



**AMRITA**  
VISHWA VIDYAPEETHAM

**NAME: KURRA VENKATA GOKUL**

**CH.SC.U4CSE24121**

**CSE-B**

## 1.BUBBLE SORT

CODE:

```
#include <stdio.h>
// CH.SC.U4CSE24121 - Bubble Sort

void bubbleSort(int a[], int n) {
    int i, j, temp, swapped;
    for (i = 0; i < n - 1; i++) {
        swapped = 0;
        for (j = 0; j < n - i - 1; j++) {
            if (a[j] > a[j + 1]) {
                temp = a[j];
                a[j] = a[j + 1];
                a[j + 1] = temp;
                swapped = 1;
            }
        }
        if (!swapped)
            break;
    }
}

int main() {
    printf("CH.SC.U4CSE24121\n");
    int a[5] = {5, 1, 4, 2, 8};
    int n = 5, i;
    bubbleSort(a, n);
    for (i = 0; i < n; i++)
        printf("%d ", a[i]);
    return 0;
}
```

## OUTPUT:

```
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>gcc bubble.c  
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>.\bubble  
CH.SC.U4CSE24121  
1 2 4 5 8  
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>|
```

### Time Complexity:

- Best Case:  $O(n)$
- Average Case:  $O(n^2)$
- Worst Case:  $O(n^2)$

### Justification:

The algorithm uses two nested loops.

- The outer loop runs  $(n - 1)$  times.
- The inner loop compares adjacent elements.
- In the best case (already sorted), the loop terminates early.  
Hence, time complexity is  $O(n^2)$  in average and worst cases.

### Space Complexity:

- Space Complexity:  $O(n)$

### Justification:

The program uses an array of size  $n$  to store elements.

Only a few extra variables ( $i$ ,  $j$ ,  $temp$ ,  $swapped$ ) are used.

No additional memory is allocated.

Therefore, space complexity is  $O(n)$ .

## 2. Insertion Sort

Code:

```
#include <stdio.h>
// CH.SC.U4CSE24121 - Insertion Sort

void insertionSort(int a[], int n) {
    int i, key, j;
    for (i = 1; i < n; i++) {
        key = a[i];
        j = i - 1;
        while (j >= 0 && a[j] > key) {
            a[j + 1] = a[j];
            j--;
        }
        a[j + 1] = key;
    }
}

int main() {
    printf("CH.SC.U4CSE24121\n");
    int a[5] = {12, 11, 13, 5, 6};
    int n = 5, i;
    insertionSort(a, n);
    for (i = 0; i < n; i++)
        printf("%d ", a[i]);
    return 0;
}
```

## Output:

```
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>gcc -o insert insert.c  
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>.\insert  
CH.SC.U4CSE24121  
5 6 11 12 13  
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>|
```

## Time Complexity:

- Best Case:  $O(n)$
- Average Case:  $O(n^2)$
- Worst Case:  $O(n^2)$

## Justification:

The outer loop runs  $(n - 1)$  times.

- In the best case, the array is already sorted and no shifting is needed.
- In the worst case, each element is compared with all previous elements.

Hence, time complexity is  $O(n^2)$ .

## Space Complexity:

- Space Complexity:  $O(n)$

## Justification:

The program stores elements in an array of size  $n$ .

Only constant extra variables ( $i, j, \text{key}$ ) are used.

No extra data structures are created.

Therefore, space complexity is  $O(n)$ .

### 3.selection sort

Code:

```
#include <stdio.h>
// CH.SC.U4CSE24121 - Selection Sort

void selectionSort(int a[], int n) {
    int i, j, min, temp;
    for (i = 0; i < n - 1; i++) {
        min = i;
        for (j = i + 1; j < n; j++) {
            if (a[j] < a[min])
                min = j;
        }
        temp = a[i];
        a[i] = a[min];
        a[min] = temp;
    }
}

int main() {
    printf("CH.SC.U4CSE24121\n");
    int a[5] = {64, 25, 12, 22, 11};
    int n = 5, i;
    selectionSort(a, n);
    for (i = 0; i < n; i++)
        printf("%d ", a[i]);
    return 0;
}
```

## Output:

```
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>.\select
CH.SC.U4CSE24121
11 12 22 25 64
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>|
```

## Time Complexity:

- Best Case:  $O(n^2)$
- Average Case:  $O(n^2)$
- Worst Case:  $O(n^2)$

## Justification:

The algorithm uses two nested loops.

- The inner loop always scans the unsorted part of the array.
- The number of comparisons remains the same in all cases.  
Hence, time complexity is  $O(n^2)$ .

## Space Complexity:

- Space Complexity:  $O(n)$

## Justification:

The input array of size  $n$  is used.

Only a few extra variables ( $i$ ,  $j$ ,  $min$ ,  $temp$ ) are used.

No additional memory is required.

So, space complexity is  $O(n)$ .

#### 4.bucket sort

Code:

```
#include <stdio.h>
// CH.SC.U4CSE24121 - Bucket Sort

void bucketSort(int a[], int n) {
    int bucket[100] = {0};
    int i;
    for (i = 0; i < n; i++)
        bucket[a[i]]++;

    int index = 0;
    for (i = 0; i < 100; i++) {
        while (bucket[i] > 0) {
            a[index++] = i;
            bucket[i]--;
        }
    }
}

int main() {
    printf("CH.SC.U4CSE24121\n");
    int a[7] = {10, 3, 5, 2, 9, 3, 1};
    int n = 7, i;
    bucketSort(a, n);
    for (i = 0; i < n; i++)
        printf("%d ", a[i]);
    return 0;
}
```



## Output:

```
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>gcc -o bucket bucket.c

C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>.\bucket
CH.SC.U4CSE24121
1 2 3 3 5 9 10
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>|
```

## Time Complexity:

- Best Case:  $O(n + k)$
- Average Case:  $O(n + k)$
- Worst Case:  $O(n + k)$

(where  $k$  is the number of buckets)

## Justification:

- One loop inserts  $n$  elements into buckets.
- Another loop scans all  $k$  buckets.  
Since  $k$  is fixed, total time depends mainly on  $n$ .

## Space Complexity:

- **Space Complexity:  $O(n + k)$**

## Justification:

The program uses:

- Input array of size  $n$
- Bucket array of size  $k$

Therefore, total space used is  $O(n + k)$ .

## 5. Max heap

Code:

```
#include <stdio.h>
// CH.SC.U4CSE24121 - Max Heap

void maxHeapify(int a[], int n, int i) {
    int largest = i;
    int left = 2 * i + 1;
    int right = 2 * i + 2;
    int temp;

    if (left < n && a[left] > a[largest])
        largest = left;
    if (right < n && a[right] > a[largest])
        largest = right;

    if (largest != i) {
        temp = a[i];
        a[i] = a[largest];
        a[largest] = temp;
        maxHeapify(a, n, largest);
    }
}

void buildMaxHeap(int a[], int n) {
    int i;
    for (i = n / 2 - 1; i >= 0; i--)
        maxHeapify(a, n, i);
}

int main() {
    printf("CH.SC.U4CSE24121\n");
    int a[6] = {3, 9, 2, 1, 4, 5};
    int n = 6, i;
    buildMaxHeap(a, n);
    for (i = 0; i < n; i++)
        printf("%d ", a[i]);
    return 0;
}
```

## Output:

```
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>gcc -o heap heap.c  
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>.\heap  
CH.SC.U4CSE24121  
9 4 5 1 3 2  
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>|
```

## Time Complexity:

- Best Case:  $O(n)$
- Average Case:  $O(n)$
- Worst Case:  $O(n)$

## Justification:

Heap is built by calling heapify from  $n/2 - 1$  to 0.

Heapify takes logarithmic time but fewer operations are needed at lower levels.

Overall heap construction takes  $O(n)$  time.

## Space Complexity:

- Space Complexity:  $O(n)$

## Justification:

The heap is stored in an array of size  $n$ .

Only recursive calls and a few variables are used.

Hence, space complexity is  $O(n)$ .

## 6.Min Heap

### Code:

```
#include <stdio.h>
// CH.SC.U4CSE24121 - Min Heap

void minHeapify(int a[], int n, int i) {
    int smallest = i;
    int left = 2 * i + 1;
    int right = 2 * i + 2;
    int temp;

    if (left < n && a[left] < a[smallest])
        smallest = left;
    if (right < n && a[right] < a[smallest])
        smallest = right;

    if (smallest != i) {
        temp = a[i];
        a[i] = a[smallest];
        a[smallest] = temp;
        minHeapify(a, n, smallest);
    }
}

void buildMinHeap(int a[], int n) {
    int i;
    for (i = n / 2 - 1; i >= 0; i--)
        minHeapify(a, n, i);
}

int main() {
    printf("CH.SC.U4CSE24121\n");
    int a[6] = {3, 9, 2, 1, 4, 5};
    int n = 6, i;
    buildMinHeap(a, n);
    for (i = 0; i < n; i++)
        printf("%d ", a[i]);
    return 0;
}
```

## Output:

```
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>gcc -o mheap mheap.c
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>.\mheap
CH.SC.U4CSE24121
1 3 2 9 4 5
```

## Time Complexity:

- Best Case:  $O(n)$
- Average Case:  $O(n)$
- Worst Case:  $O(n)$

## Justification:

Similar to max heap.

Heapify is applied to non-leaf nodes only.

Overall heap construction takes linear time  $O(n)$ .

## Space Complexity:

- **Space Complexity:  $O(n)$**

## Justification:

The array of size  $n$  stores heap elements.

Only recursion stack and few variables are used.

So, space complexity is  $O(n)$ .

## 7.BFS

### Code:

```
#include <stdio.h>
#define MAX 100
// CH.SC.U4CSE24121 - BFS

int queue[MAX], front = 0, rear = 0;

void enqueue(int x) {
    queue[rear++] = x;
}

int dequeue() {
    return queue[front++];
}

void bfs(int graph[MAX][MAX], int n, int start) {
    int visited[MAX] = {0};
    int i, node;

    enqueue(start);
    visited[start] = 1;

    while (front != rear) {
        node = dequeue();
        printf("%d ", node);
        for (i = 0; i < n; i++) {
            if (graph[node][i] == 1 && !visited[i]) {
                visited[i] = 1;
                enqueue(i);
            }
        }
    }
}

int main() {
    printf("CH.SC.U4CSE24121\n");
    int n = 4;
    int graph[MAX][MAX] = {
        {0,1,1,0},
        {1,0,1,1},
        {1,1,0,0},
        {0,1,0,0}
    };
    bfs(graph, n, 0);
    return 0;
}
```

## Output:

```
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>gcc -o bfs bfs.c  
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>.\bfs  
CH.SC.U4CSE24121  
0 1 2 3
```

## Time Complexity:

- Time Complexity:  $O(V^2)$

## Justification:

The graph is represented using an adjacency matrix.

- Each vertex is visited once.
- For every vertex, all  $V$  vertices are checked.  
Hence, time complexity is  $O(V^2)$ .

## Space Complexity:

- Space Complexity:  $O(V^2)$

## Justification:

The program uses:

- Adjacency matrix of size  $V \times V$
- Queue and visited array  
Therefore, space complexity is  $O(V^2)$ .

## 8.DFS

Code:

```
#include <stdio.h>
#define MAX 100
// CH.SC.U4CSE24121 - DFS

void dfs(int graph[MAX][MAX], int visited[], int n, int node) {
    int i;
    visited[node] = 1;
    printf("%d ", node);

    for (i = 0; i < n; i++) {
        if (graph[node][i] == 1 && !visited[i])
            dfs(graph, visited, n, i);
    }
}

int main() {
    printf("CH.SC.U4CSE24121\n");
    int n = 4, i;
    int visited[MAX] = {0};
    int graph[MAX][MAX] = {
        {0,1,1,0},
        {1,0,1,1},
        {1,1,0,0},
        {0,1,0,0}
    };
    dfs(graph, visited, n, 0);
    return 0;
}
```

Output:

```
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>gcc -o dfs dfs.c
C:\Users\kurra\OneDrive\Desktop\SEM IV\DAA>.\dfs
CH.SC.U4CSE24121
0 1 2 3
```



**Time Complexity:**

- **Time Complexity:  $O(V^2)$**

**Justification:**

DFS uses an adjacency matrix.

- For each vertex, all other vertices are checked.  
Hence, time complexity is  $O(V^2)$ .

**Space Complexity:**

- **Space Complexity:  $O(V^2)$**

**Justification:**

The program uses:

- Adjacency matrix
- Visited array
- Recursion stack

Thus, space complexity is  $O(V^2)$ .