

Slot Shifting

L J GOKUL VASAN

INTRODUCTION

- An algorithm that tries to give schedulable guarantee for Event Based tasks within the periodic task-set that is offline guaranteed.
- Event Based Tasks can be classified as
 - Soft aperiodics
 - tasks with no deadline
 - Firm aperiodics
 - tasks with deadlines but dl miss will not collapse system
 - Hard aperiodics
 - Task with deadlines but dl miss will collapse/cause fatal system failure.
- Slot shifting can only support Soft and Firm aperiodics.

Overview on Working

- The Algorithm Works in 2 phases
 1. Offline Phase
 2. Online Phase
- Fundamentals of online **Selection Function** is EDF.
- The **Decision Function** is based on uniform time intervals called **SLOTS**.

1. OFFLINE PHASE

- Basic schedulable test is done on **periodics** .
- **A layer of certainty** is applied on periodics called **Interval** notated as **I** defined with following terms.
 - **End(I)** : $deadline(J_i)$. $i \in$ all jobs with same deadline
 - **Est(I)** : $\min (est(J_i))$. $i \in$ all jobs with same deadline
 - **Start(I)**: $\max(est(I_i), end(I_{i-1}))$
 - **Length(I)**: $| I_i | \rightarrow End(I_i) - Start(I_i)$.
 - **SC(I)** : $| I_i | - \sum WCET(J_i) + \min(SC(I_{i+1}), 0)$, Spare Capacity of the interval I
 - So Basic record of interval looks like:

```
struct interval {  
    int id; /// unique ID for the interval  
    int start; /// start of interval  
    int end; /// end of interval  
    int sc; /// spare capacity of interval  
    list *jobs; /// list of jobs that belong to this interval  
    list *nxt; /// reference to next interval / NULL  
};
```

1. OFFLINE PHASE

function **create_new_interval()**

input s, e : s is the start of interval

e is the end of interval

output i : i is the created new interval

$i :=$ allocate memory for new record of type struct interval

$i.start := s$;

$i.end := e$;

$i.jobs := NULL$;

$i.next := NULL$;

return i ;

function **get_first_job_dl()**

input j : j is list of jobs

output j^0 : j^0 is the 1st job in the list / $NULL$.

if(j) return $j[0]$; else return $NULL$;

function **get_rmve_1st_job()**

input j : j is list of jobs

output j^0 : j^0 is the 1st job in list, removed from list / $NULL$

if(j)

$j^0 = \text{get_first_job_dl}(j)$;

remove_job(j, j^0); /// removes job j^0 from list j

return j^0 ;

else return $NULL$;

function **CreateInterval()**

input T : T is schedulable periodic task-set

output t, l : t is list of jobs sorted based on dl.

l is list of intervals

$st :=$ sort jobs of task-set based on deadline.

$t := st$ ///duplicate list of st on t

while(st) **//CREATE INTERVAL**

$l.next = \text{create_new_interval}(0,0)$;

$l.end = \text{get_first_job_dl}(st)$;

while($l.end == \text{get_first_job_dl}(st)$)

$l.jobs.next = \text{Get_rmve_1st_job}(st)$;

$est_i = \min(l.jobs.est)$;

$l.start = \max(est_i, l_{i-1}.end)$;

while(l) **//CREATE SC OF INTERVAL**

$l.sc = |l| - \sum WCET(l.jobs) + \min(l_{i+1}.sc, 0)$;

$l = l.next$;

return t, l ;

2. ONLINE PHASE

- Basic selection function is based on EDF on Guaranteed ready list or FCFS on !Guaranteed list.
- Selection function also manages **updating intervals and its spare capacity**.
- It checks **Firm aperiodics arrival**; if the presence is detected then **acceptance test** is run and if it can be guaranteed then **guarantee algorithm** is run.
- If **Soft aperiodics arrives**, then simply add the task to not-guaranteed list and schedule them when slot is empty after guaranteed EDF function.

2. ONLINE PHASE

Selection_function ()

input t_{prev} : t_{prev} is the previous task that was scheduled.

output t_{nxt} : t_{nxt} is the next task to be scheduled.

///1.INTERVAL UPDATE

update_sc (t_{prev}); /// update spare capacity of the current interval based on prev scheduled task.

I = Get_current_interval();

if(slot_count > I.end) /// check we are on end of interval, if so move to next interval

move_to_next_interval();

end if

///2.APERIODIC CHECK

If(aperiodic_task) ///check on aperiodics task

if(FIRM == aperiodic_task.type)

acceptance_guarantee_algorithm(aperiodic_task); /// run acceptance and guarantee algorithm

else if(SOFT == aperiodic_task.type)

add_to_notGuaranteed_list(aperiodic_task);

end if

///3.EDF

t_{nxt} = get_nxt_guaranteed_ready_task(); /// EDF on ready list

if(! t_{nxt})

t_{nxt} = get_nxt_notGuaranteed_ready_task(); ///FCFS on not guaranteed ready list

end if

return t_{nxt} ;

2. ONLINE PHASE

function Update_till_positive()

Input *l* : *l* is the interval from where sc needs update in backwards.

i_tmp = *l*;

do{

i_tmp.sc++;

i_tmp = get_prev_interval(*i_tmp*);

 if(*i_tmp* == get_current_interval())

 break;

 end if

}while(*i_tmp.sc* < 0);

i_tmp.sc++;

function Update_sc()

input *t*: *t* is the job that got scheduled

l := get_current_interval();

while(*l.job*)

 if(*l.job* == *t*) return;

l.job = *l.job.nxt*;

i_tmp = get_task_interval(*t*);

if(*i_tmp.sc* < 0)

update_till_positive(*i_tmp*);

end if

negate_sc(*l*);

2. ONLINE PHASE

FIRM APERIODIC ARRIVAL (J_A): *acceptance_guarantee_algorithm(J_A)*

- On Firm aperiodic arrival acceptance test needs to run on task to check whether it can be admitted.
- To run the **acceptance test** we traverse through 3 kinds intervals and sum there SC.
 1. $SC(I_C)$: Spare capacity of current interval.
 2. $SC(I_i) \cdot C < i \leq l$. $end(I_i) \leq dl(J_A) \wedge end(I_{l+1}) > dl(J_A)$.
 - All intervals that are between current and last interval.
 3. $Min[SC(I_{l+1}), dl(J_A) - start(I_{l+1})]$
 - Required spare capacity in the last interval within $dl(J_A)$.

If(1+2 +3 \geq WCET(J_A)) then J_A can be accepted i.e. return true.

- Simple snippet of acceptance guarantee Algorithm :

Acceptance_guarantee_algorithm(J_A)

if(acceptance_test(J_A))

guarantee_algorithm(J_A);

else

move_to_notGuaranteed_list(J_A);

2. ONLINE PHASE

TRADITIONAL GUARANTEE ALGORITHM:

Function Guarantee_algorithm(J_A)

if($dl(J_A) < end(I_{l+1})$)

split_interval(I_{l+1}, J_A);

$I_i = I_{l+1}$;

while($I_i \neq I_c$)

$sc(I_i) = |I_i| - \sum WCET(J_j) + \min(SC(I_{i+1}), 0)$;

$I_i = get_prev_interval(I_i)$;

Complexity : $O(N.t)$ where N is the number of intervals and t is the jobs in that intervals.

- Redoing what we did offline and traversing both vertically and horizontally in a list is very cumbersome to implement.
- We tried to use this offline effort to recalculate online change in SC without having vertical traversal. Giving us $O(N)$ guarantee algorithm.

2. ONLINE PHASE : $O(N)$ GUARANTEE ALGORITHM

Function split_interval()

Input : I_{right} , split_point, curr_slot

Output : I_{left}

If ($I_{\text{right}} == \text{get_current_interval}()$)

 new_len = (split_point + start(I_{right})) - curr_slot;

else

 new_len = split_point + start(I_{right}) - start(I_{right});

$SC(I_{\text{right}}) = SC(I_{\text{right}}) - \text{new_len};$

$SC(I_{\text{left}}) = \text{new_len} + \min(0, SC(I_{\text{right}}));$

$END(I_{\text{left}}) = \text{split_point} + \text{start}(I_{\text{right}});$

$START(I_{\text{left}}) = \text{start}(I_{\text{right}});$

$START(I_{\text{right}}) = \text{split_point} + \text{start}(I_{\text{right}});$

Add_intr_to_list (I_{left});

return I_{left}

Function **Guarantee_algorithm()**

Input I_{dl} , J_A , slot_no : I_{dl} is the interval in which aperiodic deadline ,i.e. I_{t+1}
 J_A is the aperiodic task that cleared acceptance test.
slot_no is the current slot number.

Output : NONE

If ($dl(J_A) < END(I_{\text{dl}})$)

 split_point = $DL(J_A) - START(I_{\text{dl}})$;

$I_{\text{left}} = \text{split_interval}(I_{\text{dl}}, \text{split_point}, \text{slot_no})$;

$J_A.\text{Interval} = I_{\text{left}} ? I_{\text{left}} : I_{\text{dl}} ;$

$\text{delta} := \text{WCET}(J_A) ;$

$i_temp := J_A.\text{Interval}$;

while(delta)

if($SC(i_temp) > 0$)

if($SC(i_temp) \geq \text{delta}$)

$SC(i_temp) = SC(i_temp) - \text{delta}$;

 delta = 0;

else

 delta = delta - $SC(i_temp)$;

$SC(i_temp) = -\text{delta}$;

else

$SC(i_temp) += -\text{delta}$;

$i_temp = \text{PREV}(i_temp)$;