数据科学与工程算法基础 习题3

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1

由于 $\mu = E(X), \sigma^2 = E\left[(X - \mu)^2\right]$

故由Chebyshev不等式可知

$$P(|X^*| \ge c) = P\left(\left|\frac{X - \mu}{\sigma}\right| \ge c\right)$$

$$= P(|X - \mu| \ge c|\sigma|)$$

$$\le \frac{\sigma^2}{c^2 \sigma^2}$$

$$= \frac{1}{c^2}$$

2

由 X_i 独立同分布可知

$$egin{aligned} E(ar{X}) &= E\left(rac{1}{n}\sum_{i=1}^n X_i
ight) \ &= rac{1}{n}\sum_{i=1}^n E(X_i) \ &= \mu \ Var(ar{X}) &= Var\left(rac{1}{n}\sum_{i=1}^n X_i
ight) \ &= rac{1}{n^2}\sum_{i=1}^n Var(X_i) \ &= rac{\sigma^2}{n} \end{aligned}$$

故由Chebyshev不等式得

$$P\left(|\bar{X} - \mu| \ge \varepsilon\right) \le \frac{\frac{\sigma^2}{n}}{\varepsilon^2} = \frac{\sigma^2}{n\varepsilon^2}$$

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(1) 易见 $X \sim b(n, \frac{1}{2})$

于是 $E(X) = \frac{n}{2}, Var(X) = \frac{n}{4}$

现取 $\varepsilon=rac{n}{4}$,则由Chebyshev不等式可知

$$P\left(\left|X - \frac{n}{2}\right| \ge \frac{n}{4}\right) \le \frac{\frac{n}{4}}{\frac{n^2}{16}} = \frac{4}{n}$$

由二项分布的概率密度函数可知 $P(X \geq rac{3}{4}n) = P(X \leq rac{n}{4})$

故

$$egin{split} P\left(\left|X-rac{n}{2}
ight|\geqrac{n}{4}
ight)&=P\left(\{X\geqrac{3}{4}n\}\cap\{X\leqrac{n}{4}\}
ight)\ &=P(X\geqrac{3}{4}n)+P(X\leqrac{n}{4})\ &=2P(X\leqrac{n}{4})\leqrac{4}{n} \end{split}$$

也即 $P(X \leq \frac{n}{4}) \leq \frac{2}{n}$

当 $\frac{n}{4} \notin \mathbb{Z}^+$ 时,

$$P\left(X < \frac{n}{4}\right) = P\left(X \le \frac{n}{4}\right) \le \frac{2}{n}$$

当 $\frac{n}{4} \in \mathbb{Z}^+$ 时,

由于

$$P\left(X = \frac{n}{4}\right) = \binom{n}{\frac{n}{4}} \left(\frac{1}{2}\right)^{\frac{n}{4}} \left(\frac{1}{2}\right)^{n - \frac{n}{4}} = \binom{n}{\frac{n}{4}} \left(\frac{1}{2}\right)^n$$

故

$$egin{split} P\left(X < rac{n}{4}
ight) &= P\left(X \leq rac{n}{4}
ight) - P\left(X = rac{n}{4}
ight) \ &\leq rac{2}{n} - inom{n}{rac{1}{4}}igg(rac{1}{2}igg)^n \end{split}$$

(2) 取 $\delta = \frac{1}{2}$

故由Chernoff不等式可知

$$P\left(X < \frac{n}{4}\right) = P\left(X < \left(1 - \frac{1}{2}\right)\frac{n}{2}\right)$$
$$< \exp\left\{-\frac{1}{2} \cdot \frac{n}{2} \cdot \frac{1}{4}\right\}$$
$$= \exp\left\{-\frac{n}{16}\right\}$$

4

(1) 对任意 t > 0,有

$$P\left(X > (1+\delta)\mu\right) = P\left(\exp\left\{tX\right\} > \exp\left\{t(1+\delta)\mu\right\}\right)$$
 $< \frac{\prod_{i=1}^{n} E(\exp\left\{tX_{i}\right\})}{\exp\left\{t(1+\delta)\mu\right\}}$

而由 $1 - x < e^{-x}$ 可知

$$E\left(\exp\{tX_i\}\right) = p_i e^t + (1 - p_i) \ = 1 - p_i (1 - e^t) \ < \exp\left\{p_i (e^t - 1)\right\}$$

故

$$\prod_{i=1}^n E(\exp\left\{tX_i
ight\}) < \prod_{i=1}^n \exp\left\{p_i(e^t-1)
ight\} = \exp\left\{\mu(e^t-1)
ight\}$$

于是

$$P\left(X>(1+\delta)\mu
ight)<rac{\exp\left\{\mu(e^t-1)
ight\}}{\exp\left\{t(1+\delta)\mu
ight\}}=\exp\left\{\mu(e^t-1-t-t\delta)
ight\}$$

要使得该式对任意 t > 0 成立,则左式应当小于右式的最小值。

令
$$f(t) = e^t - 1 - t - t\delta$$
,由 e^x 的单调性及 $\mu > 0$ 可知

$$\mathbf{minimize}_t \exp\left\{\mu(e^t - 1 - t - t\delta)\right\} \leftrightarrow \mathbf{minimize}_t \ f(t)$$

求导易知当 $t = \ln{(1+\delta)}$ 时,f(t) 取到最小值

因此

$$egin{aligned} P(X > (1+\delta)\mu) &< \exp\left\{\mu(\exp\{\ln{(1+\delta)}\} - 1 - \ln{(1+\delta)} - \delta \ln{(1+\delta)})
ight\} \ &= \exp\left\{\mu \ln{rac{e^{\delta}}{(1+\delta)^{1+\delta}}}
ight\} \ &= \left(rac{e^{\delta}}{(1+\delta)^{1+\delta}}
ight)^{\mu} \end{aligned}$$

(2)

Lemma: 当
$$x > 0$$
 时, $\ln(1+x) \ge x/(1+\frac{x}{2})$

Proof:
$$\Rightarrow f(x) = \ln(1+x) - \frac{2x}{2+x}$$

易见
$$f(x) = \ln(1+x) - 2 + \frac{4}{2+x}$$

故
$$f'(x) = \frac{1}{1+x} - \frac{4}{(2+x)^2}$$

由此可得唯一驻点 x=0

而
$$f(0) = 0, f(2) = \ln 3 - 1 > 0$$

故对任意
$$x>0$$
, $\ln\left(1+x\right)\geq x/(1+\frac{x}{2})$

由(1)可知对任意 $\delta \in (0,1)$,有

$$P(X>(1+\delta)\mu)<\left(rac{e^{\delta}}{(1+\delta)^{1+\delta}}
ight)^{\mu}=\exp\left\{\mu(\delta-(1+\delta)\ln{(1+\delta)})
ight\}$$

而由引理可得 $(1+\delta)\ln{(1+\delta)} \geq rac{2\delta(1+\delta)}{2+\delta}$

因此

$$egin{split} P(X > (1+\delta)\mu) &< \exp\left\{\mu\left(\delta - rac{2\delta(1+\delta)}{2+\delta}
ight)
ight\} \ &= \exp\left\{-rac{\mu\delta^2}{2+\delta}
ight\} \ &< \exp\left\{-rac{\mu\delta^2}{3}
ight\} \end{split}$$

6

由 $X_i \overset{i.i.d}{\sim} b(1,p)$ 可知

$$E(ar{X}) = p, Var(ar{X}) = rac{1}{n}p(1-p)$$

故由Chebyshev定理可得

$$P(|ar{X}-p| \geq arepsilon p) \leq rac{1-p}{narepsilon^2 p}$$

因此

$$egin{split} P(|ar{X}-p| \leq arepsilon p) &= 1 - P(|ar{X}-p| \geq arepsilon p) \ &\geq 1 - rac{1-p}{narepsilon^2 p} \end{split}$$

要使得 $P(|ar{X}-p| \leq arepsilon p) \geq 1-\delta$ 成立,则有

$$1 - \frac{1-p}{n\varepsilon^2 p} \geq 1 - \delta$$

也即

$$n \geq rac{1-p}{arepsilon^2 p \delta}$$