

# Real Time Systems

Liu

# According to their timing attributes, real time applications are of 4 types

## Purely Cyclic

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every task executes periodically, I/O operations are polled, resources do not vary from period to period. eg: digital controllers

## Mostly cyclic

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most tasks execute periodically, must respond to external events asynchronously. eg: modern avionics

## Asynchronous and somewhat predictable

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tasks not periodic, resources vary from period to period. eg: radar signal processing

## Asynchronous and unpredictable

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tasks have high run time complexities

# Important Terms

- Release time: instant of time at which job becomes available for execution
- Deadline: instant of time by which execution is required to be completed
- Response time: length of time from the release time of the job to the instant when it completes
- Relative deadline: maximum allowable response time
- $\text{Deadline} = \text{release time} + \text{relative deadline}$
- Tardiness: measure of how late a job completes respective to deadline.

# Important Terms

- Release time jitter: Range between earliest release time and latest release time
- Execution time,  $e_i$ : amount of time required to complete the execution
- Laxity type: indicated whether the job is soft or hard
- Usefulness function: gives usefulness of result as a function of tardiness.

# Hard vs Soft

## **Hard Deadline**

failure to meet deadline is fatal

## **Soft Deadline**

failure to meet deadline is not fatal

Performance degrades in case of missing deadlines

# Periodic Task

- task is executed repeatedly at regular or semi regular time intervals
- Its period ( $p_i$ ) is the minimum length of all time intervals between release times of consecutive jobs
- Its execution time ( $e_i$ ) is the maximum execution time of all jobs in it
- Accuracy of the periodic task model decreases with increase in jitter in release times and variations in execution times

# Periodic Task

- Release time of first job in each task is called the phase of that task ( $\phi_i$ ).
- Hyperperiod ( $H_i$ ) is the LCM of periods of all tasks
- Utilization of a task,  $u_i = e_i/p_i$
- Total Utilization = sum of all individual utilizations

## Sporadic

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Tasks containing jobs that are released at random time instants and have hard deadlines

## Aperiodic

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Tasks containing jobs that are released at random time instants and have soft deadlines or no deadlines

## Preemptable Jobs

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its execution can be stopped any time and can be resumed from the point of suspension

## Nonpreemptable Jobs

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execution should not be interrupted



# Conditions for a valid schedule

- Every processor is assigned to at most one job at any time
- Every job is assigned at most one processor at any time
- No job is scheduled before its release time
- The total amount of processor time assigned to every job is equal its maximum or actual execution time
- All the precedence and resource usage constraints are satisfied.

# Clock Driven Approach

- Decisions on what jobs execute at what time are made at specific time instances. These instances are chosen before the system begins execution.
- All parameters of a hard real time job are fixed and known
- Schedule of the jobs is computed offline and is stored for use at runtime. Scheduler schedules according to this schedule at each scheduling decision time
- Minimal runtime overhead

# Weighted Round-Robin Approach

- Every job joins a FIFO queue when it becomes ready for execution
- If the job is not completed by the end of the time slice, it is preempted and places at the end of the queue
- Processor sharing algorithm
- Each job gets  $1/n$ th share of the processor
- More the weight of the job, more the time fraction allocated to the job

# Priority Driven Approach

- resource idles only when no job requiring the resource is ready for execution
- it tries to make locally optimal decisions
- Processor allocated to processes according to priority

# Earliest Deadline First (EDF)

- process with earliest deadline given highest priority
- processes do not share resources
- Can be preempted

# Least Slack Time First (LSF)

- slack = relative deadline – execution left
- Optimal whenever EDF is optimal

# Rate Monotonic Scheduling (RM)

- For periodic tasks only
- Task priority inversely proportional to its period