathbb{F}_2^n$ let $|x\rangle \in \mathbb{C}^{2^n}$ be the n -qubit state indexed by the bit string x , and define $|x + D_2\rangle := \frac{1}{\sqrt{|D_2|}} \sum_{y \in D_2} |x + y\rangle$ for all $x \in D_1$. Now define $C := C(D_1, D_2) := \{|x + D_2\rangle \mid x \in D_1\}$ as the CSS code from D_1 over D_2 .

Remark 2.20. The above defined $C = C(D_1, D_2)$ is an $[[n, k_1 - k_2, d]]$ code, with $d \geq 2t + 1$. A proof can be found in [?].

3 Quantum Weight Enumerators

3.1 The Shor-Laflamme and Rains enumerators

Definition 3.1. Let $V = (\mathbb{C}^2)^{\otimes n}$, and $\mathcal{E} = \{\sigma_1 \otimes \dots \otimes \sigma_n \mid \sigma_i \in \{X, Y, Z, \text{Id}\} \text{ for all } i \in \underline{n}\}$ (generating set for the pauli group on V). Further let $S \subseteq \underline{n}$ and $M_1, M_2 \in \mathbb{C}^{2^n \times 2^n} (\cong \text{End}(V))$, with $\text{Tr}(M_1) \neq 0 \neq \text{Tr}(M_2)$. Then we define the coefficients of the Shor-Laflamme enumerators [?] of M_1, M_2 as follows:

$$A_S(M_1, M_2) := \frac{1}{\text{Tr}(M_1)\text{Tr}(M_2)} \sum_{\substack{E \in \mathcal{E} \\ \text{Supp}(E)=S}} (\text{Tr}(EM_1)\text{Tr}(EM_2)) \quad (14)$$

$$B_S(M_1, M_2) := \frac{1}{\text{Tr}(M_1 M_2)} \sum_{\substack{E \in \mathcal{E} \\ \text{Supp}(E)=S}} (\text{Tr}(EM_1 EM_2)) \quad (15)$$

And for $d \in \underline{n}$ we define

$$A_d(M_1, M_2) := \sum_{\substack{S \subseteq \underline{n} \\ |S|=d}} A_S(M_1, M_2); B_d(M_1, M_2) := \sum_{\substack{S \subseteq \underline{n} \\ |S|=d}} B_S(M_1, M_2) \quad (16)$$

Definition 3.2. Again let $V = (\mathbb{C}^2)^{\otimes n}$, and $\mathcal{E} = \{\sigma_1 \otimes \dots \otimes \sigma_n \mid \sigma_i \in \{X, Y, Z, \text{Id}\} \text{ for all } i \in \underline{n}\}$. Further let $S \subseteq \underline{n}, S^c = \underline{n} \setminus S$ and $M_1, M_2 \in \mathbb{C}^{2^n \times 2^n}$ with $\text{Tr}(M_1) \neq 0 \neq \text{Tr}(M_2)$. We define the coefficient of the Rains enumerators [?] of M_1, M_2 as follows:

$$A'_S(M_1, M_2) := \frac{1}{\text{Tr}(M_1)\text{Tr}(M_2)} \text{Tr}(\text{Tr}_{S^c}(M_1)\text{Tr}_{S^c}(M_2)) \quad (17)$$

$$B'_S(M_1, M_2) := \frac{1}{\text{Tr}(M_1 M_2)} \text{Tr}(\text{Tr}_S(M_1)\text{Tr}_S(M_2)) \quad (18)$$

As well as for a $d \in \underline{n}$:

$$A'_d(M_1, M_2) := \sum_{\substack{S \subseteq \underline{n} \\ |S|=d}} A'_S(M_1, M_2); B'_d(M_1, M_2) := \sum_{\substack{S \subseteq \underline{n} \\ |S|=d}} B'_S(M_1, M_2) \quad (19)$$

Note that these definitions differ by a normalisation factor from those in the original paper [?].

Definition 3.3. Let $C \leq (\mathbb{C}^{2^n})$ be a quantum code, and $P : \mathbb{C}^{2^n} \rightarrow C$ be the projection onto C . The Shor-Laflamme enumerator polynomial of C is defined as $\sum_{i=0}^n A_i(P, P)x^{n-i}y^i$, and the Rains enumerator polynomial similarly as $\sum_{i=0}^n A'_i(P, P)x^{n-i}y^i$.

Lemma 3.4. Let $C \leq (\mathbb{C}^2)^{\otimes n} =: V$ be an additive quantum code and $G \subseteq \mathcal{E}$ be its stabilizer group (i.e. $C = V^G$). and $E, E' \in \mathcal{E} = \{\sigma_1 \otimes \dots \otimes \sigma_n \mid \sigma_i \in \{X, Y, Z, \text{Id}\} \text{ for all } i \in \underline{n}\}$. Then

- a) The projection operator on C is given by $P_C = \frac{1}{|G|} \sum_{\sigma \in G} \sigma$
- b) $\text{Tr}(EP_C) = \frac{1}{|G|} \cdot \begin{cases} 2^n = \dim(V) & \text{if } E \in G \\ 0 & \text{if } E \notin G \end{cases}$
- c) $\text{Tr}(\text{Tr}_S(E)\text{Tr}_S(E')) = 0$ if $E \neq E'$ or if there exists an $i \in S$ which satisfies that for $E = \sigma_1 \otimes \dots \otimes \sigma_n$, $\sigma_i \neq \text{Id}$, otherwise $\text{Tr}(\text{Tr}_S(E)\text{Tr}_S(E)) = 2^{n+|S|}$.
- d) $\text{Tr}(EP_C EP_C) = \begin{cases} \frac{2^n}{|G|} & \text{for } E \in G^\perp \\ 0 & \text{otherwise} \end{cases}$

Proof. a) From a simple calculation it follows that $\sigma P_C(v) = P_C(v)$ for all $v \in V, \sigma \in G$ and $P_C(v) = v$ for all $v \in V^G$. From this it follows immediately that $\text{Im}(P_C) = V^G$ and $P_C^2 = P_C$, and hence that it is the projection onto V^G .

- b) Since $\text{Tr}(X) = \text{Tr}(Z) = \text{Tr}(Y) = 0$ and $\text{Tr}(\sigma_1 \otimes \sigma_2) = \text{Tr}(\sigma_1)\text{Tr}(\sigma_2)$ it follows immediately that \mathcal{E} is an orthogonal basis with respect to the trace bilinear form (since $\langle \mathcal{E} \rangle_{\mathbb{C}} = \mathbb{C}^{2^n \times 2^n}$), and for all $E \in \mathcal{E}$ we have $E^2 = \text{Id}_V$, hence $\text{Tr}(E^2) = \dim(V) = 2^n$. From a) we get:

$$\text{Tr}(EP_C) = \frac{1}{|G|} \sum_{\sigma \in G} \text{Tr}(\sigma E) = \frac{1}{|G|} \cdot \begin{cases} 2^n = \dim(V) & \text{if } E \in G \\ 0 & \text{if } E \notin G \end{cases} \quad (20)$$

- c) Since $\text{Tr}(X) = \text{Tr}(Y) = \text{Tr}(Z) = 0$, as soon as $S \cap \text{Supp}(E) \neq \emptyset$, it follows that $\text{Tr}_S(E) = 0$, and hence $\text{Tr}_S(E)\text{Tr}_S(E') = 0$. The same applies obviously for E' . Now if $E \neq E'$ and $S \cap \text{Supp}(E) = \emptyset$, then $\text{Tr}_S(E) \neq \text{Tr}_S(E')$ and thus $\text{Tr}_S(E)\text{Tr}_S(E') \neq 2^{|S|}\text{Id}$, which in turn implies that $\text{Supp}(\text{Tr}_S(E)\text{Tr}_S(E')) \neq \emptyset$ and so its trace vanishes. On the other hand, if $E = E'$ and $\text{Supp}(E) \cap S = \emptyset$, then $\text{Tr}_S(E)\text{Tr}_S(E) = (2^{|S|})^2 \text{Id}_{S^c}$ and $2^{2|S|}\text{Tr}(\text{Id}_{S^c}) = 2^{2|S|} \dim(V_{S^c}) = 2^{n+|S|}$.
- d) First note that $E = E^{-1}$ for all $E \in \mathcal{E}$ and that $\mathcal{P}'_n = \{\pm \text{Id}\}$. This means that for $D \in \mathcal{E} \subseteq \mathcal{P}_n$: $EDE = EDE^{-1} = [E, D]D = \pm \text{Id}D$. It follows that

$$\text{Tr}(EP_C EP_C) = \frac{1}{|G|^2} \sum_{\sigma, \sigma' \in G} \text{Tr}(\underbrace{E\sigma E}_{=[E, \sigma]_\sigma} \sigma') \quad (21)$$

$$= \frac{1}{|G|^2} \sum_{\sigma, \sigma' \in G} \langle E, \sigma \rangle \text{Tr}(\sigma \sigma') = \frac{2^n}{|G|^2} \sum_{\sigma \in G} \langle E, \sigma \rangle \quad (22)$$

where $\langle \cdot, \cdot \rangle$ denotes the symplectic form onto $\{\pm 1\}$ induced by the commutator. For $E \in G^\perp$, per definition, $\langle E, \sigma \rangle = 1$ for all $\sigma \in G$, hence $\sum_{\sigma \in G} \langle E, \sigma \rangle = |G|$. On the other hand, if $E \notin G^\perp$, there exists a $\sigma_0 \in G : \langle E, \sigma_0 \rangle = -1$. The mapping $\epsilon : G \rightarrow \{\pm 1\}, \sigma \mapsto \langle E, \sigma \rangle$ is an epimorphism of groups and in particular yields that the cosets $\text{Ker}(\epsilon), \sigma_0 \text{Ker}(\epsilon)$ have the same size and give a partition of G in classes of elements with $\langle \sigma, E \rangle = 1$ and $\langle \sigma, E \rangle = -1$, which implies that $\sum_{\sigma \in G} \langle E, \sigma \rangle = 0$. \square

Theorem 3.5. Let $C \leq (C^2)^{\otimes n} =: V$ be an additive quantum code and $G \subseteq \mathcal{E}$ be its stabilizer group (i.e. $C = V^G$); let P_C be the projection onto C , and let $S \subseteq \underline{n}$. Then the coefficients are given as follows:

$$A_S(P_C, P_C) = \#\{E \in G \mid \text{Supp}(E) = S\} \quad (23)$$

$$B_S(P_C, P_C) = \#\{E \in G^\perp \mid \text{Supp}(E) = S\} \quad (24)$$

$$A'_S(P_C, P_C) = 2^{-|S|} \#\{\sigma \in G \mid \sigma|_{S^c} = \text{Id}\} \quad (25)$$

$$B'_S(P_C, P_C) = \frac{2^{|S|}}{|G|} \#\{\sigma \in G \mid \sigma|_S = \text{Id}\} \quad (26)$$

Where for $\sigma = \sigma_1 \otimes \dots \otimes \sigma_n$ and $S = \{s_1 < \dots < s_k\}$, $\sigma|_S$ denotes $\sigma_{s_1} \otimes \dots \otimes \sigma_{s_k}$.

Proof.

$$\text{First } A_S(P_C, P_C) = \frac{1}{\dim(C)^2} \sum_{\substack{E \in \mathcal{E} \\ \text{Supp}(E) = S}} \text{Tr}(P_C E)^2 = \frac{1}{(\dim(C)|G|)^2} \sum_{\substack{E \in \mathcal{E} \\ \text{Supp}(E) = S}} \sum_{\substack{\sigma, \sigma' \in G \\ =0, \sigma \neq E \text{ or } \sigma' \neq E \\ =2^{2n}, \sigma = \sigma' = E}} \underbrace{\text{Tr}(\sigma E) \text{Tr}(\sigma' E)}_{=2^{2n}, \sigma = \sigma' = E} \quad (27)$$

$$= \underbrace{\left(\frac{2^n}{|G| \dim(C)}\right)^2}_{=1, \text{see prop. 2.9}} \#\{E \in G \mid \text{Supp}(E) = S\} \quad (28)$$

Where the second and third equalities follow from lemma 3.4. The calculation for $B_S(P_C, P_C)$ is completely analogous.

Now, for the Rains enumerators:

$$A'_S(P_C, P_C) = \frac{1}{\text{Tr}(P_C)} \text{Tr}(\text{Tr}_{S^c}(P_C) \text{Tr}_{S^c}(P_C)) = \frac{1}{\dim(C)^2} \frac{1}{|G|^2} \sum_{\sigma, \sigma' \in G} \text{Tr}(\text{Tr}_{S^c}(\sigma) \text{Tr}_{S^c}(\sigma')) \quad (29)$$

$$= \frac{1}{\dim(C)^2} \frac{1}{|G|^2} \sum_{\sigma \in G} \text{Tr}(\text{Tr}_{S^c}(\sigma) \text{Tr}_{S^c}(\sigma)) = \underbrace{\frac{1}{\dim(C)^2} \frac{2^{n+|S^c|}}{|G|^2}}_{=2^{|S^c|-n}=2^{-|S|}} \#\{\sigma \in G \mid \sigma|_{S^c} = \text{Id}\} \quad (30)$$

Again, the 2. and 3. equalities follow from lemma 3.4. For B'_S it is enough to see that $B'_S = \dim(C) A'_{S^c}$, which follows directly from their definitions. \square

3.2 The relationship between the enumerators

There is a characterization of the Rains enumerators, which also sheds some light to the motivation for defining A' and B' . They are basically the Shor-Laflamme enumerators, but “averaged over” the whole unitary group, instead of over \mathcal{E} .

Theorem 3.6. Let V be a finite-dimensional complex vector space, $n = \dim V$, $S \subseteq \underline{n}$, $M_1, M_2 \in \text{End}(V)$ hermitian, and let U_S be a random, uniformly distributed unitary matrix on V_S , and $U_S \otimes \text{Id}_{S^c}$ its embedding on the full $\text{End}(V)$. Then:

$$A'_S(M_1, M_2) = 2^{|S|} E(\text{Tr}(M_1(U_S \otimes \text{Id}_{S^c})) \text{Tr}(M_2(U_S \otimes \text{Id}_{S^c})^\dagger)) \quad (31)$$

$$B'_S(M_1, M_2) = 2^{|S|} E(\text{Tr}(M_1(U_S \otimes \text{Id}_{S^c}) M_2(U_S \otimes \text{Id}_{S^c})^\dagger)) \quad (32)$$

Proof. With equation 8 we have:

$$\begin{aligned} A'_S(M_1, M_2) &= \text{Tr}(\text{Tr}_{S^c}(M_1) \text{Tr}_{S^c}(M_2)) = 2^{|S|} E(\text{Tr}(\text{Tr}_{S^c}(M_1) U_S) \text{Tr}(\text{Tr}_{S^c}(M_2) U_S^\dagger)) \\ &= 2^{|S|} E(\text{Tr}(M_1(U_S \otimes \text{Id}_{S^c})) \text{Tr}(M_2(U_S \otimes \text{Id}_{S^c})^\dagger)) \end{aligned}$$

Where the last equality follows from remark ???. Now for the B' , define $\tilde{B}_S(M_1, M_2) := 2^{|S|} E(\text{Tr}(M_1(U_S \otimes \text{Id}_{S^c}) M_2(U_S \otimes \text{Id}_{S^c})^\dagger))$. Then it follows with equation 7:

$$\begin{aligned} \tilde{B}_S(M_1, M_2) &= \tilde{B}_S(E(U_S M_1 U_S^\dagger), E(U_S M_2 U_S^\dagger)) = 2^{-2|S|} \tilde{B}_S(\text{Tr}_S(M_1) \otimes \text{Id}_S, \text{Tr}_S(M_2) \otimes \text{Id}_S) \\ &= 2^{-|S|} \text{Tr}((\text{Tr}_S(M_1) \otimes \text{Id}_S)(\text{Tr}_S(M_2) \otimes \text{Id}_S)) = B'_S(M_1, M_2) \end{aligned}$$

□

From here on, we will focus exclusively on the enumerators in the case of additive/stabilizer codes. For the remainder of this section let $C \leq \mathbb{C}^{2^n}$ be an additive quantum code, G its stabilizer group, and for $S \subseteq \underline{n}$ let $A_S := A_S(P_C, P_C)$, and similarly for A'_S, B_S, B'_S . The results of this section have been proved by Eric Rains [?] for non-additive codes as well.

Proposition 3.7. Let $S \subseteq \underline{n}$. Then the Rains and Shor-Laflamme enumerators are related by the following:

$$A'_S = 2^{-|S|} \sum_{T \subseteq S} A_T \text{ and } B'_S = 2^{-|S|} \sum_{T \subseteq S} B_T \quad (33)$$

Proof. For A'_S The crucial observation here is that for $\sigma \in G$, $\sigma|_S = \text{Id} \Leftrightarrow \text{Supp}(\sigma) \subseteq S^c$. It follows that

$$\{\sigma \in G \mid \sigma|_S = \text{Id}\} = \bigcup_{T \subseteq S^c} \{\sigma \in G \mid \text{Supp}(\sigma) = T\}$$

And this in turn implies:

$$A'_S = 2^{-|S|} \#\{\sigma \in G \mid \sigma|_{S^c} = \text{Id}\} = 2^{-|S|} \sum_{T \subseteq S} \underbrace{\#\{\sigma \in G \mid \text{Supp}(\sigma) = T\}}_{=A_T} = 2^{-|S|} \sum_{T \subseteq S} A_T$$

Now for B'_S we use theorem 3.6:

$$B'_S = \frac{2^{-|S|}}{\dim(C)} E(\text{Tr}(P_C(U_S \otimes \text{Id}_{S^c}) P_C(U_S \otimes \text{Id}_{S^c}))) = \frac{2^{-|S|}}{\dim(C) |G|^2} \sum_{g, g' \in G} E(\text{Tr}[g(U_S \otimes \text{Id}_{S^c}) g'(U_S^\dagger \otimes \text{Id}_{S^c})])$$

Using the expansion in the orthogonal basis $\mathcal{E} : U_S \otimes \text{Id}_{S^c} = \sum_{\sigma \in \mathcal{E}} 2^{-n} \text{Tr}((U_S \otimes \text{Id}_{S^c}) \sigma) \sigma$ yields:

$$\begin{aligned} B'_S &= \frac{2^{|S|-2n}}{\dim(C) |G|^2} \sum_{g, g' \in G} \sum_{\sigma, \sigma' \in \mathcal{E}} E(\text{Tr}[g \text{Tr}(\sigma(U_S \otimes \text{Id}_{S^c})) \sigma g' \text{Tr}(\sigma'(U_S^\dagger \otimes \text{Id}_{S^c})) \sigma']) \\ &= \frac{2^{|S|-2n}}{\dim(C) |G|^2} \sum_{g, g' \in G} \sum_{\sigma, \sigma' \in \mathcal{E}} \text{Tr}(g \sigma g' \sigma') E[\text{Tr}(\sigma(U_S \otimes \text{Id}_{S^c})) \text{Tr}(\sigma'(U_S^\dagger \otimes \text{Id}_{S^c}))] \end{aligned}$$

Lemma 3.4 together with equation 8 yield

$$E[\text{Tr}((U_S \otimes \text{Id}_{S^c})\sigma)\text{Tr}((U_S^\dagger \otimes \text{Id}_{S^c}^c)\sigma')] = 2^{-|S|} \text{Tr}(\sigma|_S \sigma'|_S) \text{Tr}(\sigma|_{S^c}) \text{Tr}(\sigma'|_{S^c}) = \begin{cases} 2^{2n-2|S|} \delta_{\sigma, \sigma'} & \text{if } \text{Supp}(\sigma) \subseteq S \\ 0 & \text{otherwise} \end{cases} \quad (34)$$

from which in turn follows that:

$$B'_S = \frac{2^{-|S|}}{\dim(C)} \sum_{\sigma \in \mathcal{E}, \text{Supp}(\sigma) \subseteq S} \frac{1}{|G|^2} \sum_{g, g' \in G} \text{Tr}(g\sigma g'\sigma) = \frac{2^{-|S|}}{\dim(C)} \sum_{\sigma \in \mathcal{E}, \text{Supp}(\sigma) \subseteq S} \underbrace{\text{Tr}(P_C \sigma P_C \sigma)}_{=\text{Tr}(\sigma P_C \sigma P_C)} = 2^{-|S|} \sum_{T \subseteq S} B_T$$

□

Corollary 3.8. For $d \in \underline{n}$ we get thus:

$$A'_d = 2^{-d} \sum_{i=0}^d \binom{n-i}{n-d} A_i \text{ and } B'_d = 2^{-d} \sum_{i=0}^d \binom{n-i}{n-d} B_i \quad (35)$$

Proof.

$$A'_d = \sum_{S \subseteq \underline{n}, |S|=d} A'_S = 2^{-d} \sum_{S \subseteq \underline{n}, |S|=d} \sum_{T \subseteq S} A_T = 2^{-d} \sum_{i=0}^d \sum_{T \subseteq \underline{n}, |T|=i} \sum_{T \subseteq S, |S|=d} A_T \quad (36)$$

$$= 2^{-d} \sum_{i=0}^d \sum_{T \subseteq \underline{n}, |T|=i} \binom{n-i}{n-d} A_T = 2^{-d} \sum_{i=0}^d \binom{n-i}{n-d} A_i \quad (37)$$

The same works for B'_d .

□

Corollary 3.9. The enumerators $A_d(M_1, M_2)$, $A'_d(M_1, M_2)$, $B_d(M_1, M_2)$, $B'_d(M_1, M_2)$ are invariant under equivalence of codes, i.e., under the action of $\mathcal{U}_2(\mathbb{C}) \wr S_n$, where S_n is the symmetric group on the tensor product components.

Proof. For the Rains enumerators it is immediately clear from theorem 3.6. For the Shor-Laflamme enumerators it follows from cor. 3.8 as we can explicitly write $A_d(M_1, M_2)$ in terms of $A'_d(M_1, M_2)$ and the same for the B 's. □

Example 3.10. We can return to the example of the Shor code, the associated additive code has the (\mathbb{F}_2) -Generator matrix:

$$\begin{pmatrix} \omega^2 & \omega^2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & \omega^2 & \omega^2 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & \omega^2 & \omega^2 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & \omega^2 & \omega^2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & \omega^2 & \omega^2 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & \omega^2 & \omega^2 \\ \omega & \omega & \omega & \omega & \omega & \omega & 0 & 0 & 0 \\ 0 & 0 & 0 & \omega & \omega & \omega & \omega & \omega & \omega \end{pmatrix}$$

Using this we can easily get the Shor-Laflamme enumerator polynomials for this code, as it is the Hamming weight enumerator of the additive code

$$x^9 + 9x^7y^2 + 27x^5y^4 + 75x^3y^6 + 144xy^8 \quad (38)$$

The Rains enumerator polynomial is a little bit harder to calculate. It is:

$$x^9 + \frac{9}{2}x^8y + \frac{45}{4}x^7y^2 + \frac{147}{8}x^6y^3 + \frac{171}{8}x^5y^4 + 18x^4y^5 + \frac{93}{8}x^3y^6 + \frac{45}{8}x^2y^7 + \frac{9}{4}xy^8 + \frac{1}{2}y^9 \quad (39)$$

This is obtained using cor. 3.8, or equivalently, a relationship between both enumerator polynomials which follows from it, as will be seen in 3.11. Note that the coefficients of the rains enumerator polynomial need not be integers, the reason for this becomes clear after examining theorem 3.5

3.3 The MacWilliams transform

Lemma 3.11. The enumerator polynomials are related by the following:

$$A'(x, y) = A(x + \frac{y}{2}, \frac{y}{2}) \quad (40)$$

$$B'(x, y) = B(x + \frac{y}{2}, \frac{y}{2}) \quad (41)$$

Proof. Using cor. 3.8 we get:

$$A'(x, y) = \sum_{d=0}^n A'_d x^{n-d} y^d = \sum_{d=0}^n \sum_{i=0}^d 2^{-d} \binom{n-i}{n-d} A_i x^{n-d} y^d \quad (42)$$

$$= \sum_{i=0}^n \sum_{d=i}^n 2^{-d} \binom{n-i}{n-d} A_i x^{n-d} y^d \stackrel{j:=n-d}{=} \sum_{i=0}^n \sum_{j=0}^{n-i} 2^{-(n-j)} \binom{n-i}{j} A_i x^j y^{n-j} \quad (43)$$

$$= \sum_{i=0}^n y^i 2^{-i} \sum_{j=0}^{n-i} 2^{-(n-i-j)} \binom{n-i}{j} A_i x^j y^{n-i-j} = \sum_{i=0}^n A_i \left(\frac{y}{2}\right)^i \left(x + \frac{y}{2}\right)^{n-i} = A(x + \frac{y}{2}, \frac{y}{2}) \quad (44)$$

The calculation for B' is literally the same, replacing A with B . \square

Theorem 3.12. We get analogies of the MacWilliams identity for Quantum codes: Let $C \leq \mathbb{C}^{2^n}$ be an additive quantum code on n qubits, and $A(x, y), B(x, y)$ be the corresponding Shor-Laflamme enumerator polynomials, $A'(x, y), B'(x, y)$ the Rains enumerator polynomials. Then:

$$A'(x, y) = \dim(C) B'(y, x) \quad (45)$$

$$A(x, y) = \dim(C) B\left(\frac{x+3y}{2}, \frac{x-y}{2}\right) \quad (46)$$

Proof. For the first one we again note that immediately from its definition it is clear that $A'_i = \dim(C) B'_{n-i}$. It follows then

$$A'(x, y) = \sum_{d=0}^n A'_d x^{n-d} y^d = \sum_{d=0}^n \dim(C) B'_{n-d} x^{n-d} y^d = \dim(C) B'(y, x) \quad (47)$$

Using lemma 3.11 we can translate this identity to the Shor-Laflamme enumerators:

$$A(x, y) = A'(x-y, 2y) = \dim(C) B'(2y, x-y) = \dim(C) B\left(2y + \frac{x-y}{2}, \frac{x-y}{2}\right) = \dim(C) B\left(\frac{x+3y}{2}, \frac{x-y}{2}\right) \quad (48)$$

\square

4 Applications of the weight enumerators: Bounds

Having introduced weight enumerators for quantum error correcting codes, and having found out a few properties of these, we can go on to see one of the most important applications of the weight enumerators, namely that of finding bounds for the distance or dimension of codes with specific parameters. This is analogous to classical coding theory and is the reason why these enumerators for quantum codes were introduced in the first place by Shor and Laflamme. To do this, we first have to find out what the relationship is between the code distance and the weight enumerators.

4.1 The relationship between the code distance and the enumerators

Lemma 4.1. Let C be a quantum code on n qubits, and let v be a random unit vector of C , uniformly distributed. Further let P be the orthogonal projection onto C , and K be the dimension of C , $K = \dim_{\mathbb{C}} C$. Then

$$E_{v \in C}(|v\rangle\langle v|) = \int_C |v\rangle\langle v| dv = \frac{P}{K} \quad (49)$$

Proof. We start by identifying that the left hand side is well defined. This is the case since the unit sphere on \mathbb{C}^K is compact (with respect to the natural topology), and is a submanifold of \mathbb{R}^{2K} , which allows us to define an finite integral on C , which can be chosen to be normed, see [?]. Now, an integral of an operator can be defined using the scalar product - see [?], and thus the integral is well defined. Now, since C is isometric to \mathbb{C}^K , we will assume without loss of generality that $C = \mathbb{C}^K$ and thus it suffices to show that for $\tilde{P} := \int_C |v\rangle\langle v| dv$ it holds that $\tilde{P} = \frac{1}{K} \text{Id}_{\mathbb{C}^K}$. To show this we note that the unitary group $\mathcal{U}_K(\mathbb{C})$ acts on $\text{End}(\mathbb{C}^K)$ via conjugation, and this action defines a representation of $U_K(\mathbb{C})$, which we will denote by π . Now, let μ denote the measure induced on \mathbb{C}^K by the probability measure of v and the isometry between C and \mathbb{C}^K , and let dU denote the Haar-measure on the unitary group $\mathcal{U}_K(\mathbb{C})$. It follows that

$$\tilde{P} = \int_{\mathbb{C}^K} |v\rangle\langle v| d\mu(v) = \int_{\mathbb{C}^K} U^\dagger |v\rangle\langle v| U d\mu(v) = \int_{\mathcal{U}_K(\mathbb{C})} \int_{\mathbb{C}^K} U^\dagger |v\rangle\langle v| U d\mu(v) dU \quad (50)$$

Where in the first equality we used the uniformity of the probability measure on v and on the second one we used the normalization of the Haar measure. Let $|v_0\rangle \in \mathbb{C}^K$, $\langle v_0|v_0\rangle = 1$. Then because of the normalization of the μ measure it holds that

$$\tilde{P} = \int_{\mathcal{U}_K(\mathbb{C})} \int_{\mathbb{C}^K} U^\dagger |v\rangle\langle v| U d\mu(v) dU = \int_{\mathcal{U}_K(\mathbb{C})} U^\dagger |v_0\rangle\langle v_0| U dU = \int_{\mathcal{U}_K(\mathbb{C})} (\pi(U))(|v_0\rangle\langle v_0|) dU \quad (51)$$

Since the only fix point of the representation π is I_K , and the mapping $A \mapsto \int_{\mathcal{U}_K(\mathbb{C})} (\pi(U))(A) dU$ is an orthogonal projection to an invariant subspace of the representation π (Reynolds operator), it follows that $\tilde{P} \in \mathbb{C} I_K$. That $\tilde{P} = \frac{1}{K} I_K$ can be seen by explicitly calculating the orthogonal projection: $\tilde{P} = \frac{\langle |v_0\rangle\langle v_0|, I_K \rangle}{\langle I_K, I_K \rangle} I_K = \frac{1}{K} I_K$. \square

Lemma 4.2. Let C be a Quantum Code on n qubits, $M \in \text{End}(\mathbb{C}^{2^n})$ be any operator on those n qubits, and let v be a random unit vector of C , uniformly distributed. Further let $K := \dim_{\mathbb{C}} C$ and P be the orthogonal projection onto C . Then

$$E[|\text{Tr}(M|v\rangle\langle v|) - \frac{1}{K} \text{Tr}(MP)|^2] = \frac{1}{K^2(K+1)} (K \text{Tr}(MPM^\dagger P) - |\text{Tr}(MP)|^2) \quad (52)$$

If we write the term $|a|^2$ under the modulus as aa^* and multiplying out yields:

$$\begin{aligned} & E[|\text{Tr}(M|v\rangle\langle v|) - \frac{1}{K} \text{Tr}(MP)|^2] \\ &= E[|\text{Tr}(\langle v|M|v\rangle)|^2 + \frac{1}{K^2} |\text{Tr}(MP)|^2 - \frac{1}{K} \text{Tr}(M|v\rangle\langle v|)^* \text{Tr}(MP) - \frac{1}{K} \text{Tr}(M|v\rangle\langle v|) \text{Tr}(MP)^*] \\ &\quad \underbrace{=}_{\text{lemma 4.1}} E[|\text{Tr}(\langle v|M|v\rangle)|^2 - \frac{1}{K^2} |\text{Tr}(MP)|^2] \end{aligned}$$

Now consider only the first integral:

$$\begin{aligned} & \int_{\mathbb{C}^K} |\text{Tr}(\langle v|M|v\rangle)|^2 dv = \int_{\mathbb{C}^K} |\langle v|M|v\rangle|^2 dv = \int_{\mathbb{C}^K} \langle v|M|v\rangle \langle v|M^\dagger|v\rangle \\ &= \int_{\mathbb{C}^K} \text{Tr}(\langle v|M|v\rangle \langle v|M^\dagger|v\rangle) = \int_{\mathbb{C}^K} \text{Tr}(|v\rangle\langle v|M|v\rangle \langle v|M^\dagger) =: \int_{\mathbb{C}^K} \varphi_M(|v\rangle\langle v|, |v\rangle\langle v|) \end{aligned}$$

Where the mapping φ_M is the bi-linear functional defined by $\varphi_M(A, B) := \text{Tr}(AMB M^\dagger)$. Using the universal property of the tensor product, there exists exactly one linear mapping $\psi_M : \text{End}(\mathbb{C}^K) \otimes \text{End}(\mathbb{C}^K) \rightarrow \mathbb{C}$, so that $\varphi_M(A, B) = \psi_M(A \otimes B)$ for all $A, B \in \text{End}(\mathbb{C}^K)$.

Theorem 4.3. For both, the Rains and Shor-Laflamme enumerators A'_i, B'_i, A_i, B_i of a quantum Code C of dimension K it holds that $B'_i \geq A'_i \geq 0$, where equality $B'_i = A'_i$ holds for $i \in \underline{d}$ if and only if C has minimal distance d . The same is true for A_i, B_i .

Proof. First for the Shor-Laflamme enumerators A_i, B_i : Consider that by definition, for any hermitian operator M , it holds that $A_i(M, M) = A_i(M, M^\dagger) = \sum_{\substack{E \in \mathcal{E} \\ \text{wt}(E)=i}} |\text{Tr}(EM)|^2$. Since the orthogonal projection P onto C is hermitian, it follows immediately that $A_i \geq 0$. Now for the relationship between A_i, B_i consider the hermitian

operator M_v for any normed $|v\rangle \in C$: $M_v := |v\rangle\langle v| - \frac{1}{K}P$. Now let v be a uniformly distributed, random unit vector from C . From our previous observation, and lemma 4.2, we get that

$$0 \leq E[A_i(M_v, M_v)] \stackrel{\text{def.}}{=} E[\frac{1}{K^2} \sum_{\substack{\sigma \in \mathcal{E} \\ \text{wt}(\sigma)=i}} \text{Tr}(\sigma M_v)^2] = \frac{1}{K^2} \sum_{\substack{\sigma \in \mathcal{E} \\ \text{wt}(\sigma)=i}} E[\text{Tr}(\sigma |v\rangle\langle v| - \frac{1}{K}\sigma P)^2] \quad (53)$$

$$= \frac{1}{K^2(K+1)} (\frac{1}{K} \sum_{\substack{\sigma \in \mathcal{E} \\ \text{wt}(\sigma)=i}} \text{Tr}(\sigma P \sigma^\dagger P) - \frac{1}{K^2} \sum_{\substack{\sigma \in \mathcal{E} \\ \text{wt}(\sigma)=i}} |\text{Tr}(\sigma P)|^2) = \frac{1}{K^2(K+1)} (B_i(P, P) - A_i(P, P)) \quad (54)$$

Where in order to use lemma 4.2 we examine, that the expressions under the modulus are real, since the operators are hermitian. Further we note that $A_i = B_i$ iff $E[\text{Tr}(\langle v|\sigma|v\rangle) - \frac{1}{K}\text{Tr}(\sigma P)] = 0$ for all $\sigma \in \mathcal{E}$ with $\text{wt}(\sigma) = i$. Since this condition is linear, it holds for the span of $\{\sigma \in \mathcal{E} \mid \text{wt}(\sigma) = i\}$ which includes all i -qubit errors, i.e. all unitary operators whose support is exactly i qubits. The condition $E(\langle v|U_i|v\rangle - \frac{1}{K}\text{Tr}(U_i P)) = 0$ for all i -qubit errors U_i can be seen, using lemma 4.1, as the variance of the random variable $\langle v|U_i|v\rangle = \text{Tr}(\langle v|U_i|v\rangle) = \text{Tr}(U_i|v\rangle\langle v|)$. Thus, this variance is zero if and only if the random variable is constant, which means that for all $|v_0\rangle, |v'_0\rangle \in C$ with $\langle v_0|v_0\rangle = \langle v'_0|v'_0\rangle = 1$ it holds that $\langle v_0|U_i|v_0\rangle = \langle v'_0|U_i|v'_0\rangle$. But this is exactly definition 2.18. This proof transfers to the rains enumerators directly thanks to theorem 3.6. \square

With this results we can see the usefulness of the weight enumerators. Since any quantum code must have a weight enumerator, and we have conditions it has to follow, we can get bounds on the minimal distance or on the dimension for a given number of qubits using these conditions. This proceeding is inspired by classical coding theory, which also gives a theorem for additive codes:

Definition 4.4. Let $n \in \mathbb{N}, j \in \underline{n}$. We define the Krawtchouk polynomials $P_j(x, n)$ as follows:

$$P_j(x, n) := \sum_{s=0}^j (-1)^s 3^{j-s} \binom{x}{s} \binom{n-x}{j-s} \quad (55)$$

Theorem 4.5 (LP bounds for additive codes). Let C be an additive $[[n, k, d]]$ code and $A_i, i = 0 \dots n$ be the Shor-Laflamme weight enumerator of C . Further let the associated \mathbb{F}_4 code have no vectors of weight 1. Further let $P_j(x, n)$ denote the Krawtchouk polynomials. Then, the coefficients of the weight enumerator have to satisfy the following

$$A_0 = 1, A_1 = 0, A_j \geq 0 \text{ for all } j \in \{2, 3, \dots, n\} \quad (56)$$

$$A_0 + A_1 + \dots + A_n = 2^{n-k} \quad (57)$$

$$A'_j := \frac{1}{2^{n-k}} \sum_{r=0}^n P_j(r, n) A_r \text{ for all } j \in \{0, 1, \dots, n\} \quad (58)$$

$$A_j := A'_j \text{ for all } j \in \{0, \dots, d-1\}, A_j \leq A'_j \text{ for all } j \in \{d, \dots, n\} \quad (59)$$

$$\sum_{j=0}^{\frac{n}{2}} A_{2j} = 2^{n-k-1} \text{ or } 2^{n-k} \quad (60)$$

$$\frac{1}{2^{n-k-1}} \sum_{r=0}^n P_j(2r, n) A_{2r} \geq A'_j \text{ for all } j \in \underline{n} \quad (61)$$

A proof of this can be found in [?], and can be proved with classical coding theory. It will be omitted here since an analogous result for general quantum codes will be proven later. Note that the requirement of no vectors of weight 1 is not really much of a restriction: It can be proved (see [?]) that if an $[[n, k, d]]$ code with vectors of weight 1 in the associated \mathbb{F}_4 exists, then also an $[[n-1, k, d]]$ code exists. In particular, for finding bounds we can safely make this assumption.

With this we can get bounds on the distance d for a given pair n, k (the maximal number for which the system of equations on theorem 4.5 have a solution), or bounds on the number of encoded qubits k , by maximizing $2^k = \sum_{i=0}^n A_i$ with the rest of the constraints given by theorem 4.5 - a problem which falls under what is known as linear programming, and can also easily be solved. This bounds hold for additive codes, but they are not strict for general quantum error correcting codes. The following is an example of a non-additive code that is strictly better than any additive code.

Example 4.6. According to theorem 4.5 there is no additive $[[5, k, 2]]$ code for $k > 2$, which means that code space dimension for an additive code with $n = 5$ qubits and minimal distance $d = 2$ is bounded by $2^2 = 4$. There is, however, a $((5, 6, 2))$ code - which means there is strictly better than the $((5, 4, 2))$ bound for additive codes. This code can be given explicitly, for example, by giving a projection operator onto it:

$$P := \frac{1}{16} [3(I_2 \otimes I_2 \otimes I_2 \otimes I_2 \otimes I_2) + (I_2 \otimes Z \otimes Y \otimes Y \otimes Z)_{\text{cyc}} + (I_2 \otimes X \otimes Z \otimes Z \otimes X)_{\text{cyc}} - (I_2 \otimes Y \otimes X \otimes X \otimes Y)_{\text{cyc}} + 2(Z \otimes X \otimes Y \otimes Y \otimes X)_{\text{cyc}} - 2(Z \otimes Z \otimes Z \otimes Z \otimes Z)_{\text{cyc}}]$$

Here, the subscript cyc denotes that it is the sum of all cyclic shifts of that element. A simple calculation lets us see that this is indeed a projection operator ($P^2 = P$) and since $\text{Tr}(P) = 6$ it does project onto a 6-dimensional subspace of \mathbb{C}^{2^5} . That it has minimal distance two can be seen from its weight enumerator, $A(x, y) = x^5 + \frac{10}{6}xy^4 + \frac{8}{3}y^5$ using theorem 4.3. For further properties of the code, see [?]

The main reason for introducing the Shor-Laflamme/Rains enumerators is generalizing the Hamming weight enumerator from additive codes to general quantum codes. In this way, we can get a result analogous to theorem 4.5 which holds for any quantum code.

Theorem 4.7 (LP bound for general quantum codes). Let C be a quantum $((n, K, d))$ code and let $\sum_{i=0}^n A_i x^{n-i} y^i$ be its (Shor-Laflamme) weight enumerator. Then the A_i satisfy the following set of linear equations and inequalities:

$$A_0 = 1, A_i \geq 0 \text{ for all } i \in \underline{n} \quad (62)$$

$$B_i := \frac{1}{K 2^n} \sum_{r=0}^n P_i(r, n) A_r \quad (63)$$

$$A_i = B_i \text{ for all } i \in \{0, \dots, d-1\} \quad (64)$$

$$A_i \leq B_i \text{ for all } i \in \{d, \dots, n\} \quad (65)$$

Proof. $A_0 = 1$ follows directly from the definition. Equations 62, 64 and 65 all follow from theorem 4.3. Equation 63 is, of course, the MacWilliams transform (theorem 3.12). This can be seen immediately if we expand the terms on the MacWilliams transform:

$$A(x, y) = \dim(C) B\left(\frac{x+3y}{2}, \frac{x-y}{2}\right) \Leftrightarrow B(x, y) = \frac{1}{K} A\left(\frac{x+3y}{2}, \frac{x-y}{2}\right) \quad (66)$$

As can be seen with a simple variable substitution. Now, expanding the right hand side of the last equation yields:

$$\begin{aligned} \frac{1}{K} A\left(\frac{x+3y}{2}, \frac{x-y}{2}\right) &= \frac{1}{K} \sum_{r=0}^n \left(\frac{x+3y}{2}\right)^{n-r} \left(\frac{x-y}{2}\right)^r A_r = \frac{1}{2^n K} \sum_{r=0}^n (x+3y)^{n-r} (x-y)^r A_r \\ &= \frac{1}{2^n K} \sum_{r=0}^n \sum_{j=0}^{n-r} \binom{n-r}{j} x^j (3y)^{n-r-j} \sum_{i=0}^r \binom{r}{i} (-1)^i x^i y^{r-i} A_r \\ &= \frac{1}{2^n K} \sum_{r=0}^n \sum_{i=0}^r \sum_{j=0}^{n-r} \binom{r}{i} \binom{n-r}{j} (-1)^i 3^{n-r-j} x^{i+j} y^{n-(i+j)} A_r \end{aligned} \quad (67)$$

From this we can read B_d as the coefficient of $x^{n-d} y^d$ to be the Krawtchouk polynomial.

4.2 Shadow enumerators

Theorem 4.7 gives us a good start on finding bounds on the minimal distance or code space dimension for general quantum error correcting codes. The motivation for this theorem came from the hamming weight enumerators of classical coding theory. There is another enumerator on binary codes which has not been discussed, that (as we will see) can also be generalized for quantum error correcting codes and gives even tighter bounds, namely, the shadow enumerator.

Definition 4.8. Let C be an additive quantum code on n qubits, and S be the associated stabilizer group. We define the shadow of C to be the set $S(C)$ of all $\sigma \in \mathcal{E}$ with $\langle \sigma, \sigma' \rangle \equiv \text{wt}(\sigma') \pmod{2}$ for all $\sigma' \in S$, where $\langle \cdot, \cdot \rangle$ denotes the trace inner product. The (hamming) weight enumerator of $S(C)$ we call the shadow enumerator of C , and denote it by $S(x, y)$.

Just as for the enumerator of the dual code, there is a simple relationship relating the weight enumerator to the shadow enumerator, as the MacWilliams identity. For additive codes this is also very simple:

Proposition 4.9. Let C be an additive quantum code with associated stabilizer group G , then the (Shor-Laflamme) weight enumerator is related to the shadow enumerator in the following way:

$$S(x, y) = A\left(\frac{x+3y}{2}, \frac{y-x}{2}\right) \quad (68)$$

Proof. If S has no elements of odd weight, then the assertion follows from the fact that $S(C) = G^\perp$ in this case, as can be seen from the definition of both. Since for a code with only even-weighted elements, the weight enumerator fulfils $A(x, y) = A(x, -y)$ we get the required identity from the MacWilliams transform. Hence, we assume G has elements of odd weight.

Define $G_0 \subset G$ as the subset of all elements of even weight. We know that G is self-orthogonal with respect to the symplectic form given by the commutator. Applying this we get the following inclusions:

$$G_0 \subset G \subset G^\perp \subset G_0^\perp \quad (69)$$

From the self-orthogonality we also get that G_0 is in fact subspace of G . Let $A^{(0)}, B^{(0)}$ be the (Shor-Laflamme) weight enumerators of G_0 . Now, consider $S(C) = \{\sigma \in \mathcal{E} \mid \langle \sigma, \sigma' \rangle \equiv wt(\sigma') \pmod{2} \text{ for all } \sigma' \in G\} \subseteq G_0^\perp$. In fact, since for any $\sigma \in S(C), \sigma' \in G$ it has to hold that $\langle \sigma, \sigma' \rangle = 0$ if and only if $\sigma' \in G_0$, it turns out that $S(C)$ is exactly $G_0^\perp \setminus G^\perp$. Using the combinatorial interpretation of the weight enumerators (as hamming weight enumerators), we can thus express the shadow enumerator:

$$S(x, y) = B^{(0)}(x, y) - B(x, y) = A^{(0)}\left(\frac{x+3y}{2}, \frac{x-y}{2}\right) - A\left(\frac{x+3y}{2}, \frac{x-y}{2}\right) \quad (70)$$

On the other hand we see that $A(x, y) + A(x, -y) = 2A^{(0)}(x, y)$, as exactly the elements of even weight cancel out like this. This yields the required identity.

As already mentioned, shadow enumerators can be generalized for any quantum code, just as the hamming weight enumerators were. The original article in which these were introduced, [?], includes a detailed discussion of the steps taken to do this generalization, and the rationale behind it. Here, we will just present the finished generalization along with the important accompanying results.

Definition 4.10. Let M be a hermitian operator acting on \mathbb{C}^{2^n} , and let $M = \sum_{\sigma \in \mathcal{E}} c_\sigma \sigma$ be the decomposition in the basis \mathcal{E} . We define

$$M^{(-)} := \sum_{\sigma \in \mathcal{E}} (-1)^{wt(\sigma)} c_\sigma \sigma \quad (71)$$

Definition 4.11. Let C be a quantum code on n qubits and P be the orthogonal projection operator onto C . We define the shadow enumerator coefficients

$$S_d := B_d(P, P^{(-)}) \quad (72)$$

and similarly the shadow enumerator polynomial as

$$S(x, y) := \sum_{d=0}^n S_d x^{n-d} y^d \quad (73)$$

Note that for any hermitian operator M , $M^{(-)}$ is also hermitian, and thus, the shadow enumerator is well defined.

We also get a theorem analogous to prop. 4.9.

Theorem 4.12. Let C be a K -dimensional quantum code on n qubits, P be the orthogonal projection operator onto C , and $A(x, y), S(x, y)$ be the (Shor-Laflamme) weight enumerator and shadow enumerator respectively. Then they are related by the following MacWilliams transform:

$$KS(x, y) = A\left(\frac{x+3y}{2}, \frac{y-x}{2}\right) \quad (74)$$

Proof. Consider the polynomial $W(x, y) := \sum_{d=0}^n A_d(P, P^{(-)})$. Note that though the MacWilliams transform, theorem 3.12, was stated explicitly for the projection operator onto a quantum code, the proof does hold the same for the enumerator coefficients with respect to any two hermitian operators. Thus, by the MacWilliams transform we get that $KS(x, y) = W(\frac{x+3y}{2}, \frac{x-y}{2})$. It is therefore sufficient to prove that $W(x, y) = A(x, -y)$. Writing the coefficients explicitly we get the desired result:

$$\begin{aligned} A_d(P, P^{(-)}) &= \frac{1}{\text{Tr}(P)\text{Tr}(P^{(-)})} \sum_{\substack{\sigma \in \mathcal{E} \\ \text{wt}(\sigma)=d}} \text{Tr}(P\sigma)\text{Tr}(P^{(-)}\sigma) \\ &= \frac{1}{\text{Tr}(P)\text{Tr}(P^{(-)})} \sum_{\substack{\sigma \in \mathcal{E} \\ \text{wt}(\sigma)=d}} (-1)^d \text{Tr}(P\sigma)\text{Tr}(P\sigma) = (-1)^d A_d(P, P) \end{aligned}$$

In the case of stabilizer codes we defined the shadow enumerators via a combinatorial identity, which also restricted many of their properties. Now we have properly generalized shadow codes and identified the relationship to the (Shor-Laflamme) weight enumerators via a MacWilliams transform. We still have, however, no new information from these in the general case. While these do not have to be integers in the general case, we can prove that they have to be nonnegative, which adds indeed new restrictions to the set of equations and inequalities. This is proven in the following theorem:

Theorem 4.13. Let C be a quantum code on n qubits, and let $S(x, y) = \sum_{d=0}^n S_d x^{n-d} y^d$ be its shadow enumerator polynomial. Then

$$S_d \geq 0 \text{ for all } d \in \{0, \dots, n\} \quad (75)$$

Proof. It suffices to show that for the orthogonal projection operator P onto C , $P^{(-)}$ is positive semi-definite. Indeed, for two hermitian, positive semi-definite operators M, N with $\text{Tr}(MN) \neq 0$ we always have that

$$B_d(M, N) = \frac{1}{\text{Tr}(MN)} \sum_{\substack{E \in \mathcal{E} \\ \text{wt}(E)=d}} \underbrace{(\text{Tr}(MENE^{-1}))}_{\geq 0} \quad (76)$$

Now, to show that $P^{(-)}$ is positive definite, we can use the fact that for any hermitian operator M , $M^{(-)} = Y^{\otimes n} \bar{M} Y^{\otimes n}$, where $\bar{\cdot}$ denotes the complex conjugate. Since M is hermitian as are all $E \in \mathcal{E}$, the coefficients c_E of the basis decomposition $M = \sum_{E \in \mathcal{E}} c_E E$ are all real, and thus it suffices to show this holds for $M \in \mathcal{E}$. For $M \in \mathcal{E}$, we have $M^{(-)} = (-1)^{\text{wt} M} M$. Since the products of elements in \mathcal{E} can be computed 'qubit-wise', i.e. diagonally on the tensor product components, we can see what happens for each of I_2, X, Y, Z : $Y I_2 Y = Y^2 = I_2$, $Y Y Y = Y^2 Y = Y$, so I_2 and Y are invariant under conjugation with Y . $Y X Y = i^2 X Z X X Z = -X$, $Y Z Y = i^2 X Z Z X Z = -Z$, so X and Z have both change signs with the conjugation. On the other hand, $X = \bar{X}$ and $Z = \bar{Z}$, but $Y = -\bar{Y}$, so that the equality holds for all cases. Since $Y^{\otimes n}$ is unitary, this means that $P^{(-)}$ is positive semi-definite, since P is as a projection, which completes the proof. \square

With this result we can now refine theorem 4.7 to include the new inequalities, giving tighter bounds.

Theorem 4.14 (tighter LP bound for general quantum codes). Let C be a quantum $((n, K, d))$ code and let $\sum_{i=0}^n A_i x^{n-i} y^i$ be its (Shor-Laflamme) weight enumerator. Then the A_i satisfy the following set of linear equations and inequalities:

$$A_0 = 1, A_i \geq 0 \text{ for all } i \in \underline{n} \quad (77)$$

$$B_i := \frac{1}{K 2^n} \sum_{r=0}^n P_i(r, n) A_r \quad (78)$$

$$A_i = B_i \text{ for all } i \in \{0, \dots, d-1\} \quad (79)$$

$$A_i \leq B_i \text{ for all } i \in \{d, \dots, n\} \quad (80)$$

$$S_i := \frac{1}{K 2^n} \sum_{r=0}^n (-1)^r P_i(r, n) A_r \quad (81)$$

$$S_i \geq 0 \text{ for all } i \in \{0, \dots, n\} \quad (82)$$

Proof. This is just theorem 4.7 enhanced by theorems 4.12 and 4.13. \square