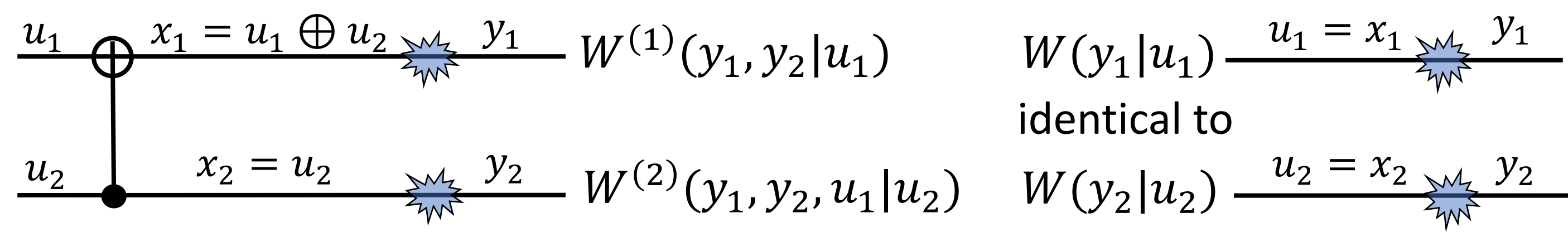


# Improved Logical Error Rate via List Decoding of Quantum Polar Codes

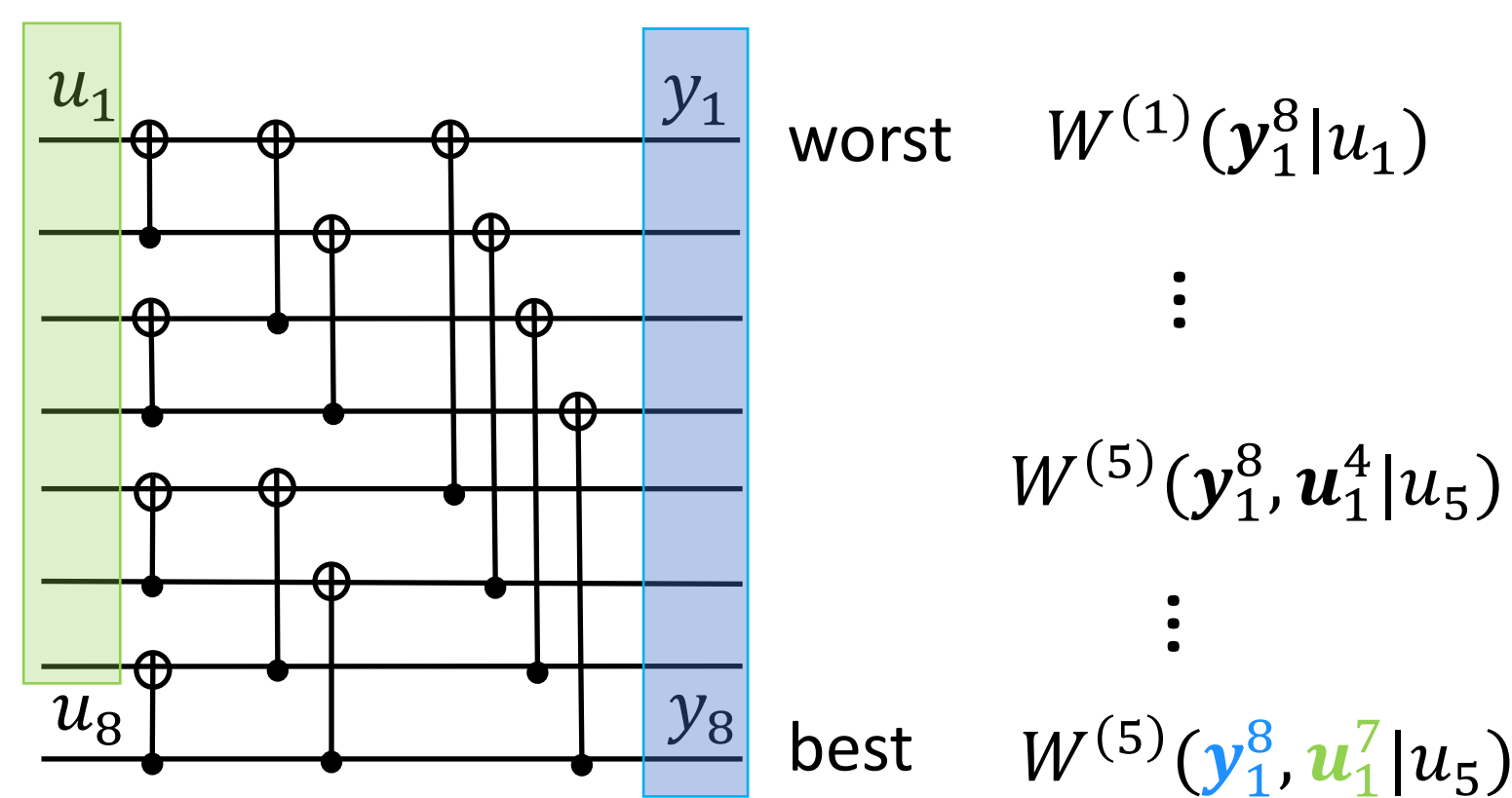
Anqi Gong, Joseph M. Renes Institute for Theoretical Physics, ETH Zürich, 8093 Zürich, Switzerland

## Polar Codes and Decoding [1,2]



**Channel Polarization:**  $I(W^{(1)}) < I(W) < I(W^{(2)})$

**Successive Cancellation Decoding**



For  $i$  in  $1, \dots, N$  do:

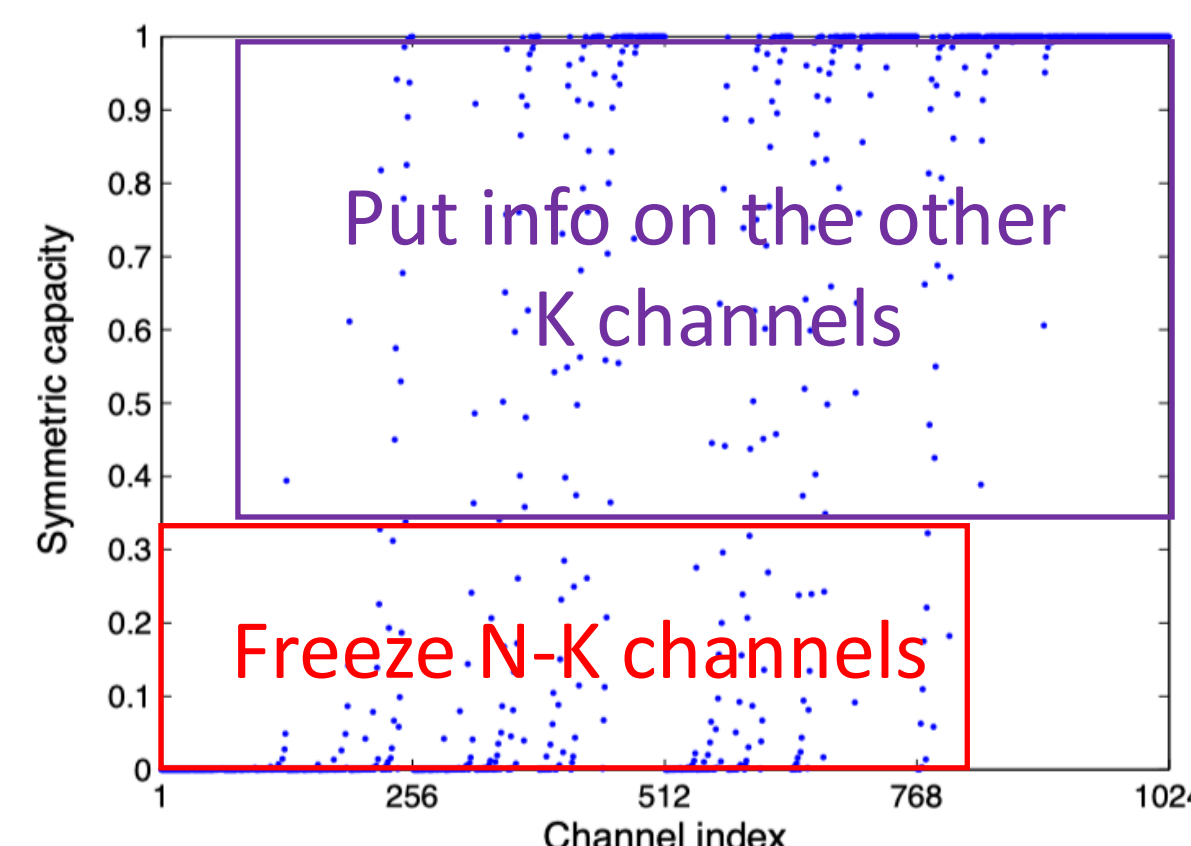
if  $i \in \mathcal{A}^c$ : /\*  $u_i$  is a frozen bit \*/

$\hat{u}_i \leftarrow u_i$

else: /\*  $i \in \mathcal{A}$ ,  $u_i$  is an info bit \*/

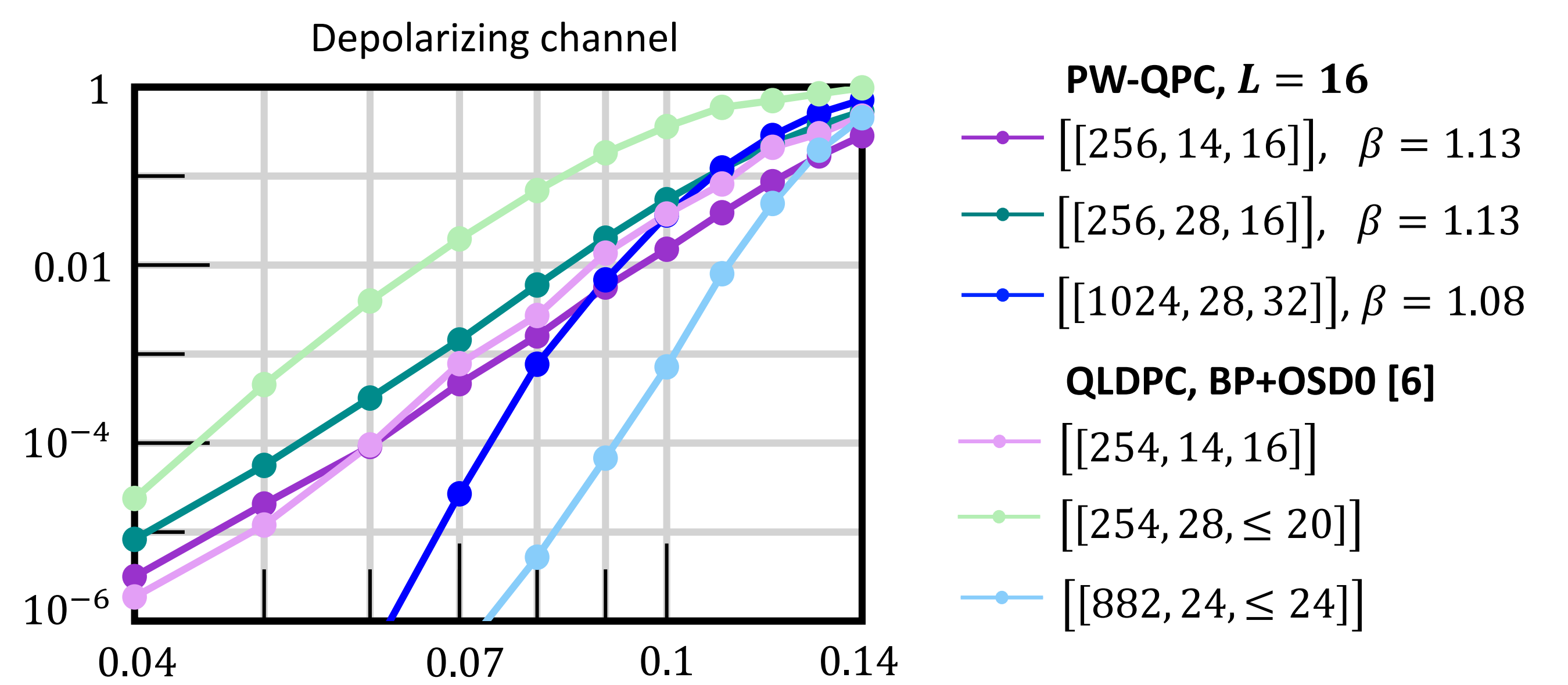
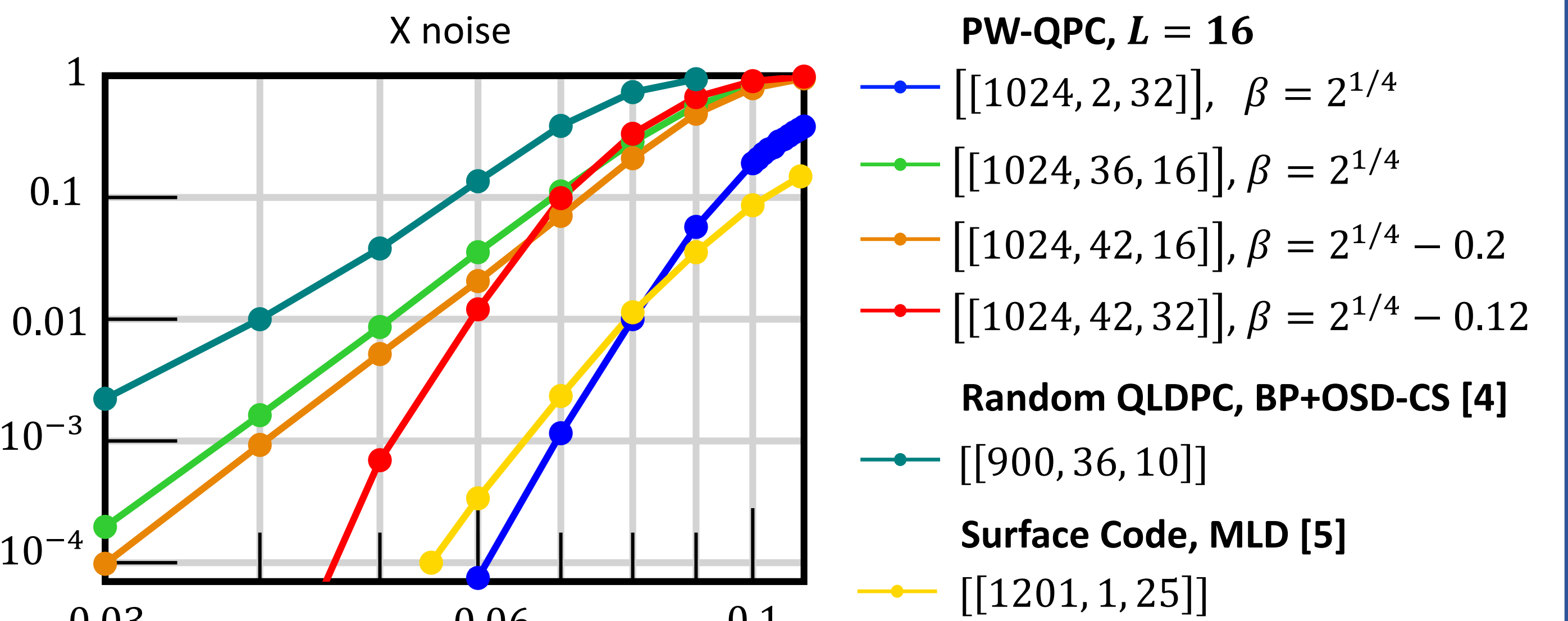
$\hat{u}_i \leftarrow \arg\max_{u_i \in \{0,1\}} W^{(i)}(\mathbf{y}, \hat{\mathbf{u}}_1^{i-1} | u_i)$

$\mathcal{O}(N \log N)$  Time  
 $\mathcal{O}(N)$  Space

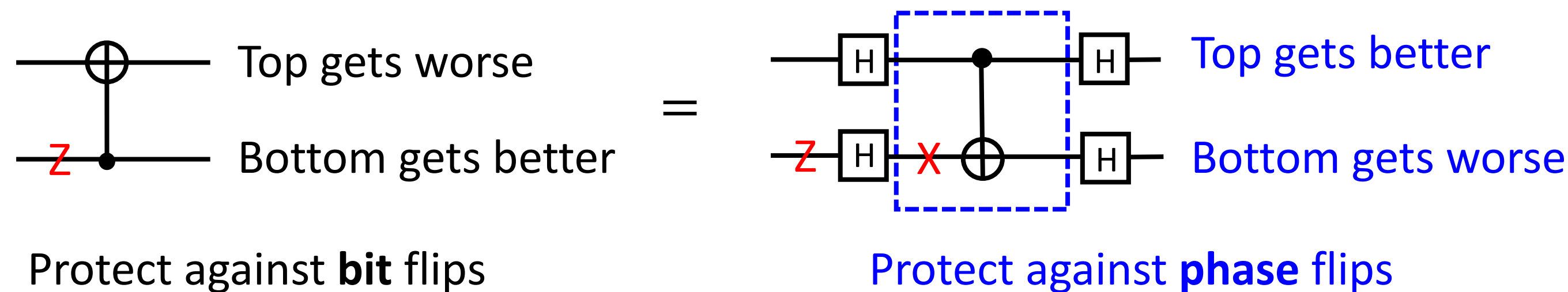


List decoding [2] always keeps a list of  $L$  decisions  $\hat{\mathbf{u}}_1^i$ , and their path metrics  $\text{PM}(\hat{\mathbf{u}}_1^i) = \Pr[\hat{\mathbf{u}}_1^i | \mathbf{y}]$  at each step.

## Comparison to other Codes



## Polarization Weight Construction



CSS constraint: no qubit is simultaneously frozen in both X and Z basis

## Polarization Weight (PW) Construction

Want to define Polarization Weight of row 6

$$6 = 110_2 = 1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$$

$$\text{PW}(6) = 110_\beta = 1 \times \beta^2 + 1 \times \beta^1 + 0 \times \beta^0$$

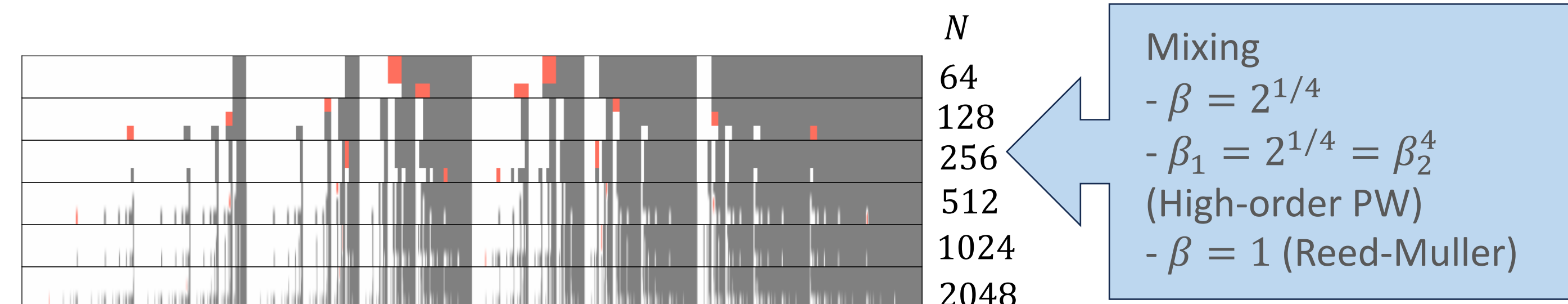
Satisfies CSS constraint.  
No entanglement assistance needed.  
Gives good codes.

Ex: want an  $[[8, 2]]$  quantum code Against bit flips Against phase flips

	Binary	PW $\beta = 2^{1/4}$	PW $\beta = 2^{1/4}$
$ 0\rangle$ Row 0	000	0	3.603
$ 0\rangle$ Row 1	001	1	2.603
$ 0\rangle$ Row 2	010	1.189	2.414
$ 0\rangle$ Row 3	011	2.189	1.414
$ 0\rangle$ Row 4	100	1.414	2.189
$ +\rangle$ Row 5	101	2.414	1.189
$ +\rangle$ Row 6	110	2.603	1
$ +\rangle$ Row 7	111	3.603	0

freeze 3 worst

freeze 3 worst

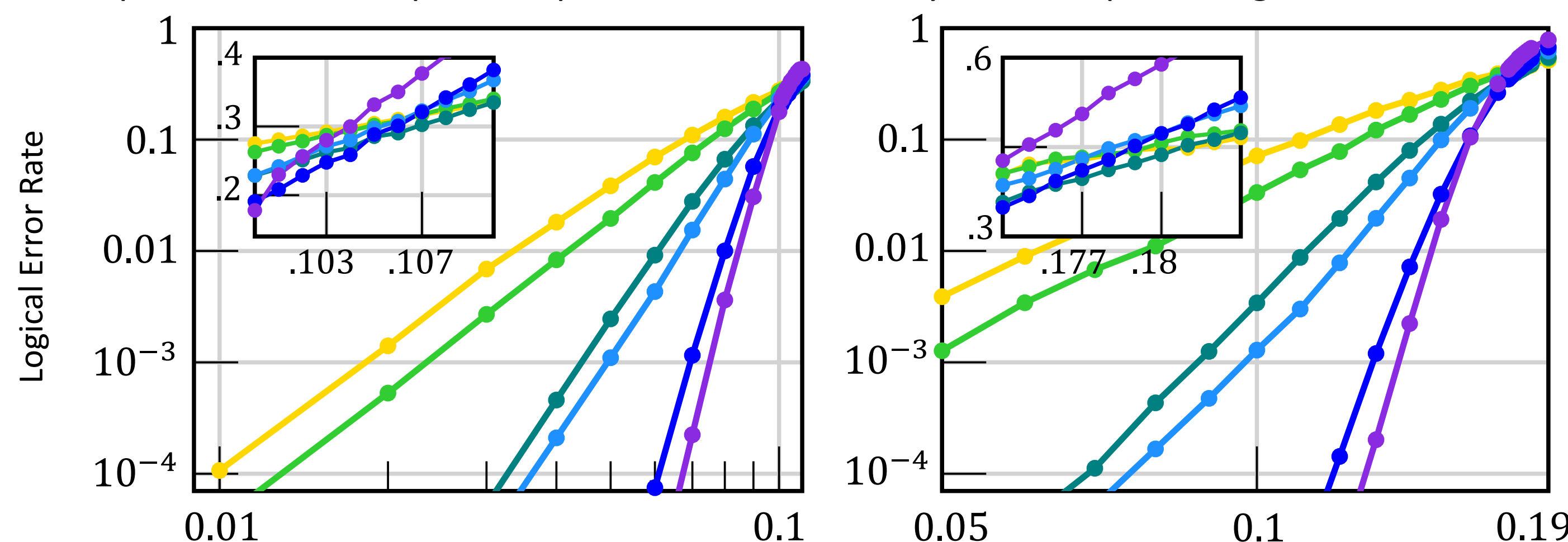


A family of  $[[N, 2]]$  PW-QPC ( $\beta = 2^{1/4}$ ) using list decoder of size  $L$

$N, L$	64, 4	256, 8	1024, 16
	128, 8	512, 16	2048, 32

Independent bit- and phase-flip channel, X noise only

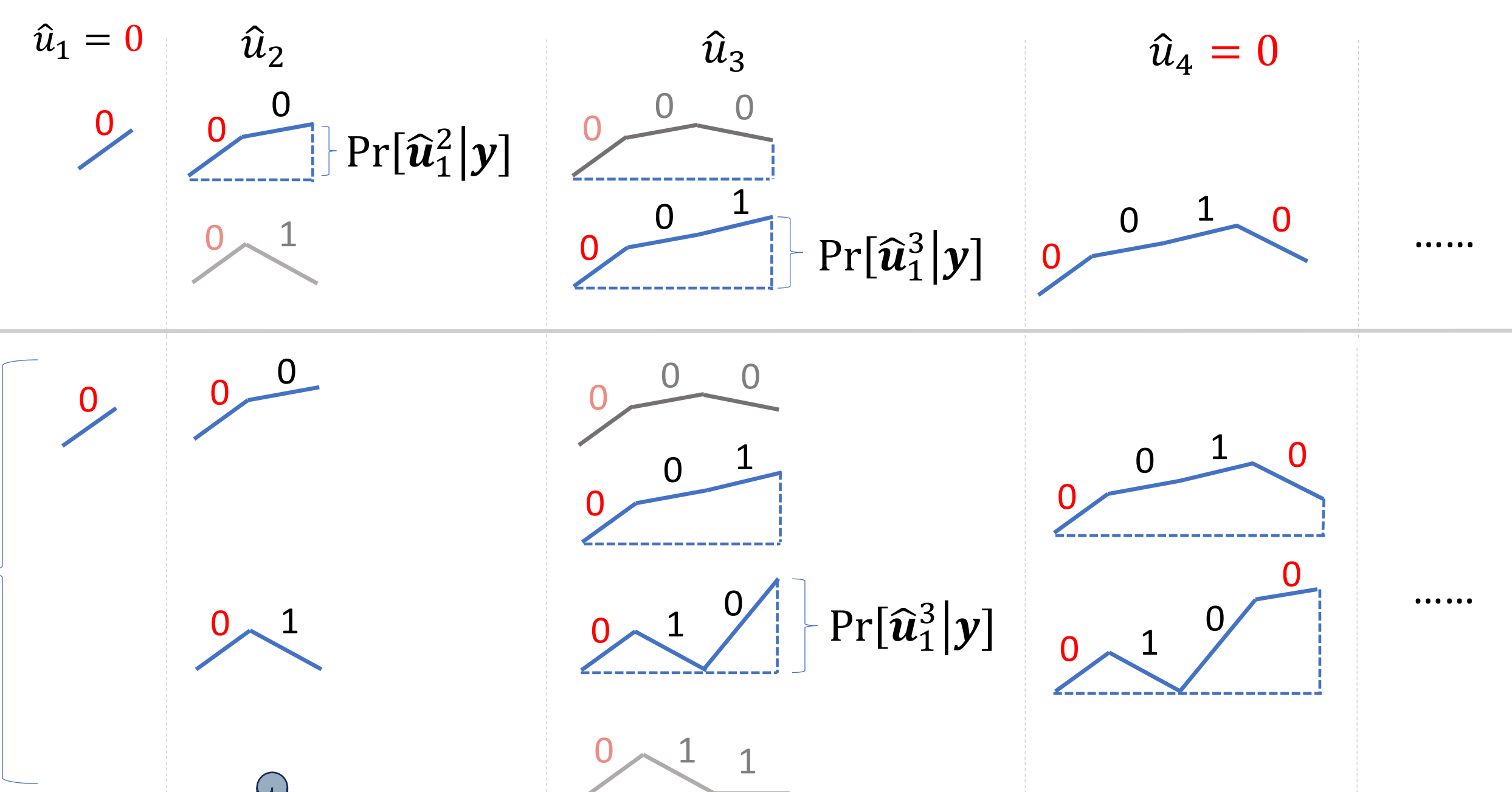
Depolarizing channel



## List Decoding and Degeneracy

Following a local minimum each time does not necessarily lead to global minimum.

Successive Cancellation Decoding ( $L = 1$ )



Time  $\mathcal{O}(LN \log N)$   
Space  $\mathcal{O}(LN)$

Conversion  
- Put syndrome  $\mathbf{s}$  at  $\mathbf{u}_{\mathcal{A}^c}$  (both size  $N - K$ )  
- Treat all-zero as the noisy codeword  
- Decode  $\mathbf{u}_{\mathcal{A}}$

Codeword list decoder

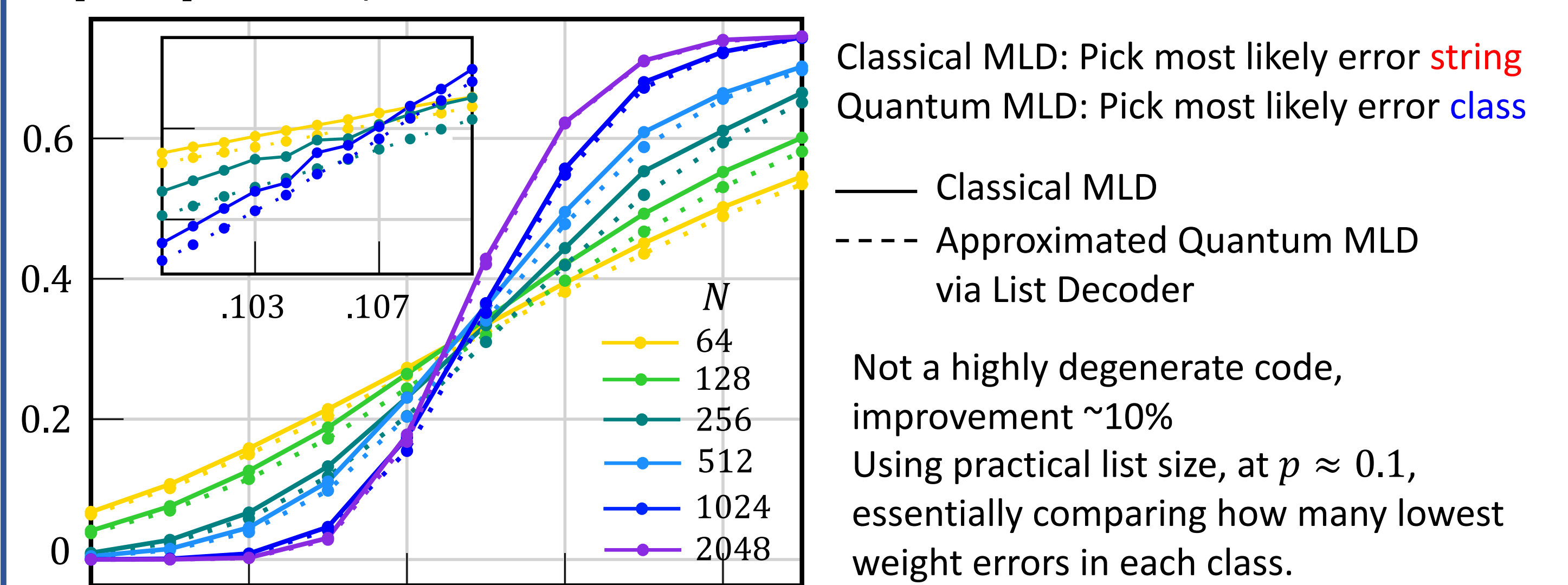
Syndrome list decoder

**What's on the list?** (with high probability)

- Classical:  $\hat{\mathbf{u}}_1^i$ 's such that  $\Pr[\hat{\mathbf{u}}_1^i | \mathbf{y}]$  are the largest

- Quantum: Given syndrome  $\mathbf{s}$ , most likely error strings compatible with  $\mathbf{s}$

$[[N, 2]]$  PW-QPC ( $\beta = 2^{1/4}$ ),  $L = 128$ , X noise



Not a highly degenerate code, improvement  $\sim 10\%$   
Using practical list size, at  $p \approx 0.1$ , essentially comparing how many lowest weight errors in each class.

## References

- [1] E. Arkan, "Channel Polarization: A Method for Constructing Capacity-Achieving Codes for Symmetric Binary-Input Memoryless Channels", IEEE Transactions on Information Theory, 2009.
- [2] I. Tal and A. Vardy, "List Decoding of Polar Codes", IEEE Transactions on Information Theory, 2015.
- [3] G. He et al, "Beta-Expansion: A Theoretical Framework for Fast and Recursive Construction of Polar Codes", IEEE Global Communications Conference, 2017.
- [4] J. Roffe et al, "Decoding across the quantum low-density parity-check code landscape", Physical Review Research, 2020.
- [5] S. Bravyi, M. Suchara, and A. Vargo, "Efficient algorithms for maximum likelihood decoding in the surface code", Physical Review A, 2014.
- [6] P. Panteleev and G. Kalachev, "Degenerate Quantum LDPC Codes With Good Finite Length Performance", Quantum, 2021.