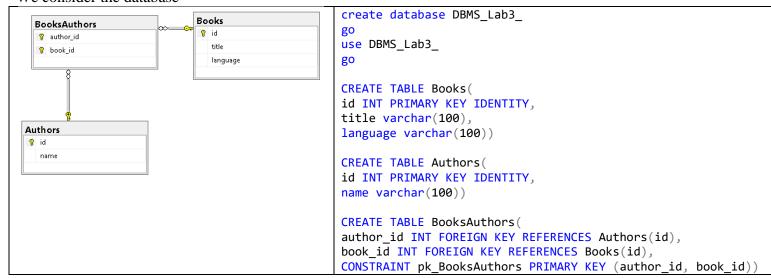
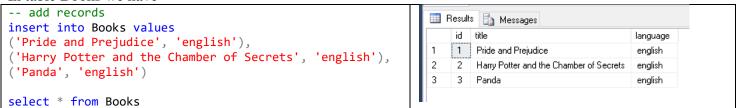
Concurrency problems – a simple example

We consider the database



In table Books we have



Please, put in each file use DBMS Lab3

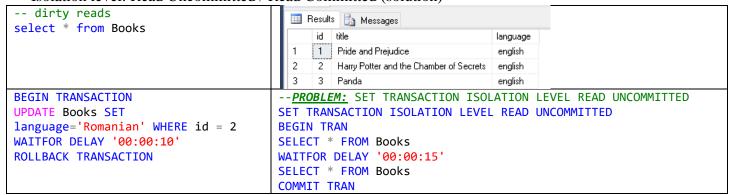
In what follows, we will work with the table **Books** and

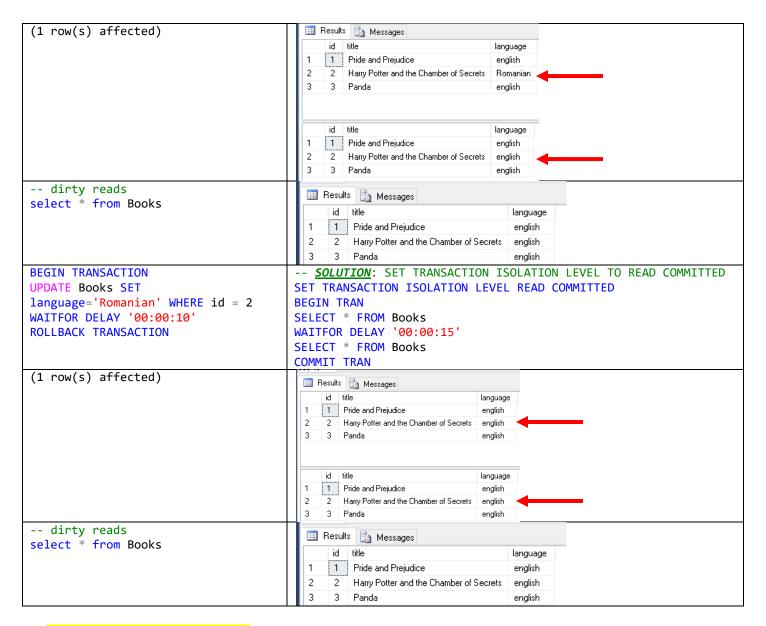
T1=Transaction 1 starts first and finish first (the first column form the table(s))

T2=Transaction start immediately after T1 and finish after T1 (the second column form the table(s))

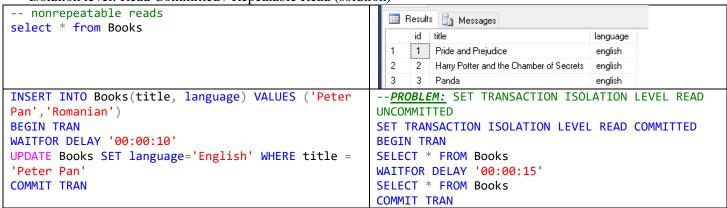
1. DIRTY READS – T1: update + delay + rollback, T2: select + delay + select -> see the update in the first select, even if it is rollback then (the order in the execution of the operations is: update – select – rollback – select)

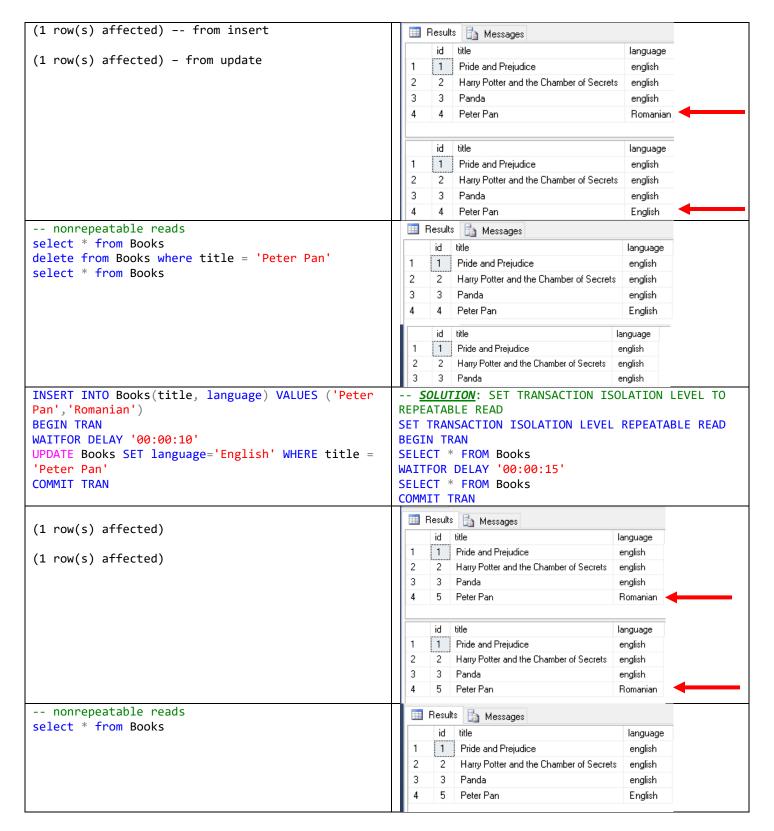
Isolation level: Read Uncommitted / Read Committed (solution)



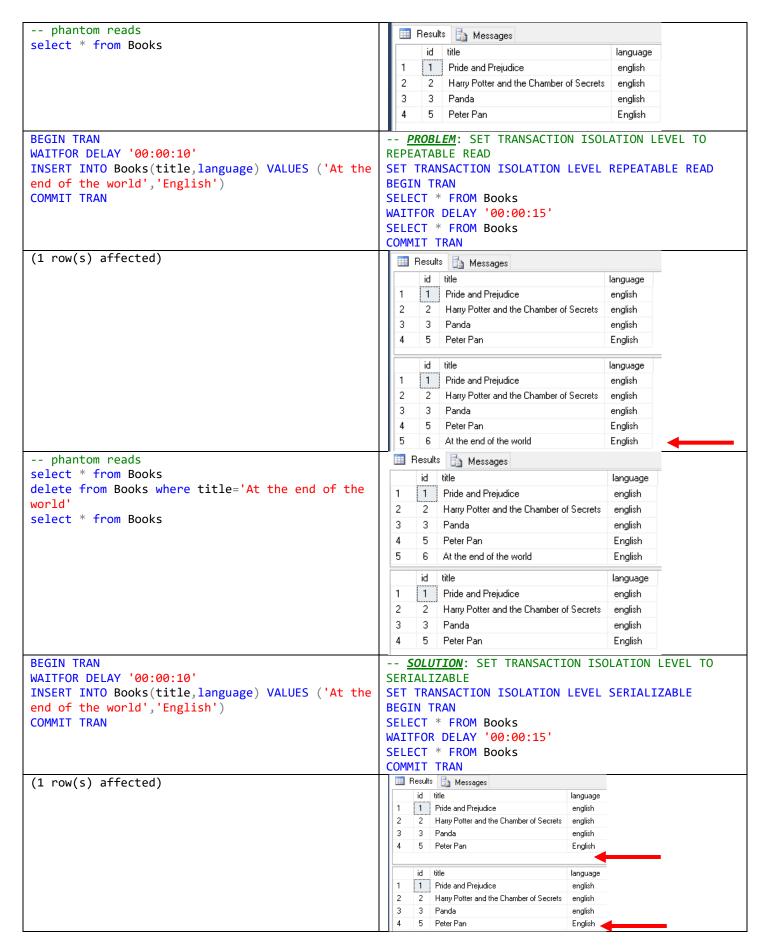


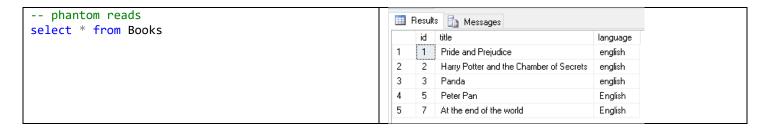
2. NON-REPEATABLE READS – T1: delay + update + commit, T2: select + delay + select -> see the value inserted before the transaction from the first select of T2 + see the update of the value inserted before the transaction, from the second select of T2 (the order in the execution of the operations is: select – update – select) Isolation level: Read Committed / Repeatable Read (solution)





3. PHANTOM READS – T1: delay + insert + commit, T2: select + delay + select -> see the inserted value only at the second select from T2 (the order in the execution of the operations is: select – insert – select) Isolation level: Repeatable Read / Serializable (solution)





DEADLOCK - T1: update on table A + delay + update on table B
 T2: update on table B + delay + update on table A

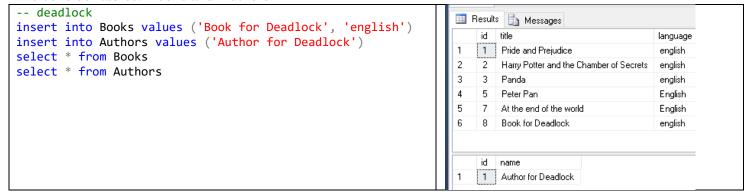
T1- update on table A ->	delay	Try to update table	Table B is blocked in	One of the blocked
exclusive lock on table A		В	T2	transactions, T1 or
				T2, will be chosen
T2 - update on table B ->	delay			as a deadlock victim
exclusive lock on table B		Try to update table	Table A is blocked in	and terminates with
		A	T1	an error. The other
				transaction wins and
				update both table A
				and table B

The only solution is to decide which of the 2 transactions to win, by using the DEADLOCK_PRIORITY, that can be set (LOW, NORMAL, HIGH, or from -10 (-5) to 10 (5)). Implicit is NORMAL (0).

The victim transaction is chosen like this:

- The transaction with the lowest DEADLOCK_PRIORITY
- 2. If both of the transactions have the same DEADLOCK_PRIORITY, the victim is the one, less expensive at ROLLBACK
- 3. If both of the transactions have the same DEADLOCK_PRIORITY and the same cost, the victim is chosen randomly

We consider tables: Books and Authors

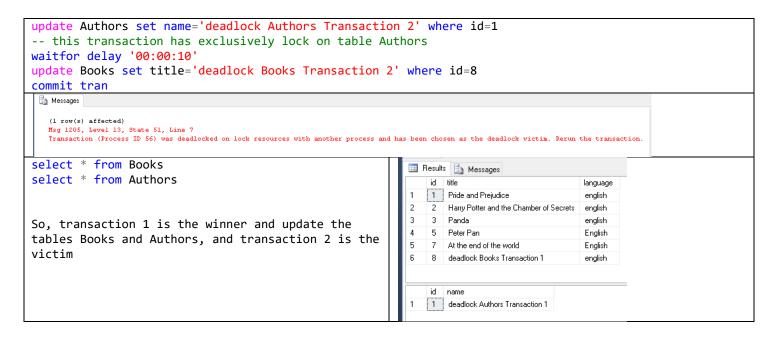


Deadlock example:

```
-- transaction 1
begin tran
update Books set title='deadlock Books Transaction 1' where id=8
-- this transaction has exclusively lock on table Books
waitfor delay '00:00:10'
update Authors set name='deadlock Authors Transaction 1' where id=1
commit tran
-- transaction 2
begin tran

(1 row(s) affected)

(1 row(s) affected)
```



If in transaction 2, we set DEADLOCK_PRIORITY to HIGH, or, if in transaction 1 we set DEADLOCK PRIORITY to LOW, the winner will be transaction 2 and the victim transaction 1.

