

**National University of Singapore**

**School of Computing**

**CS3202: Software Engineering Project II**

**TEAM 05: Flying Cockroach**

Semester 1, AY2014/2015

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Consultation Day/Hour: Monday 6-6.30pm

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# SPA

Static Program Analyser (SPA) is a program to answer queries about an input SIMPLE program. In this report, we will be describing the design and implementation decisions made during the development of the SPA during CS3201 and CS3202.

## Architecture

The architecture for the prototype consists of 3 main components: the Code Parser, the PKB and the Query Processor. Both the Code Parser and the Query Processor are dependent on PKB but not dependent on each other. Code Parser parses the code and stores design abstractions in each of the 8 tables in the PKB. After Query Parser has parsed the query, the Query Evaluator consults the PKB API to answer queries.



Figure 1

## Interaction

/\* **DINDA**

PLEASE REMOVE THIS COMMENT

Please change the diagrams as necessary to factor in the changes we made in CS3202.

Look at pg 67 of handbook.

Also, include more UML diagrams that were used in project planning or testing plan!

Draw UML diagrams that you found useful. For each diagram that you draw, explain how you used it (e.g., in project planning, communication, test planning or in other situations), and comment on the value a diagram added to your project.

\*/

CodeParser works by evaluating each line of the given source code. It creates AST Node, set the pointers accordingly; set the tables and the appropriate databases in PKB.

The attributes in PKB (the tables) will then be used by Query evaluator to answer queries. Testing for CodeParser is done by checking the content of each table, whether it has set the values properly, and check the content of each node in the AST, whether it matches the correct AST.



Figure 2

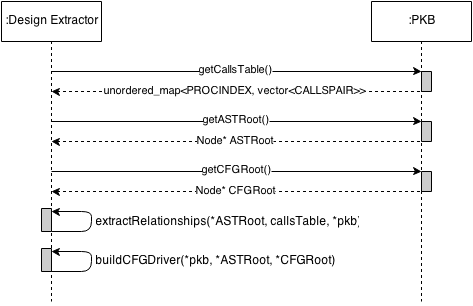


Figure 3 shows the sequence diagram of query evaluation process. This diagram was useful in demarcating the responsibilities of each PQL group member. For example, QueryEvaluator directly assumes that the Query it receives is valid and syntactically correct. Therefore it is the responsibility of QueryParser to validate each query before passing it to the evaluator.



Figure 3

This diagram also helps to keep track of the dependencies between components. This is especially useful during debugging process of integration testing. When QueryProcessor fails to return the correct result, the team knows that the errors could come from at least three places, i.e. QueryParser, QueryEvaluator, and PKB.

# Summary of Achievements

## Basic SPA

All the required functionalities for Iteration 1-3 were implemented. These include:

1. All the grammar rules for the SIMPLE source program as outlined on Page 103 of the Handbook.
2. All the grammar rules for the PQL as described on Page 98-101 of the Handbook.
3. All the relationships are extracted and can be queried using PQL. These include:

|  |  |
| --- | --- |
| **Relationship** | **Implemented using** |
| Calls and Calls\* | CodeParser, PKB, Query Evaluator |
| Modifies | CodeParser, Design Extractor, PKB, Query Evaluator |
| Uses | CodeParser, Design Extractor, PKB, Query Evaluator |
| Parent and Parent\* | CodeParser, PKB, Query Evaluator |
| Follows and Follows\* | CodeParser, PKB, Query Evaluator |
| Next and Next\* | DesignExtractor, PKB, Query Evaluator |
| Affects and Affects\* | PKB, Query Evaluator |

## Bonus Features

We have extended the features of SPA beyond its call of duty. The features will be explained in more detail below:

* + 1. Flexible CodeParser

According to the handbook, the source code that is to be tested against our SPA, is defined to be in standardized format and neatly arranged. They have regular and consistent spacing, indentation, tabs, and endline characters. Please refer to the example below:

procedure Orchid {

while idx {

y = z\*3 + 2\*x;

call Tulip;

idx = idx - 1; }

z = z + y + idx; }

procedure Lily {

while v {

y = x + y; }

x = y + x; }

However, our SPA source code parser is very flexible. It is able to detect erroneous and inconsistent spaces, tabs, endline characters and erroneous close curly brackets. A flexible CodeParser is, thus, one of our bonus features, as it was not required for us to have in the basic SPA. To demonstrate this ability, please refer to the example below (which will have the same extracted information as the code fragment above):

procedure Orchid

{

while idx {

y = z \*3 + 2\* x ;

call Tulip ;

idx = idx -1;

}

z=z+y+idx;

}

procedure Lily{

while v{

y =x + y ;}

x = y + x;}

### Flexible QueryParser

The queries that are to be tested against basic SPA are defined to be in a standard format. They have regular and consistent lowercase/uppercase command, spacing, and characters. Please refer to the sample standard query below:

if ifstat; Select ifstat such that Follows\* (ifstat, 17)

assign a; Select a such that Modifies (a, "idx") and Uses (a, "idx")

assign a; while w; Select a such that Modifies (a, "idx") and Uses (a, "idx") and Follows (15, a) and Parent\* (w, a)

However, our QueryParser is very flexible and is able to detect queries with inconsistent/extra spaces and lowercase/uppercase clauses. As it was not required for us to have such flexibility for CS3202’s basic SPA, we have this advanced feature. To demonstrate this ability, please refer to the example below (which will have the same information/result as the example above):

If ifstat ; Select ifstat such that Follows\* ( ifstat,17)

assign a;Select a such that Modifies(a,"idx") and Uses(a,"idx")

assign a ; while w ; Select a such that Modifies (a, "idx" ) and Uses ( a , "idx") and Follows (15, a) and Parent\* ( w , a )

### Highly Organised Repository

In our shared repository, our code, report files and directories are very well organised. We do this by following standard naming conventions and following a structured hierarchy such that each team member has fast and easy access.



Appropriate milestones and issues are also tracked and reported regularly to achieve any goals or/and objectives.





Every commit related to certain issues or milestones was tagged and linked to the issue for easy tracking and reference in the future. It was also used to monitor the progress of any issues and the SPA system as a whole. Github features such as milestones, issue tracker with assignees, milestones, and issue category (bugs, testing, documentation, enhancement, etc) were highly useful for our project management process.





Overall, our code and API are well-documented and in sync with each other. Every detail of the code is described clearly, showing and explaining all methods, attributes, and inheritance diagram whenever applicable.



Overall, our progress in building this project was conscientiously and thoroughly tracked and documented so that any new addition to the team could understand the project. Also, by posting and tracking issues, we can determine who tackled what.

Our project is currently hosted on GitHub, and the relevant links are as follows:

* Repository : <https://github.com/yulonglong/Static-Code-Analyzer>
* Issue tracker : <https://github.com/yulonglong/Static-Code-Analyzer/issues>
* Milestones : <https://github.com/yulonglong/Static-Code-Analyzer/milestones>
* Commits : <https://github.com/yulonglong/Static-Code-Analyzer/commits/master>

# Project Plan

## 3.1. Schedule For Whole Project

|  |  |  |  |
| --- | --- | --- | --- |
| Iteration | Major Task | Target Completion | Completed |
| 0 | Change query parser to use regular expression checking | Week 3 | Week 2 |
| 0 | Improve internal PKB data structures | Week 3 | Week 2 |
| 1 | Implement Calls, Modifies, and Uses in PKB | Week 5 |  |
| 1 | Extend Query Parser to support with clause | Week 5 |  |
| 1 | Extend Query Evaluator to support with clause | Week 5 |  |
| 2 | Implement Next in PKB | Week 7 |  |
| 2 | Extend Query Parser to support multiple clauses | Week 7 |  |
| 2 | Extend Query Evaluator to support multiple clauses | Week 7 |  |
| 3 | Implement Affects in PKB | Week 10 |  |
| 3 | Extend Query Parser to support Affects and tuple results | Week 10 |  |
| 3 | Extend Query Evaluator to support Affects and tuple result | Week 10 |  |
| 4 | Implement Sibling in PKB | Week 12 |  |
| 4 | Extend Query Parser to support Sibling | Week 12 |  |
| 4 | Extend Query Evaluator to support Sibling | Week 12 |  |

## 3.2. Comments & Problems

1. Encountering difficulties when inter modules

API and .h in advance compilation error

1. Very hard to trace errors whenever we have wrong answer

debug mode by checkpoints state of program, which module is giving a bug

1. coordinate

Miscommunication. Frequent discussion, testing and informing

# 4. Components

## 4.1. Code Parser

Code Parser’s main functions are to read in the source code, build the AST, and set the tables (VarTable, ProcTable, TypeTable, Follows, Parent, Modifies, Uses) in PKB according to the input source code.

To build the AST, Code Parser depends on the implementation of node, which is the node structure being used to build AST. A node can have many children, thus we decided to use NODE\_PTR\_LIST (C++ vector) for dynamic storage of the children pointers, rather than array with fixed size.

Code Parser keeps track of every relevant parent of each node being built, by storing pointers to their parents. For example when there is a while statement, a pointer to the while statement will be stored and then link all the children to their parent node.

Code Parser works by tokenizing the source code as string line by line, detecting the tokens in each line, generating the types, program lines, and setting the tables accordingly.

For an assignment statement, Code Parser will convert the expression from infix to postfix, and then create the expression tree.

Code Parser does its validation by keeping track of the curly brackets (i.e. “{“ and “}” ). It keeps track of the number currently present open curly bracket, “{“. When Code Parser encounters an open curly bracket, it will push it to a stack. When it encounters a closed curly bracket, it will pop from the stack.

When Code Parser reaches the end of the source code, it will return invalid if the stack is not empty, or if Code Parser is trying to pop from an empty stack. It means there is a mismatch in the number of curly brackets.

When the stack is empty, Code Parser will accept a line which defines a procedure. If it encounters any other statement while the stack is empty, it will return invalid.

Example:

Source code (source1.txt):

procedure Mini {

A1 = 29;

a1 = 31;

i = 51; }

The following describes how the Code Parser parses the above SIMPLE program:

1. Code Parser starts reading the source code at line 1, it will check whether the stack is empty. If the stack is empty, it will be expecting a procedure declaration.
2. It then parses procedure Mini, creates an ASTNode, sets it as root, and pushes the curly bracket “{“ into the stack. Insert “Mini” into the ProcTable.
3. At line 2, Code Parser tokenizes and checks the type of statement. Since it starts with a variable, it detects the statement as an assignment statement.
4. Code Parser will check if there exists a semicolon at the end of the line because it is compulsory to have a semi colon at the end of an assignment statement. If it exists, it will create a node containing “=”, and link “A1” as the first child. The expression on the right hand side will be converted into a postfix expression, and then build the expression tree.
5. Code Parser will link the root of the expression tree as the second child of “=”, and then link the “=” to its parent, which in this case, is “procedure Mini”.
6. Code Parser will also set the VarTable, Modifies, Uses, Follows, and Parent accordingly, in this case, it will set Modifies (line 1, and variable A1), and insert A1 into the VarTable.
7. At line 3, it detects that it is an assignment statement, and repeat step 4 to 6.
8. At line 4, it detects that it is an assignment statement, and repeat step 4 to 6. In addition, Code Parser detects a closed curly bracket. Therefore, Code Parser will pop the stack.
9. End of source code is reached. It will now check whether the stack is empty.
10. Since the stack is empty, and there is no violation of the rules stated earlier. Code Parser has built the AST successfully and stored the design abstractions in the relevant tables.

CodeParser’s Unit testing is done by checking the content of each table, whether it has set the tables properly, and by checking the contents of each node in the AST, whether it matches the expected AST.

## 4.2. PKB

/\***YOLIM**

PLEASE REMOVE THIS COMMENT

Follow guidelines in Handbook Section 10.2 to analyze, justify and document detailed design decisions. Pay attention to clarity of the description (check hints in Section 10.2).

If you applied design patterns, document them in this section:

a) Explain the design problem and pattern you applied

b) Document expected benefits and costs of applying a design pattern

c) Document the actual benefits and costs of a design pattern that you experienced in the project after applying it.

\*/

## 4.3. Design Extractor

The main role of the design extractor is to extract relationships about the SIMPLE program that could not be extracted in the one-time parsing done by the Code Parser. This includes:

* 1. Extracting information about Modifies and Uses for procedures and for program lines that are calls statements
  2. Building the Control Flow Graph (CFG) from the AST and subsequently storing it in the PKB and storing the Next relationship in the PKB

### 4.3.1. Extracting Relationships

The following shows the steps required to extract relationships like Modifies and Uses for multiple procedures in a single SIMPLE program:

1. Obtain the Calls Table from the PKB. An example of the calls Table is shown below where procedure with index 0 calls procedure with index 1 from program line 3 and procedure with index 4 calls procedure with index 3 from program line 21.

0 🡪 (1, 3) (4, 5)  
1 🡪 (2, 8) (3, 10)  
2 🡪  
3 🡪  
4 🡪 (3, 21)

This translates into a Calls Tree as follows where the nodes represent procedure indices and the edges represent the program line at which there is a calls statement.



1. Run Depth First Search (DFS) on the Calls table (Note: There can be cycles in the graph) in order to obtain a topological sort order of the procedure indices in a queue (below) which starts on the left. Each Queue contains a procedure index and a vector of program lines in which the procedure is called either directly or indirectly.

\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_

(2, [3, 8]) (3, [3, 10]) (1, [3]) (3, [5, 21]) (4, [5]) (0, [ ])  
\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_

1. Starting from the head of the queue, find all the variables that are modified and all the variables that are used in the procedure. For each of the program lines in the Queue Item, set these program lines (p) to modify and use the respective variables. For each of the program lines (p), if they are contained in another container statement (c), then set these program lines (c) to modify and use the respective variables too.

### 4.3.2. Building CFG

Given the AST, the following shows the pseudo code to build the CFG. Then, the CFG is traversed and the Next relationships are stored in the PKB’s Next table.

Here are the steps taken to build a CFG:

1. Maintain a currCFGNode pointer.
2. Create a CFG Root.
3. Iteratively traverse each of the procedure type nodes in AST.
   1. For each node, update the currASTNode and currCFGNode. Then create a CFG for stmtLst using the currASTNode.

Here are the steps taken to create a CFG for stmtLst:

1. Iteratively, traverse the children AST nodes and do the following:
   1. If type is ASSIGN, create CFG for Assign.
   2. If type is CALL, create CFG for Call.
   3. If type is WHILE, create CFG for While.
   4. If type is IF, create CFG for If.

To create CFG for Assign or Call, just create a new node and attach to existing CFG.

To create CFG for While, the steps are:

1. Create a new node and attach to existing CFG.
2. Save a pointer to the currCFGNode. This is the toNode for the backwards pointer of the while loop in the CFG that is to be set later.
3. Create the CFG for stmtLst.
4. Find the fromNode for the backwards pointer by running DFS on the currently created CFG.
5. Create a link from the fromNode to the toNode in the CFG.
6. Update currCFGNode.

To create CFG for If, the steps are:

1. Create a new node and attach to existing CFG.
2. For the then:stmtLst, create CFG for stmtLst.
3. Store a pointer to the currCFGNode. This is a leaf node that is to be connected to a dummy -1 node later.
4. For the else:stmtLst, create a CFG for stmtLst.
5. Store a pointer to the currCFGNode. This is another leaf node.
6. Create a dummy -1 node that is connected to both these leaf nodes.
7. Update currCFGNode.

After building the CFG, traverse the CFG and set the Next relationship in the PKB’s Next table using the appropriate method. This would allow the Query Evaluator to easily access the Next information without traversing the CFG.

## 4.4. Query Processor

Query processor consists of three parts: query processor (controller), query parser, and query evaluator.

### 4.4.1 Query Processor

Query Processor is a façade class for the whole component. The following shows the steps it takes:

1. Query Processor calls QueryParser to create a Query object from the given query string.
2. Query Processor then passes the Query object to the QueryEvaluator.
3. Query Evaluator will compute all necessary relations and return the results in the form of a list of integers.
4. Query Processor transforms the result into the correct display format and returns the answer to the user.

**Result Projection**

Once the results for each relation are computed by the Query Evaluator, the Query Processor does the following steps to project the correct result:

1. Create a table containing tuples of synonyms used during the Query Evaluation.

To create this table, we iterate through each relation and do one of the following

1. If none of the synonyms in the relations are already in the table, we set the new table to be the Cartesian product of the current table and the new synonyms
2. If one or both of the synonyms in the relation are already in the table, we set the new table to be the Join of the current table and the new synonyms
3. If there are no synonyms in the relation(the relation evaluates to a Boolean value), we do nothing to the table is its true, and clear the table if its false.
4. Augment the table to have synonyms selected by the query but not used in any relation. This is done by setting the new table to be the Cartesian product of the current table and the unused synonyms.
5. Create a set of strings to store the final answers. We use a set to prevent duplicate answer from appearing.
6. For each tuple in the table, create an answer based on the selected synonyms of the query.
7. Return all answers

### 4.4.2 Query Parser

Query parser has two major functionalities: query validation and query parsing, and they are implemented as functions in the **QueryParser** class. The controller calls query validator to check if the given query is syntactically correct. If it is, query controller will then parse the query by calling the query parser.

**Query Validation (Grammar check)**

**Query validation** is done using regular expression method using the grammar rules written in the handbook.

We have considered the option to validate while parsing the query (tokens), but decided to use regular expression to validate the query first then parse afterwards, because of the following reasons:

1. The code is much neater and simpler
2. It is simpler to define the book grammar rule in regex, hence less prone to mistakes
3. It is faster to detect errors rather than parsing and validating (especially if the error is towards the end of the query)
4. Query parsing becomes simpler if we already know the exact possible format of the string query that needs to be parsed.

All types of queries have been defined in a static string, following the grammar rules in the textbook. The strings are then used to validate the queries using regular expression, where the definitions earlier are used. It is very convenient because the grammar rules format in the book is very similar to a regex grammar rules.

As an example, consider this valid query

assign a; while w; Select a such that Follows(w, a) pattern a (“x”, \_”x+y”\_)

Query validator will first break down the query into statements, separated by semicolon. The query above will be broken down into three statements:

1. assign a
2. while w
3. Select a such that Follows(w, a) pattern a (“x”, \_”x+y”\_)

The validator will then use regular expression to check the validity of the statement and retrieve the tokens. Valid declaration statements (design-entity) will be stored into a map with the synonym as the key and its type as the value. This map is called the synonym map and will be used later by the parsing function. In this example, the synonym-map will look as follows:

|  |  |
| --- | --- |
| Synonym | Type |
| BOOLEAN | BOOLEAN |
| a | ASSIGN |
| w | WHILE |

This map enables easy look up when the query evaluator evaluates the query. Note that the BOOLEAN type always exists in the synonym table because user can use “BOOLEAN” in his select statement without any declaration.

Select statement will be broken down into such-that, with, or pattern clauses, whose parameters will be checked against the grammar rule.

Consider statement number 3 from the example above. Query validator will use regular expression to check the validity of the statement and retrieve the appropriate tokens.

If the regular expression matching fails, (for example, the number of arguments in the clause is not exactly two) the validator will instantly terminate and declare the query invalid. In the case where statement is valid, all the tokens from the select statement will be stored in a vector, *selectStatement*. This vector will be accessed by the parsing function later on. For efficiency, the *selectStatement* vector will only contain relevant tokens from the statement. Therefore, the unnecessary syntactic punctuation will be removed.

From the example above, the value of *selectStatement[] (dynamic array, c++ vector)* will be:

|  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| Select | a | such | that | Follows | w | a | pattern | a | x | \_”x+y”\_ |

**Query Parsing (Semantic check)**

The parser processes the *selectStatement* vector from the earlier. The *selectStatement* vector will be processed to construct a Query object with the following structure.

|  |
| --- |
| **Query** |
| *vector<string>* selected-synonym  *vector<Relationship>* relationship-vector  *map<string, synonym-type>* synonym-table |

Since each query can contain many select clauses, the selected synonyms and clauses are stored inside a vector for scalability purposes. All the synonyms present in the *selected-synonym* vector will be detected and validated once again on whether they have been declared earlier. If it is not declared, the query is invalid and **QueryParser** will indicate and return invalid. The string vector will then be stored in a **Query** class.

The synonym map that was created earlier will also be included in the **Query** object. Both such-that and pattern clauses will be stored as another object, **Relationship**, as the following.

|  |
| --- |
| **Relationship** |
| *enum* relationship-type  *enum* token-type  *string* token1  *string* token2  *TokenType* token1-type  *TokenType* token2-type  *string* pattern-synonym |

**Relationship** arguments will also be semantically checked to determine if they are valid. For example, both arguments in Follows clause have to be a statement (i.e. stmt, while, if, assign, call). If the arguments contain a constant synonym (e.g. constant c; while w; Select w such that Follows (c,w), then **QueryParser** will detect the error and return invalid.

The types of the tokens/arguments are detected by **QueryParser** and store them in the **Relationship** class. For example, if the token is “Second”, then its type will be IDENTIFIER, if the token is 1, then its type will be INTEGER, and if the token is w, then its type will be SYNONYM.

From the valid example above, the select-statement vector will be processed to produce the following.

|  |  |
| --- | --- |
| **Query** | |
| selected-synonym | a |
| relationships | [rel1, rel2] |
| synonym-table | map1 |

|  |  |
| --- | --- |
| **rel1** | |
| relationship-type | FOLLOWS |
| token1 | w |
| token2 | a |
| token1-type | SYNONYM |
| token2-type | SYNONYM |

|  |  |
| --- | --- |
| **rel2** | |
| relationship-type | PATTERN |
| token1 | “x” |
| token2 | \_”x+y”\_ |
| token1-type | IDENTIFIER |
| token2-type | UNDERSCOREEXPR |
| pattern-synonym | a |

|  |  |
| --- | --- |
| **map** | |
| Synonym | Type |
| a | ASSIGN |
| w | WHILE |
| BOOLEAN | BOOLEAN |

When the controller calls the parsing function, the function will return a query object. This object will then be passed to query evaluator.

**With-clause**

For with-clause, Query Parser with detect the conditions whether they are valid.

For example:

procedure p,q; Select q such that Calls (p,q) with p.procName=”Second”

p.procName = ”Second” will be parsed into two parts, left-hand-side and right-hand-side. Right-hand-side includes “Second” and will be stored in token2/argument-2 in a relationship object. Left-hand-side includes p.procName and will be stored in token1/argument-1 if token is valid. Synonym p will be checked against the map whether it exists. Since the attribute name is procName, it will also be checked on whether it is of type procedure.

Left-hand-side and right-hand-side will be checked on whether they are of the same type, either character strings or integers. If they are of different type, return invalid.

If it passes both validations, then the query is valid and stored in the relationship below:

|  |  |
| --- | --- |
| **rel2** | |
| relationship-type | WITH |
| token1 | p |
| token2 | “Second” |
| token1-type | SYNONYM |
| token2-type | IDENTIFIER |

“procName” is not stored because it is known that a synonym of type procedure can only have “procName” as its attribute name.

However, for special cases when the synonym is of type call. There will be another attribute to represent whether call refers to the procedure name (c.procName) or statement number (stmt#). **QueryEvaluator** will then check the additional attribute before evaluating the clause.

**Multiple Clauses**

For multiple clauses (such that, pattern, with), Query Parser works by keeping the previous clause keyword (e.g. “pattern”, or “with”), and use it to detect the clause type when it encounters “and”.

For example:

assign a; Select a pattern a(\_,\_) and a(“x”,\_”x+y”\_);

Parser will detect the first clause and store in the relationship class accordingly

|  |  |
| --- | --- |
| **rel1** | |
| relationship-type | PATTERN |
| token1 | \_ |
| token2 | \_ |
| token1-type | UNDERSCORE |
| token2-type | UNDERSCORE |
| pattern-synonym | A |

It also keeps track of the last relationship-type. Therefore, when Query Parser reaches the word “and”, it knows that it will be parsing a pattern clause again, translating the “and” keyword into “pattern”. Query Parser will then validate and parse accordingly with respect of the clause type, in this case it is “pattern”.

|  |  |
| --- | --- |
| **rel2** | |
| relationship-type | PATTERN |
| token1 | “x” |
| token2 | \_”x+y”\_ |
| token1-type | IDENTIFIER |
| token2-type | UNDERSCOREEXPR |
| pattern-synonym | a |

Both relationships will then be stored in a **Query** class and passed to **QueryEvaluator**.

### 4.4.3 Query Evaluator

**General Query Evaluation**

Our general query evaluation is as follows:

1. Given 2 synonyms s1, s2 and relation(s1,s2), check if s1 and s2 are already evaluated
2. If they are already evaluated, fetch them from storage and check if relation(s1,s2) is true
3. If they are not fetch them from pkb of the respective synonym type and check if relation(s1,s2) is true
4. Returns all pairs of (s1,s2) such that relation(s1,s2) is true

**Query Optimization**

To optimize and facilitate a faster running time for each query, we decided to order the relationships in our query object. Below is a summary of which relationships we will evaluate first and last.

**Relationships in order of priority**

|  |  |
| --- | --- |
| 1. **With relationships with only 1 synonym** | With c.value = 3 |
| 1. **Pattern relationships with no wildcards** | Pattern a(“x”, \_”x+y”\_) |
| 1. **Relationships with no synonyms** | Modifies(5, “x”), Follows( 1, 4), Parent(\_,4) |
| 1. **Direct Relationships with only 1 synonym** | Modifies(a, “x”), Follows( s, 4), Parent(s, \_) |
| 1. **Indirect Relationships with only 1 synonym** | Follows\*(s1,5), Parent\*(s1,8) |
| 1. **Direct Relationships with 2 synonyms** | Uses(s, v), Next(n, n2) |
| 1. **Indirect Relationships with 2 synonyms** | Next\*(n1, n2) |
| 1. **Affects\* Relationships with 2 synonyms** | Affects\*(a1,a2) |

**Query Clause Evaluation Examples**

stmt s1; Select s1 such that Follows(s1, 4)

As s1 first appears in Follows(s1, 4), QE will fetch all valid statement numbers by calling PKB method getAllStmts(Type Stmt) and check that for each value of s1 does Follows(s1,4) evaluate to true.

variable v; Select v such that Modifies(5,v) with v.varName = “x”

Here, we have 2 relationship clauses: Modifies and with. As with has a higher priority, we move it to the top and evaluate it first. We then store “x” as the evaluated answer for v. Hence when we move on to evaluate Modifies(5,v), we do not pull all valid variables but only the previously evaluated “x”.

**Alternative Query Data Representation**

Previously in our earlier reports we mentioned that we utilized a table to store the information each relationship has with each synonyms. For instance, Modifies(s, “v”) contains the synonym s, and hence we will store in our table “s is linked to Modifies(s,”v”)”.

An alternative representation of this will be to draw a graph and consider all synonyms as vertices and relationships as edges. To illustrate, consider the following example

Select BOOLEAN such that Modifies(s,v) and Follows(s,a) and Next(s, n)

The output graph would be:

Modifies

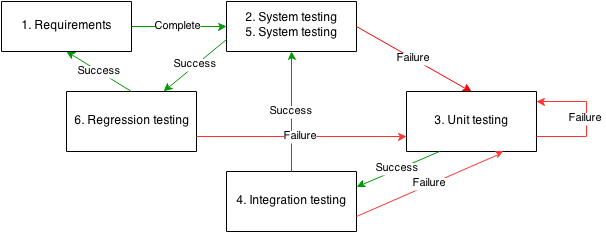
Follows

Next

# 5. Testing

**5.1. Overview**

Throughout 4 iterations, the team followed the following testing cycle:

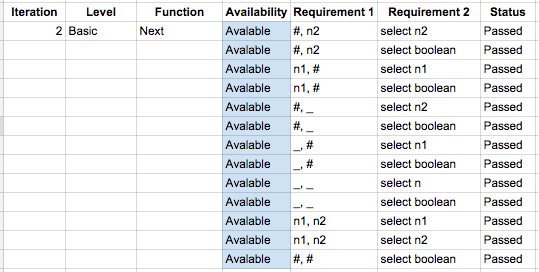


The cycle is started with the creation of system testing scripts from the current iteration's requirements by the test lead. These test scripts will become executable requirements, which will be referenced by other team members. As the development progresses, each team member is responsible for his/her own unit testing. As soon as two or more components pass the unit testing, integration testing will be carried out. The success of this integration testing will prompt another run of system testing. Upon successful completion of system testing, the testing cycle is concluded by regression testing. This is done by running all of the past and current test cases to make sure that the new feature does not interfere with other features.

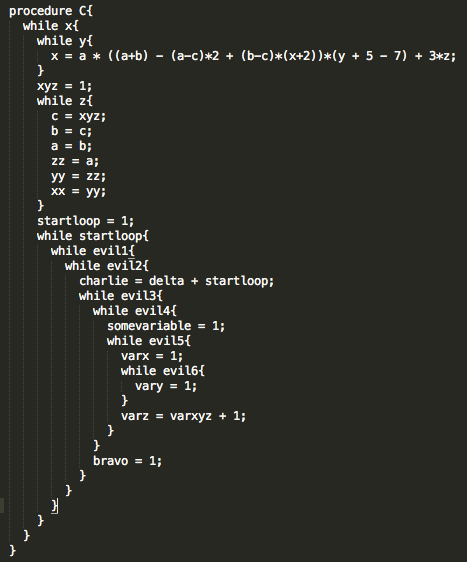
**5.2. Testing Standards**

To ensure a thorough test coverage, the following testing standards are observed.

1. Feature-centric testing



1. Stress testing



assign a1,a2,a3,a4,a5,a6,a7,a8;

Select <a1,a2,a3,a4,a5,a6,a7,a8> such that Follows\*(a1,a2) and Follows\*(a3,a4) and Follows\*(a5,a6) and Follows\*(a7,a8)

1. Optimization testing

procedure p1, p2; stmt s1, s2; prog\_line n1, n2; variable v;

Select BOOLEAN such that Calls\* (p1, p2) and Next\* (n1, n2) and Uses (s1, v) and Affects\*(1, 30)

procedure p1, p2; stmt s1, s2, s3, s4, s5, s6, s7;

Select s7 such that Calls\*(p1, p2) and Modifies (s1, v1) and Uses (s2, v2) and Modifies (s2, v2) and Affects\* (s2, s3) and Affects\* (s3, s4) and Parent\* (s5, s6) and Modifies (s7, v) with s7.stmt# = 29

1. Positive, negative, and disruptive testing
2. Separation of concerns

# 6. Coding Standards

Our entire team adopted similar coding standards to maintain a common coding quality. Some of the coding standards are listed below:

1. Indentation and whitespace
   1. Use it to indicate different code segments
2. Comments to enhance understanding and communication
3. Descriptive variable declarations
   1. Always start with lower case
   2. Use CamelCase
   3. Use only letters and numbers
4. Informative function naming conventions
   1. All getters start with “get”
   2. All setters start with “set”
   3. All functions that start with “is” returns a Boolean value
5. Keep it simple and effective
   1. Avoid complex code fragments
6. Refactoring

**Standards between abstract APIs and concrete APIs**

The correspondence between the abstract APIs and the concrete APIs was improved by:

1. Ensuring that the abstract APIs provide an interface for the concrete APIs to be built upon.
2. Making abstract APIs as comprehensive as possible by:
   1. Offering an extensive description and responsibility of the abstract API
   2. Specifying the complete parameters needed for the function

# 7. Project Evaluation

/\***YOLIM** \*/

1. How would you improve your SPA if more time was available?

2. What would you done differently if you were to start project again?

3. Comment on the experience gained in this project in respect to: a) working in the team,

b) incremental development,

c) complexity of the SPA problem and program solution,

d) what did work well?

e) what did not work well?

f) what did you learn in this project course?

4. Comment on the tools used for the project a) Were the recommended tools useful?

b) What other tools did you use (if any), and in what ways were they useful?

c) What were the problems you faced when using each tool?

d) In which areas would you like to have had more tool support?

5. What management lessons have you learned?

Time Management

We learnt the importance of doing your own allocated work properly and on time so as not to slow down the progress of the entire team. Not only that, we learnt that although the iterations were different in their requirements, they were similar in the duties each of the team members had to execute. For example, Adinda was to write the test cases for the upcoming iteration at the beginning. Then, each of the other team members was supposed to work on their individual components requirements for that iteration. Following that, Kester, the designated tester, ran these test cases. After that, debugging was done by the respective team members as necessary.

6. What advice would you give to the students who will take this course in the future?

It is important to delegate clear duties from the start and be disciplined enough to do one’s own duties well for each iteration. Moreover, students should start Autotester testing quite a few days before the end of the iteration to allow sufficient time to debug.

7. Suggest how we could improve this project course.

This course could be improved by asking the students to submit their code at various midpoints so that the progress could be monitored more closely. Also, I suggest that there could be a greater percentage for another individual component in the final grade which could be more fair to each team member in terms of effort they put in. This component could either be code related or related to testing or related to testing us on the new relationships for the extension.

8. Discuss any other experiences.

-

9. Comment on Handbook

The Handbook is concise but it could provide more test case examples to give a more comprehensive explanation. Also, it could include more explanation on special corner cases or rules that we might miss.

# API

Please view our Doxygen at:

[www.comp.nus.edu.sg/~kester/cs3202](http://www.comp.nus.edu.sg/~kester/cs3202)