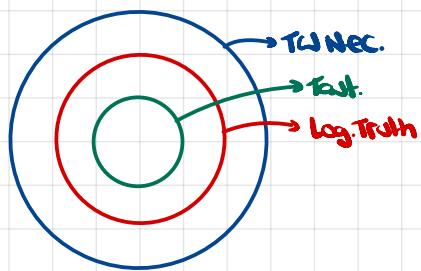


1.

Euclid - Geometry	S-A
Russell - the, both, neither	G-B
Grice - Implicature	4-C
Gödel - no complete axiomatization of arithmetic in FOL	2-F

4.



2.

- A: All FOL sentences are wffs.
 B: All wffs are expressions.
 C: All FOL sentences are expressions.

- A: $\forall x (\text{Sentence}(x) \rightarrow \text{WFF}(x))$
 B: $\forall x (\text{WFF}(x) \rightarrow \text{Expr}(x))$

- C: $\forall x (S(x) \rightarrow E(x))$

S(α)
 $S(\alpha) \rightarrow L(\alpha)$
 $L(\alpha) \rightarrow E(\alpha)$
 $L(\alpha)$
 $E(\alpha)$
 $\forall x (S(x) \rightarrow E(x))$

3.

$\exists x (\text{CNF}(x) \wedge \text{NNF}(x))$
 $\forall x (\text{DNF}(x) \rightarrow \neg \text{NNF}(x))$

$\exists x (\text{CNF}(x) \wedge \neg \text{DNF}(x))$

$\text{CNF}(\alpha) \wedge \text{NNF}(\alpha)$
 $\text{NNF}(\alpha)$
 $\neg (\neg \text{NNF}(\alpha))$
 $\neg \text{DNF}(\alpha)$
 $\text{CNF}(\alpha)$
 $\text{CNF}(\alpha) \wedge \neg \text{DNF}(\alpha)$
 $\forall x (\text{CNF}(x) \wedge \neg \text{DNF}(x))$

- $\exists x (T(x) \wedge L(x)) \quad T$
 $\forall x (L(x) \rightarrow N(x)) \quad T$
 $\exists x (T(x) \wedge N(x))$

All axioms are necessary truths. F
 No theorem is a necessary truth. F

No theorems are axioms

$\forall x (A(x) \rightarrow N(x))$
 $\forall x (T(x) \rightarrow \neg N(x))$

$\forall x (T(x) \rightarrow \neg A(x))$

$T(\alpha)$
 $T(\alpha) \rightarrow \neg N(\alpha)$
 $N(\alpha) \rightarrow N(\alpha)$
 $\neg N(\alpha)$
 $\neg A(\alpha)$
 $\forall x (T(x) \rightarrow \neg A(x))$

6.

$$\forall x (AI(x) \rightarrow \neg S(x)) \quad T$$

$$\exists x (S(x) \wedge \neg V(x)) \quad F$$

$$\forall x (AI(x) \rightarrow \neg V(x))$$

Counterexample:

γ is an unsound argument that is valid.

γ has inconsistent premises.

$$\begin{array}{l} \forall x \neg S(x) \wedge V(x) \\ AI(y) \end{array}$$

$$\text{Modus } AI(y) \wedge V(y)$$

$$\begin{array}{l} \exists x (AI(x) \wedge V(x)) \\ \exists x \neg (\neg AI(x) \vee \neg V(x)) \\ \exists x \neg (AI(x) \rightarrow \neg V(x)) \\ \neg \forall x AI(x) \rightarrow \neg V(x) \end{array}$$

7.

Abrraham can fool all of the people all of the time.

$$\forall x \forall y ((\text{Time}(y) \wedge \text{Person}(x)) \rightarrow \text{Can Fool}(abraham, x, y))$$

Abrraham can't fool ...

$$\neg \forall x \forall y ((\text{Time}(y) \wedge \text{Person}(x)) \rightarrow \text{Can Fool}(abraham, x, y))$$

$$\neg \forall x \forall y (\neg (\text{Time}(y) \wedge \text{Person}(x)) \vee \text{Can Fool}(abraham, x, y))$$

$$\exists x \exists y \neg (\neg (\text{Time}(y) \wedge \text{Person}(x)) \vee \text{Can Fool}(abraham, x, y))$$

$$\exists x \exists y ((\text{Time}(y) \wedge \text{Person}(x)) \wedge \neg \text{Can Fool}(abraham, x, y))$$

8. There is a smaller than every minute

$$\exists x \neg (\forall y (x+y \rightarrow \neg \text{Tet}(y)) \rightarrow \neg \exists z \text{Smaller}(x, z))$$

II. One man in his time plays many parts.

many +

satisfying $P(x, y)$

satisfy $Q(y)$

many +

satisfying $f = 1$

satisfy $Q(y)$

Invariant under expansion (φ) $\leftrightarrow (\varphi_u \rightarrow \varphi_{ut})$

" " contraction (φ) $\leftrightarrow (\varphi_u \rightarrow \varphi_{ut})$

$$\exists x (\forall y (x+y \rightarrow \text{Smaller}(x, y)) \rightarrow \exists z \text{Tet}(z))$$

$$\exists x (\forall y (x+y \rightarrow \text{Smaller}(x, y)) \rightarrow \exists z \text{Tet}(z))$$

$$\exists x (\neg (\forall y (x+y \rightarrow \text{Smaller}(x, y))) \vee \exists z \text{Tet}(z))$$

$$\exists x (\forall y (A(x, y) \rightarrow B(x, y)) \rightarrow \exists z C(z))$$

$$\exists x (\exists y (\forall z (x+z \wedge \neg \text{Smaller}(x, z)) \vee \exists z \text{Tet}(z)))$$

$$\exists x (\neg (\forall y (A(x, y) \rightarrow B(x, y))) \vee \exists z C(z))$$

$$\exists x \exists y \exists z ((x+y \wedge \neg \text{Smaller}(x, y)) \vee \exists z \text{Tet}(z))$$

$$\exists x (\exists y \neg (\neg A(x, y) \vee B(x, y)) \vee \exists z C(z))$$

$$\exists x \exists y \exists z ((x+y \rightarrow \text{Smaller}(x, y)) \rightarrow \exists z \text{Tet}(z))$$

$$\exists x \exists y \exists z ((A(x, y) \rightarrow B(x, y)) \rightarrow \exists z C(z))$$

TRE

$$\exists x (\forall y (x+y \rightarrow \text{Smaller}(x, y)) \text{ TKE}$$

$$\exists x (\exists y \neg (x+y \rightarrow \text{Smaller}(x, y)) \text{ TKE} \\ \neg (x+y \vee \text{Smaller}(x, y)) \\ (x+y \wedge \neg \text{Smaller}(x, y))$$

IN EXP.

\rightarrow IN CON.

$$\forall x \forall y \ x = y$$

$$\exists x \forall y \ x = y \mid x + y$$

$$\forall x (\exists y (x + y \rightarrow \text{SameSize}(x, y)) \rightarrow \neg \exists z \text{Tel}(z))$$

$$\forall x (\neg \exists y (x + y \rightarrow \text{SameSize}(x, y)) \vee \neg \exists z \text{Tel}(z))$$

$$\forall x (\forall y \neg (x + y \rightarrow \text{SameSize}(x, y)) \vee \neg \exists z \text{Tel}(z))$$

$$\forall x (\forall y \neg (x + y \rightarrow \text{SameSize}(x, y)) \vee \forall z \neg \text{Tel}(z))$$

$$\forall x \forall y \forall z ((x + y \rightarrow \text{SameSize}(x, y)) \rightarrow \neg \text{Tel}(z))$$

When is $\forall x (\exists y (x + y \rightarrow \text{SameSize}(x, y)) \rightarrow \neg \exists z \text{Tel}(z))$ true?

• $\forall x \exists y (x + y \rightarrow \text{SameSize}(x, y))$ false.

$\forall x (\neg \exists y (x + y \rightarrow \text{SameSize}(x, y)))$ true

$\forall x (\forall y \neg (x + y \rightarrow \text{SameSize}(x, y)))$

$\forall x \forall y (x + y \neg \rightarrow \text{SameSize}(x, y))$

↳ Inv Exp.

• $\neg \exists z \text{Tel}(z)$ true

When is $\forall x \forall y \forall z ((x + y \rightarrow \text{SameSize}(x, y)) \rightarrow \neg \text{Tel}(z))$ true?

• $\forall x \forall y \forall z (x + y \rightarrow \text{SameSize}(x, y))$ false

$\forall x \forall y \forall z \neg (x + y \rightarrow \text{SameSize}(x, y))$

$\forall x \forall y \forall z \neg (x + y \vee \text{SameSize}(x, y))$

• $\forall x \forall y \forall z \neg \text{Tel}(z)$ true

$$\exists x \neg (\forall y (x+y \rightarrow \neg \text{Tet}(y)) \rightarrow \exists z \text{Smaller}(x, z))$$

$$\exists x \neg (\neg \forall y (x+y \rightarrow \neg \text{Tet}(y)) \vee \exists z \text{Smaller}(x, z))$$

$$\exists x (\exists y \neg (x+y \rightarrow \neg \text{Tet}(y)) \vee \exists z \text{Smaller}(x, z))$$

$$\exists x (\forall y (x+y \rightarrow \neg \text{Tet}(y)) \wedge \exists z \text{Smaller}(x, z))$$

$$\exists x \forall y \exists z ((x+y \rightarrow \neg \text{Tet}(y)) \wedge \text{Smaller}(x, z))$$

$$\exists x \forall y \exists z \neg (\neg (x+y \rightarrow \neg \text{Tet}(y)) \vee \neg \text{Smaller}(x, z))$$

$$\exists x \forall y \exists z \neg ((x+y \rightarrow \neg \text{Tet}(y)) \rightarrow \neg \text{Smaller}(x, z))$$

TRUE if

$$x+y \rightarrow \neg \text{Tet}(y) \quad \text{and} \quad \text{Smaller}(x, z)$$

$$x+y \vee \neg \text{Tet}(y) \quad \text{and} \quad \text{Smaller}(x, z)$$

There is one object and it isn't a tet. There is another object and the sole object is smaller than it.

Always false. Incarnation.