1 Introdução

O entrelaçamento é um ingrediente fundamental em informação quântica e por isso é importante testar a sua presença em estados preparados por um computador quântico. No caso de dois qubits esse teste pode ser feito recorrendo à versão da desigualdade de Bell devida a Clauser-Horne-Shimony-Holt (CHSH). A desigualdade de Bell pode ser extendida para um número n de qubits superior a 2 recorrendo aos chamados polonómios de Mermim.

1.1 Sistemas de dois qubits

Relembremos o paradoxo de EPR na versão de Bohm : um pião π^0 decai em repouso, num eletrão, e^- , e na sua antipartícula, o positrão, e^+ :

$$\pi^0 \to e^- + e^+$$
 .

O estado conjunto do eletrão e positrão tem de ter spin igual ao do pião, ou seja, é um estado singleto :

$$\frac{1}{\sqrt{2}}(|\uparrow\downarrow\rangle - |\downarrow\uparrow\rangle) \tag{1}$$

A conservação de momento linear impõe que o eletrão e o positrão se afastem em direções opostas. O eletrão é detetado no laboratório A (de Alice), o positrão é detetado no laboratório B (de Bob).

Alice
$$\stackrel{e^-}{\longleftarrow} \stackrel{\pi^0}{\longleftarrow} \stackrel{e^+}{\longrightarrow}$$
 Bob

Alice mede σ na direção \vec{a} ou $\vec{a'}$, Bob mede σ na direção \vec{b} e $\vec{b'}$. A escolha entre as duas direções possíveis é arbitrária mas o valor medido - que designaremos por a ou a' no caso de Alice, e por b ou b' no caso de Bob - tem necessariamente o valor +1 ou -1 em qualquer dos casos:

$$a, a' = \pm 1$$
 , $b, b' = \pm 1$,

o mesmo sucedendo com qualquer dos produtos ab, ab', a'b ou a'b'.

No texto que se segue seguir-se-á a abordagem das referências [1, 2].

O polinómio de Bell devido a Clauser-Holt-Horne-Shimony (CHSH) define-se como

$$B_{\text{CHSH}} = ab + ab' + a'b - a'b' \tag{2}$$

$$= a(b+b') + a'(b-b')$$
 (3)

O princípio do realismo local (LR, iniciais de *Local Realism*) assume que o resultado de uma medida num sistema não pode influenciar instantaneamente o resultado da medida num segundo sistema e que os valores das quantidades físicas têm uma realidade física independentemente de serem medidos ou não.

Assim, de acordo com o realismo local, um dos termos da expressão (3) é nulo e o outro é ± 2 , ou seja o valor máximo de B_{CHSH} é $\langle B_{\text{CHSH}} \rangle^{\text{LR}} = 2$.

Em Mecânica Quântica as quantidades a,a',b e b' são operadores (neste caso, operadores de spin de Pauli) cujos valores próprios são ± 1 e satisfazem as igualdades

$$a^2 = a'^2 = b^2 = b'^2 = 1$$
. (4)

O operador de Bell escreve-se como

$$B_{\text{CHSH}} = a \otimes b + a \otimes b' + a \otimes b' + a \otimes b'$$
 (5)

Se combinarmos as expressões(4) e (5) podemos concluir que

$$B_{\text{CHSH}}^2 = 4\mathbb{1}_a \otimes \mathbb{1}_b - [a, a'] \otimes [b, b'] \tag{6}$$

Se admitirmos o realismo local, todos os observáveis comutam, o que conduz ao resultado já obtido anteriormente : $\langle B_{\rm CHSH} \rangle^{\rm LR} = \sqrt{\langle B_{\rm CHSH}^2 \rangle^{\rm LR}} = \sqrt{4} = 2$.

Em Mecânica Quântica, contudo, os observáveis a, a' e b, b' não comutam. O valor absoluto máximo dos valores próprios dos comutadores ocorrem quando as duas direções alternativas são perpendiculares e os observáveis a, a' e b, b' são representados por operadores de Pauli que obedecem às relações $[\sigma_j, \sigma_k] = 2 i \epsilon_{jkl} \sigma_l$ onde ϵ_{jkl} é o símbolo de Levy-Civita. Uma vez que σ_l tem valores próprios ± 1 , o valor máximo absoluto dos valores próprios dos comutadores é 2. Assim, o valor máximo de $\langle B_{\text{CHSH}} \rangle$ previsto pela Mecânica Quântica é

$$\langle B_{\text{CHSH}} \rangle^{\text{QM}} = \sqrt{\langle B_{\text{CHSH}}^2 \rangle^{\text{QM}}} = \sqrt{8} = 2\sqrt{2}$$
.

Assim, o realismo local conduz à desigualdade $\langle B_{\rm CHSH} \rangle \leq 2$ cuja violação foi demonstrada experimentalmente, escolhendo um estado de dois qubits para os quais a Mecânica Quântica prevê $\langle B_{\rm CHSH} = 2\sqrt{2}$.

A violação de uma desigualdade de Bell pode ser caraterizada pela razão R(B) associada a um polinómio de Bell:

$$R(B) = \frac{\langle B \rangle^{\text{QM}}}{\langle B \rangle^{\text{LR}}}$$

O valor de R(B) é uma medida de não localidade pois se R(B) > 1 o estado não pode ser descrito pelo realismo local. É importante notar que entrelaçamento e não localidade são conceitos distintos porque há estados entrelaçados que não verificam a condição R(B) > 1, ou seja o entrelaçamento é uma condição necessária mas não suficiente para violar uma desigualdade de Bell. Contudo, estados que produzem valores máximos de R(B) são fortemente entrelaçados. ¹ No caso de dois qubits o valor máximo de R(B) é

$$R(B_{\text{CHSH}}) = \frac{\langle B_{\text{CHSH}} \rangle^{\text{QM}}}{\langle B_{\text{CHSH}} \rangle^{\text{LR}}} = \sqrt{2} .$$

1.2 Sistema de 3 qubits

Várias desigualdades de Bell têm sido propostas para sistemas com um número n de qubits superior a 2 [4, 5]. Para n ímpar, a quantidade R é maximizada usando as chamadas desigualdades de Mermin [6, 4, 1, 2]. Para n qubits, as desigualdades são definidas para os observáveis a_1, a_2, \dots, a_n , todos eles com valores próprios ± 1 . Os polinómios de Mermim, M_n , podem ser obtidos definindo $M_1 \equiv 1$ e construindo M_n recursivamente a partir de M_{n-1} [7, 2]. O polinómio de Mermim M_2 é o polinómio B_{CHSH} . Para 3 qubits temos

$$M_3 = (a_1 a_2 a_3' + a_1 a_2' a_3 + a_1' a_2 a_3) - (a_1' a_2' a_3')$$

$$\tag{7}$$

¹O grau de entrelaçamento de dois qubits pode ser quantificado através da medida da entropia de Von Neumann do operador densidade. Se o número de qubits entrelaçados for superior a dois, há diferentes medidas possíveis do grau de entrelaçamento [3].

cujo quadrado é

$$M_3^2 = 4 \quad \mathbb{1}_1 \otimes \mathbb{1}_2 \otimes \mathbb{1}_3 - ([a_1, a_1'] \otimes [a_2, a_2'] \otimes \mathbb{1}_3 + + [a_1, a_1'] \otimes \mathbb{1}_2 \otimes [a_3, a_3'] + \mathbb{1}_1 \otimes [a_2, a_2'] \otimes [a_3, a_3'])$$
(8)

A condição de que todas as variáveis comutam em condições de realismo local conduz-nos ao limite superior $\langle M_3 \rangle^{\text{LR}} = \sqrt{4} = 2$. O valor máximo de $\langle M_3 \rangle^{\text{QM}}$ é $\langle M_3 \rangle^{\text{QM}} = \sqrt{4+4\times 3} = 4$. Assim, o valor máximo de R para n=3 é

$$R(M_3) = \frac{\langle M_3 \rangle^{\text{QM}}}{\langle M_3 \rangle^{\text{LR}}} = \frac{4}{2} = 2$$

Em geral, para n qubits, um estado do tipo GHZ

$$|\Psi_{\text{GHZ}}\rangle = \frac{1}{\sqrt{2}} \left(|00\cdots 0\rangle + e^{i\phi} |11\cdots 1\rangle \right)$$

produz a máxima violação da desigualdade de Mermim [6, 4, 1]. Os casos em que o número de qubits é n=2 e n=3 correspondem também a estados em que o entrelaçamento é máximo. Contudo, o mesmo não sucede para $n\geq 4$ [2] e há até sugestões de utilização de machine learning para classificar o entrelaçamento de estados [8].

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4º problema para avaliação

O objetivo deste problema é o de testar o grau de entrelaçamento dos qubits de um computador quântico da IBM disponível on-line, reproduzindo a abordagem do artigo de Daniel Alsina e José Ignacio Latorre [1] que se junta em anexo.

Leia o artigo e consulte o site da IBM (https://www.ibm.com/quantum-computing/) para saber como submeter um circuito quântico. Registe-se na plataforma da IBM.

- 1. Considere um sistema de 3 qubits e o polinómio de Mermim M_3 .
 - (a) Verifique que se escolhermos $a_i = \sigma_x$ e $a_i' = \sigma_y$ (como sucede no artigo de Daniel Alsina e José Ignacio Latorre [1]) o estado que maximiza a violação da desigualdade de Mermin ou seja, o estado que maximiza $\langle M_3 \rangle^{\rm QM}$ é $|\Phi\rangle = \frac{1}{\sqrt{2}}(|000\rangle + i|111\rangle)$. Verifique também que, se escolhermos $a_i = \sigma_y$ e $a_i' = \sigma_x$, o estado a escolher é $|\Psi\rangle = \frac{1}{\sqrt{2}}(|000\rangle |111\rangle)$.
 - (b) Proceda de forma semelhante ao implementado no artigo [1] no caso de 3 qubits, mas adapte o procedimento à escolha $a_i = \sigma_y$ e $a'_i = \sigma_x$. Assim,
 - i. Escreva os circuitos necessários para testar a desigualdade de Mermim, $\langle M_3 \rangle \leq 2$, num computador quântico da IBM, no caso $a_i = \sigma_y$ e $a_i' = \sigma_x$.²
 - ii. Submeta os circuitos da alínea anterior no computador quântico da IBM. Repita um número de vezes suficientemente elevado para poder verificar a violação da desigualdade de Mermim, $\langle M_3 \rangle \leq 2.0$. Note que para concluir se há ou não violação da desigualdade é essencial calcular a incerteza associada ao valor médio $\langle M_3 \rangle$. Use o procedimento sugerido no artigo para calcular a incerteza.
 - iii. Num computador quântico o conhecimento sobre as propriedades do sistema físico é importante. Leia atentamente a informação disponível no site da IBM sobre "backend configuration". Tente criar dois circuitos que desempenham a mesma tarefa, um que lhe pareça otimizado no sentido de diminuir os erros possíveis (o seu circuito "ótimo") e um circuito que maximiza esses erros (um circuito "péssimo"). Indique claramente a informação em que se baseou e os critérios que utilizou. Compare o valor de $\langle M_3 \rangle$ e respetiva incerteza para os dois circuitos.

²Inclui-se um pacote para edição de circuitos em I⁴TEX , apenas para o caso de escolherem esse formato. Mas não é necessário, uma resolução manuscrita não tem qualquer penalização.

Experimental test of Mermin inequalities on a five-qubit quantum computer

Daniel Alsina and José Ignacio Latorre

Departament Física Quàntica i Astrofísica, Universitat de Barcelona, Diagonal 645, 08028 Barcelona, Spain and Institut de Ciències del Cosmos (ICCUB), Martí i Franquès 1, 08028 Barcelona, Spain (Received 25 May 2016; published 11 July 2016)

Violation of Mermin inequalities is tested on the five-qubit IBM quantum computer. For three, four, and five parties, quantum states that violate the corresponding Mermin inequalities are constructed using quantum circuits on superconducting qubits. Measurements on different bases are included as additional final gates in the circuits. The experimental results obtained using the quantum computer show violation of all Mermin inequalities, with a clear degradation of the results in the five-qubit case. Though this quantum computer is not competitive to test Mermin inequalities as compared to other techniques when applied to a few qubits, it does offer the opportunity to explore multipartite entanglement for four and five qubits beyond the reach of other alternative technologies.

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Quantum physics can be discriminated from classical physics using Bell-type inequalities [1]. In particular, the violation of Bell inequalities for two qubits has been extensively verified since they were first checked in atomic physics experiments [2,3]. Later on, the improvement of quantum optics techniques as well as other technologies such as nitrogen-vacancy (NV) centers has made it possible to eliminate many of the loopholes in the experimental verification of two-qubit Bell inequalities [3].

An extension of Bell inequalities to a larger number of particles corresponds to the set of Mermin inequalities [4]. Such inequalities should be maximally violated by Greenberger-Horne-Zeilinger (GHZ)-type states [5]. The experimental verification of multipartite Mermin inequalities faces the problem of a good control of three or more qubits, including the generation of entangled states, and the possibility of performing different measurements on each one. Violation of Mermin inequalities has been reported for three qubits [6] and four qubits [7], where all qubits are made out of photons, and for up to 14 qubits with a quantum computer based on ion traps [8].

In the case of superconducting qubits, violation of the Clauser-Horne-Shimony-Holt (CHSH) inequality was achieved in Ref. [9], whereas the GHZ construction and the three-qubit Mermin inequality violation was demonstrated in Ref. [10]. For a general review of theoretical and experimental progress in Bell inequalities, see Ref. [11].

The construction of the first prototypes of quantum computers allows for the possibility of experimenting with quantum states containing more than two qubits. In particular, IBM has opened the use of its five-qubit quantum computer to the community [12]. We here report results on the use of this quantum computer to test the violation of Mermin inequalitites for three, four, and five superconducting qubits.

I. MERMIN POLYNOMIALS

Local realism can be tested using Mermin polynomials. The technique to generate them is explained for example in Ref. [13]. The Mermin polynomial for three qubits is

$$M_3 = (a_1 a_2 a_3' + a_1 a_2' a_3 + a_1' a_2 a_3) - (a_1' a_2' a_3'), \tag{1}$$

where a_i and a_i' correspond to two different settings for the measurement of each qubit i. Each measurement can take

the values $\{-1,1\}$. Classical theories obey local realism (LR) which translates into a bound for the expectation value of the Mermin polynomial, $\langle M_3 \rangle^{LR} \leqslant 2$. Instead, for quantum mechanics (QM) the observables a_i and a_i' are built out of linear combinations of Pauli matrices. Each measurement is expressed as a Kronecker product of the three local measurements and the expectation value for $\langle M_3 \rangle$ is the maximum eigenvalue of the resulting 8×8 matrix. In this case, the maximum possible eigenvalue, and therefore the quantum bound, is $\langle M_3 \rangle^{QM} \leqslant 4$. We briefly construct circuits to check the violation of the classical bound on this inequality.

The Mermin polynomial that will be experimentally checked for four-qubits is

$$M_{4} = -(a_{1}a_{2}a_{3}a_{4}) + (a_{1}a_{2}a_{3}a'_{4} + a_{1}a_{2}a'_{3}a_{4} + a_{1}a'_{2}a_{3}a_{4} + a'_{1}a_{2}a_{3}a_{4}) + (a_{1}a_{2}a'_{3}a'_{4} + a_{1}a'_{2}a_{3}a'_{4} + a_{1}a'_{2}a'_{3}a_{4} + a'_{1}a_{2}a_{3}a'_{4} + a'_{1}a_{2}a'_{3}a_{4} + a'_{1}a'_{2}a_{3}a_{4}) - (a_{1}a'_{2}a'_{3}a'_{4} + a'_{1}a_{2}a'_{3}a'_{4} + a'_{1}a'_{2}a_{3}a'_{4} + a'_{1}a'_{2}a'_{3}a_{4}) - (a'_{1}a'_{2}a'_{3}a'_{4}),$$

with a classical bound of $\langle M_4 \rangle^{\rm LR} \leqslant 4$ and a quantum bound of $\langle M_4 \rangle^{\rm QM} \leqslant 8\sqrt{2}$.

In the five-qubit case, the Mermin polynomial reads

$$M_{5} = -(a_{1}a_{2}a_{3}a_{4}a_{5}) + (a_{1}a_{2}a_{3}a'_{4}a'_{5} + a_{1}a_{2}a'_{3}a_{4}a'_{5} + a_{1}a'_{2}a_{3}a_{4}a'_{5} + a'_{1}a_{2}a_{3}a_{4}a'_{5} + a_{1}a_{2}a'_{3}a'_{4}a_{5} + a_{1}a'_{2}a_{3}a'_{4}a_{5} + a'_{1}a_{2}a_{3}a'_{4}a_{5} + a_{1}a'_{2}a'_{3}a_{4}a_{5} + a'_{1}a_{2}a'_{3}a_{4}a_{5} + a'_{1}a'_{2}a_{3}a_{4}a_{5}) - (a_{1}a'_{2}a'_{3}a'_{4}a'_{5} + a'_{1}a_{2}a'_{3}a'_{4}a'_{5} + a'_{1}a'_{2}a_{3}a'_{4}a'_{5} + a'_{1}a'_{2}a'_{3}a_{4}a'_{5} + a'_{1}a'_{2}a'_{3}a'_{4}a_{5}),$$
(3)

with a classical bound of $\langle M_5 \rangle^{LR} \leq 4$ and a quantum bound of $\langle M_5 \rangle^{QM} \leq 16$.

II. CIRCUIT IMPLEMENTATION

There are a number of technical issues associated with the specific implementation of the IBM five-qubit quantum computer. This quantum computer is based on superconducting flux qubits that live on a fridge with a temperature of about 15 mK, where only one of the qubits can be used to act as

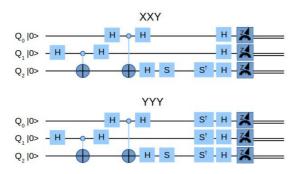


FIG. 1. The two circuits used for the three-qubit Mermin inequality. The first circuit corresponds to the $\sigma_x\sigma_x\sigma_y$ experiment, and the second circuit corresponds to the $\sigma_y\sigma_y\sigma_y$ experiment. The S^\dagger gates make the difference between a σ_x measurement and a σ_y measurement.

the target qubit of any controlled-NOT (CNOT) gate. In the test of Mermin inequalitites, only GHZ-like states have to be created. This requires the use of a Hadamard gate on a control qubit followed by CNOTs targeted to the rest. In order to implement this kind of action we need to operate CNOT gates targeted to other qubits. This can be done using the relation $\text{CNOT}_{1\rightarrow 2} = (H_1 \otimes H_2) \text{CNOT}_{2\rightarrow 1} (H_1 \otimes H_2),$ where H_1 and H_2 are Hadamard gates on qubits 1 and 2, whereas $\text{CNOT}_{1\rightarrow 2}$ is the controlled-NOT gate which is controlled by qubit 1.

In our choice of settings, the needed GHZ-like states have relative phases, as in the case of three qubits, where $|\phi\rangle=1/\sqrt{2}(|000\rangle+i|111\rangle)$. These phases are implemented using S and T gates, which are one-qubit gates that mutiply the $|1\rangle$ term with $\pi/2$ and $\pi/4$ phases, respectively. Measurements can only be done on the σ_z basis, but they can be simulated in another basis with the help of additional gates, namely an H gate for σ_x and an S^\dagger gate followed by an H gate for σ_y .

Another relevant issue to be considered is that not all of the qubits are equally robust in the present quantum computer, some have relaxation and decoherence times larger than others, although all of them are of the order of $T=O(100~\mu s)$. We adapt our circuits to minimize the number of gates on the qubits that behave more poorly. For example, gates that implement phases that can be put freely in any qubit are allocated to the most robust ones.

Figures 1 and 2 represent the three circuits for the three-, four-, and five-qubit Mermin inequalities. In principle we need to perform as many experiments as the number of terms in the Mermin inequalities (1), (2), and (3). However due to our limited access to the computer and the symmetry of particle exchange of the states and the inequalities, only one experiment for a term representative of each number of primes (a'_i) was run. In our choice of settings, the number of primes amounts to the number of σ_v measurements, whereas the nonprimes (a_i) correspond to σ_x measurements. We thus have two experiments for three qubits, five experiments for four qubits, and three experiments for five qubits. Each experiment was run 8192 times, the maximum available, except for the three-qubit experiments, which were run only 1024 times. When computing the expected value of the whole polynomial, each experiment was given the corresponding weight. In the errors discussion we compare results obtained when using the

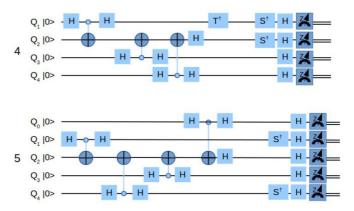


FIG. 2. Two of the circuits used for the four-qubit and five-qubit Mermin inequalities. The first circuit corresponds to the $\sigma_y \sigma_y \sigma_x \sigma_x$ experiment, whereas the second corresponds to the $\sigma_x \sigma_y \sigma_x \sigma_x \sigma_y$ experiment. The S^{\dagger} gates make the difference between a σ_x measurement and a σ_y measurement. In order to change from σ_x to σ_y , one has to add an S^{\dagger} gate, or remove it to do the opposite. With this technique one can obtain all the circuits needed to test the inequalities.

symmetry with results obtained without using it, computing all the terms, for the three-qubit case.

III. RESULTS

We now give a more detailed discussion of the results for the three-qubit case and an abridged one for the four- and five-qubit cases, as much of it is basically the same.

In order to check the violation of the inequality, one has to choose the settings and the corresponding state that maximally violate it. One possibility is to choose the settings $a_i = \sigma_x$ and $a'_i = \sigma_y$ for all the qubits. The state that maximizes the quantum violation in this case is $|\phi\rangle = 1/\sqrt{2}(|000\rangle + i|111\rangle)$.

The three-qubit Mermin inequality has four terms as shown in Eq. (1). In principle, four different circuits are needed, one for each term. The state will be the same for all of them, but the settings change. However, one can use the symmetry of the state and the inequality to reduce the number of measurements needed if there is limited access to the experimental setting as is our case. All the terms that have the same number of primes (a_i') are represented by the same circuit by symmetry. We then considered only two different experiments, with 1024 runs each, the $\sigma_x \sigma_x \sigma_y$ experiment and the $\sigma_y \sigma_y \sigma_y$ experiment. The results are shown in Table I.

TABLE I. Table of detailed results for the two three-qubit experiments. In bold are results of even parity; in italic are results of odd parity. Counts for each result are expressed in probabilities computed out of 1024 runs. Computation of the expected value of XXY gives $\langle XXY \rangle = 0.715$ and of YYY gives $\langle YYY \rangle = -0.710$. The combination $3\langle XXY \rangle - \langle YYY \rangle$ gives $\langle M_3 \rangle_{\rm exp} = 2.85 \pm 0.02$.

Result XXY 001 010 011 100 101 111 0.229 0.042 0.024 0.194 0.043 0.203 0.231 0.033 **Probability** Result YYY 000 001 010 011 100 101 110 111 Probability $0.050\ 0.188\ 0.188\ 0.028\ 0.258\ 0.026\ 0.041\ 0.221$

TABLE II. Table of results. LR corresponds to the local realism bound for each Mermin inequality, QM corresponds to the quantum bound, and Expt. is the experimental result.

	LR	QM	Expt.
Three qubits	2	4	2.85 ± 0.02
Four qubits	4	$8\sqrt{2}$	4.81 ± 0.06
Five qubits	4	16	4.05 ± 0.06

Eight probabilities for each term were obtained. In order to translate these probabilities to the expected values that appear in the inequality, one has to arrange the results in two groups according to the parity of the number of 1 (which represents the value -1.) The expected value of the term is obtained by summing all the probabilities of the results of even parity and subtracting the results of odd parity. The correctly weighted sum of the expected values of each term gives the final result $\langle M_3 \rangle_{\rm exp} = 2.85 \pm 0.02$.

In the case of four qubits, the use of the settings $a_i = \sigma_x$ and $a_i' = \sigma_y$ implies that the state that maximizes the quantum violation is $|\phi\rangle = 1/\sqrt{2}(e^{i\pi/4}|0000\rangle + |1111\rangle)$. With these settings and this state, five experiments were performed, one for each term with different numbers of primes (2), with 8192 runs for each experiment. A result of $\langle M_4 \rangle_{\rm exp} = 4.81 \pm 0.06$ was obtained.

In the case of five qubits, the use of the settings $a_i = \sigma_x$ and $a_i' = \sigma_y$ implies that the state that maximizes the quantum violation is $|\phi\rangle = 1/\sqrt{2}(|00000\rangle + |11111\rangle)$. With these settings and this state, three experiments were performed, one for each term with different numbers of primes (3), with 8192 runs for each experiment. A result of $\langle M_5 \rangle_{\rm exp} = 4.05 \pm 0.06$ was obtained. This is clearly a poor violation, which is still compatible with local realism. Improvement of the quantum computer is needed to obtain more accurate results, a summary of the results is presented in Table II.

The results obtained from the IBM quantum computer are subject to different kinds of errors. The stability of the quantum computer is still poor and the same experiments run at different times provided results that differed more than the expected behavior of statistical fluctuations. As an example, one month after the original runs, the three-qubit experiment was run again to compare results. This time, a result of $\langle M_3 \rangle_{\rm exp} = 2.57 \pm 0.02$ was obtained, clearly showing the previous point. An additional run was done computing separately the four terms of Eq. (1), without assuming any symmetry, and a similar result was obtained, $\langle M_3 \rangle_{\rm exp} = 2.57 \pm 0.02$, showing that it is safe to assume the symmetry of parity exchange.

We may get an estimation of the statistical error as a dispersion around the mean. We may, as well, treat the results as a multinomial distribution, using the expression $\delta p = \sqrt{p(1-p)/N}$, which for N=8192 gives $\delta p=O(10^{-2})$. The different Mermin inequalities for three, four, and five qubits require a different number of experiments to be done, which are considered as independent. We may then add in quadrature its errors, which is the figure we associate with the explicit results. In this sense, the five-qubit result obtained with the present quantum computer does not have sufficient statistical significance to discard local realism.

Furthermore, some of the issues related to the elimination of loopholes cannot be addressed. Experiments suffer from errors related to stability, loss of coherence, and lack of full fidelity of the quantum gates. This is clearly seen as the violation of Mermin inequalities deteriorate progressively as the numbers of qubits, and gates used in the experiment, increase. We may think of the experimental verification of Mermin inequalitites as a test of the overall fidelity of the whole Mermin circuit.

IV. CONCLUSIONS

Experimental verification of Mermin inequalities for three, four, and five qubits has been tested on a five-qubit IBM quantum computer. Results do show violation of local realism in all cases, with a clear degradation in quality as the number of qubits (and needed gates) increases. Nonetheless, this produces the first experimental violation of four- and five-qubit Mermin inequalities with superconducting qubits, though the statistical significance of the second one is still poor. It should be noted however that, in the case of the four-qubit inequality, the result shows generic nonlocality but does not provide evidence for genuine four-particle nonlocality, because this would only be implied by $M_4 > 8$. [14]. It can be argued that the measurements of Mermin polynomials for many qubits can be used as a figure of merit to assess the fidelity of a quantum computer.

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