74.343 Operating Systems

Final Examination

April 26th, 2002 – 9:00AM University Centre 220-344

Paper # 630 Instructors: Eskicioglu and Graham

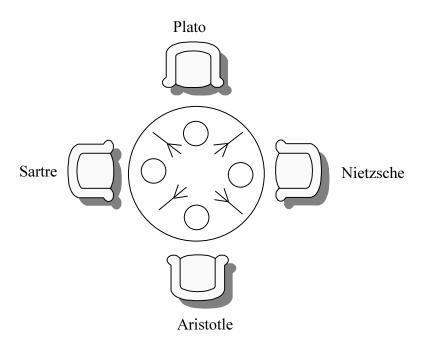
Instructions: Time Allotted - 2 hours.

No aids permitted (or necessary)

Total Marks - 65.

Name:					
Student N	lumbe	r:			
			Question	Value	Score
			1	10	
			2	8	
			3	10	
			4	6	
			5	8	
			6	6	
			7	8	
			8	9	
				Total:	
10] 1. I	Definiti	ions: (Please define <i>concisely</i> !)			
	(i)	Counting Semaphore			
	(ii)	Virtual memory			
					
	(iii)	Interrupt Handler			
	<i>(</i> ')	T1 1:			
	(iv)	Thrashing			
	(17)	Non blocking Possiva			
	(v)	Non-blocking Receive			
	_				

[8] 2. The "dining philosophers" problem is a classic synchronization problem. Each philosopher has his/her own plate of spaghetti but the forks are shared. There is one fork between each pair of philosophers and a philosopher must have both forks to be able to eat. Each "philosopher" switches between thinking and eating at unpredictable times. When a philosopher is thinking, no forks are required but when a philosopher wants to eat she/he must acquire both forks (the one to the left and the one to the right) before beginning. Once a philosopher has finished eating he/she returns both forks to the table. Pseudo code a solution to the dining philosophers problem using binary semaphores to guard each fork. Specifically, you should pseudo code the routine called "Eat(int p)" which will be executed by each philosopher with the parameter 'p' specifying the number of the philosopher invoking the "Eat()" routine. Assume there are N philosophers. Be sure to show your declarations (Hint: use an array of semaphores for the forks). Watch out for deadlocks! A picture of the dining philosophers scenario is provided below for reference:





[10] 3. Consider a paged memory management system that uses the clock algorithm for replacement. There are 5 frames in the memory and the pages initially loaded into them are shown in the first row of the following table. For each page frame there is also a reference bit, as required by the clock algorithm (0 means unreferenced, 1 means referenced). The column labeled "Clock" indicates the number of the frame that the clock pointer points to (frame 3 is where the "clock" points initially). Given the sequence of page references shown in the first column, apply the clock algorithm to complete the table. Be sure to indicate when the faults occur by placing a "Yes" in the rightmost column.

Reference	Clock	Frame	1	Frame	2	Frame	3	Frame	4	Frame	5	Fault?
Initial	3	Page 5	1	Page 7	0	Page 3	1	Page 12	0	Page 8	0	N/A
Page 7												
Page 5												
Page 4												
Page 10												
Page 12												
Page 2												
Page 12												
Page 9												
Page 17												
Page 9												

		Page 17												
		Page 9												
[6]	4.	P ₁ _Alloc(1	sequence R _b), P ₃ _A	e: P ₂ _A lloc(R _c) h	lloc(as al	R _b), P _{2_} ready been	Allo n co	$c(R_c)$, P_3 mpleted (a	_All and a	5 resources oc(R _a), Practices assuming all AG) that re-	2_Al l res	loc(R _e), ources we	P ₁ _A ere fro	Alloc(R _d), ee before
	L													
		Is the syste	em deadlo	cked? Wh	y or	why not?								

is the system deadlocked. Why of why not:

[8]	5.	Given a hard disk with 1024 cylinders (C=[01023]), 4 heads (H=[03]), and 64 sectors (S=[063]),
		fill in the table to show the order in which the sequence of disk I/O requests shown below would be
		serviced under the NEAREST-FIRST (i.e. seek next to the nearest disk location) algorithm. You
		should assume that only the next three requests are available while processing each preceding request.
		Assume also that the disk read/write heads are originally positioned at cylinder 128 and that all heads
		move in unison. The requests are in the format: <cylinder, head,="" sector="">.</cylinder,>

<10,3,42> <200,2,18> <118,1,60> <453,1,6> <1001,2,8> <621,0,0> <8,3,8> <29,2,6>

Current (Cylinder	Visible Requests	Selected Request	# Cylinders Crossed
128				
				$\Sigma =$
poi	ked List Allo	ext.	ter to the first data block an	
2. Ind	ked List Allonter to the neexed Allocatile system: ow many bloom	ocation: There is a poin ext. ion: There is a two leve	ter to the first data block and lindex where each block portequired to read blocks 500	oints to 100 other blocks.
point 2. Index For each f	ked List Allonter to the neexed Allocatile system: ow many bloom	ocation: There is a poin ext. ion: There is a two leve	ter to the first data block and lindex where each block portequired to read blocks 500	oints to 100 other blocks.
point 2. Index For each f	ked List Allonter to the neexed Allocatile system: ow many bloom	ocation: There is a poin ext. ion: There is a two leve	ter to the first data block and lindex where each block portequired to read blocks 500	oints to 100 other blocks.
poin 2. Indexed List:	ked List Allonter to the neexed Allocatile system: ow many bloom	ocation: There is a poin ext. ion: There is a two leve	ter to the first data block and lindex where each block portequired to read blocks 500	oints to 100 other blocks.
poin 2. Indexed List:	ked List Allonter to the neexed Allocatile system: ow many bloom	ocation: There is a poin ext. ion: There is a two leve	ter to the first data block and lindex where each block portequired to read blocks 500	oints to 100 other blocks.
poin 2. Indexed List:	ked List Allonter to the neexed Allocatile system: ow many bloom	ocation: There is a poin ext. ion: There is a two leve	ter to the first data block and lindex where each block portequired to read blocks 500	oints to 100 other blocks.
poin 2. Indexed List:	ked List Allonter to the neexed Allocatile system: ow many bloom	ocation: There is a poin ext. ion: There is a two leve	ter to the first data block and lindex where each block portequired to read blocks 500	oints to 100 other blocks.
poin 2. Indexed for each for e	ked List Allonter to the neexed Allocate ile system: fow many bloom many blo	ocation: There is a point ext. ion: There is a two level ock read operations are ock read operations are ssume only one request	required to read blocks 10 arrives and is processed at a	oints to 100 other blocks. through 600? Explain!
poin 2. Indexed for each for e	ked List Allonter to the neexed Allocate ile system: fow many bloom many blo	ocation: There is a poin ext. ion: There is a two level ock read operations are ock read operations are	required to read blocks 10 arrives and is processed at a	oints to 100 other blocks. through 600? Explain!
poin 2. Indexed for each for e	ked List Allonter to the neexed Allocate ile system: fow many bloom many blo	ocation: There is a point ext. ion: There is a two level ock read operations are ock read operations are ssume only one request	required to read blocks 10 arrives and is processed at a	oints to 100 other blocks. through 600? Explain!

Indexed: ____

[6]

[8] 7. Consider the following system snapshot using the data structures in the Banker's algorithm, with resources A, B, C, and D, and processes P_0 to P_4 .

		Allocation			Maximum				Available				Need			
	A	В	C	D	A	В	C	D	A	В	C	D	A	В	C	D
									3	3	1	1				
\mathbf{P}_0	2	0	0	1	6	3	1	2								
\mathbf{P}_1	1	0	0	0	1	7	5	0								
P ₂	2	2	5	3	2	3	5	6								
P ₃	0	5	3	2	1	6	4	2								
P ₄	0	1	1	5	1	6	5	6								

i.	How many resources of each type (A, B, C, D) are there? Explain how you know this.	- \
		- - - -
		- - -
ii.	Fill in the current "Need" matrix (above)	
iii.	Is the system currently in a safe state? Why or why not?	_
		-
		- -
		- -
		-
iv.	If a request from P0 arrives for additional resources of [1, 3, 1, 0] can the resources granted? Explain your answer by reference to the system state.	be
		- - -
		-
		-
		- - -
		-
		- - -
		-

	gies fo	or the following			for the Round Robin and Priority sche numbers mean higher priority and that th
	<u>ob</u>	Arrival Time	<u>Duration</u>	<u>Priority</u>	
A E		0 2	6 4	2 5	Taka wata aftha assisal times afiab
(3	2	4	Take note of the arrival times of jobs
Ι)	7	3	1	
	C D				
A	verag	ge wait time for	Round Rol	Time =	→
	-	•			
	A				
	В				
RIORITY	_				
	С				