

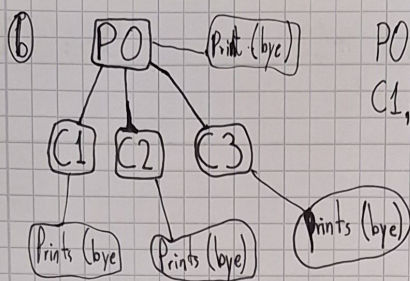
12.1 @ and ⑥ have been submitted as text files.

12.2 operating system processes

① Annotate the assembler instructions:

- `addi sp, sp, -16`: This subtracts 16 from the stack pointer `sp`, effectively allocating 16 bytes on the stack for the current function stack frame.
- `sd ra, 8(sp)`: Saves the return address `ra` to location `sp+8`.
- `sd 0, 0(sp)`: Saves frame pointer `0` to memory location `sp`.
- `addi s0, sp, 16`: Adds 16 to the stack pointer `sp` and stores result in `s0`.
- `jal fork`: Jumps to `fork` function and saves the return address in `ra`.
- `bne a0, zero, done`: This branches to the `done` label if `a0` \neq 0.
- `jal execvp`: Jumps to `execvp` function which replaces the current process image with a new one.
 (Note: `execvp` receives the return value of the `fork` as an argument.)
- `la a0, bye`: Loads address of the "bye" string into `a0`.
- `jal puts`: Jumps to `puts` function, prints the `a0` pointed string to the standard output.
- `mv a0, zero`: Moves the value 0 into `a0`.
- `ld ra, 8(sp)`: Loads the return address `ra` from `sp+8`.
- `ld s0, 0(sp)`: Loads the frame pointer `s0` from memory location at `sp`.
- `addi sp, sp, 16`: This instruction adds 16 to the stack pointer `sp`, deallocating the current function's stack frame.
- `ret`: Returns address stored in `ra`.





P0 - initial parent process that starts the program

C1, C2, C3 - Child processes created by the fork calls in P0

- ⑦ The program prints "bye" four times due to three forks creating eight processes. The "exec" call replaces these processes with the "date" command, preventing them from reaching the "puts" call. The remaining four processes, including the parent and the fork's children reach the "puts" call and print "bye".

Problem 12.3: Pre- and post conditions

a) After lines: ① $Z = X \cdot 2$

② $Y = Z + 2 = X \cdot 2 + 2$

③ $X = Y \cdot 3 = ((X \cdot 2) + 2) \cdot 3$

So the strongest post condition is $((X \cdot 2) + 2) \cdot 3 = X$ and $(2 \cdot X + 2) = Y$.

b) After lines: ① $X = X + 2$

② $Y = X \cdot 4 = (X + 2) \cdot 4$

③ $X = Y - 4 = ((X + 2) \cdot 4) - 4$

Plugging in the postcondition $(X > 10) \wedge (Y < 20)$ we can plugin $X = ((X + 2) \cdot 4) - 4$ which gives us $X > \frac{3}{2}$ and we do the same for $Y < 20$ which gives us $X < 3$ so the weakest pre conditions are $X > \frac{3}{2}$ and $X < 3$ aka. $1.5 < X < 3$.

c) For this we must work backwards to find it:

Given the post condition $7 \leq Y < 25$.

Starting from the last line: $Y := Y - 2$ we need $9 \leq Y < 27$

Now we go to the conditional statements.

Else:

$$Y := X + 6 \quad (x \geq 12)$$

we need: $3 \leq x < 21$

IF-Completed:

$$Y := 3 \cdot x - 9 \quad (\text{when } x < 12)$$

we need: $6 \leq x < 12$

Combine: $3 \leq x < 21$

Moving up we get the line: $X := 3 \cdot Y - 2$

To ensure: $3 \leq x < 21$

we need: $\boxed{5/3 \leq Y < 23/3}$

Thus the weakest precondition has been found to be $\boxed{5/3 \leq Y < 23/3}$

THE END