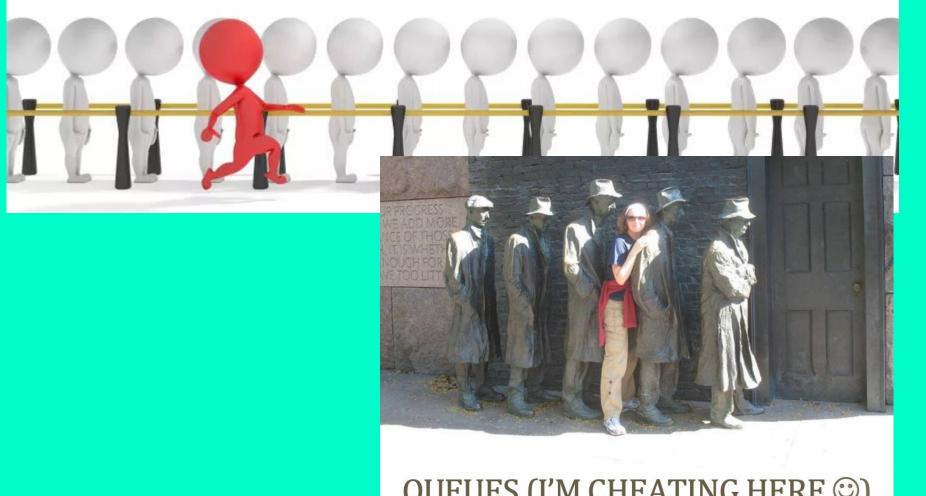


Lecture 6

MIC



QUEUES (I'M CHEATING HERE ©)

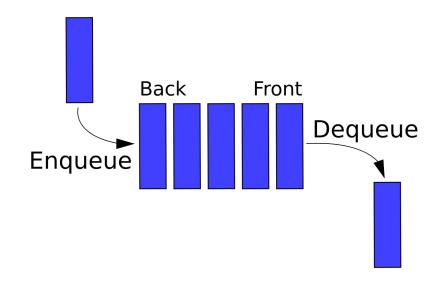
THE QUEUE ADT

A **Queue** is a collection of objects inserted and removed according to the

First In First Out (FIFO) principle. Think of a queue of people to Rubios

QUEUE OPERATIONS

Enqueue (insert) and Dequeue (remove) are the two main
operations





QUESTION

When using enqueue operation to place the following items in a queue:

enqueue(10)

enqueue(20)

enqueue (30)

enqueue(0)

enqueue (-30)

The output when dequeuing from the queue is:

A: 10, 20, 30, 0, -30

B: -30, 0, 10, 20, 30

C: 30, 10, 20, 0, -30

D: -30, 0, 30, 20, 10

E: 0, 30, -30, 10, 20

IMPLEMENTATION. ARRAYS. 0(1)

Main update methods:

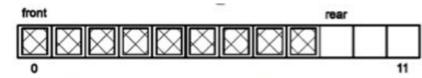
- Enqueue (e)
- Dequeue()

Additional useful methods

- Peek(): Same as dequeue, but does not remove the element
- Empty(): Boolean, True when the queue is empty
- Size(): Returns the size of the queue

(INCORRECT) ATTEMPTS TO IMPLEMENT IT

REGULAR ARRAY: DEQUEUE



(a) Queue.front is always at 0 - shift elements left on dequeue().

```
def dequeue():
    # potential issue if empty
    # for now, assume not empty

elem = array[front]
    # You code is here #
    return elem
```

Select the correct code to delete from below:

```
A: front = front + 1
```

for i in range(rear):
 array[i] = array[i+1]

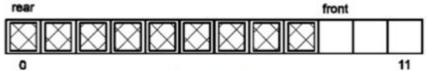
rear = rear - 1

```
B: rear = rear - 1
```

D: None of these are correct



REGULAR ARRAY: ENQUEUE



(b) Queue. rear is always at 0 - shift elements right on enqueue().

B: array[front] = elem

D: None of these are correct

def enqueue(elem):

potential issue if full

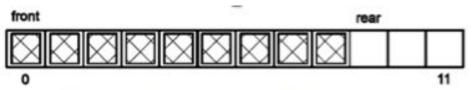
Your code is here

front = front + 1

for now, assume not empty

ISSUES

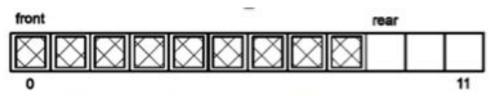
Dequeue:



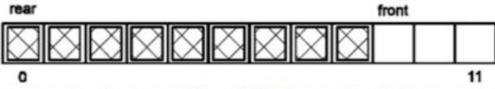
(a) Queue.front is always at 0 - shift elements left on dequeue().

ISSUES

Dequeue:



(a) Queue.front is always at 0 - shift elements left on dequeue().

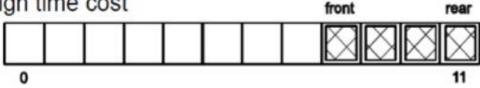


Enqueue:

(b) Queue. rear is always at 0 - shift elements right on enqueue().

REGULAR ARRAY

 Neither of those solutions is very good as they both involve moving all the existing data elements, which has high time cost

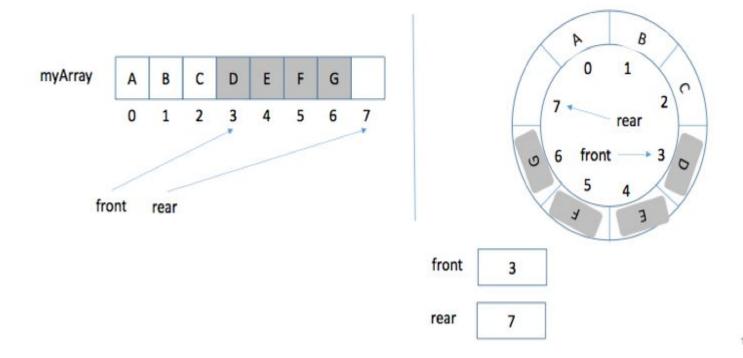


 Idea: Instead of moving data elements to a <u>fixed</u> position for *front* when removing, let *front* advance through the array

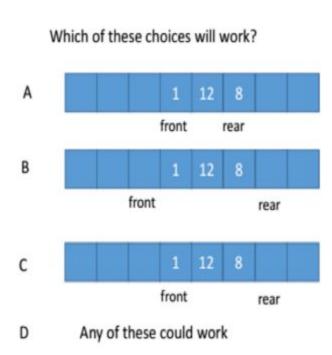
Hmmm....what do we do when we now add an element to that queue at the rear? What happens when we remove several elements, and *front* catches up with *rear*?...

Making a linear array appear circular

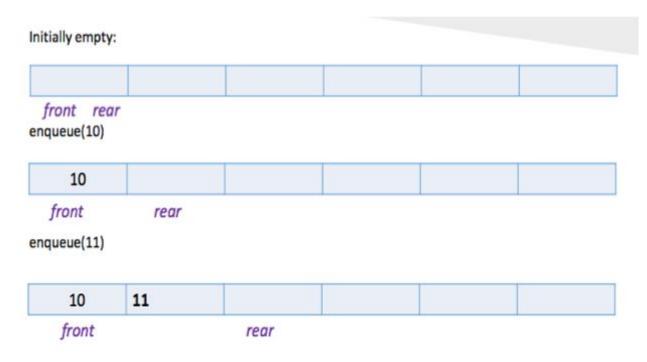
front==head rear == tail



Design decisions: Where do front and rear point?







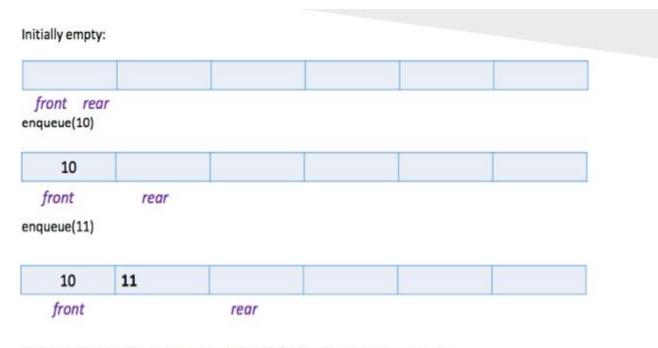
What should be the value of front after the next dequeue?

A. 0

B. 1

C. 2

D. 5



What should be the value stored at arr[0] after the next dequeue?

A. 10

B. 0

C. null

D. It doesn't matter



What is the value of rear after this enqueue?

- A. 5
- B. 0
- C. 1
- D. 2
- E. Other

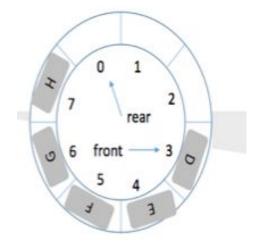
```
0 1
7 rear
0 6 front 3 0
5 4
```

```
def dequeue():
    size = size - 1
    elem = array[front]
    # Your code #
    return elem
```

Select the correct code to insert from below:

B: rear = rear - 1 if (rear < 0): rear = len(array) - 1

D: None of these are correct



rear = rear + 1

```
def enqueue(elem):
    #Your code is here#
    size = size + 1
```

Select the correct code to insert from below:

```
rear = rear + 1
                                                        for i in range(rear):
A:
       if (rear == len(array)):
                                                            array[i] = array[i+1]
            rear = 0
                                                        array[rear] = elem
        array[rear] = elem
                                                        front = front - 1
        array[rear] = elem
B:
                                                  D: None of these are correct
```

COMPLEXITY OF AN ARRAY BASED QUEUES

Method	Running Time
size	O(1)
isEmpty	O(1)
first	O(1)
enqueue	O(1)
dequeue	O(1)